

Optimistic Crash Consistency

Vijay Chidambaram

Thanumalayan Sankaranarayanan Pillai

Andrea Arpaci-Dusseau

Remzi Arpaci-Dusseau



WISCONSIN
UNIVERSITY OF WISCONSIN-MADISON



Crash Consistency Problem

Single file-system operation updates **multiple** on-disk data structures

System may **crash** in **middle** of updates

File-system is partially (incorrectly) updated

Performance **OR** Consistency

Performance **OR** Consistency

Crash-consistency solutions **degrade** performance

Users **forced** to **choose** between high performance and strong consistency

- Performance differs by **10x** for some workloads

Many users choose performance

- ext3 default configuration did not guarantee crash consistency for many years
- Mac OSX `fsync()` does not ensure data is safe

“The Fast drives out the Slow even if the Fast is wrong”

Ordering and Durability

Crash consistency is built upon ordered writes

File systems **conflate** ordering and durability

- Ideal: $\{A, B\} \rightarrow \{C\}$ (made durable later)
- Current scenario
 - $\{A, B\}$ durable
 - $\{C\}$ durable

Inefficient when **only** ordering is required

Can a file system provide
both
high performance
and strong consistency?

Is there a middle ground between:
high performance but **no** consistency
strong consistency but **low** performance?

Our solution

Optimistic File System (OptFS)

Journaling file system that provides performance **and** consistency by **decoupling** ordering and durability

Such decoupling allows OptFS to **trade freshness** for performance while maintaining crash consistency

Results

Techniques: checksums, delayed writes, etc.

OptFS provides strong consistency

- Equivalent to ext4 data journaling

OptFS improves performance significantly

- **10x** better than ext4 on some workloads

New primitive **osync()** provides ordering among writes at high performance

Outline

Introduction

Ordering and Durability in Journaling

Optimistic File System

Results

Conclusion

Outline

Introduction

Ordering and Durability in Journaling

- Journaling Overview
- Realizing Ordering on Disks
- Journaling without Ordering

Optimistic File System

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Journaling Overview

Before updating file system, **write note describing update**

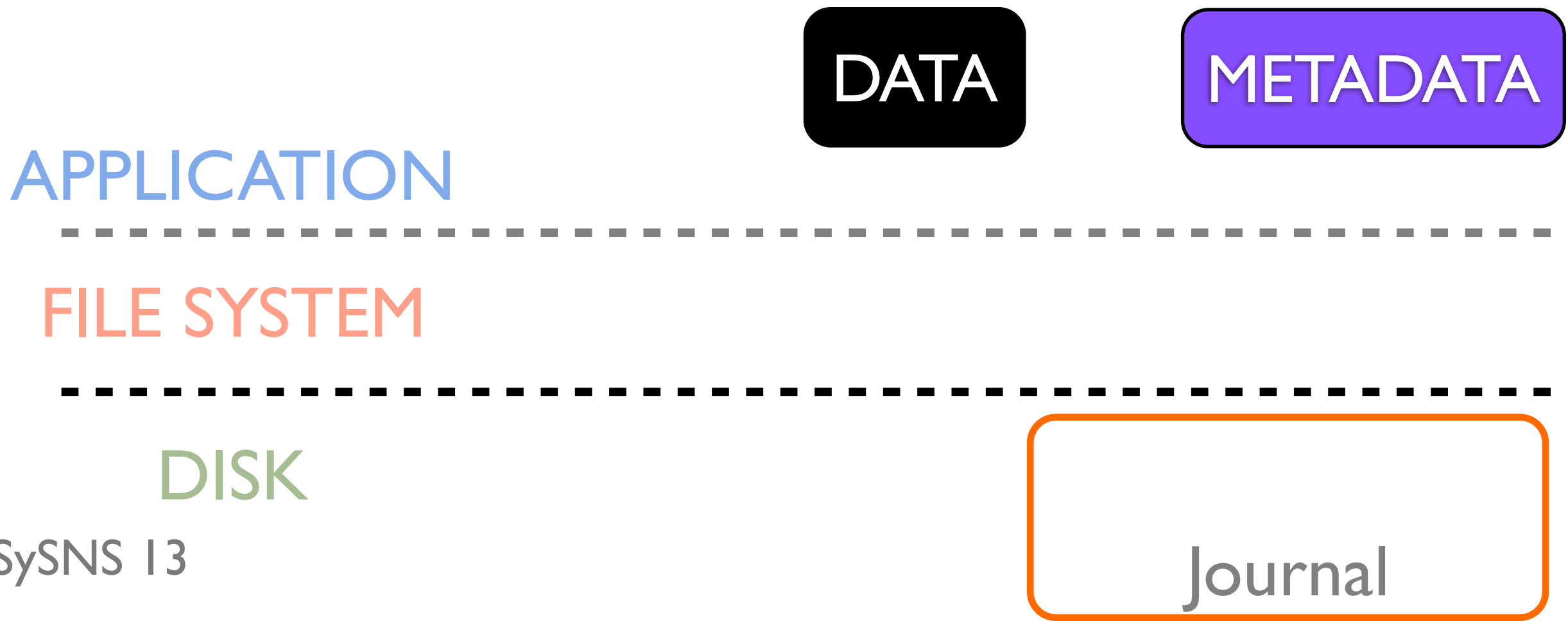
Make sure note is safely on disk

Once note is safe, **update** file system

- If interrupted, read note and **redo** updates

Journaling Overview

Workload: Creating and writing to a file
Journaling protocol (ordered journaling)

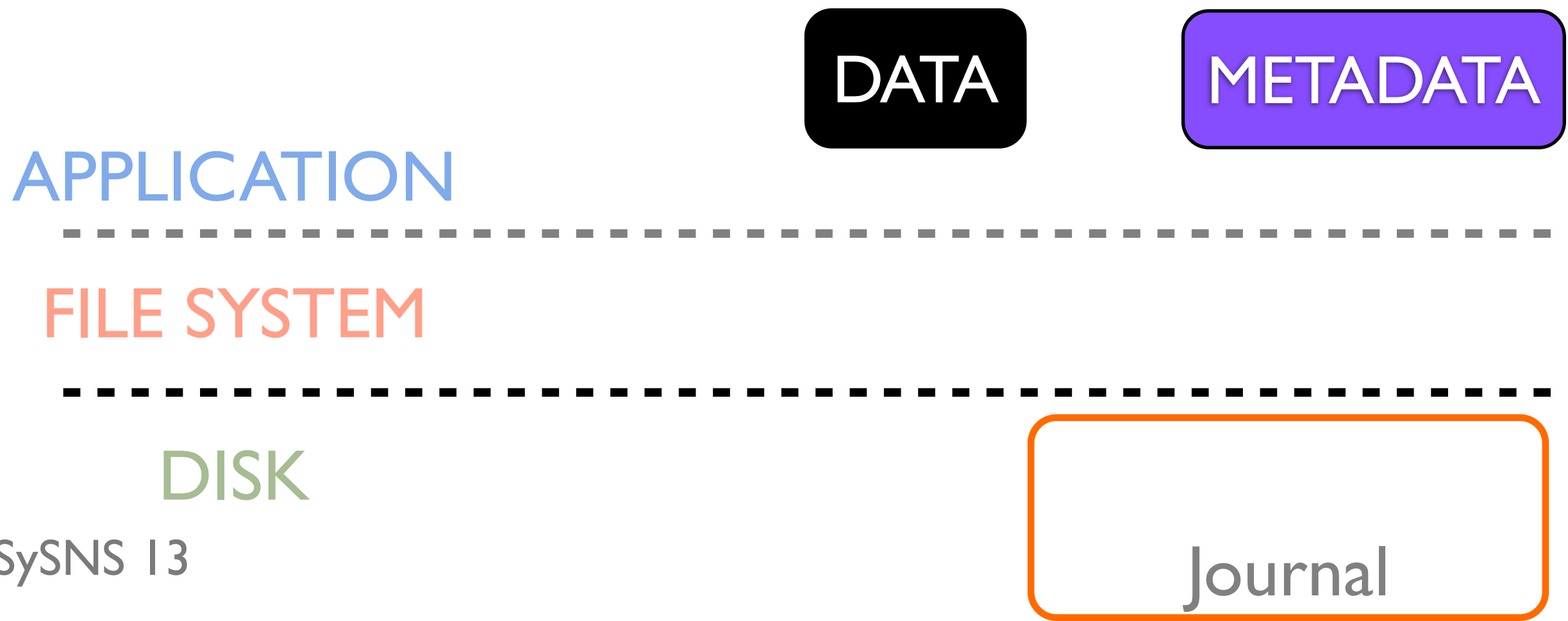


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Workload: Creating and writing to a file

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- Data write (D)

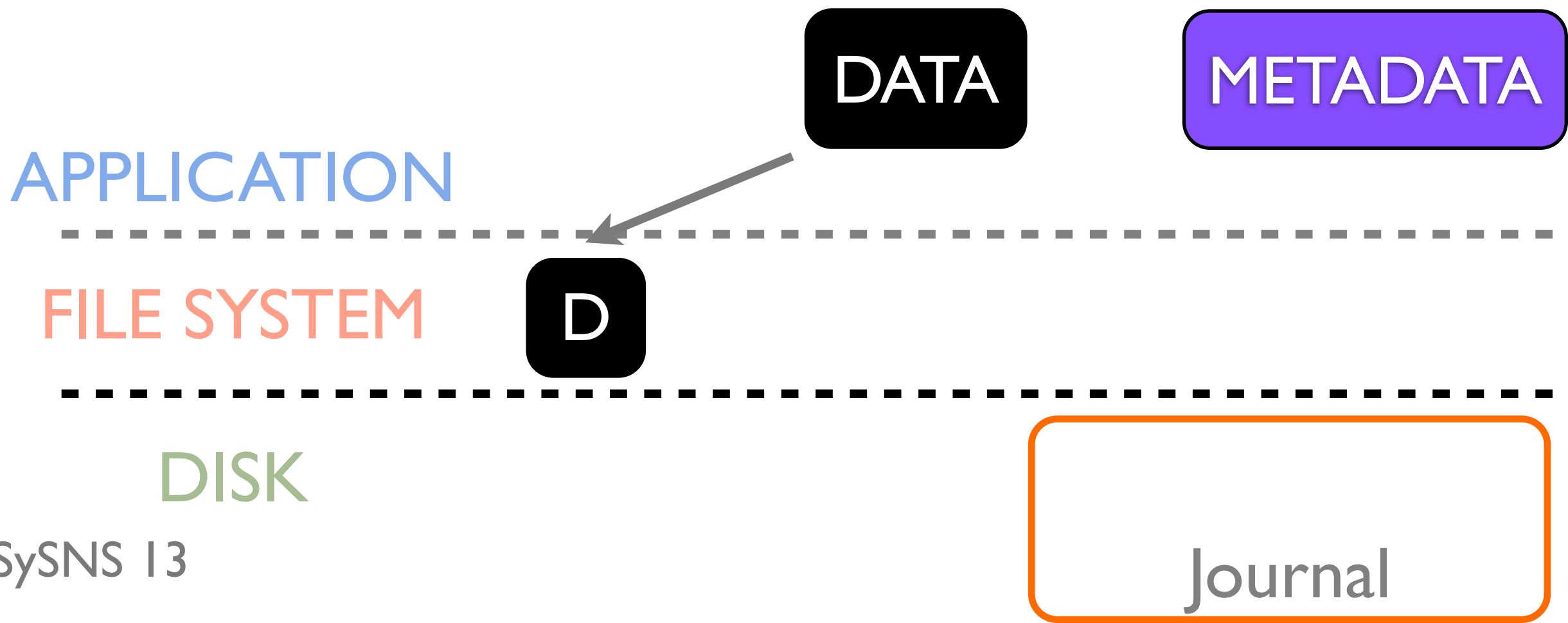


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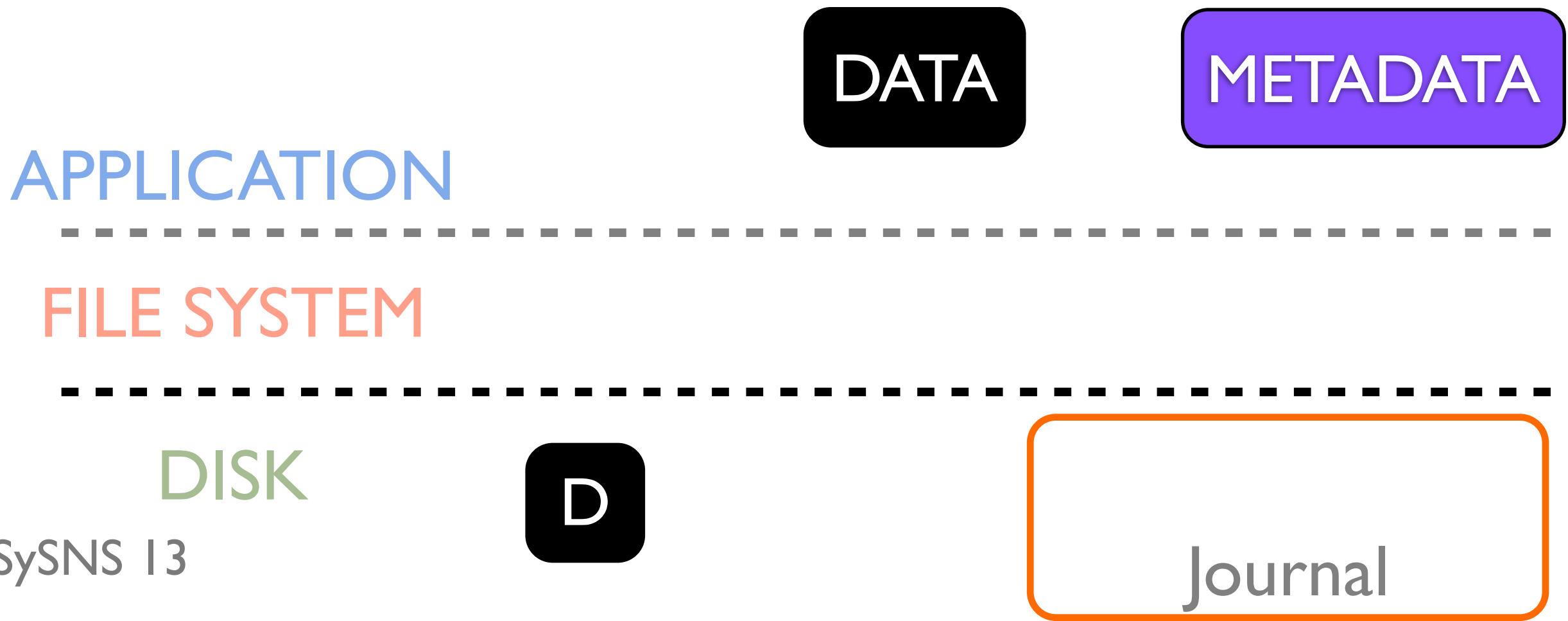


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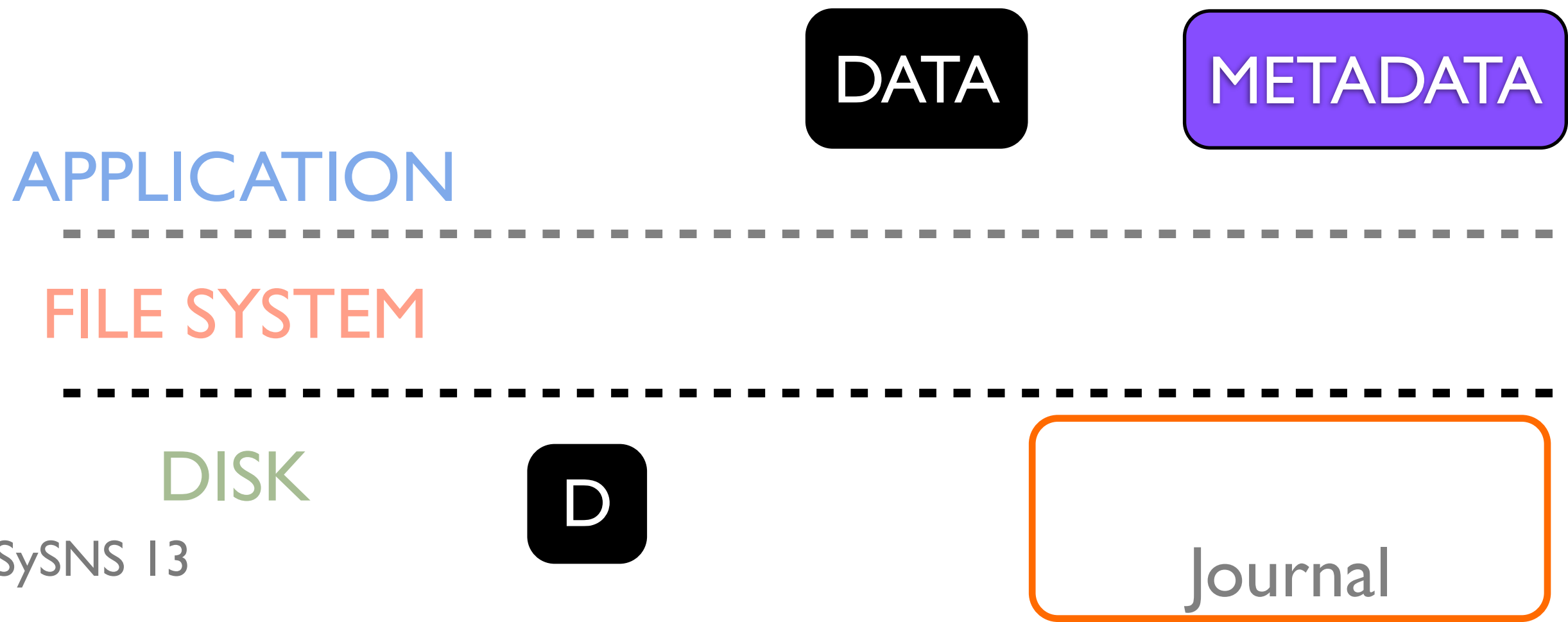


