

Today 537: Part 3

⇒ CPU virtualization

→ mechanisms

(Limited Direct
Execution)

LDE

→ policies: scheduler

which process
should run?

last time:

→ Shortest Job First (SJF)
(assumes knowledge
of run time)

→ Round Robin:

A	B
---	--------------

 not RR

ABAB

↳

time slice (quantum)

=> responsive (interactive)

Today: develop real sched
Policy => (classic Unix scheduler)
=> Multi-level Feedback Queue
(MLFQ)

Later today: (virtual Memory)

Later later today: Project
2a
(Processes, Shell)

Life ↔ Work

OS sched Policy: MLFQ

Problem: Don't know very
much about processes!

like to learn: short jobs, longer running?

how? \Rightarrow measure : using past
to predict future

many queues : job is on
one queue @
any given
time
(might
change
over
time)

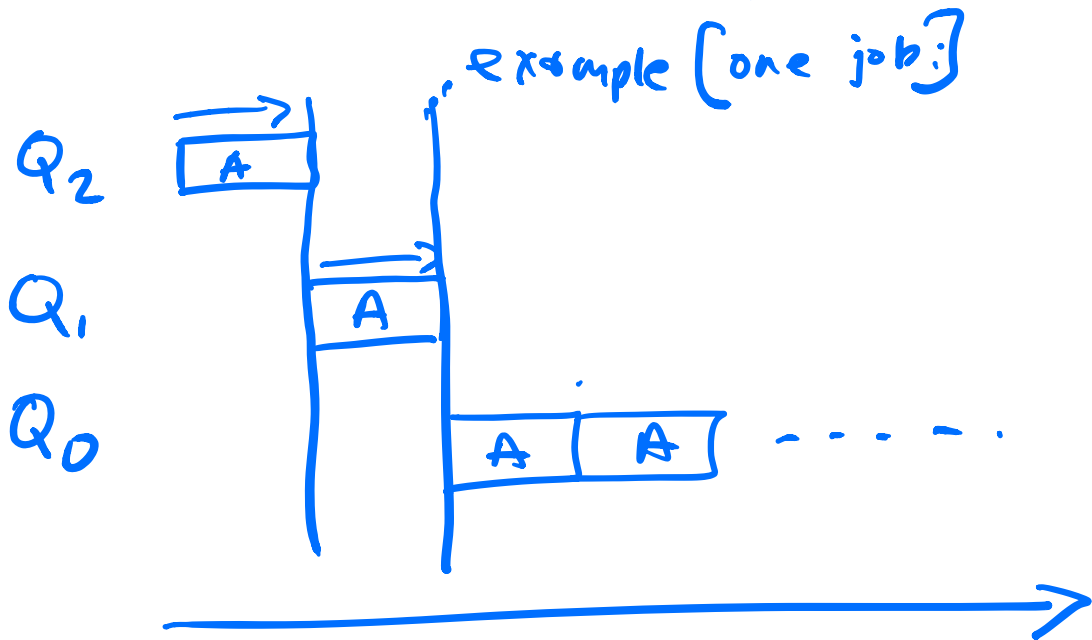
priority:
highest Q_2
middle Q_1 B, C
low Q_0 D

each queue has priority

Rules:

- 1) if $\text{Priority}(A) > \text{Priority}(B)$
 \Rightarrow A runs (B doesn't)
- 2) if $\text{Pri}(A) == \text{Pri}(B)$
 \rightarrow Round Robin between
them
- 3) Start: Highest Priority

4) if process uses time slice
 @ given priority,
 => at end of time slice,
 move down one level



Time

Example 2:

long-running job : A

short running
 $C_1,$
 $C_2, \dots,$
 C_N



C_1, \dots, C_N : τ large # of short jobs!

=> never ending

=> what happens to A?
(long running)?

=> (Starvation)

How to ensure long-running jobs make progress?

