

After an easy breach, hackers leave “TIPS WHEN RUNNING A SECURITY COMPANY”

DDoS protection firm Staminus apparently stored customers' credit card data in the clear.

by Sean Gallagher - Mar 11, 2016 11:35am CST

A Web security company's systems **are offline this morning** after an apparent intrusion into the company's network. The servers and routers of Staminus Communications—a Newport Beach, California-based hosting and distributed denial of service (DDoS) protection company—went offline at 8am Eastern Time on Thursday in what a representative **described in a Twitter post** as "a rare event [that] cascaded across multiple routers in a system wide event, making our backbone unavailable."

That "rare event" appears to have been intentional. A data dump of information on Staminus' systems includes customer names and e-mail addresses, database table structures, routing tables, and more. The data was posted to the Internet this morning, and a Staminus customer who wishes to remain anonymous confirmed his data was part of the dump. The authors of the dump claim to have gained control of Staminus' routers and reset them to factory settings.

- Use one root password for all the boxes
- Expose PDU's [power distribution units in server racks] to WAN with telnet auth
- Never patch, upgrade or audit the stack
- Disregard PDO [PHP Data Objects] as inconvenient
- Hedge entire business on security theatre
- Store full credit card info in plaintext
- Write all code with wreckless [sic] abandon

public key

cryptography

CS642

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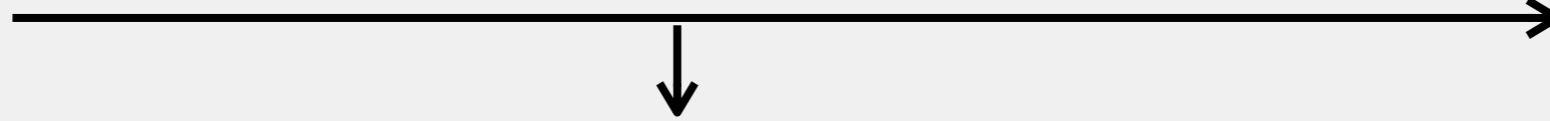
today

- * Hybrid encryption
- * Digital signatures, certificates
- * TLS overview
- * Passwords



Alice

email



Bob



Eve

- * Security goals?
/ Confidentiality, integrity, authenticity
- * Symmetric encryption: fast, hard to distribute keys
- * Public key encryption: slow, easy to distribute public keys