Journaling Overview

Workload: Creating and writing to a file

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- Data write (D)
- Logging Metadata (JM)

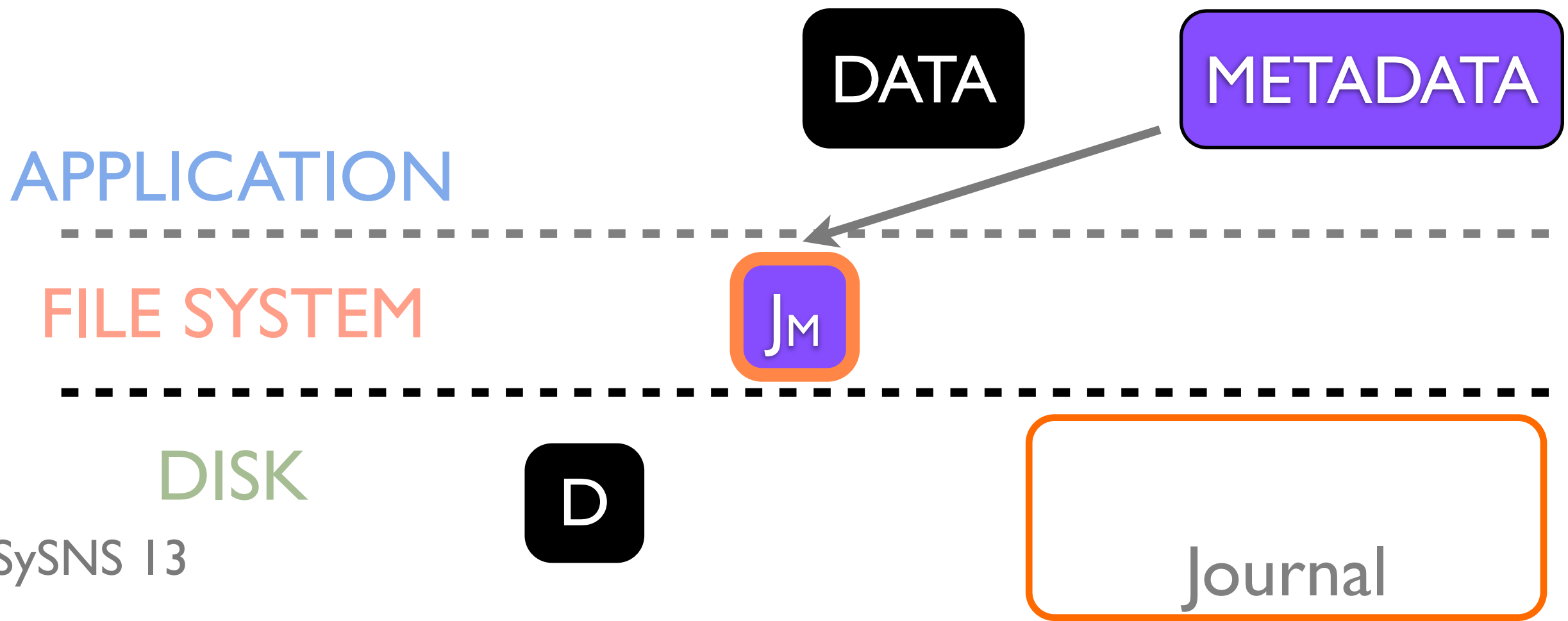


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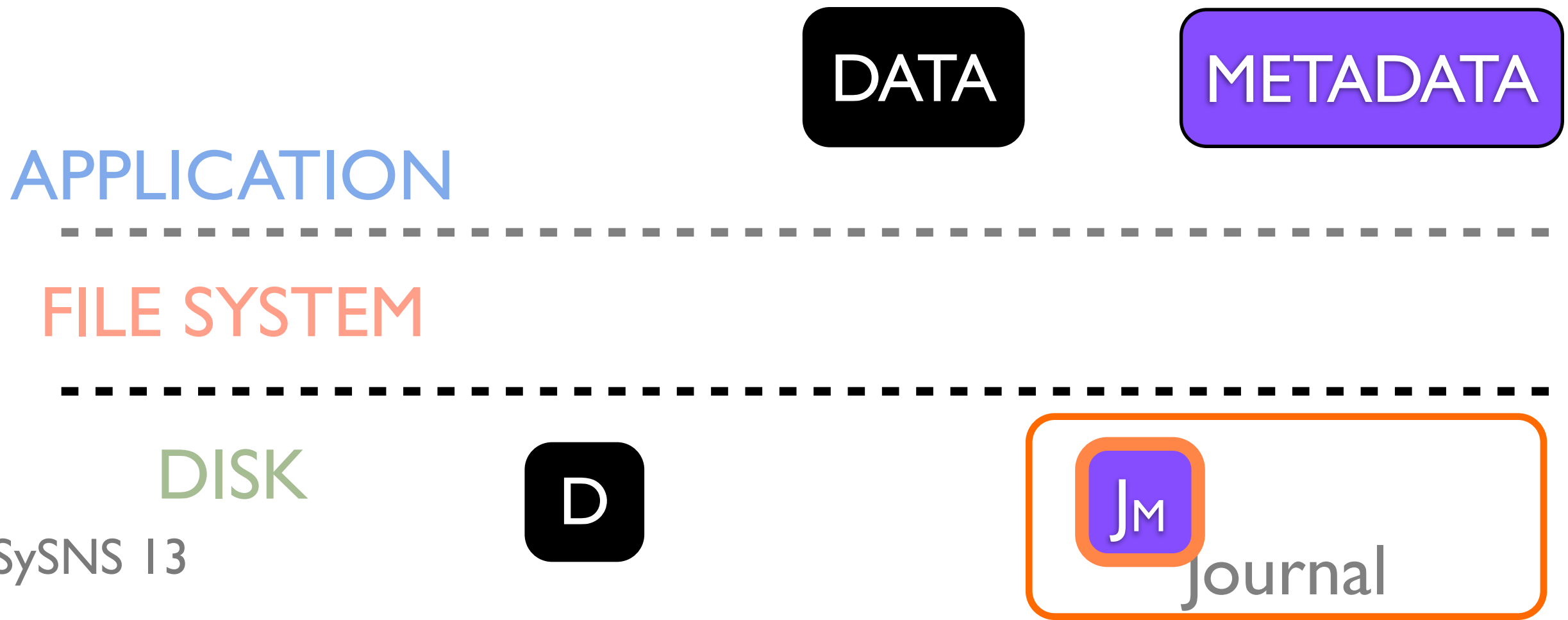


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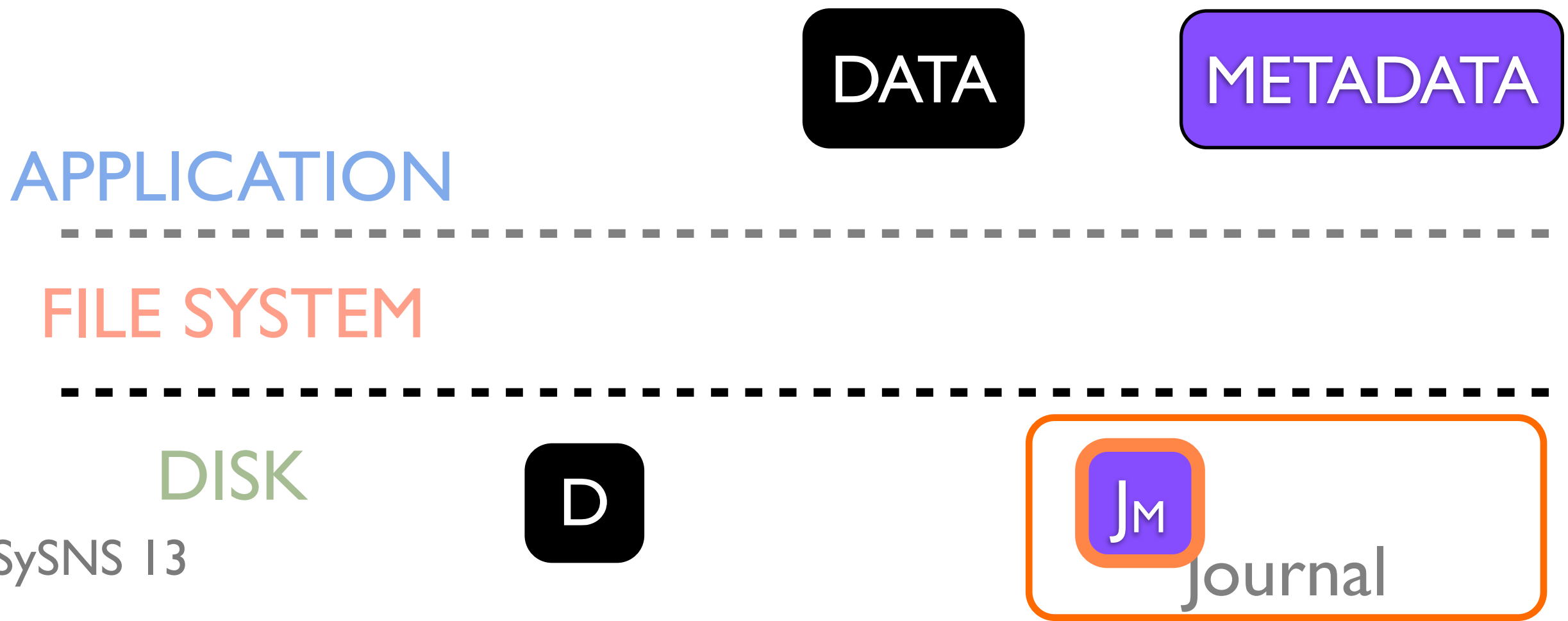


Journaling Overview

Workload: Creating and writing to a file

Journaling protocol (ordered journaling)

- Data write (D)
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- Logging Commit (Jc)

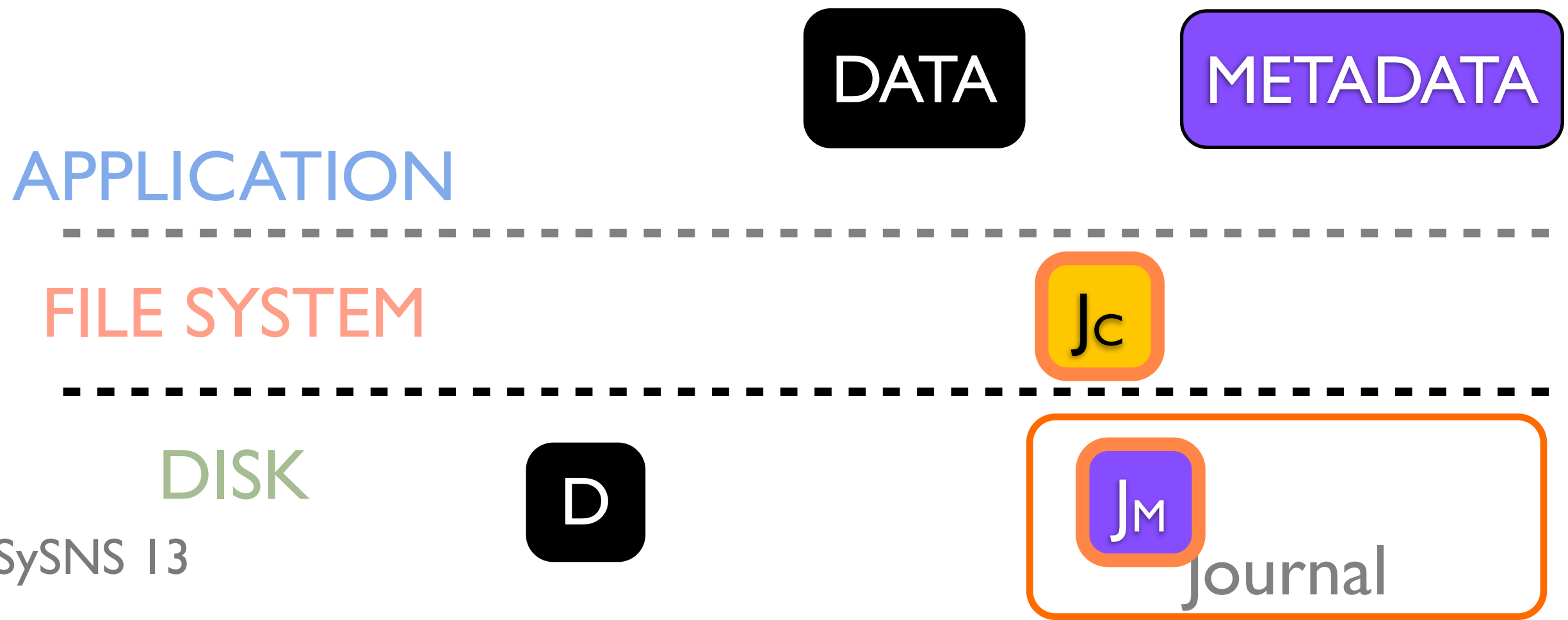


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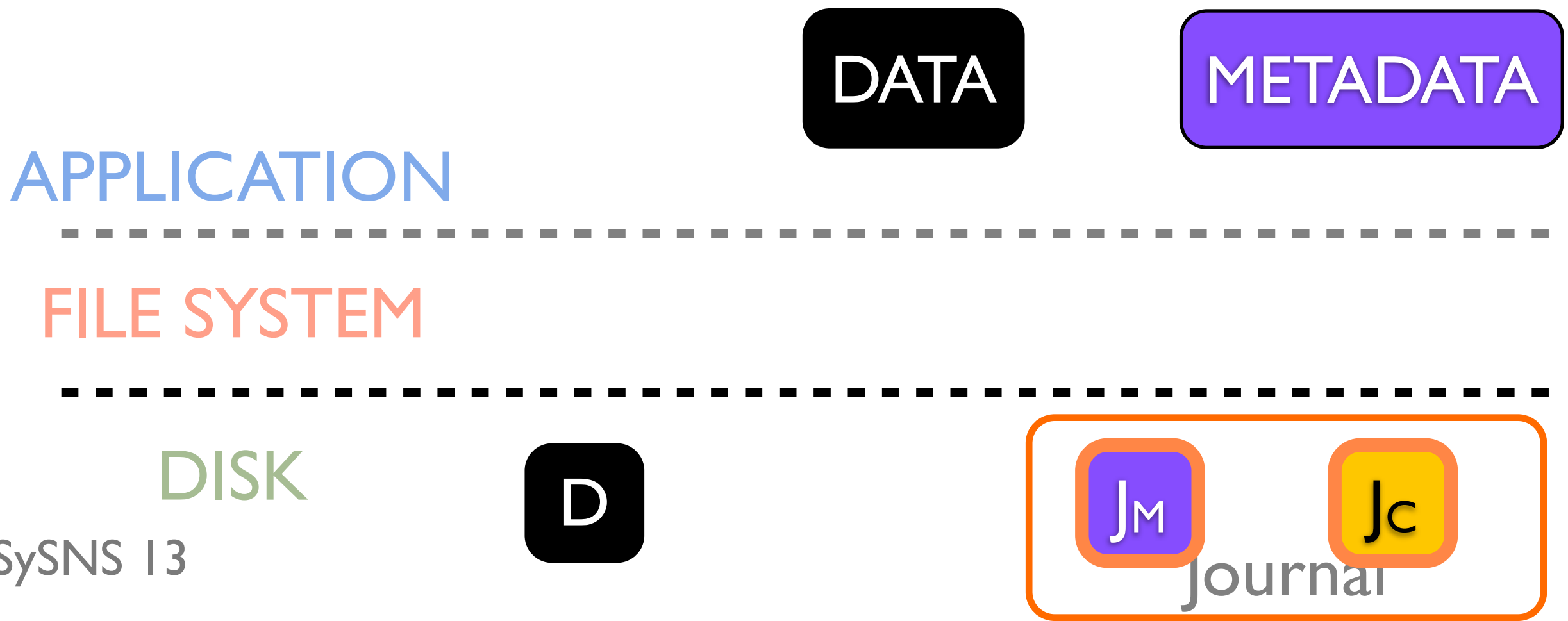