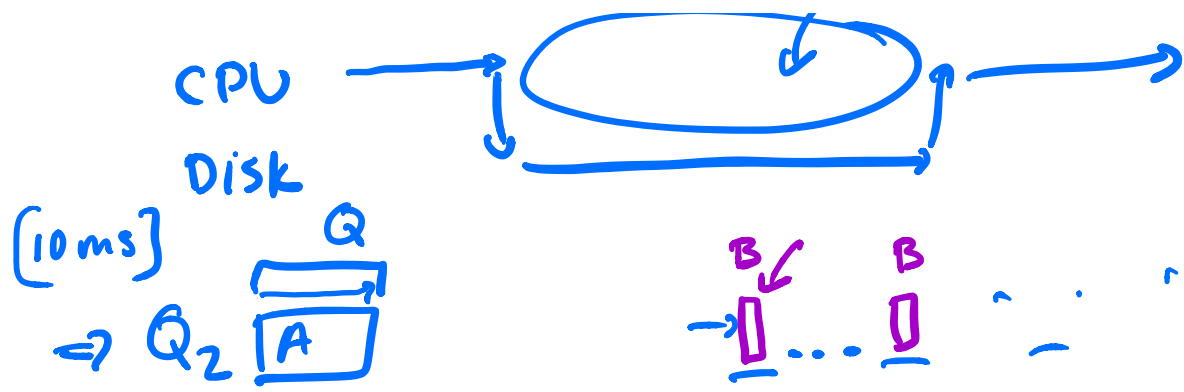
General idea: long-running need to move up

Rule: Every T seconds, move all jobs to highest priority

=> nature of job might change
between interactive and batch

new worry: (I/O)

/ run something else here



A : long running
 B : I/O job

naive rule for I/O:
 if job runs for less than timeslice,
stay at same level

⇒ "Gaming" the scheduler

[B] | [B] ...