hybrid encryption

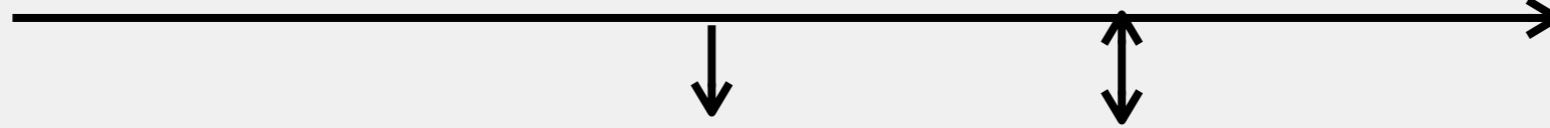


Alice

pk_A, sk_A

pk_B

$F(pk_B, x), E_x(M)$



Bob

pk_B, sk_B

pk_A



Eve



Mallory

$x \leftarrow \{0, 1\}^k$

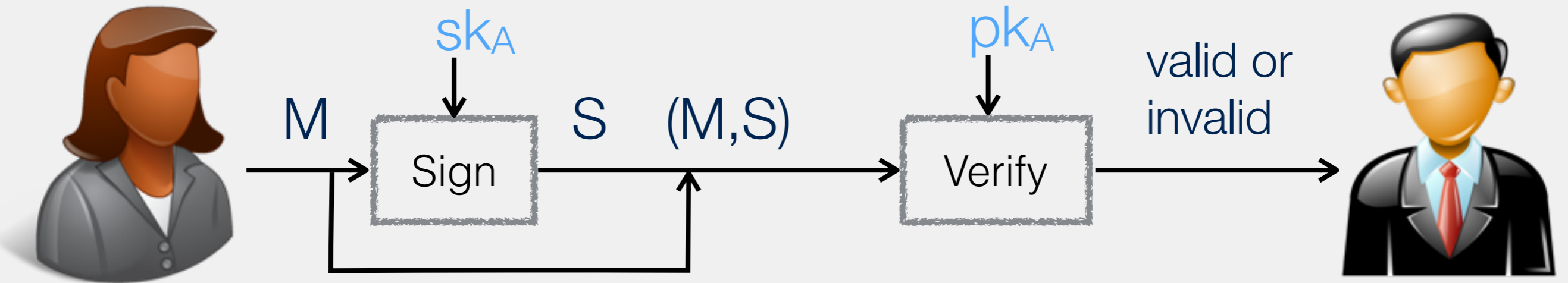
$F(pk_B, x), E_x(M)$

random key for this message

Authenticated encryption scheme

Encrypt under Bob's pubkey

hybrid encryption



Alice
 pk_A, sk_A

Trapdoor permutation
 $F_{pk}: X \rightarrow X$
 $F^{-1}_{sk}: X \rightarrow X$

Hash Fn
 $H: \{0,1\}^* \rightarrow X$

Bob
 pk_A

think-*pair*-share

Sign(sk_A, M):
 $d = H(M)$
 $S = F^{-1}(sk_A, d)$

Verify(pk_A, M, S):
 $d' = F(pk_A, S)$
 if $d' = H(M)$:
 ret VALID
 else:
 ret INVALID

digital signatures



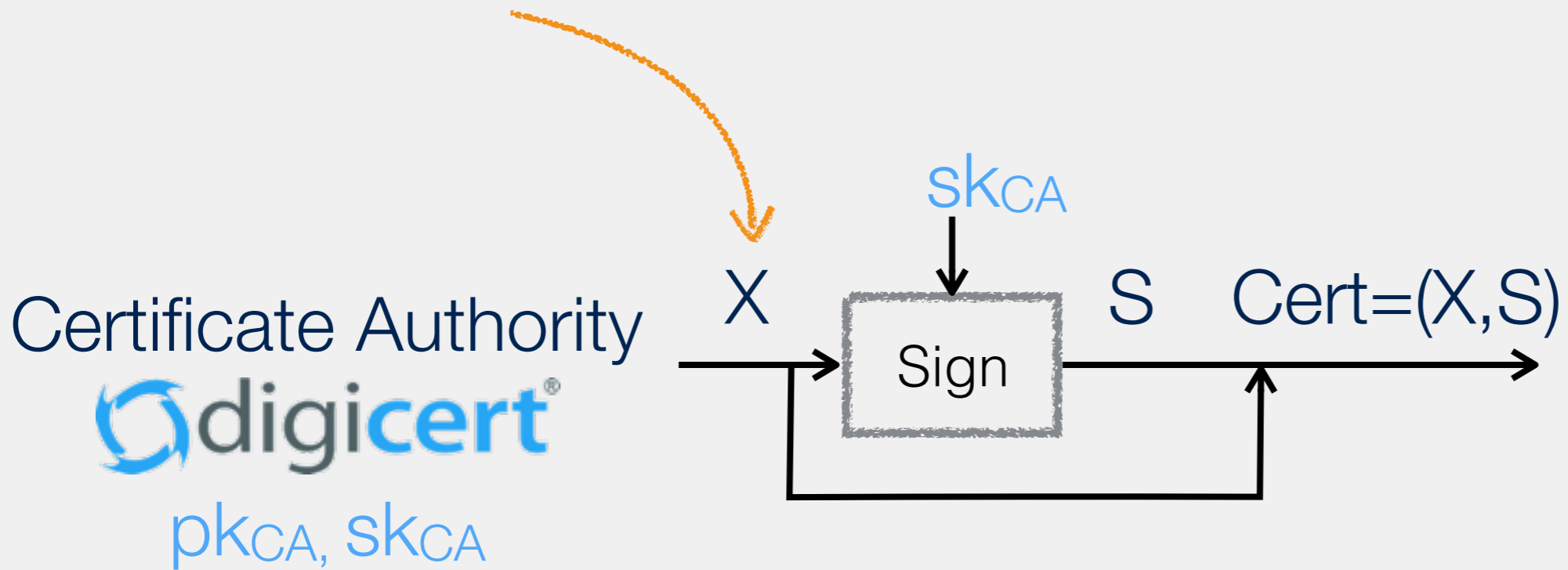
- * Problem: How does a client get the public key for a website?

certificates



Domain: *.facebook.com

Pubkey: 04 DB D1 77 ...



cert signing



Safari is using an encrypted connection to www.facebook.com.