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Journaling Overview

Workload: Creating and writing to a file

Journaling protocol (ordered journaling)

- Data write (D)
- Logging Metadata (JM)
- Logging Commit (Jc)
- Checkpointing (M)

DATA

METADATA

APPLICATION

FILE SYSTEM

DISK

D

JM

Jc

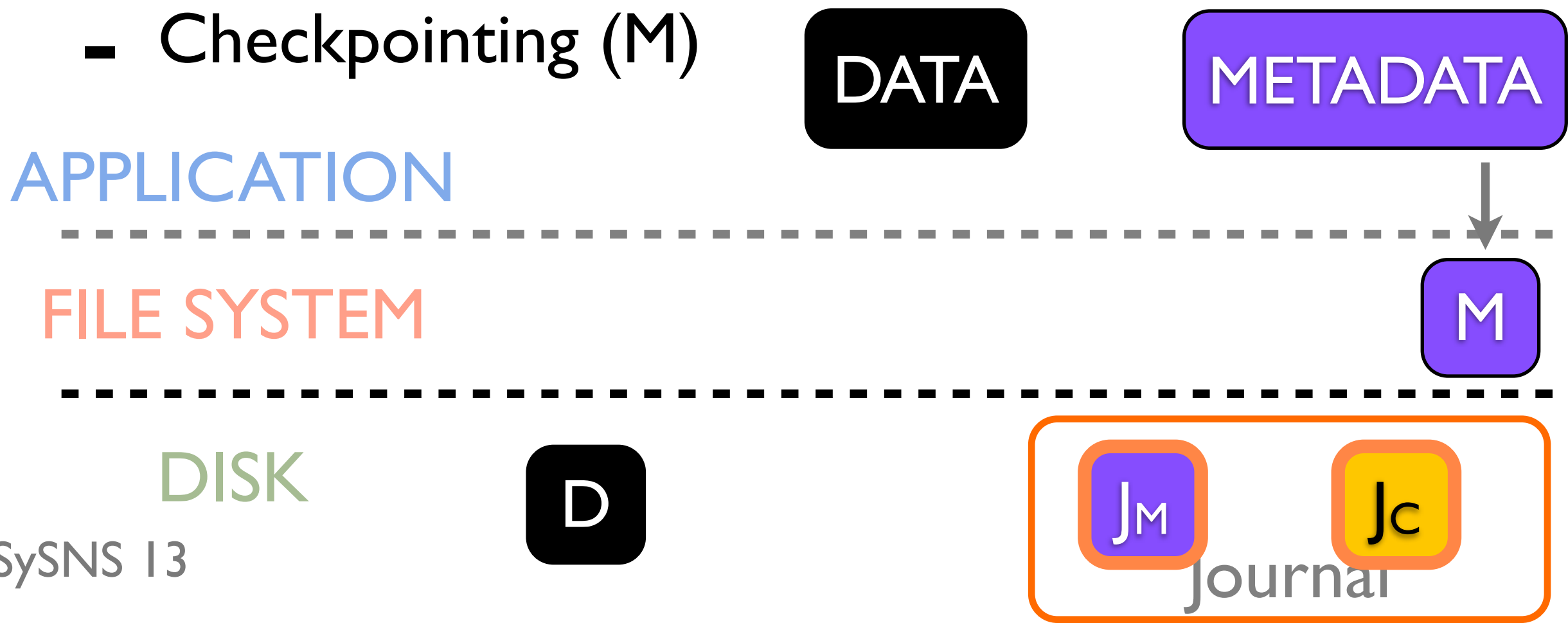
Journal

Journaling Overview

Workload: Creating and writing to a file

Journaling protocol (ordered journaling)

- Data write (D)
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DATA

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DISK

D

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Journal

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- Journaling Overview
- **Realizing Ordering on Disks**
- Journaling without Ordering

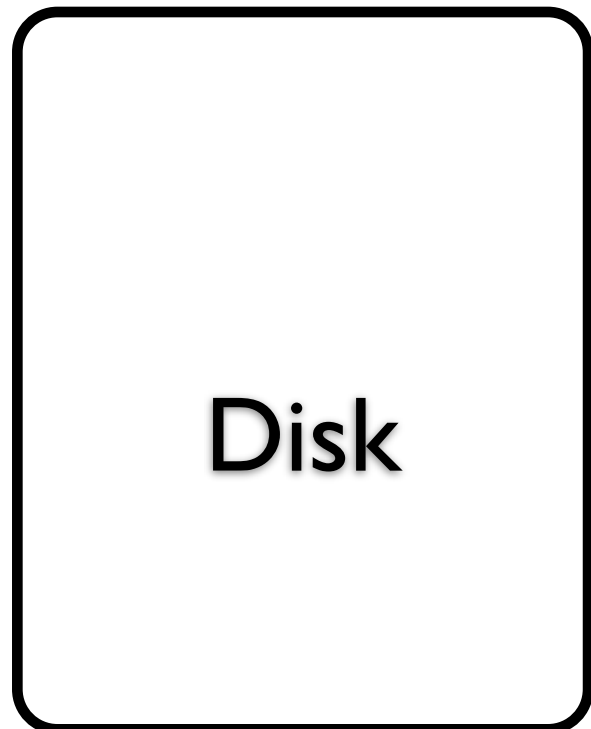
Optimistic File System

Results

Conclusion

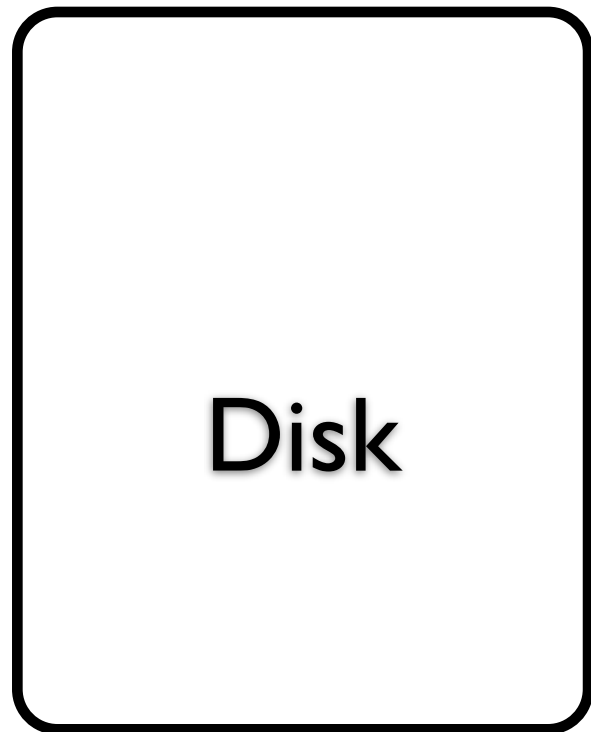
How Writes are Ordered

Original
Disks



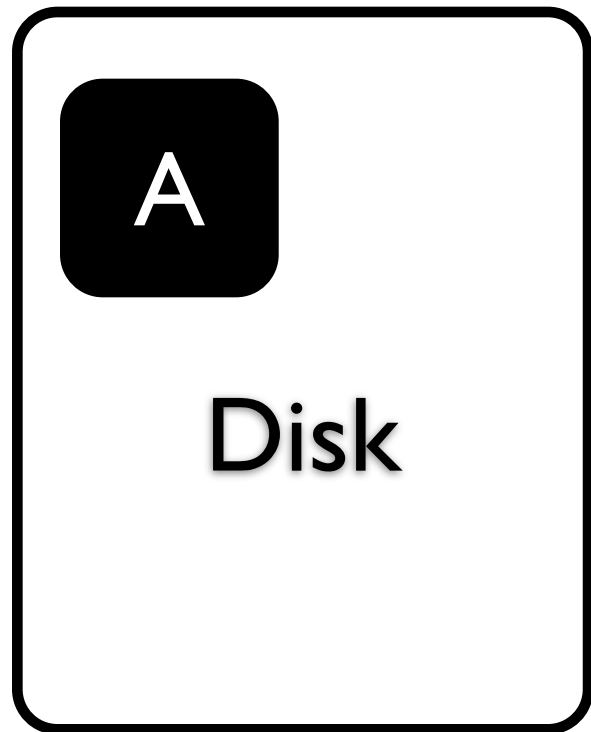
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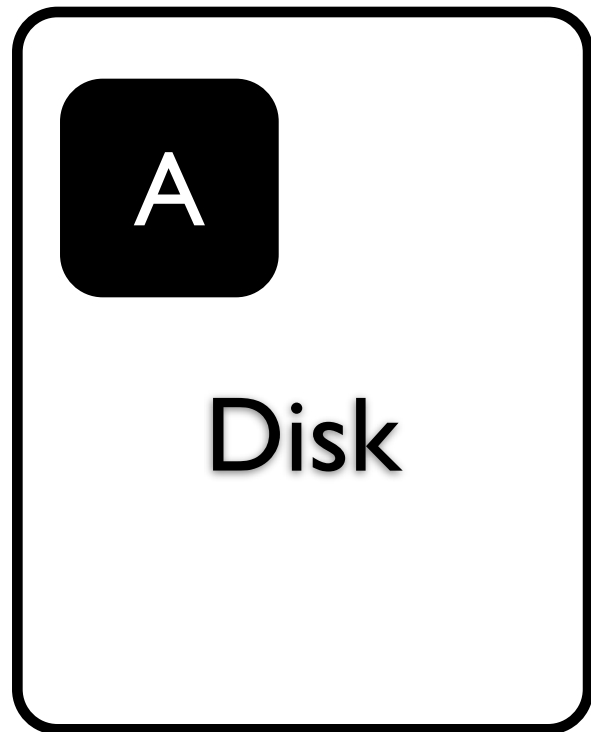
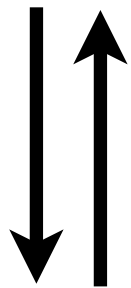
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Original
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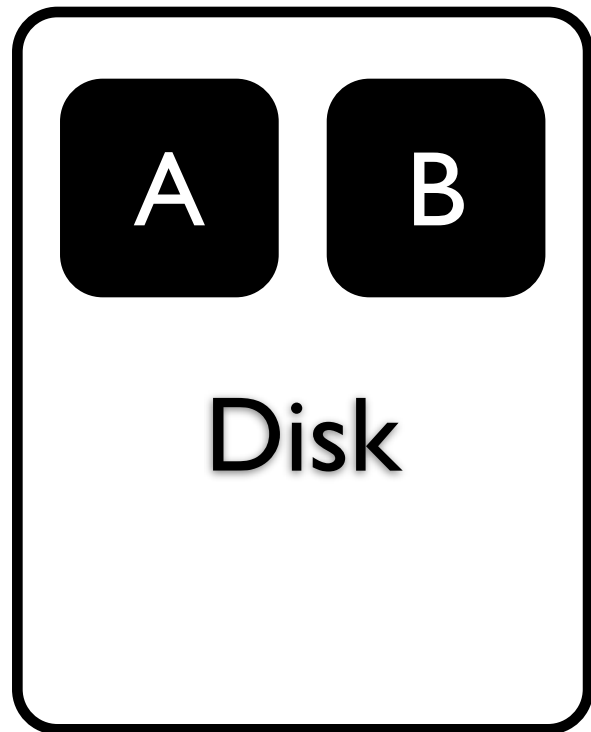
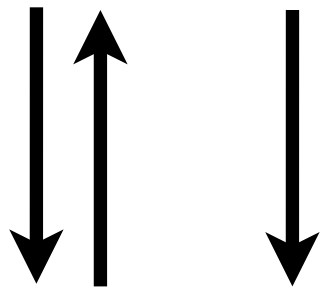
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Original
Disks



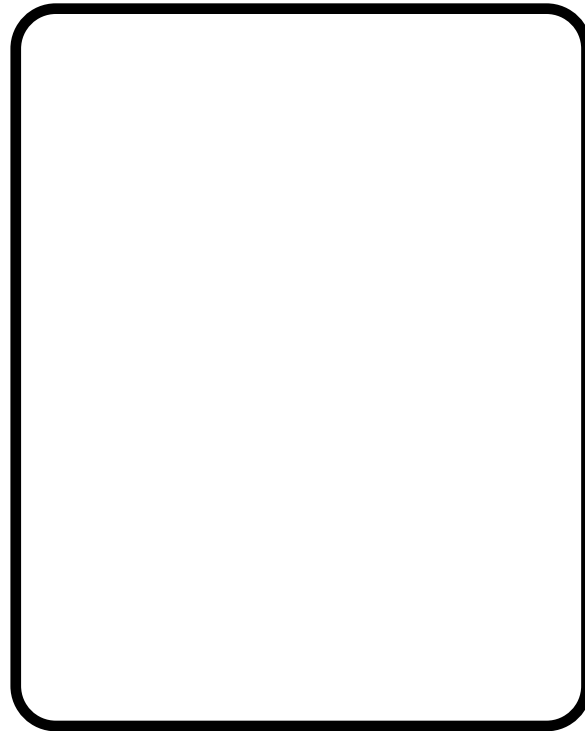
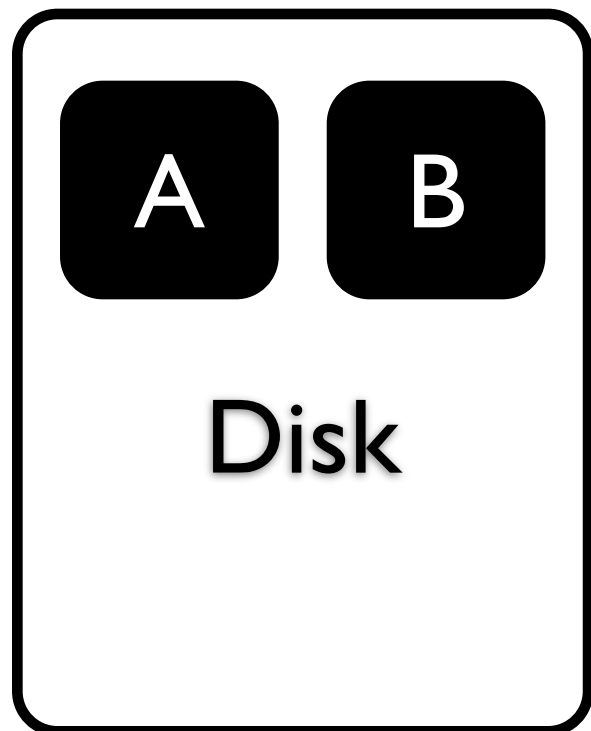
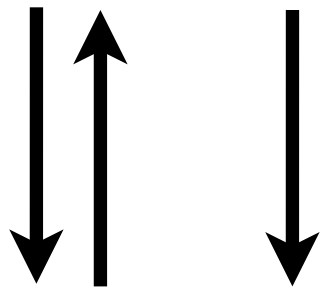
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Original
Disks



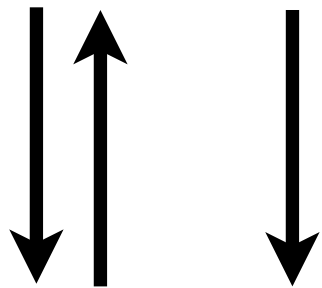
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Original
Disks