better accounting : if job uses up quantum, moves down

Parameters: ML FQ

Q_{N-1} 

⋮

ML


Q_0

ML

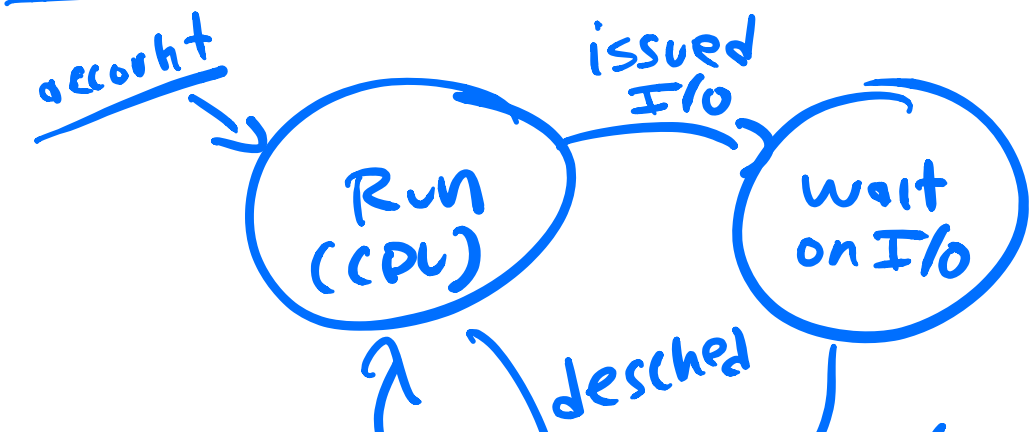
N QUEUES

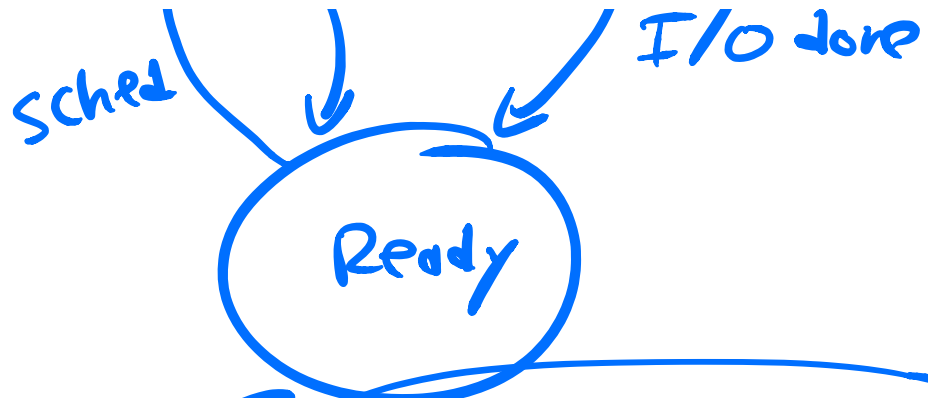
Q_i quantum length (per queue)

T : reset period

⇒ How to configure:
⇒ mystery
(use defaults) 

⇒ new: ML





Break:

Best Sports-themed
Movie

- | | |
|-----------------------------------|-------------------------------|
| => Waterboy | => Caddyshack |
| => Sandlot | => Dodgeball |
| => <u>Miracle</u> ← | => Shaolin Soccer |
| => Space Jam | => Blades of Glory |
| => Remember the Titans | => Major League |
| | => Bills of Fury |