Encryption with a digital certificate keeps information private as it's sent to or from the ht

- DigiCert High Assurance EV Root CA
- ↳ DigiCert SHA2 High Assurance Server CA
- ↳ *.facebook.com



*.facebook.com

Issued by: DigiCert SHA2 High Assurance Server CA

Expires: Friday, December 30, 2016 at 6:00:00 AM Central Standard Time

✔ This certificate is valid

▶ Trust

▼ Details

Subject Name

Country US

State/Province CA

Locality Menlo Park

Organization Facebook, Inc.

Common Name *.facebook.com

Public Key Info

Algorithm Elliptic Curve Public Key (1.2.840.10045.2.1)

Parameters Elliptic Curve secp256r1 (1.2.840.10045.3.1.7)

Public Key 65 bytes : 04 D8 D1 DD 35 BD E2 59 B6 FB 9B 1F 54 15 8C DB BF 4E 58 BD 47 BE B8 10 FC 22 E9 D2 9E 98 F8 49 2A 25 FB 94 46 E4 42 99 84 50 1C 5F 01 FD 14 25 31 5C 4E D9 64 FD C5 0C B3 46 D2 A1 BC 70 B4 87 8E

Key Size 256 bits

Key Usage Encrypt, Verify, Derive

Signature 256 bytes : AA 91 AE 52 01 8C 60 F6 02 B6 94 EB AF 6E EB DD 3C C8 E1 6F 17 AB B8 28 80 EC DC 54 82 56 24 C1 16 08 E1 C2 C8 3E 3C 0F 53 18 40 7F DF 41 36 93 95 5F B1 D9 35 43 5E 94 60 F9 D6 A7 83 6A 7D C7 B4 F6 0B 90 76 F8 B4 0A C1 31 0D 16 18 B5 CB 71 5C F9 93 02 21 AA BB 40 FD EE 0A 1B A9 F2 C3 0E 25 13 63 67 A2 42 EB 79 EA 5F 8F FB D8 BB 76 8C 5F 61 CA 2C BE 01 44 09 AF 36 1E A9 F7 40 1C A4 B3 65 78 42 68 04 F0 4B 0C 7F 1F D9 13 F6 0A 3B 35 79 73 69 C7 3C 70 E5 5D 06 98 EA 88 D5 DD 6B E6 66 62 57 CF AF D0 FB 67 9B E0 C8 20 3A B9 B6 4F 39 7A 5F C4 FD A0 46 8C BC C7 44 A7 B3 AB 52 49 DB 86 97 ED 2E BC 80 56 95 9F D2 63 84 57 E7 92 15 32 E4 75 C5 81 52 CB 3B 26 E1 5D 4B FD E0 39 5E 81 06 AF CC 7E 77 D1 9D 9A 06 6F EF F7 FC E2 86 5A 16 5A C2 04 DE 80 E3 78 1F 0F FC 7F DF

Let's Encrypt is a new Certificate Authority:
It's free, automated, and open.

[Get Started \(Public Beta\)](#)

<http://letsencrypt.org>

certificates



- * What does having a trusted TLS certificate prove?
 - / That someone paid at least \$0
 - / Proved to an intermediate CA that they controlled a given domain name for at least 5 minutes
 - / If TLS established, proves they know the corresponding private key to the pub key in cert

think-*pair*-share
- * What could possibly go wrong?
 - / Any CA secret key in chain could be **compromised**
 - / Server **secret key** could be compromised
 - / Typo-squatting domain (gmal.com)
 - / **Malicious** root CA key installed on client
 - / DNS **chicanery** during verification process

DigiNotar

- * Dutch CA DigiNotar compromised in 2011
- * Attackers generated fake certificates
- * Twitter.com was redirected to fake site
- * Attackers eavesdropped with man-in-the-middle attacks
/ Iranian govt eavesdropping on dissidents

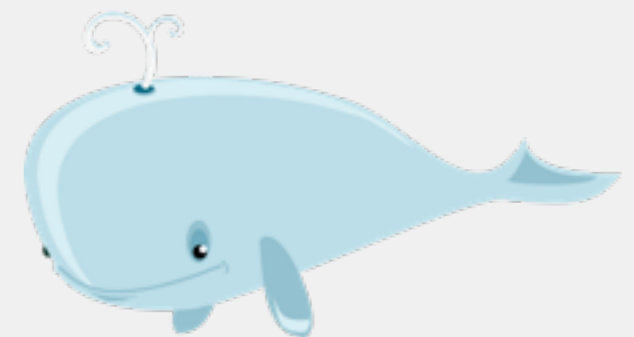
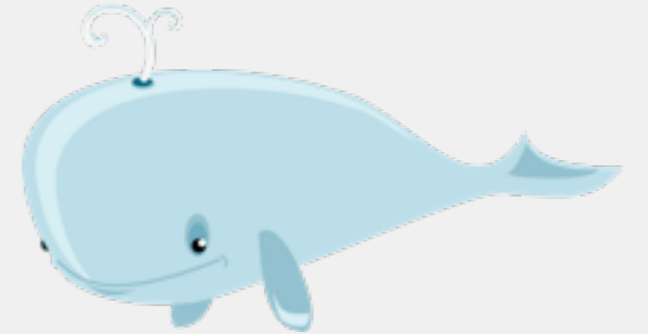