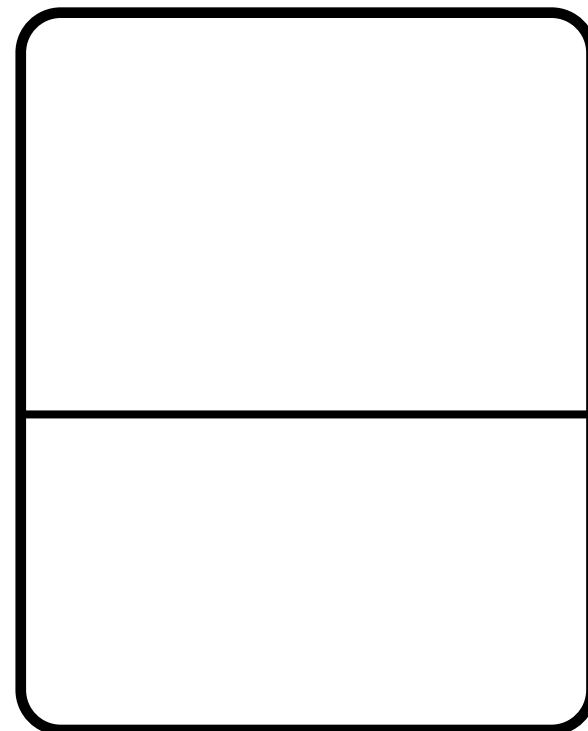
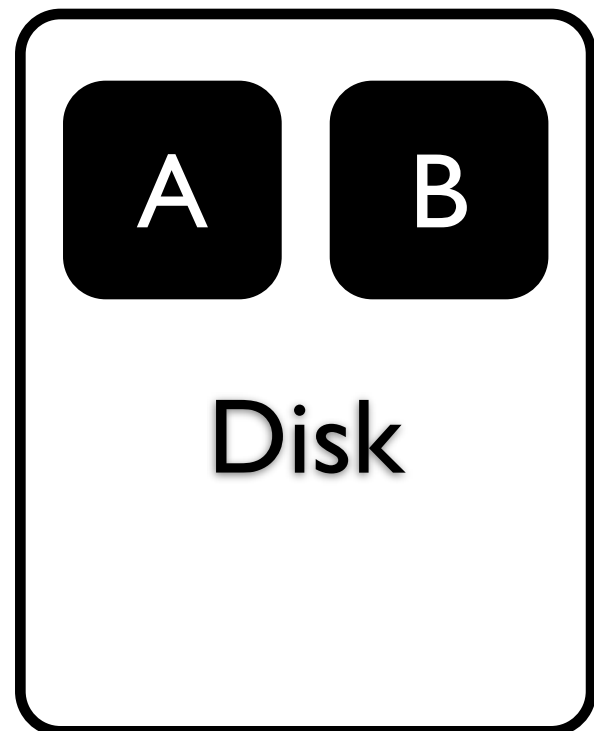