Rocky

original
Karate Kid

II
IIT
IV
V
—

Bad news ^{good?} : class tomorrow

→ X X 3 4 5 6 7 knowledge
lack of 8 9 10 11 12 13 14 knowledge

Virtual Memory : multi-cores } later
Linux CFS


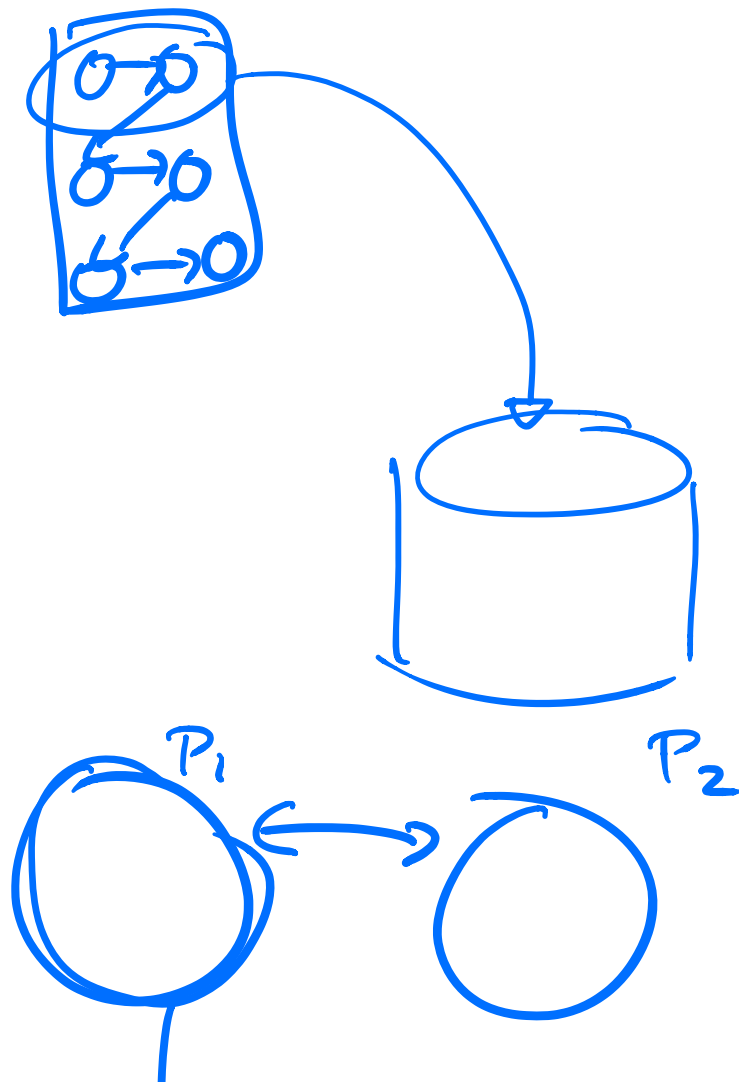
Process : running program

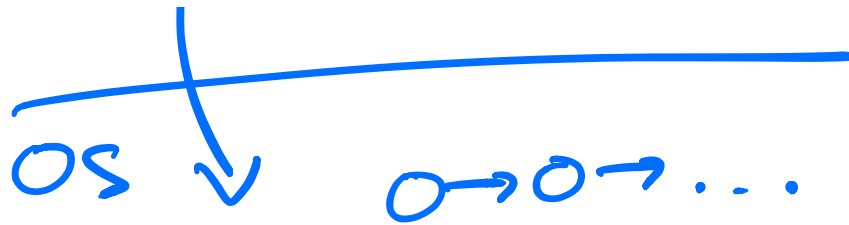


memory
=> virtual
address

goals illusion of large memory = ease
 => private (protected) memory
 => efficient memory

Space
 [32-bit AS,
 64-bit AS]



efficiency : limited
Direct Execution

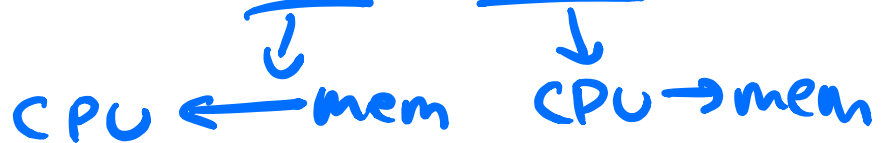
(loads, stores,
inst fetch)

Mechanisms : [s/w ⇒ software]
→ h/w (hardware) → OS support

Memory Accesses:

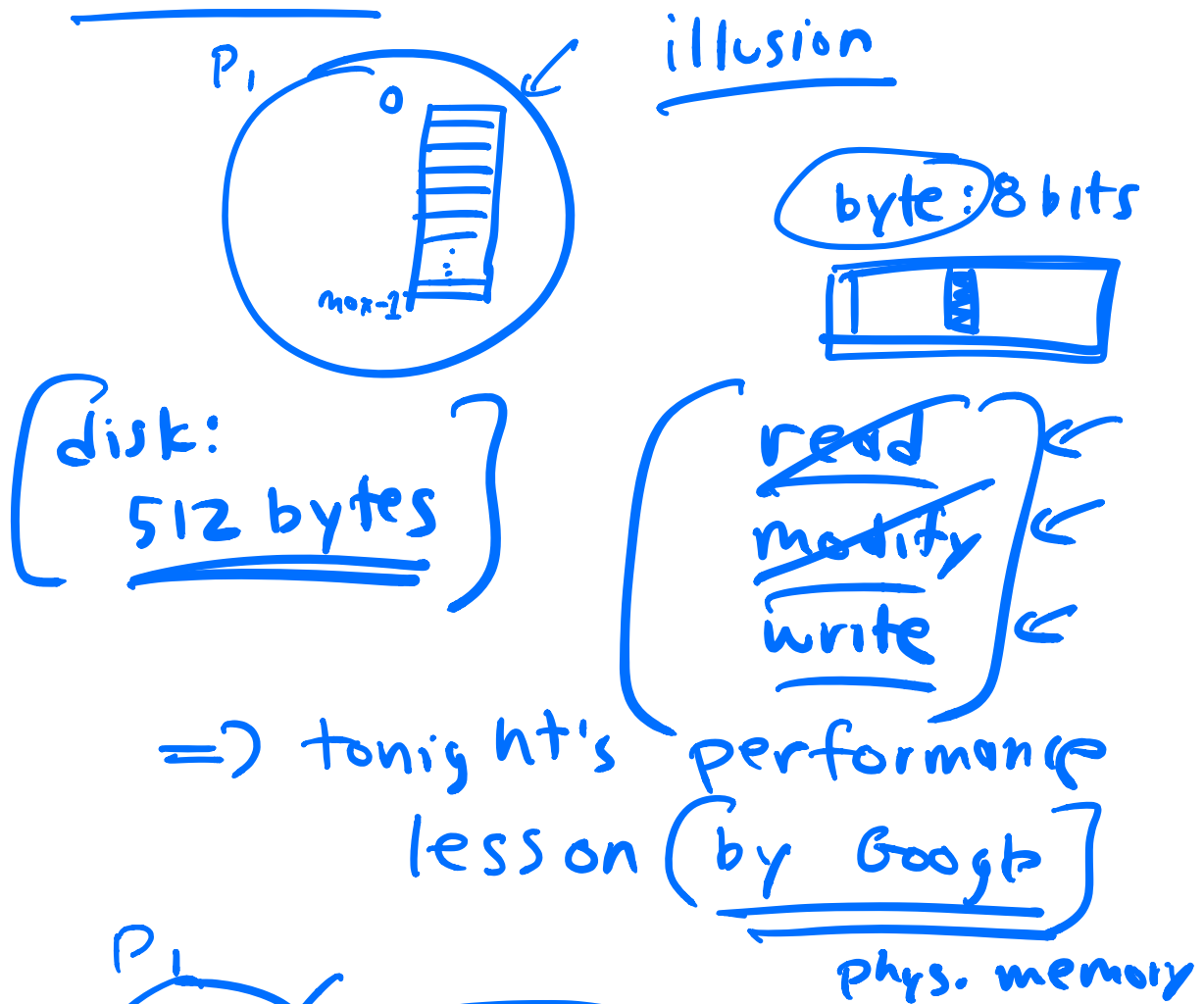
→ instruction fetch

→ explicit loads, stores



need: interpose

Address Translation :