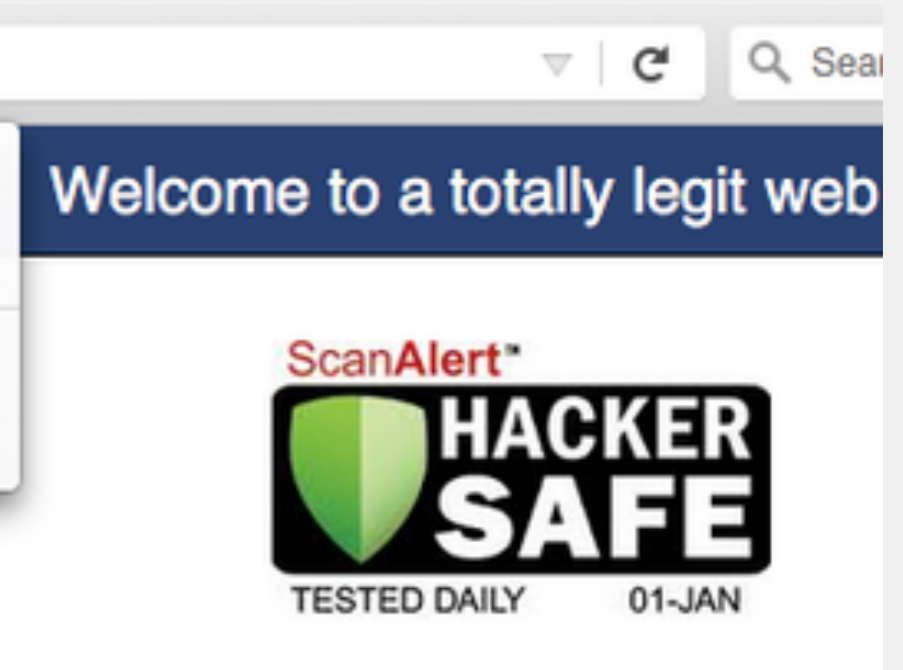
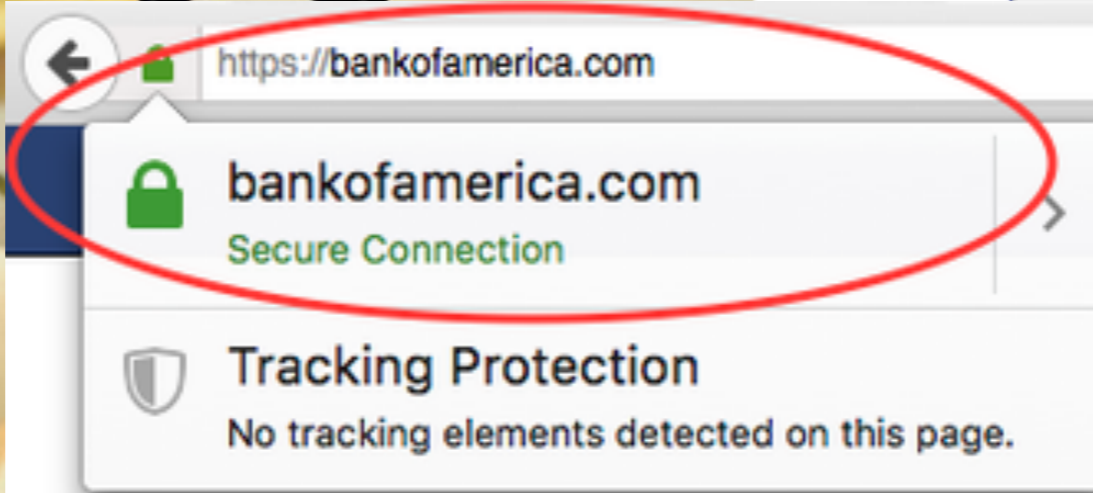
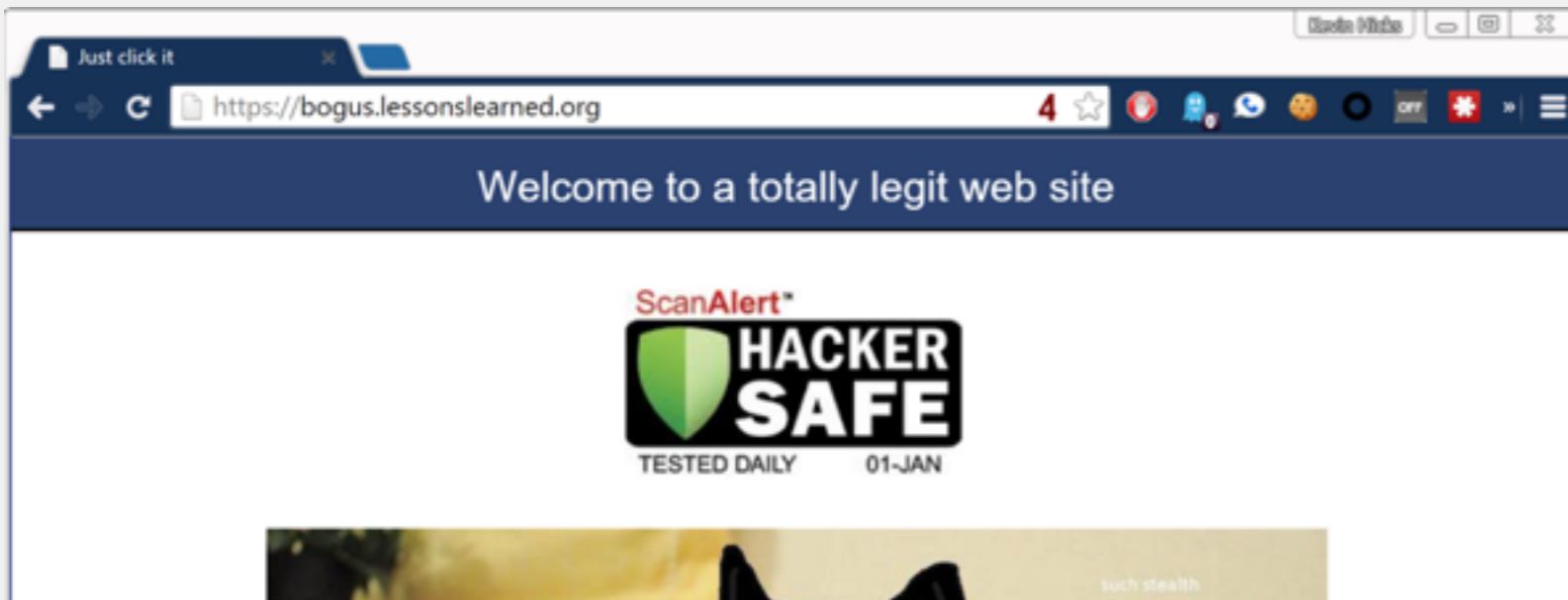
- * How did compromise occur?
- * DigiNotar had crappy security
 - / Out-of-date antivirus software
 - / Poor software patching
 - / Weak passwords
 - / No auditing of logs
 - / Poorly designed local network