How Writes are Ordered

Original
Disks



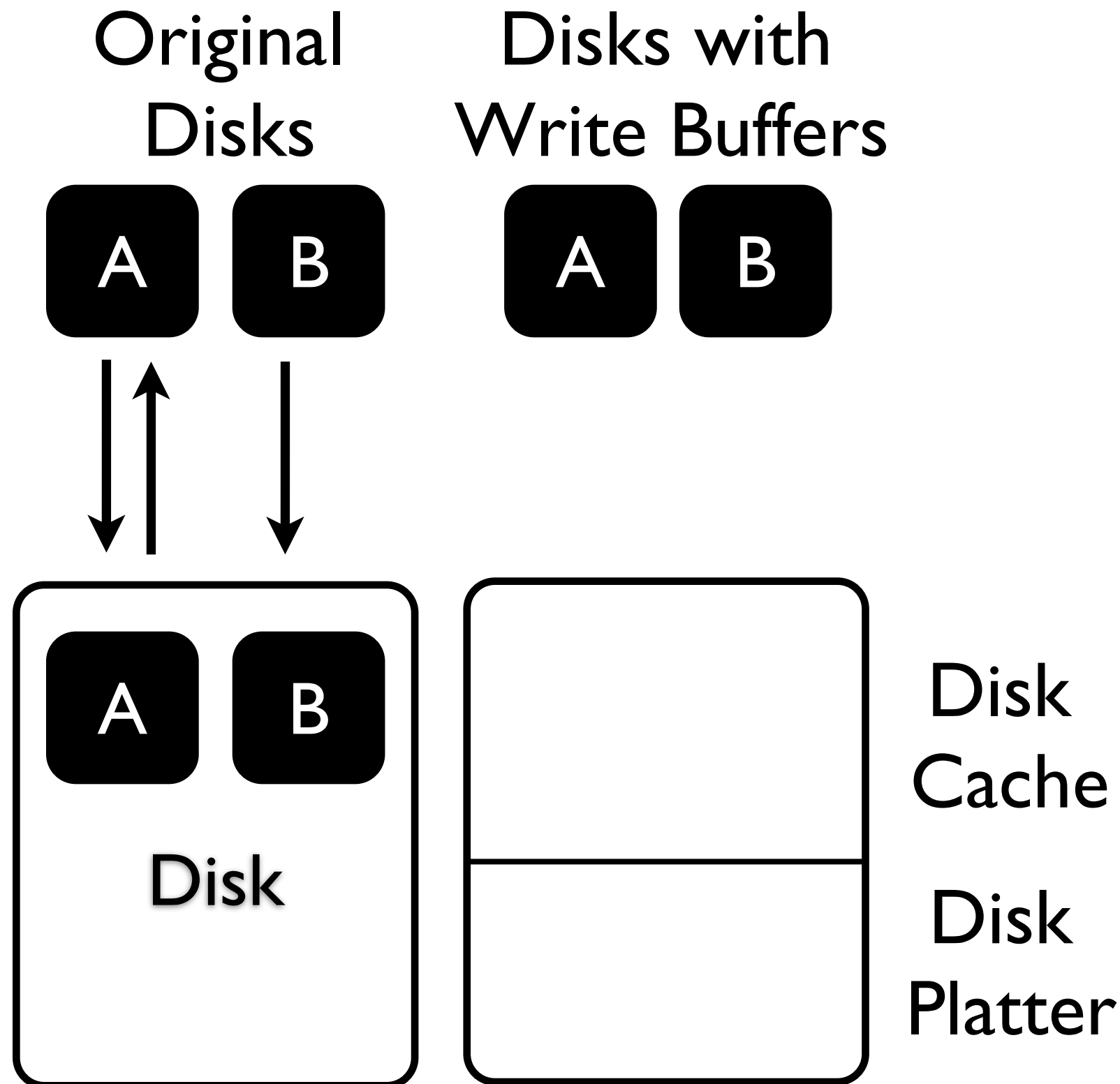
Disks with
Write Buffers



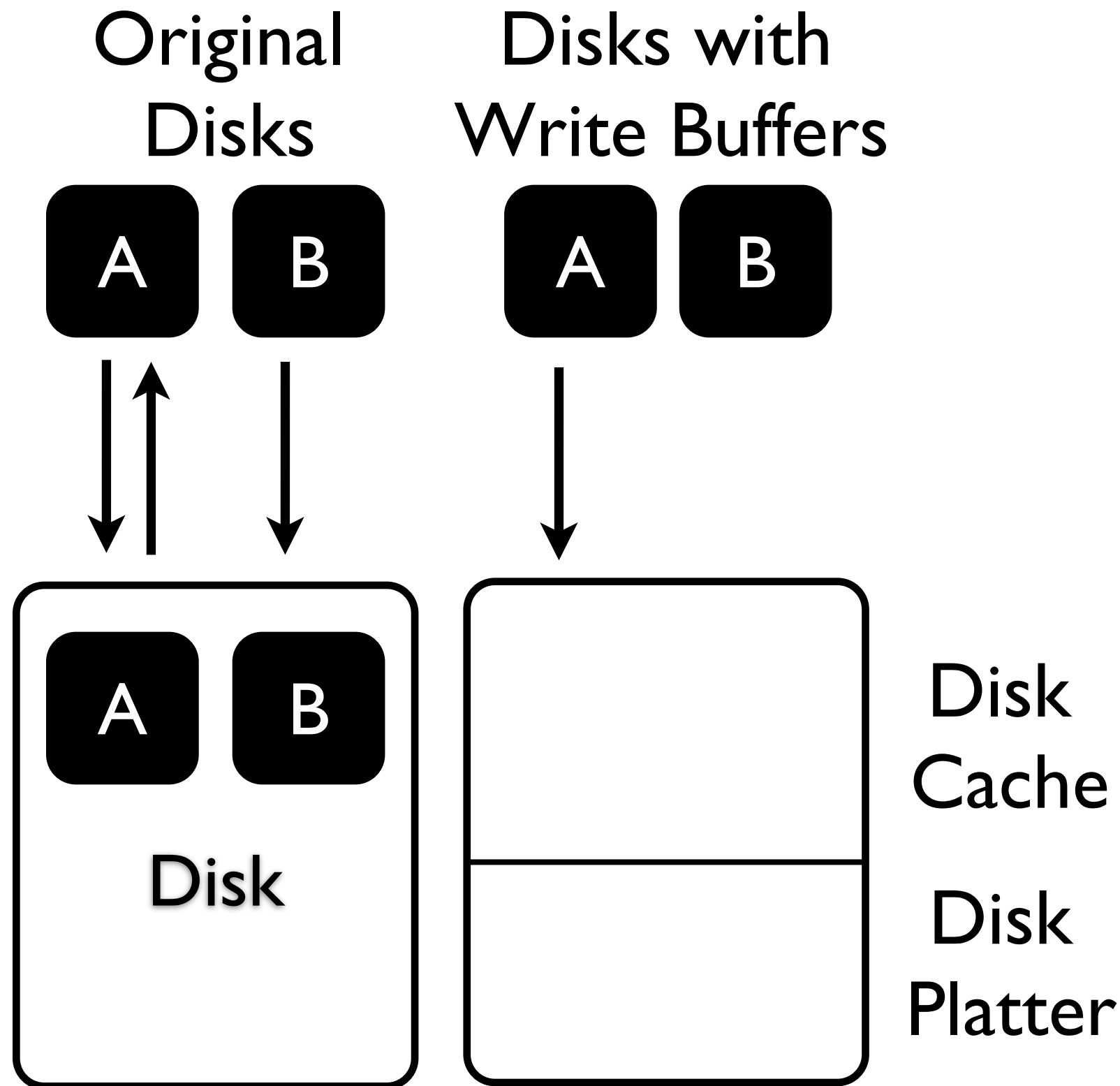
Disk
Cache

Disk
Platter

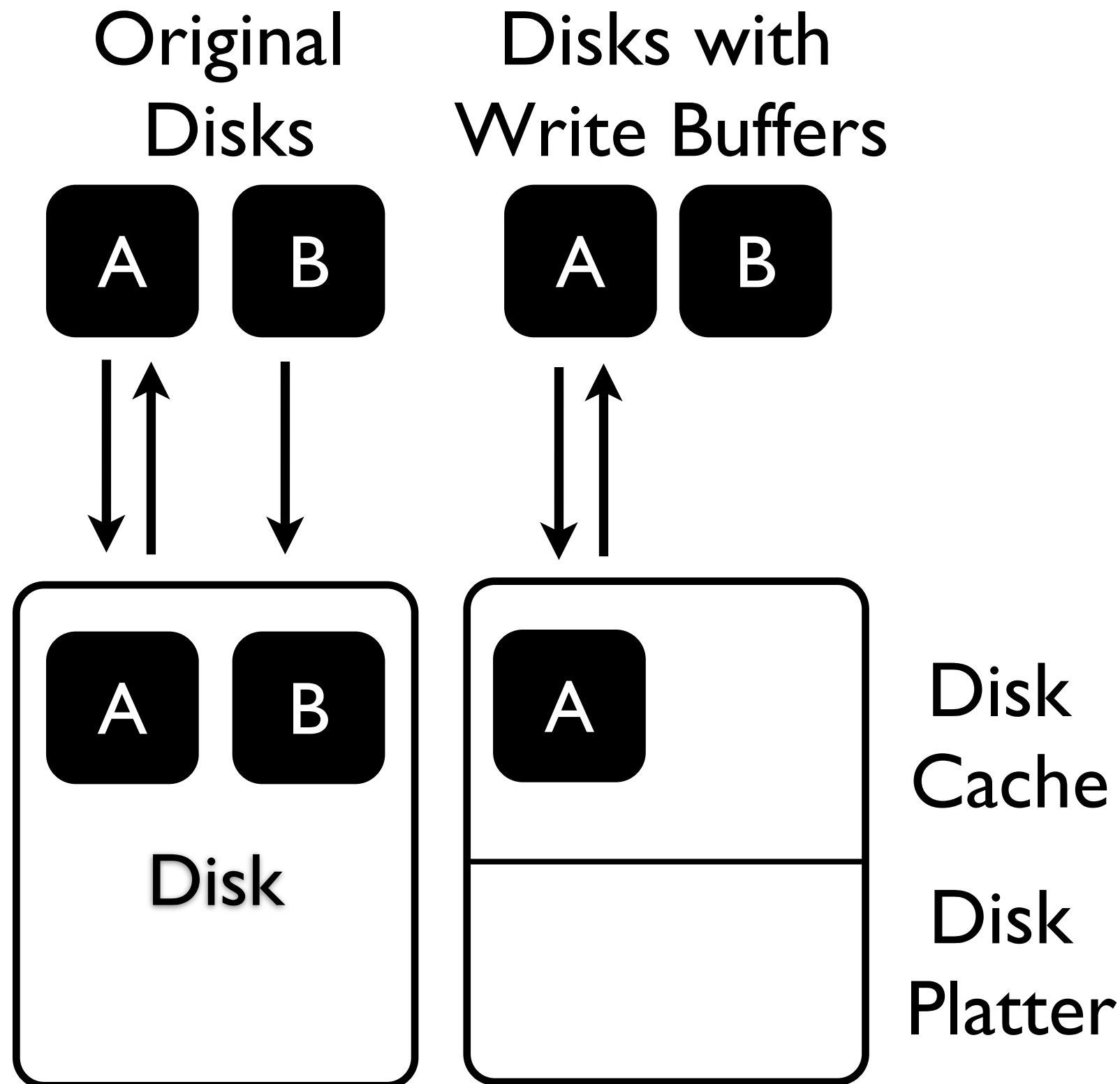
How Writes are Ordered



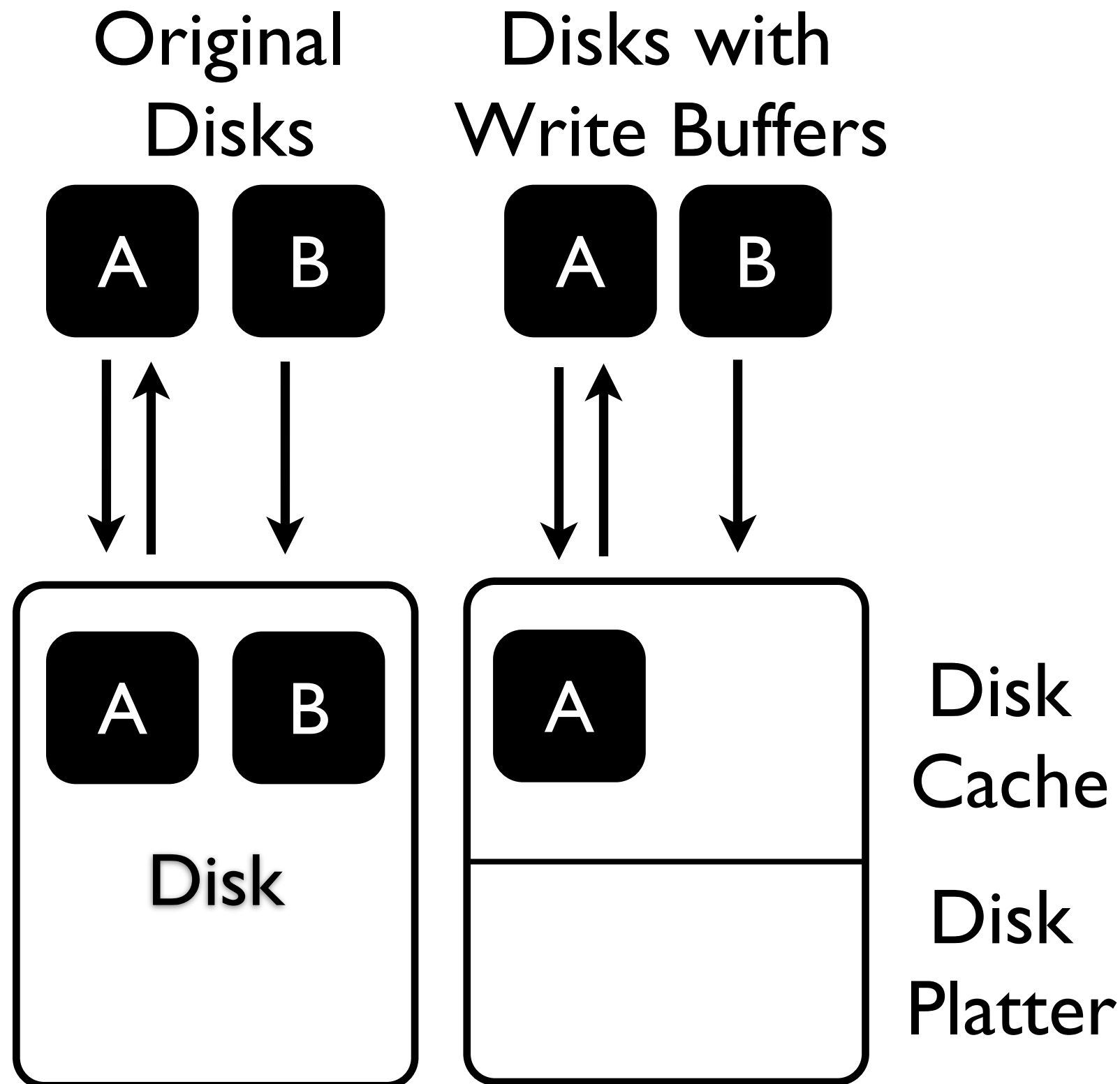
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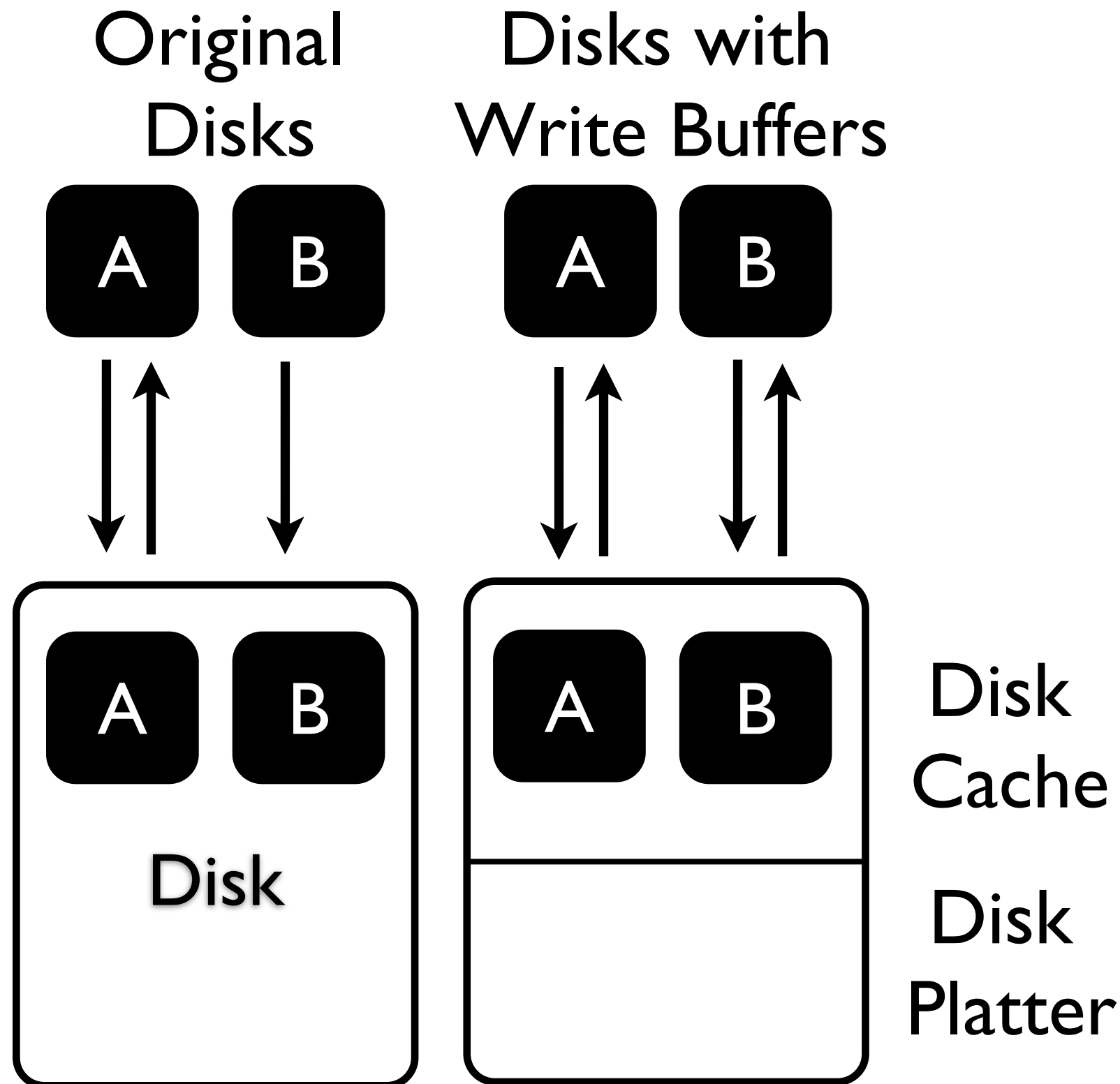
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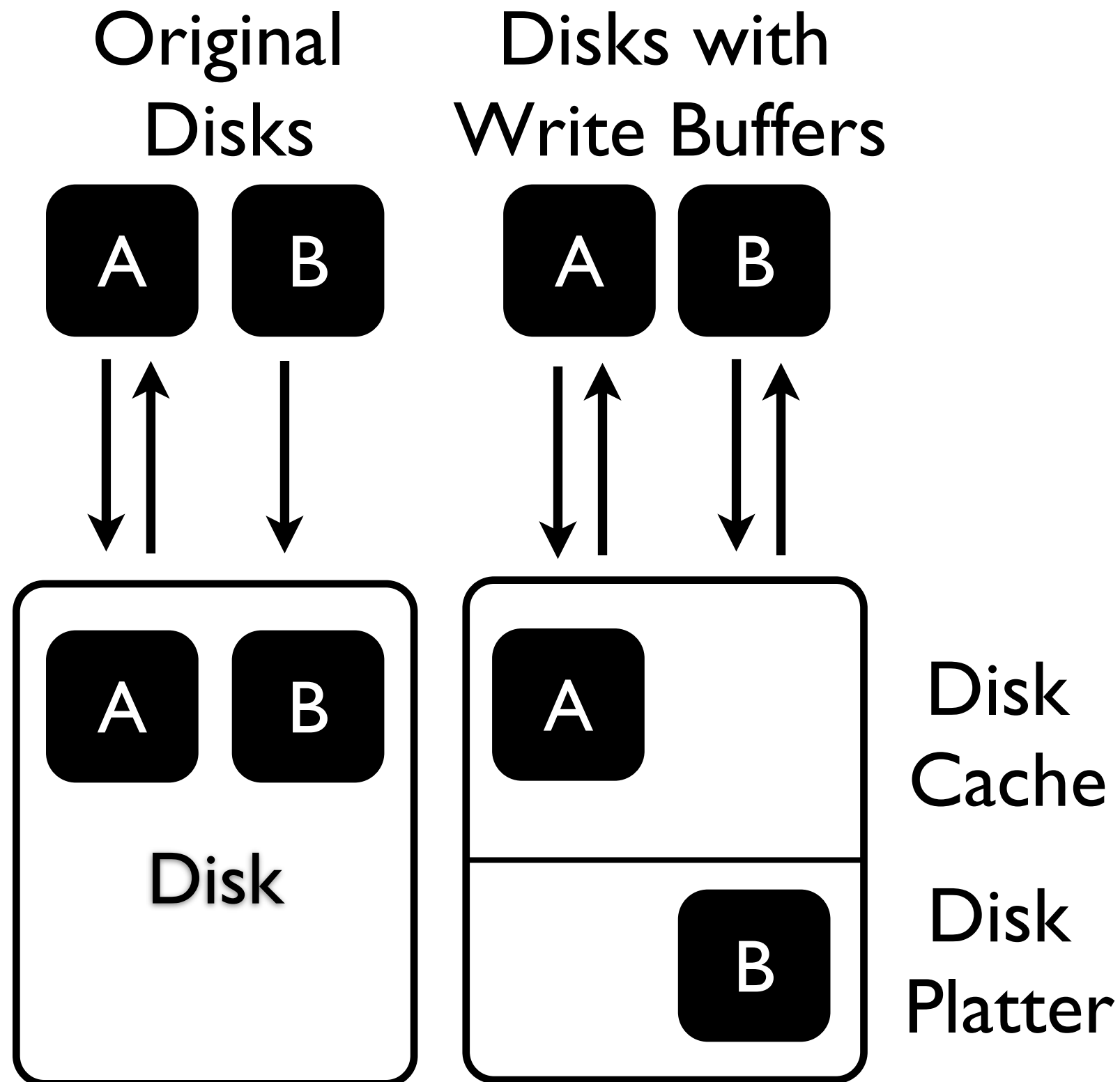
How Writes are Ordered



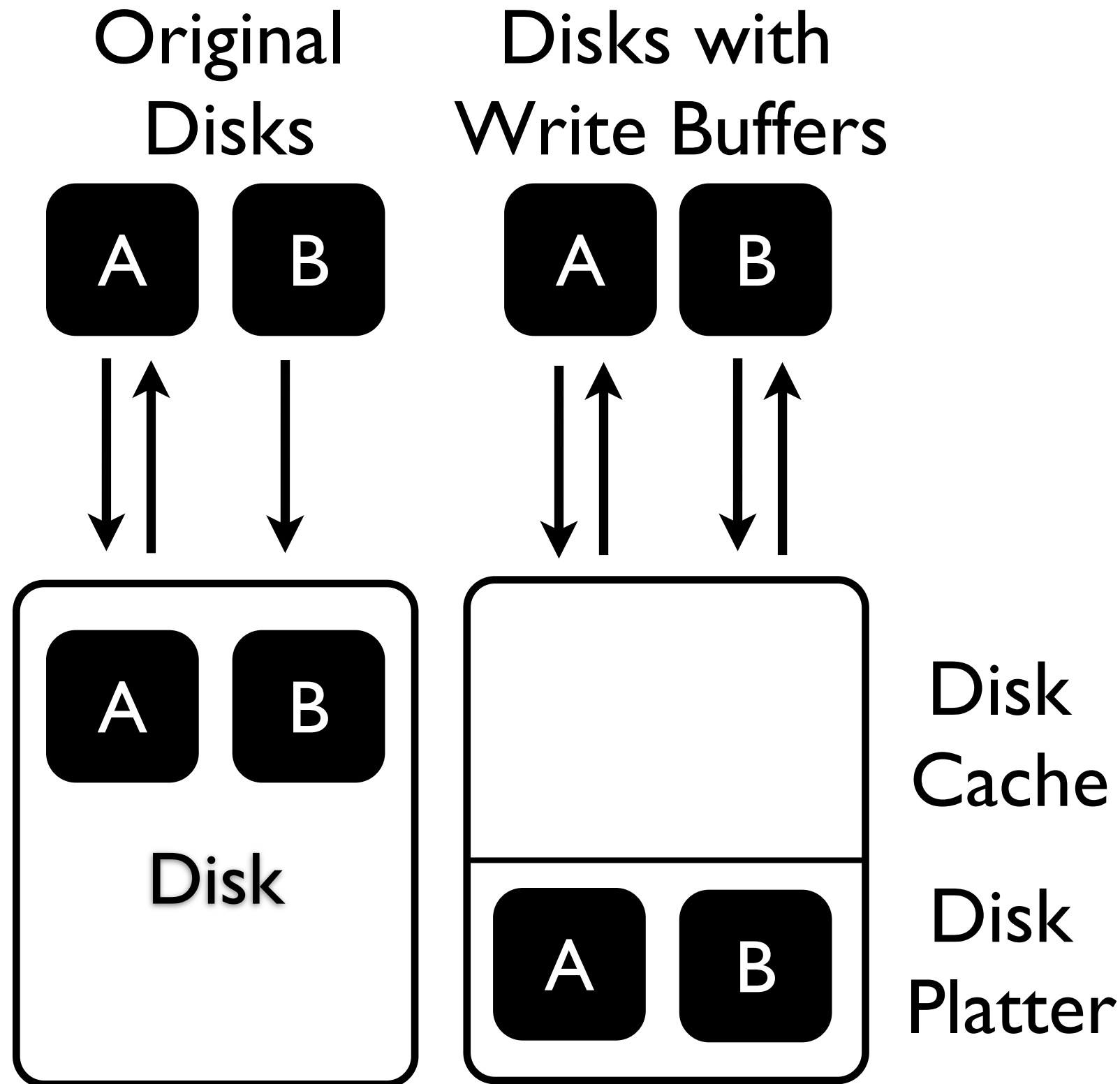
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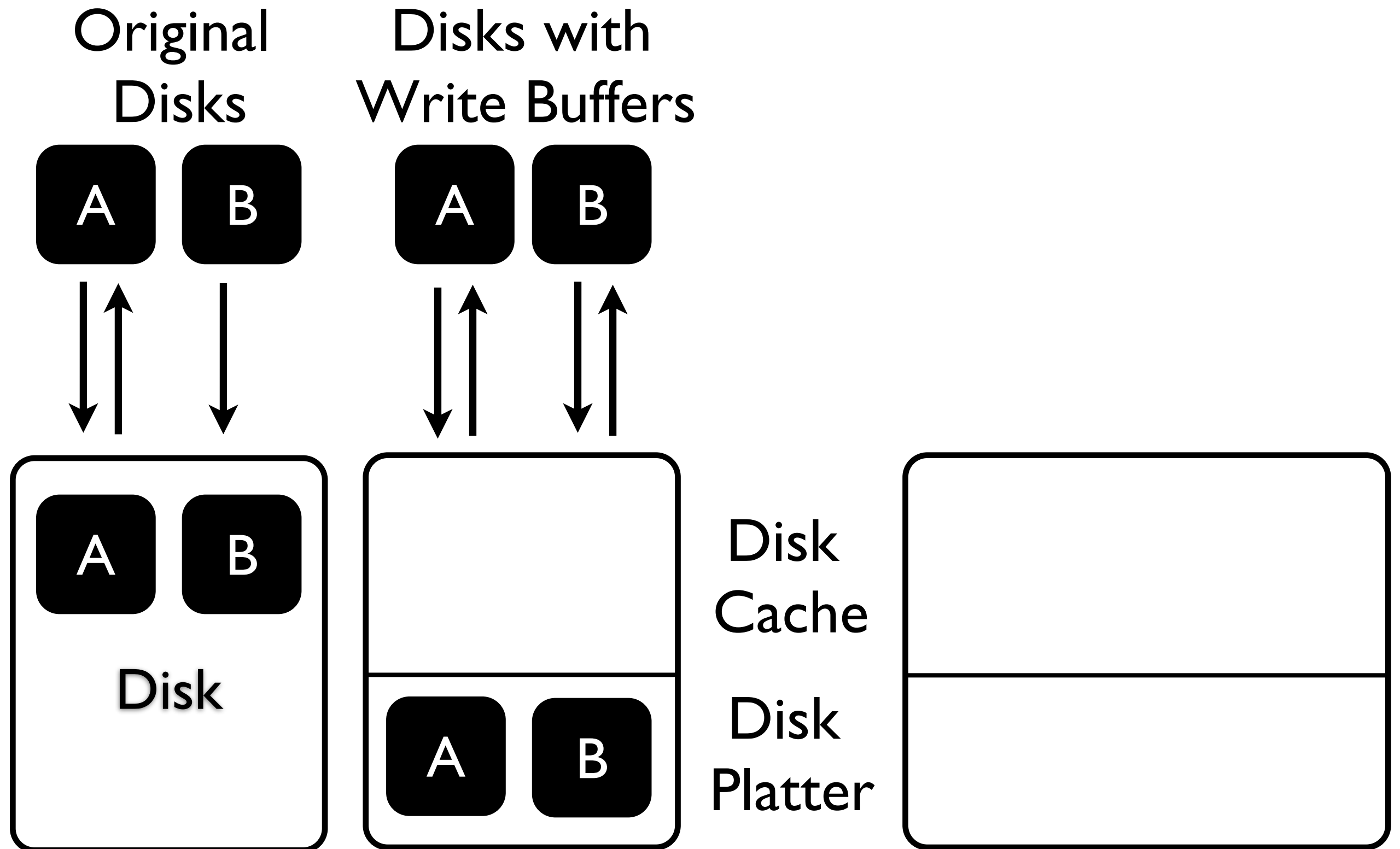
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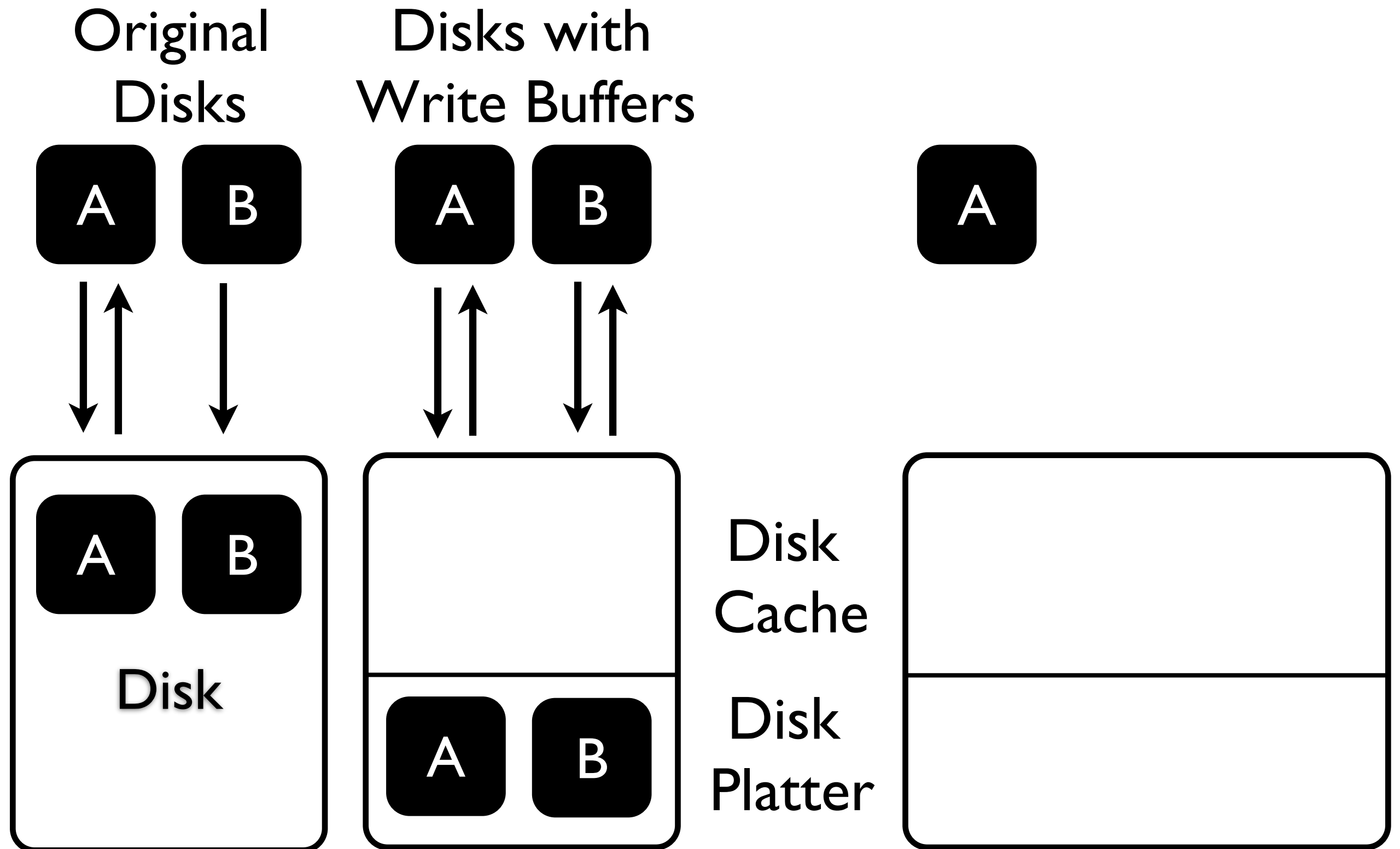
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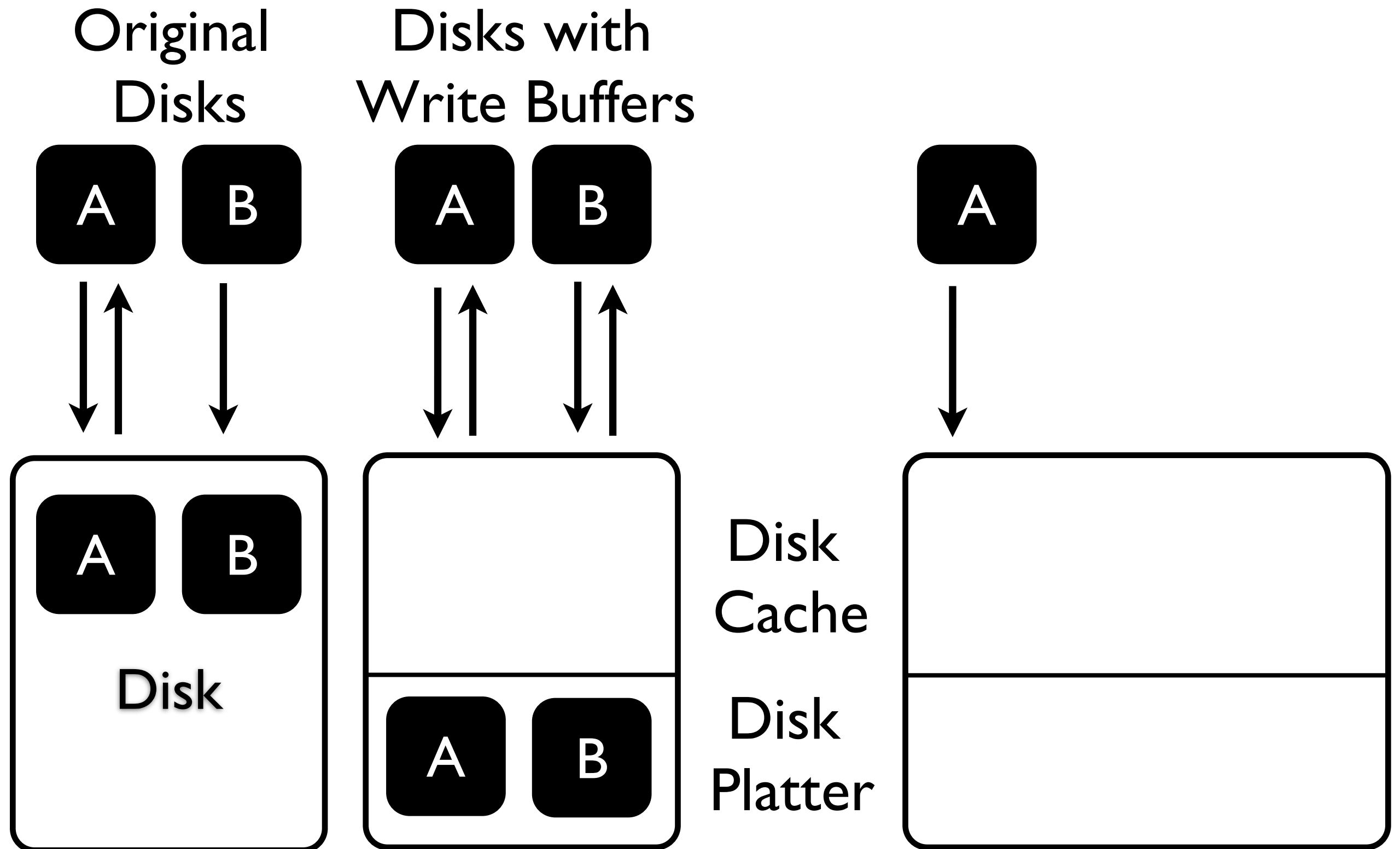
How Writes are Ordered