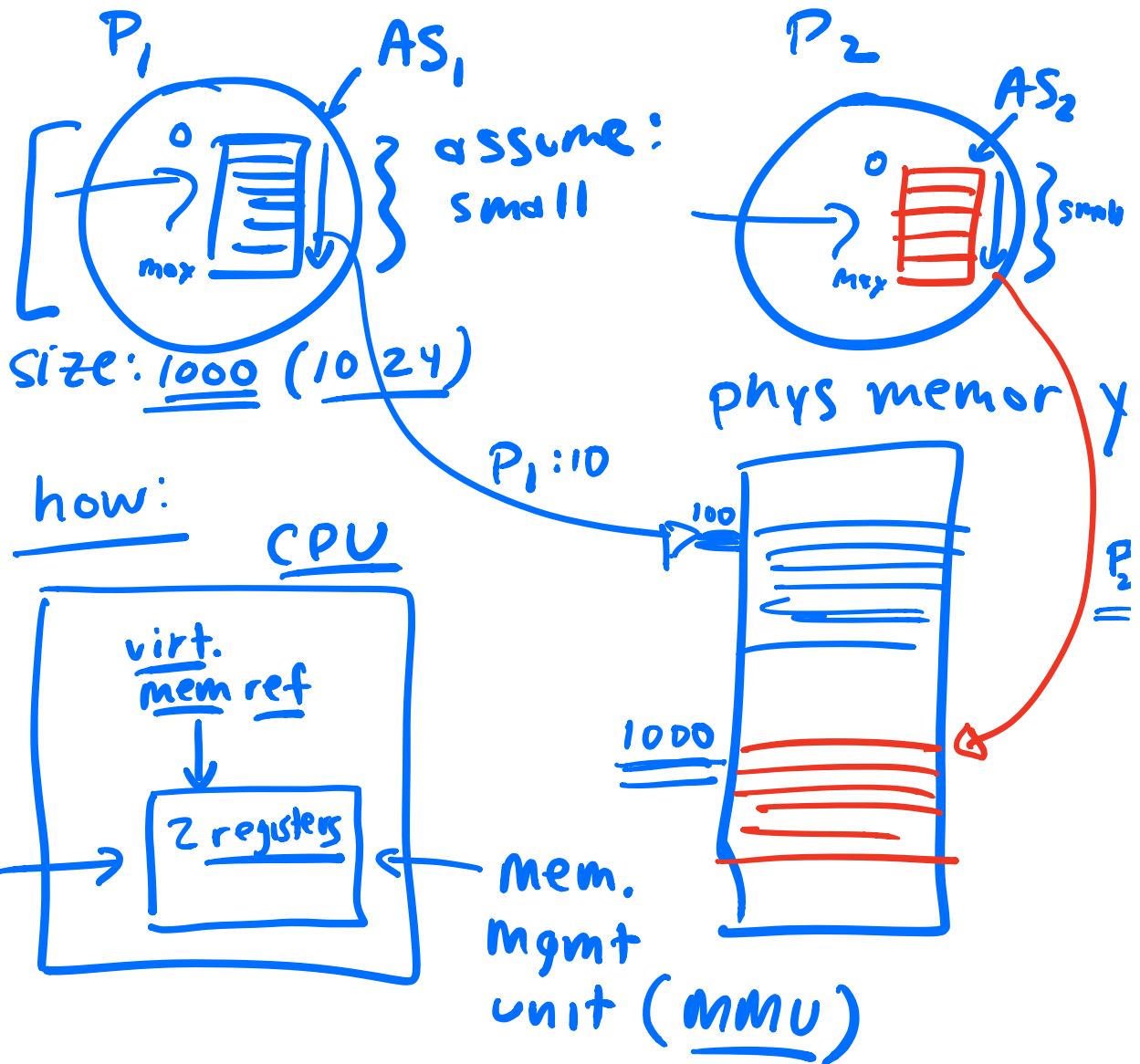


Address Translation

=> on every mem. reference,

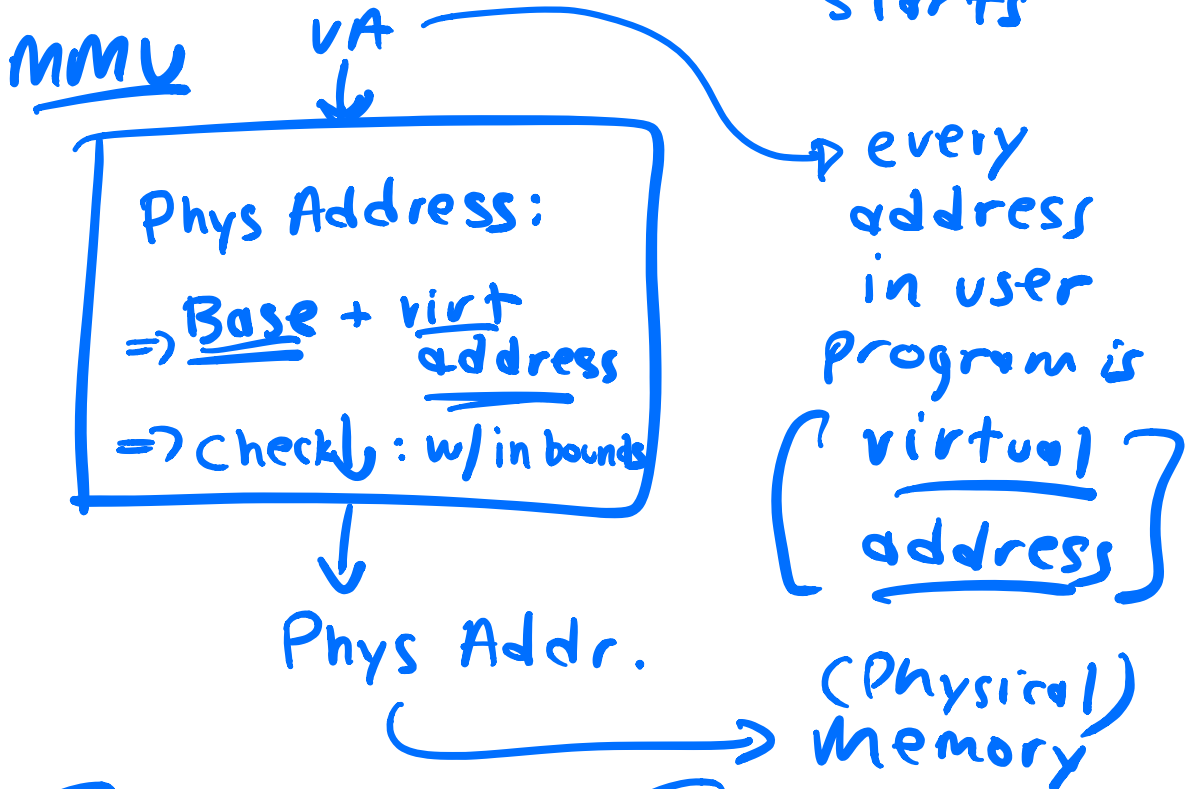
translate virtual memory addr:
 \Rightarrow phys memory addr