DigiNotar

eDellRoot

- * Dell shipped several computer systems with a self-signed root CA certificate preinstalled
/The cert also contained the CA **secret key**
- * Intended purpose: something to do with automated support software
- * If certificate removed, automatically **reinstalls** on reboot





eDellRoot



blog.com

ClientHello, MaxVersion, Nonce_C, Supported ciphersuites

ServerHello, Version, Nonces_S, SessionID, Ciphersuite

Certificate = (pk_S, domain name, signature, cert chain)

E(pk_S, PMS)

MS <- HMAC(PMS, "master secret" || Nc || Ns)

K1,K2 <- HMAC(MS, "key expansion" || Ns || Nc)

..... Change to symmetric cipher

ChangeCipherSpec, Finished,
HMAC(MS, "Client finished" || H(transcript))

ChangeCipherSpec, Finished,
HMAC(MS, "Server finished" || H(transcript'))

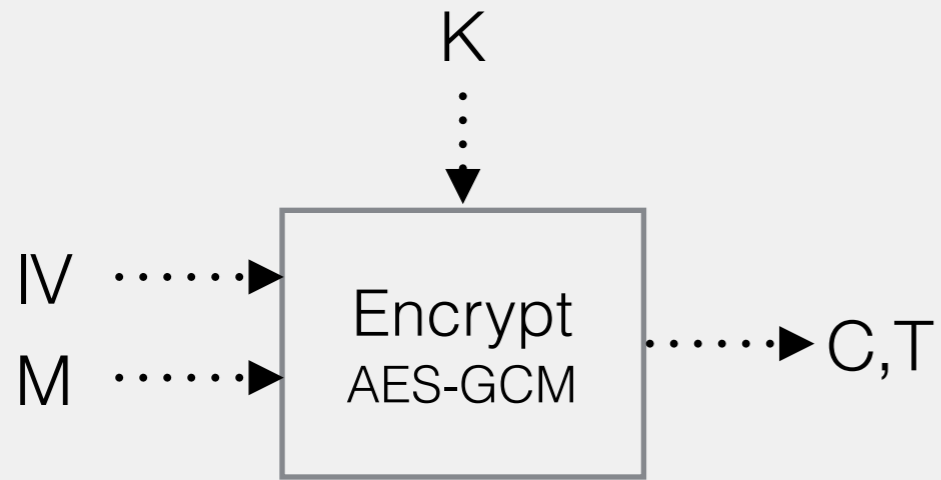
Exchange info using E_{k1}, E_{k2}

tls

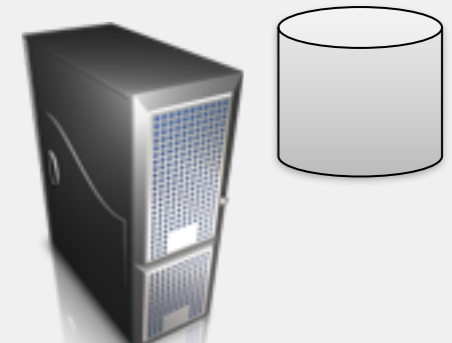
INTERMISSION

passwords

pw use cases

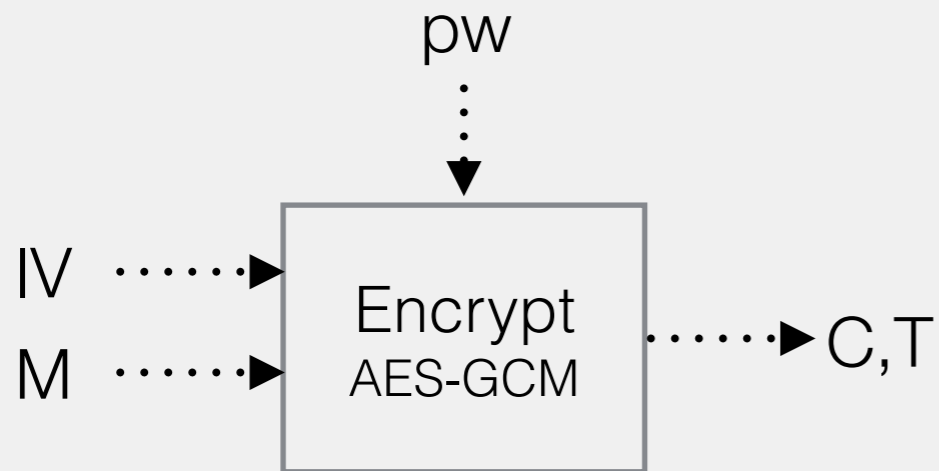


Create account:
...username,pw...



[server, desktop, or web service]