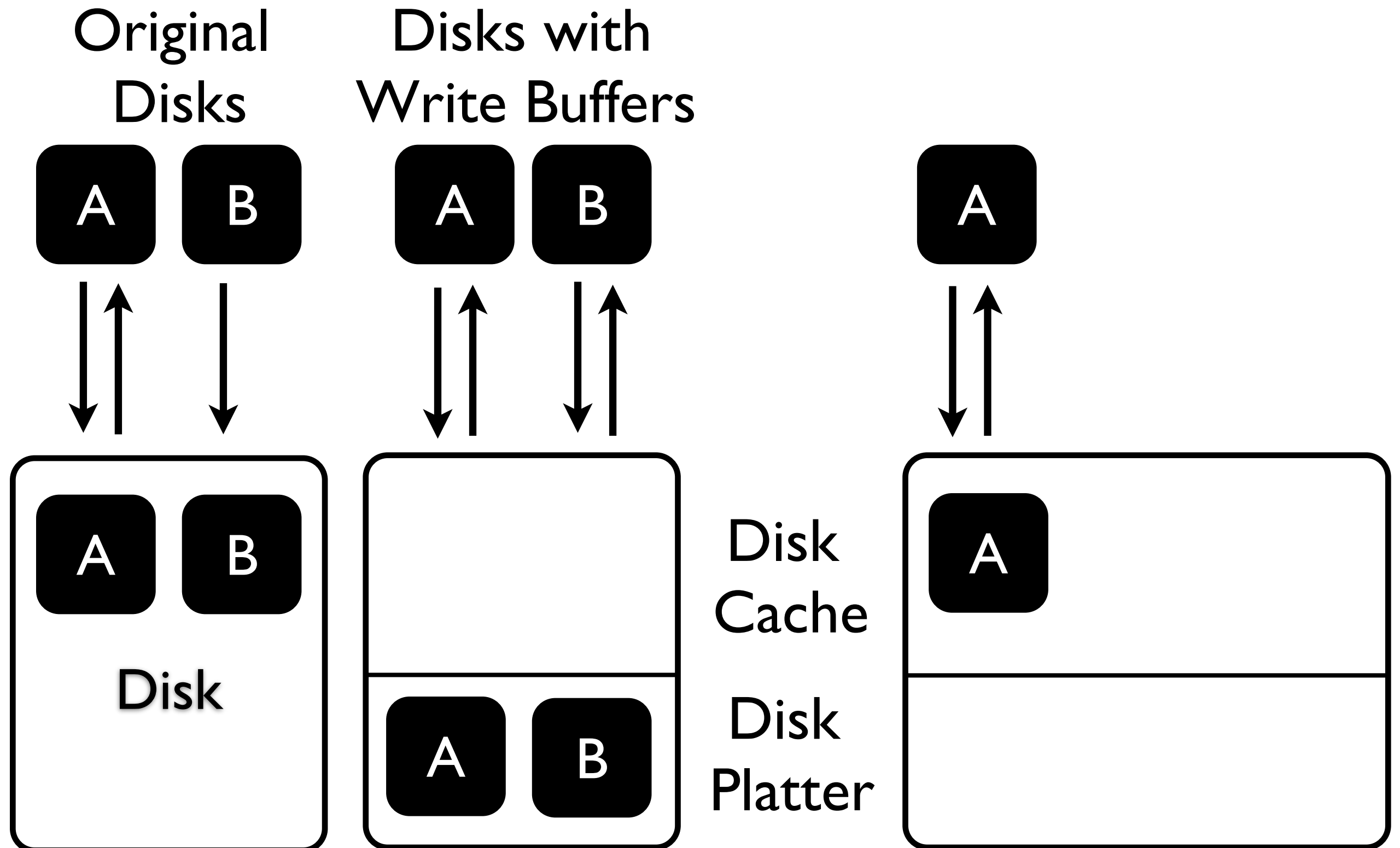
How Writes are Ordered



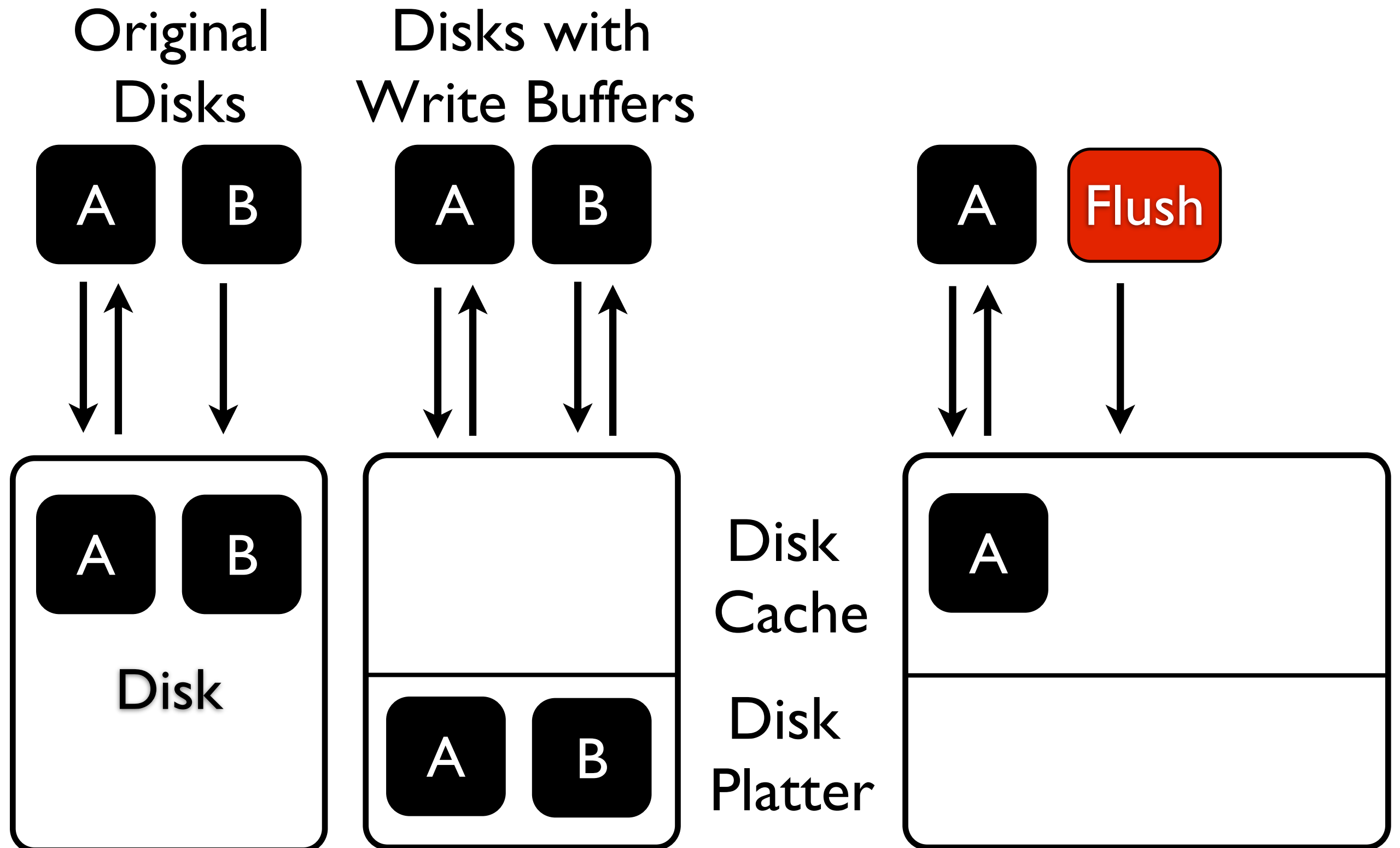
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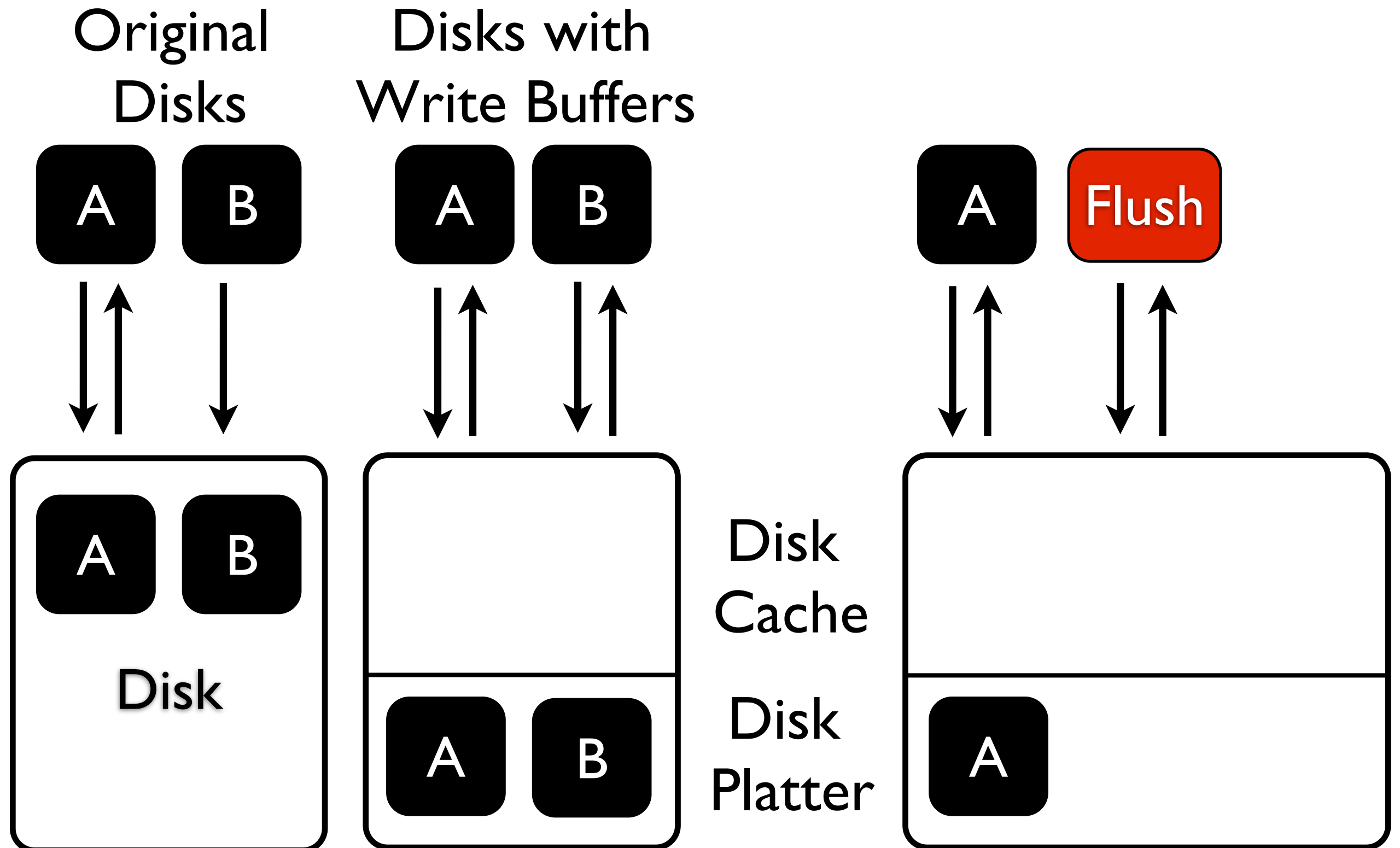
How Writes are Ordered