Mechanism #1:
Dynamic Relocation ("or Base/Bounds")



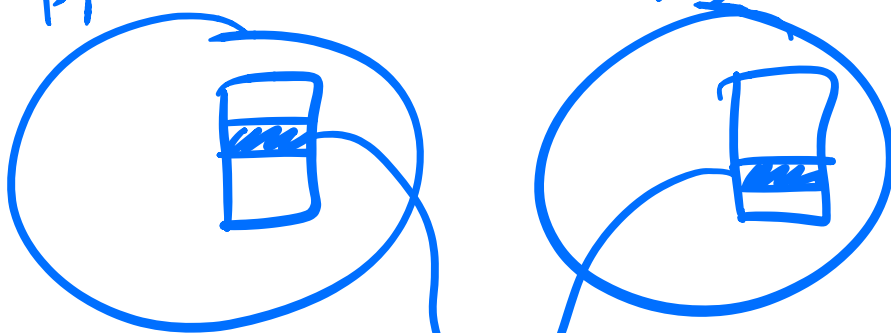
(2 registers / CPU)

- base: : address in phys mem where AS of currently running process starts
- bounds :



Modern OS: [LATER]

sharing of memory





=> Lesson #2:
(sometimes)
When?

Base: Translation

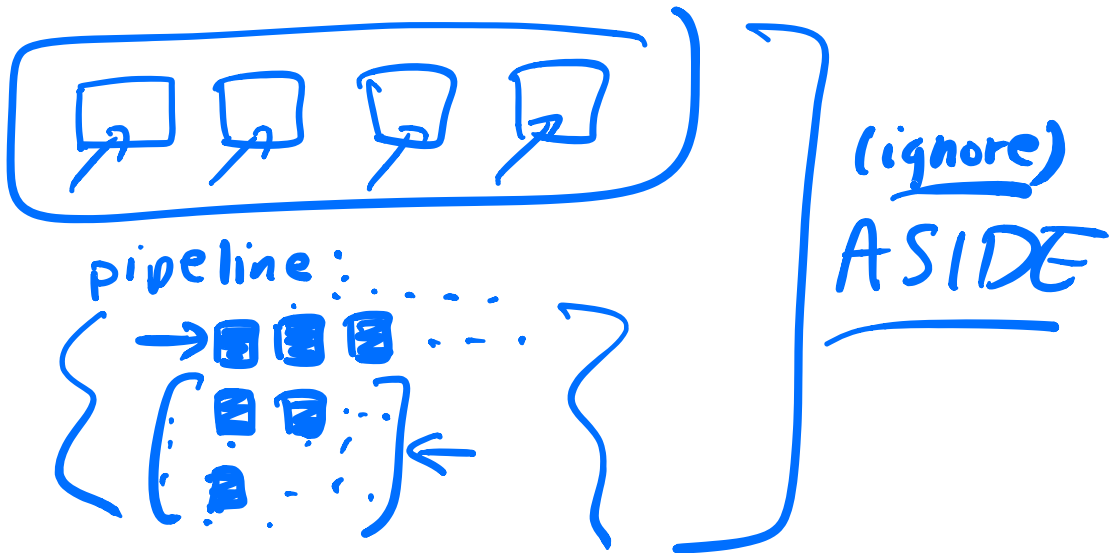
Bounds: Protection
(check w/in Addr. Space)