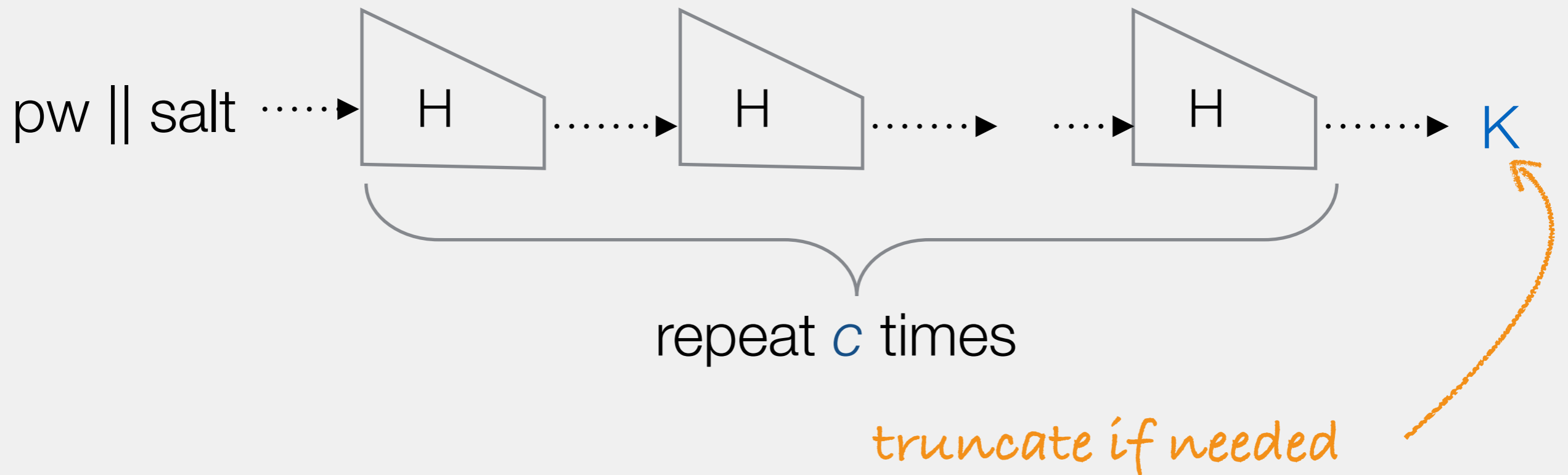
How does the server
store the pw?



Password-based symmetric encryption

[password-based key derivation function]

PBKDF(pw, salt):



pbkdf

pw-based encryption

Enc(pw, M, R):

salt || R' = R

K = PBKDF(pw, salt)

C = Enc'(K, M, R')

Return (salt, C)

Dec(pw, C):

salt || C' = C

K = PBKDF(pw, salt)

M = Dec'(K, C')

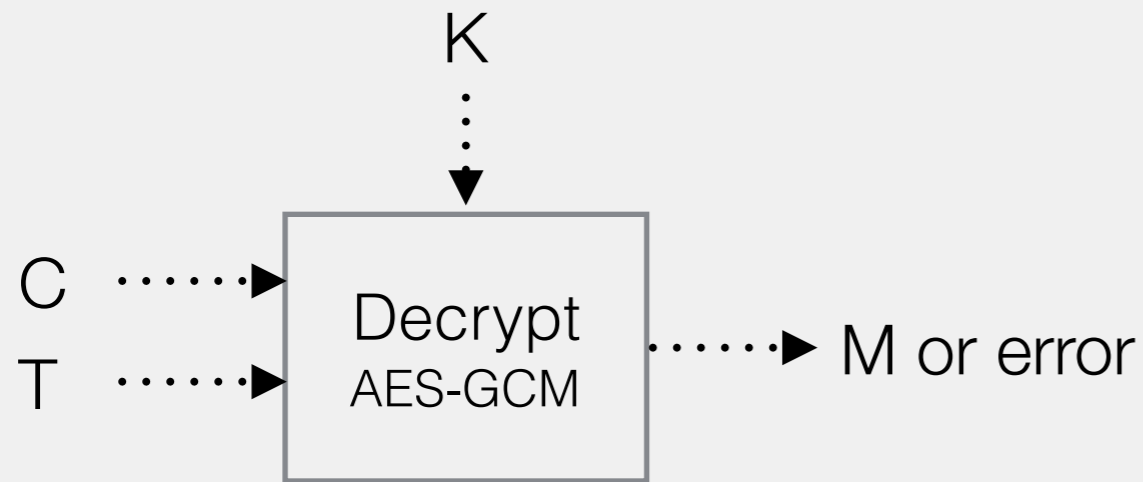
Return M

Enc'/Dec' is some authenticated encryption scheme,
like AES-GCM

PBKDF + symmetric encryption → pw-based encryption

Attacks?