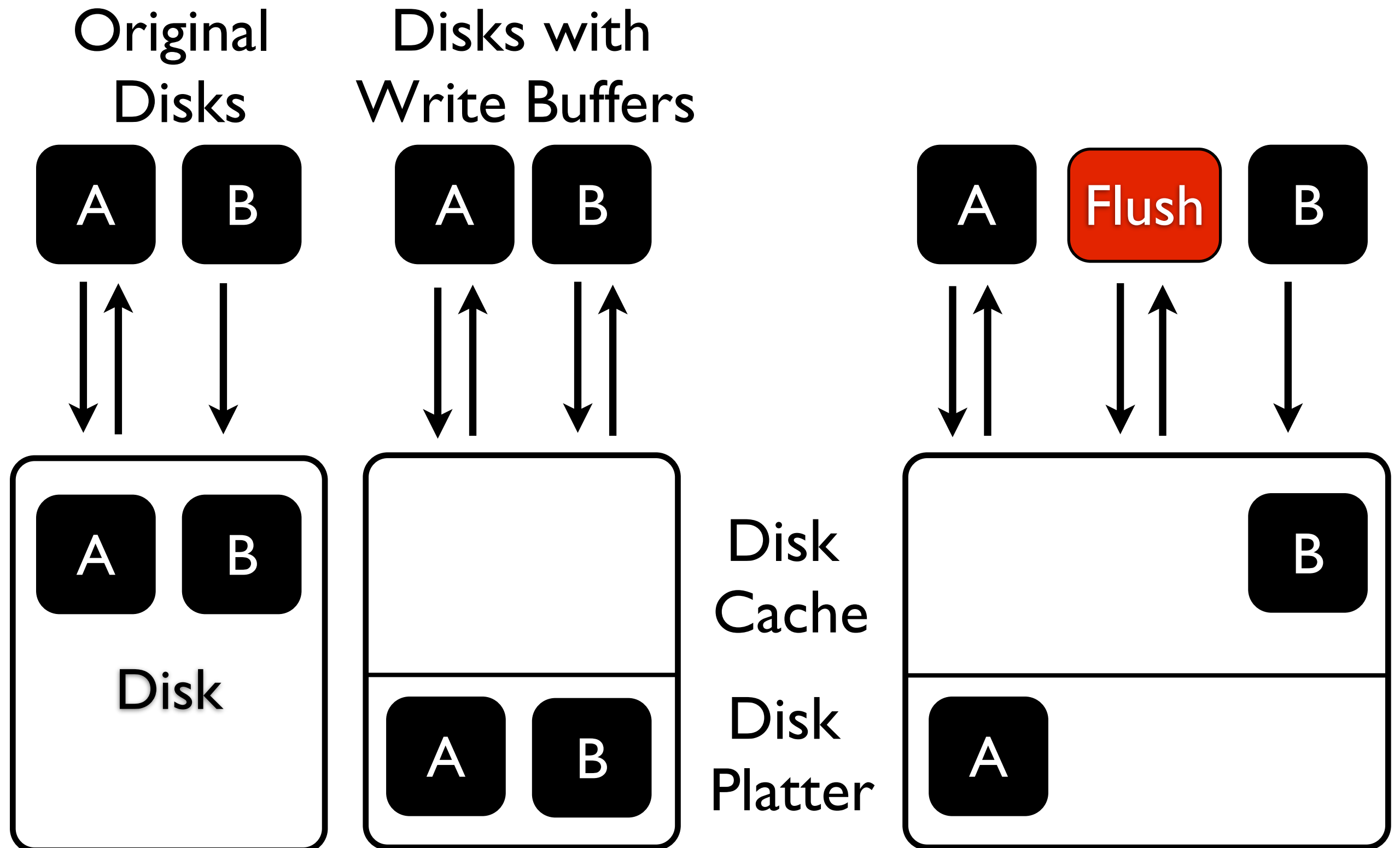
How Writes are Ordered



How Writes are Ordered



How Writes are Ordered



Journaling with Flushes

APPLICATION

FILE SYSTEM

DISK CACHE

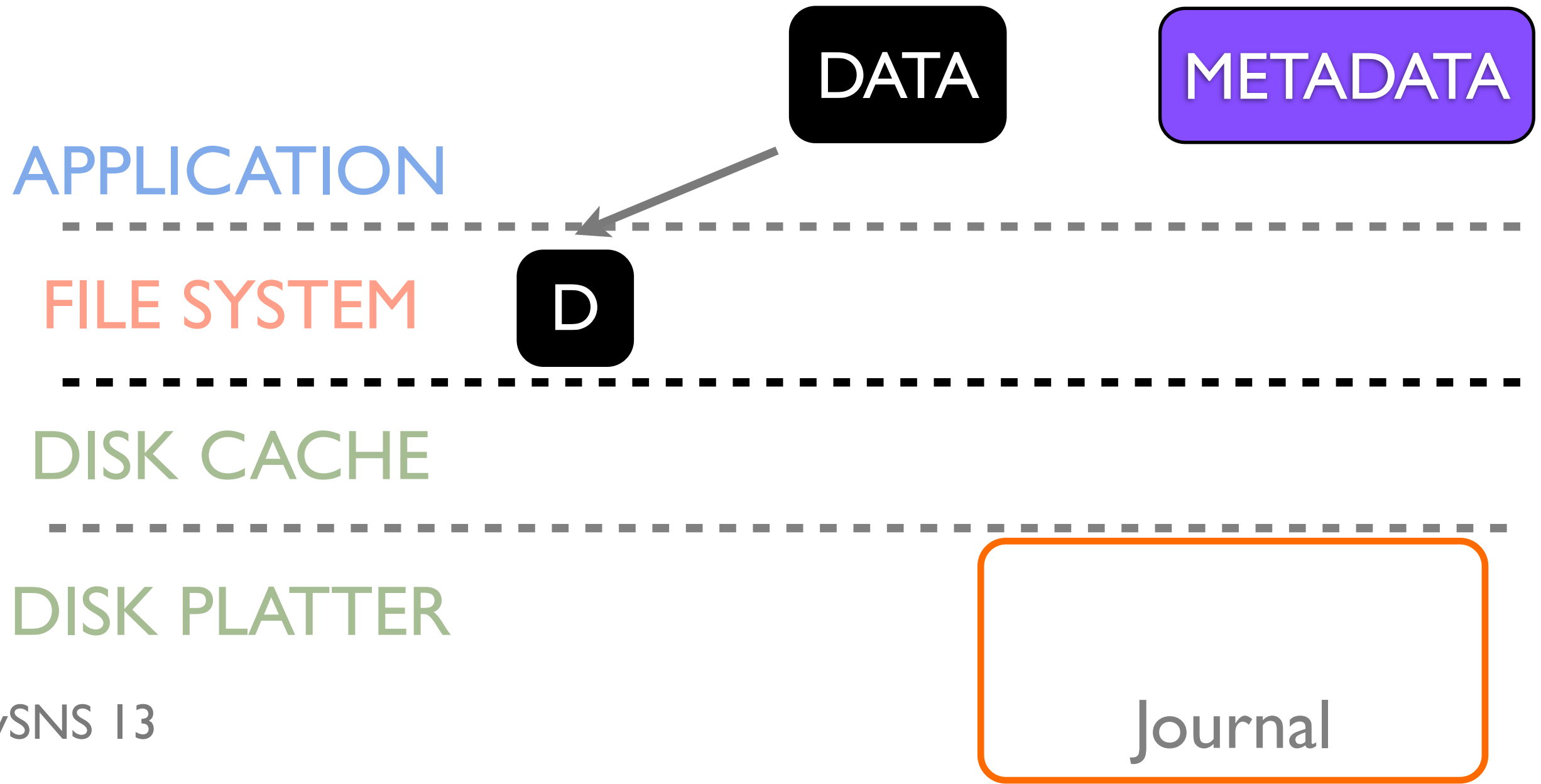
DISK PLATTER

Journal

Journaling with Flushes

Journaling protocol

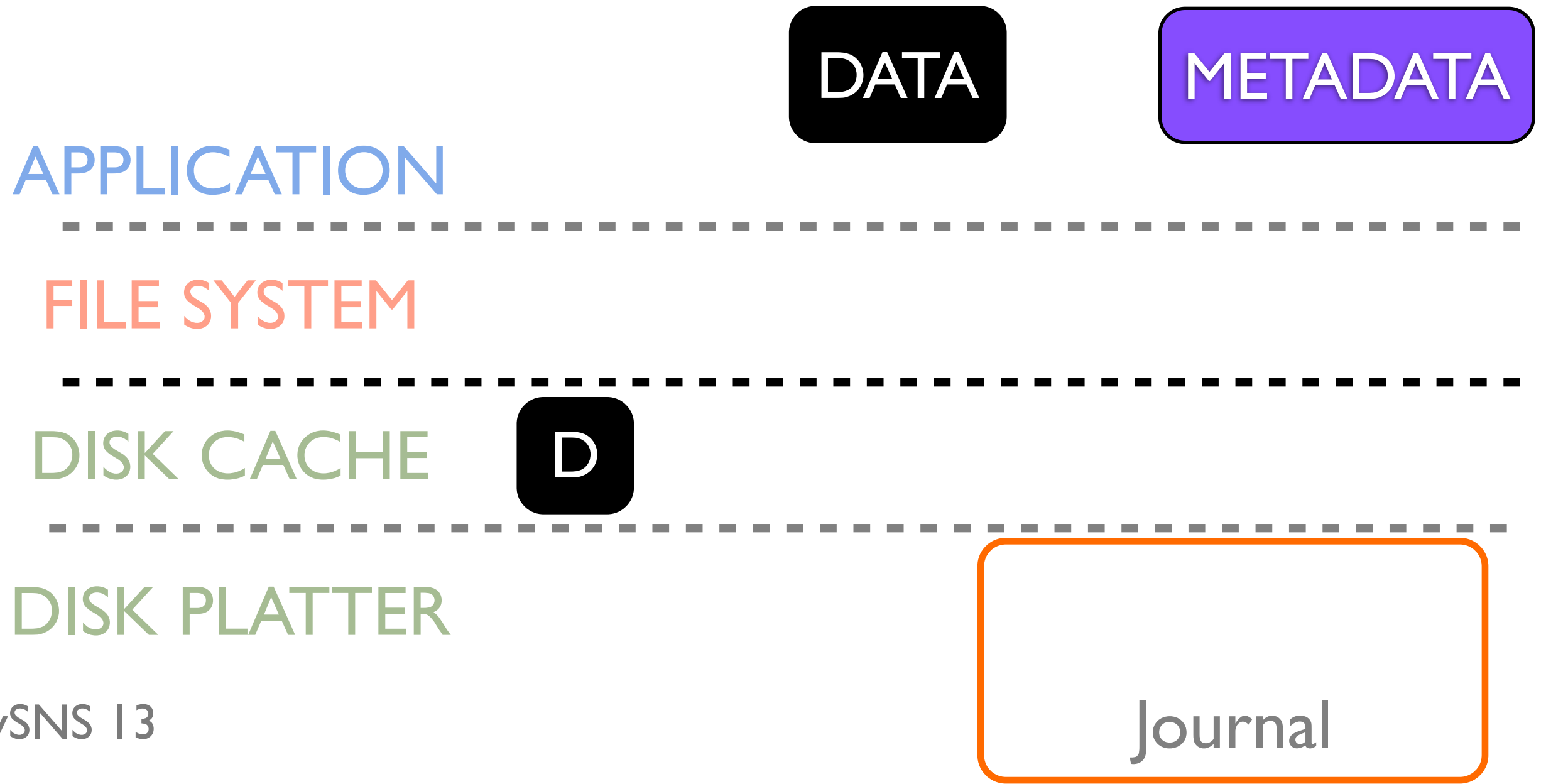
- Data write (D)



Journaling with Flushes

Journaling protocol

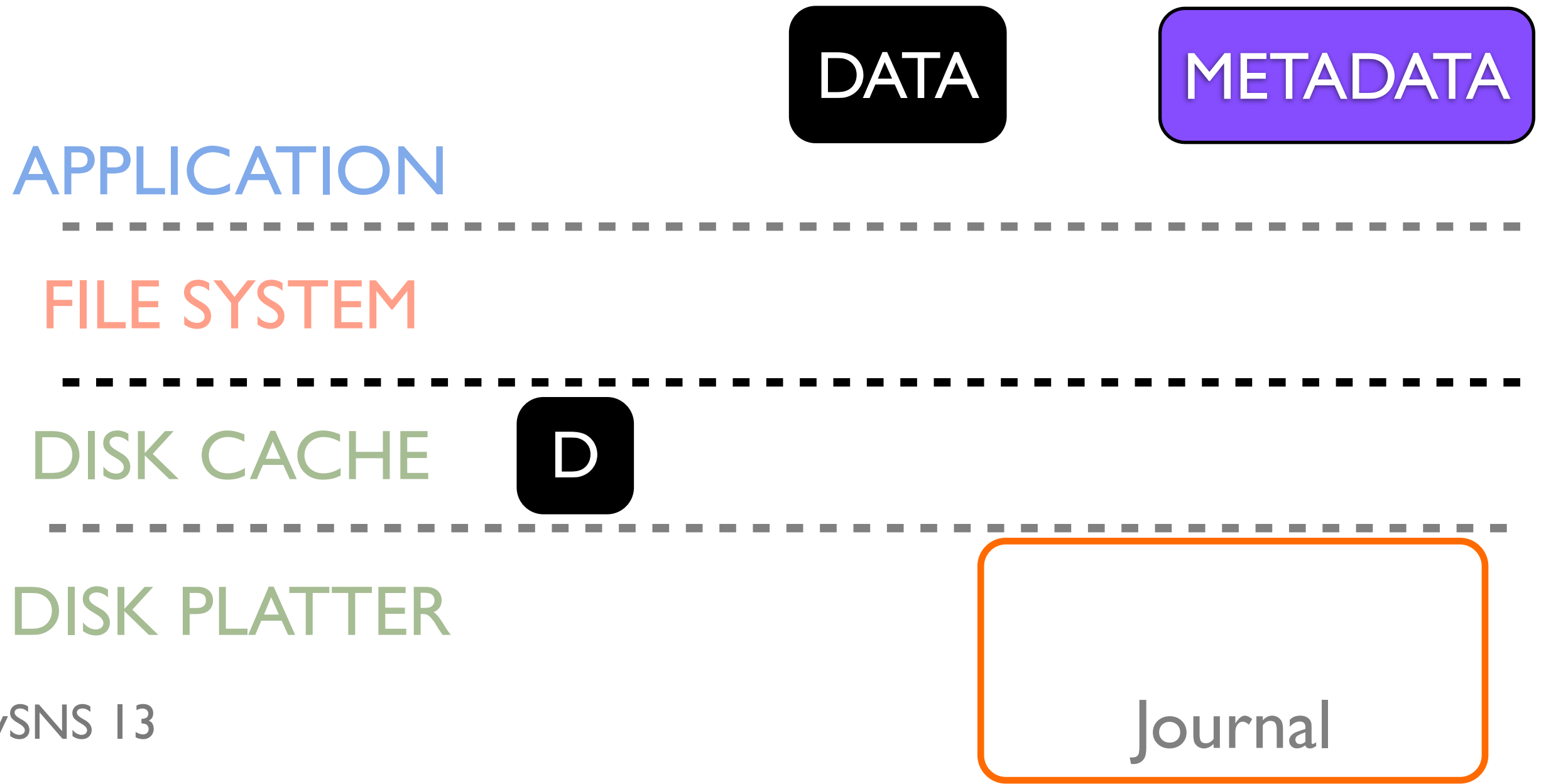
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Journaling with Flushes