CPU (abstractly):

```
while (1) {  
  Fetch (PC)  
  Decode  
  Execute ← limit  
}
```

MMU: bounds

if OK: go ahead
if not OK: raise exception

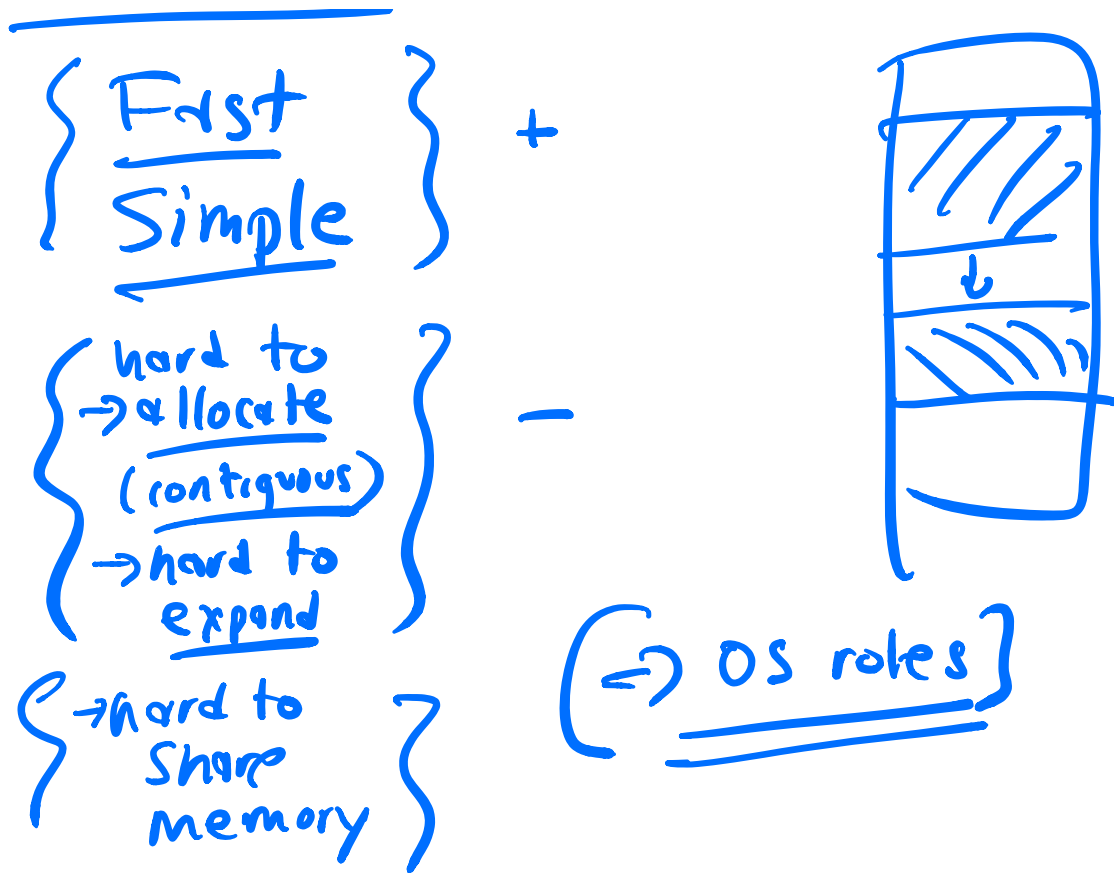


not OK: h/w raises exception
(illegal memory access)

=> H/w: OS → Process: bad mem access
@ boot: OS set up exception handlers

=> OS: [kill process]

Base/Bounds :



Question #2 : (Best non-sports
themed movie)