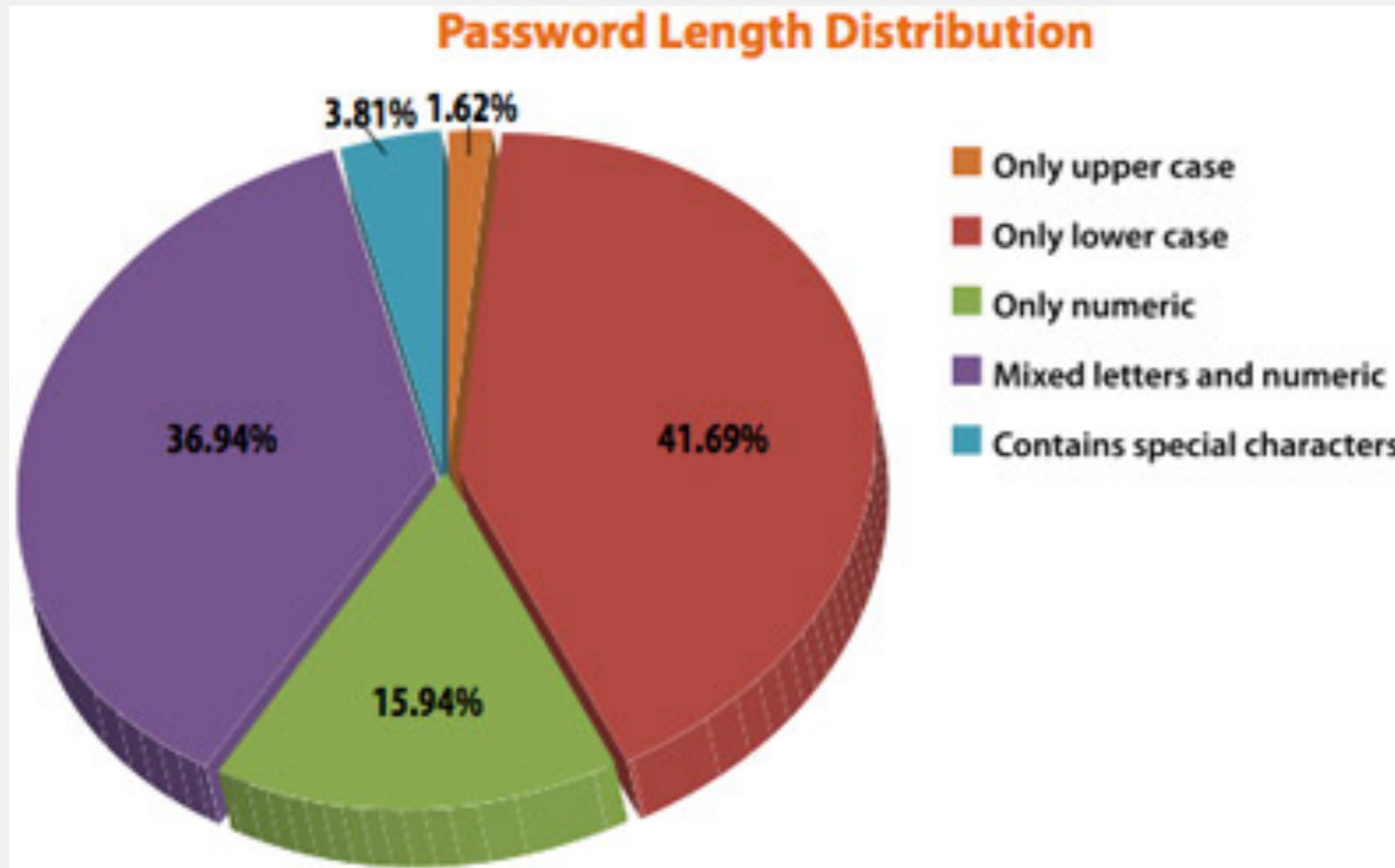
dictionary attack



DictionaryAttack(D,C,T):
for pw^* in D :
 $M^* = \text{Dec}(pw^*, C, T)$
 if $M^* \neq \text{error}$:
 return pw, M^*

- * Given an authenticated encryption output (C, T) , dictionary D of possible password
- * Enumerate D in order of likelihood
- * Test each candidate password

pw distribution



From an Imperva study of released RockMe.com password database (2010)

Facebook's Password Onion

```
$cur = 'password'  
$cur = md5($cur)  
$salt = randbytes(20)  
$cur = hmac_sha1($cur, $salt)  
$cur = remote_hmac_sha256($cur, $secret)  
$cur = scrypt($cur, $salt)  
$cur = hmac_sha256($cur, $salt)
```



- * Hybrid encryption
- * Digital signatures
- * Certificates, problems
- * Password-based key derivation
/ Dictionary attacks
- * Exit slips
/ 1 thing you learned
/ 1 thing you didn't understand

recap