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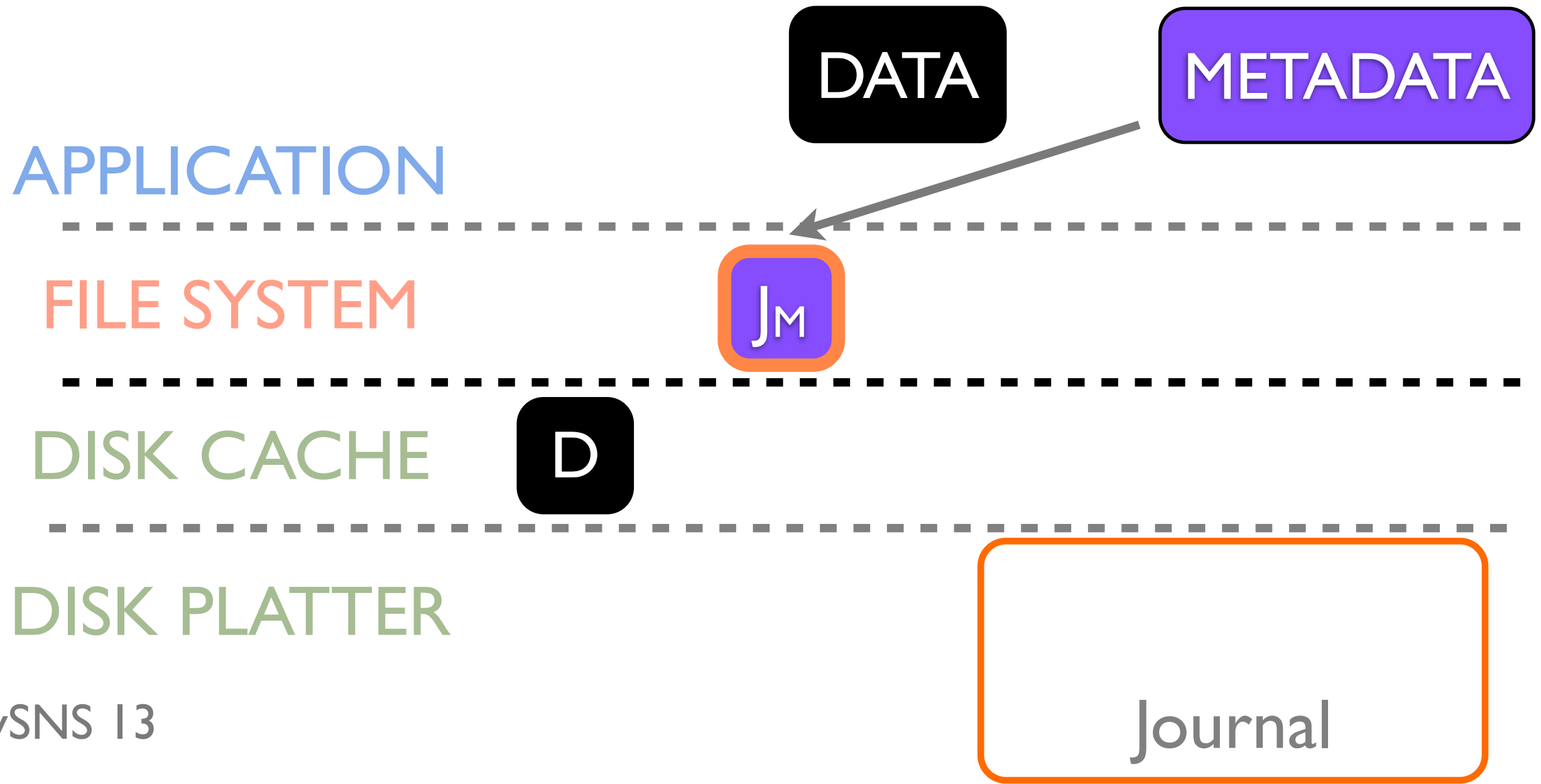
- Data write (D)
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Journaling with Flushes

Journaling protocol

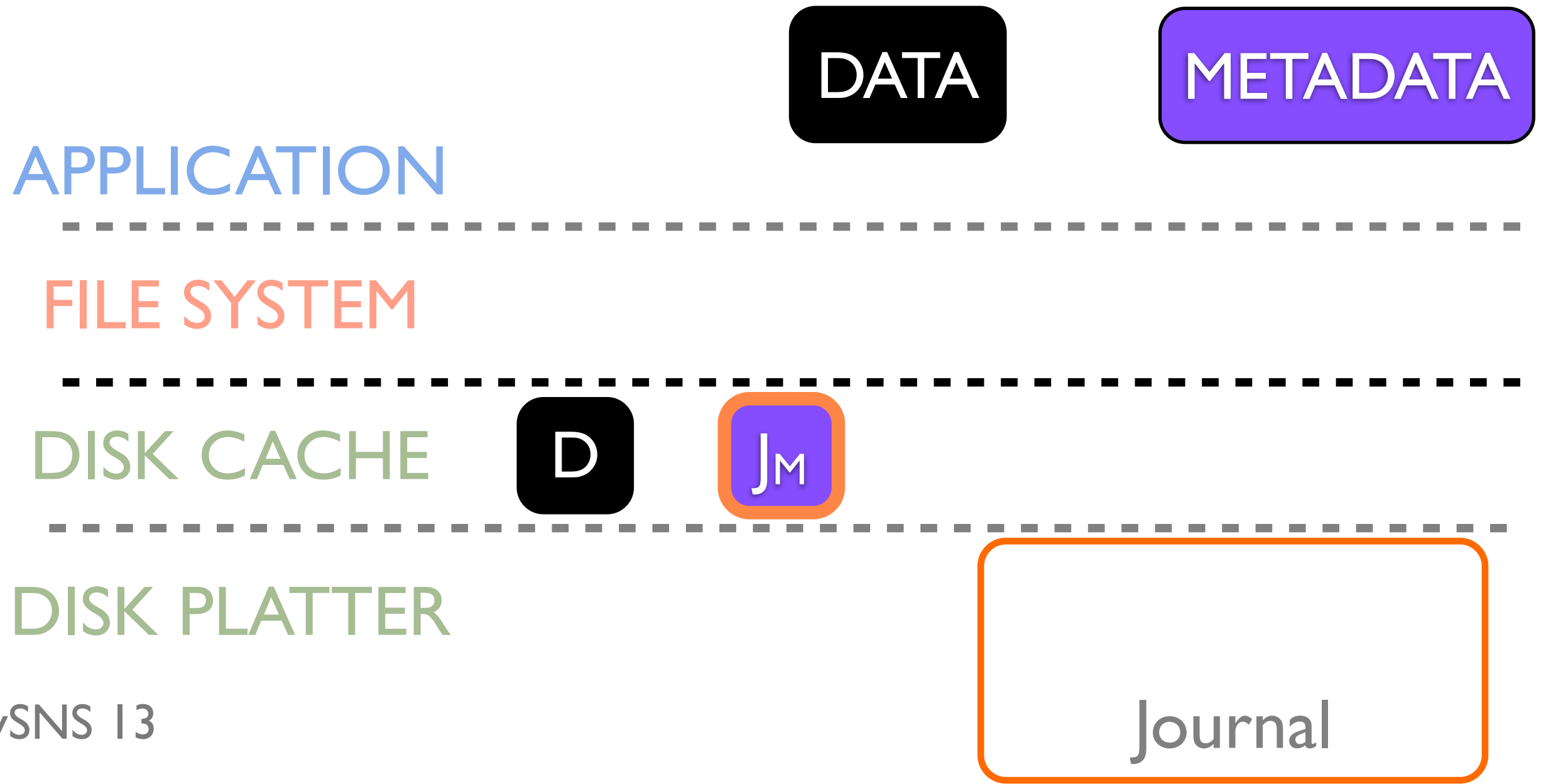
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Journaling with Flushes

Journaling protocol

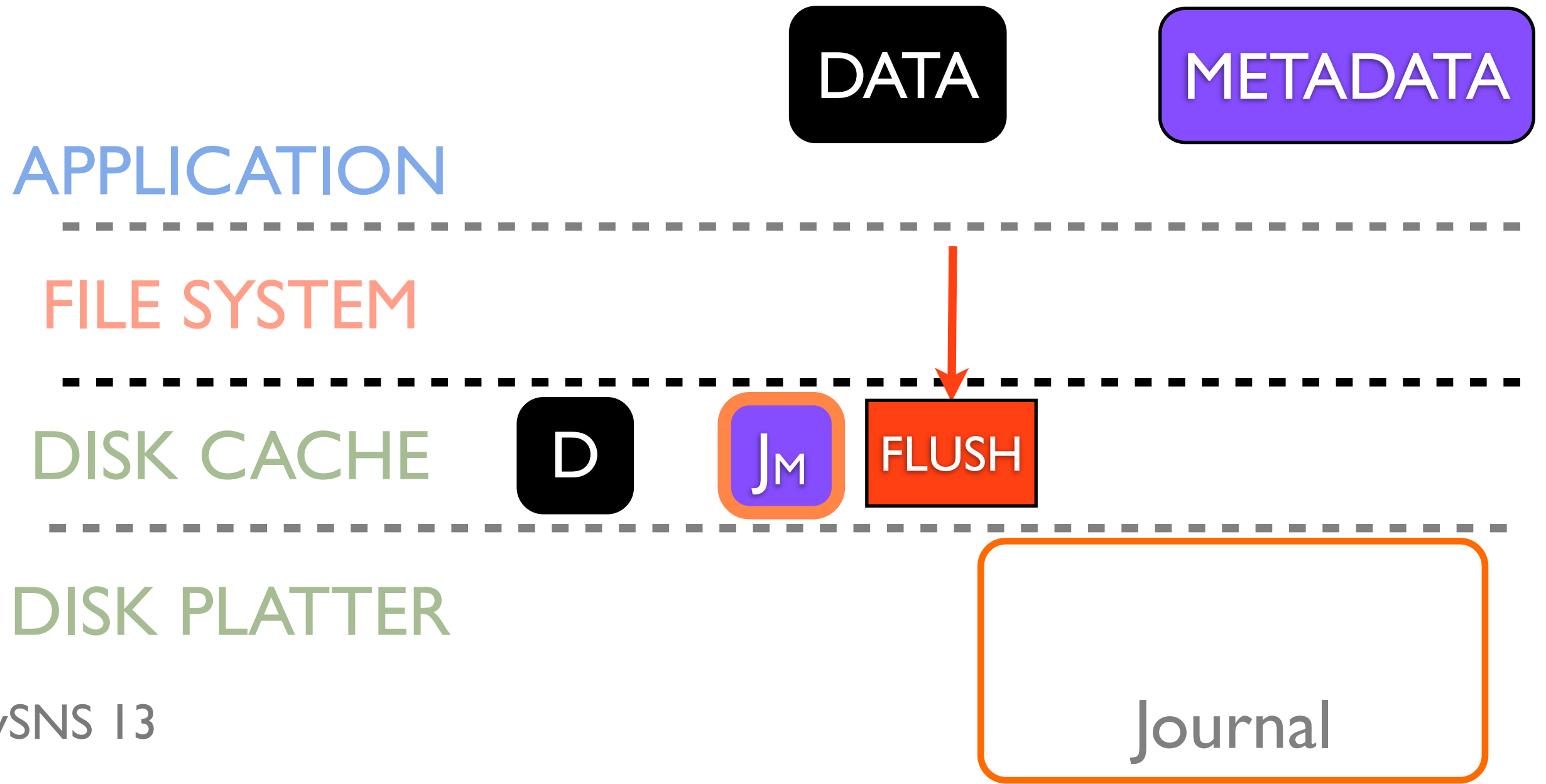
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Journaling with Flushes

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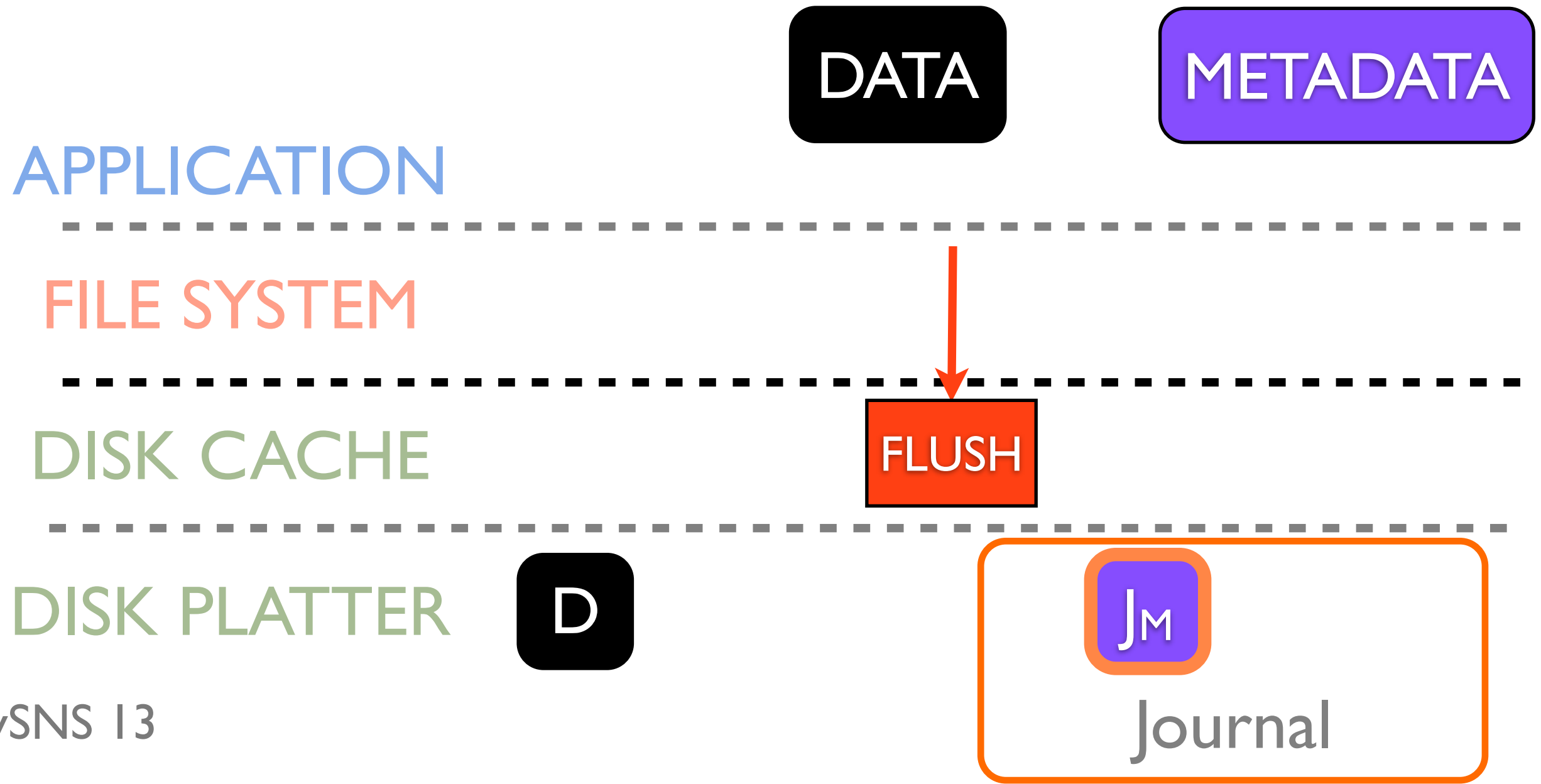
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Journaling with Flushes

Journaling protocol

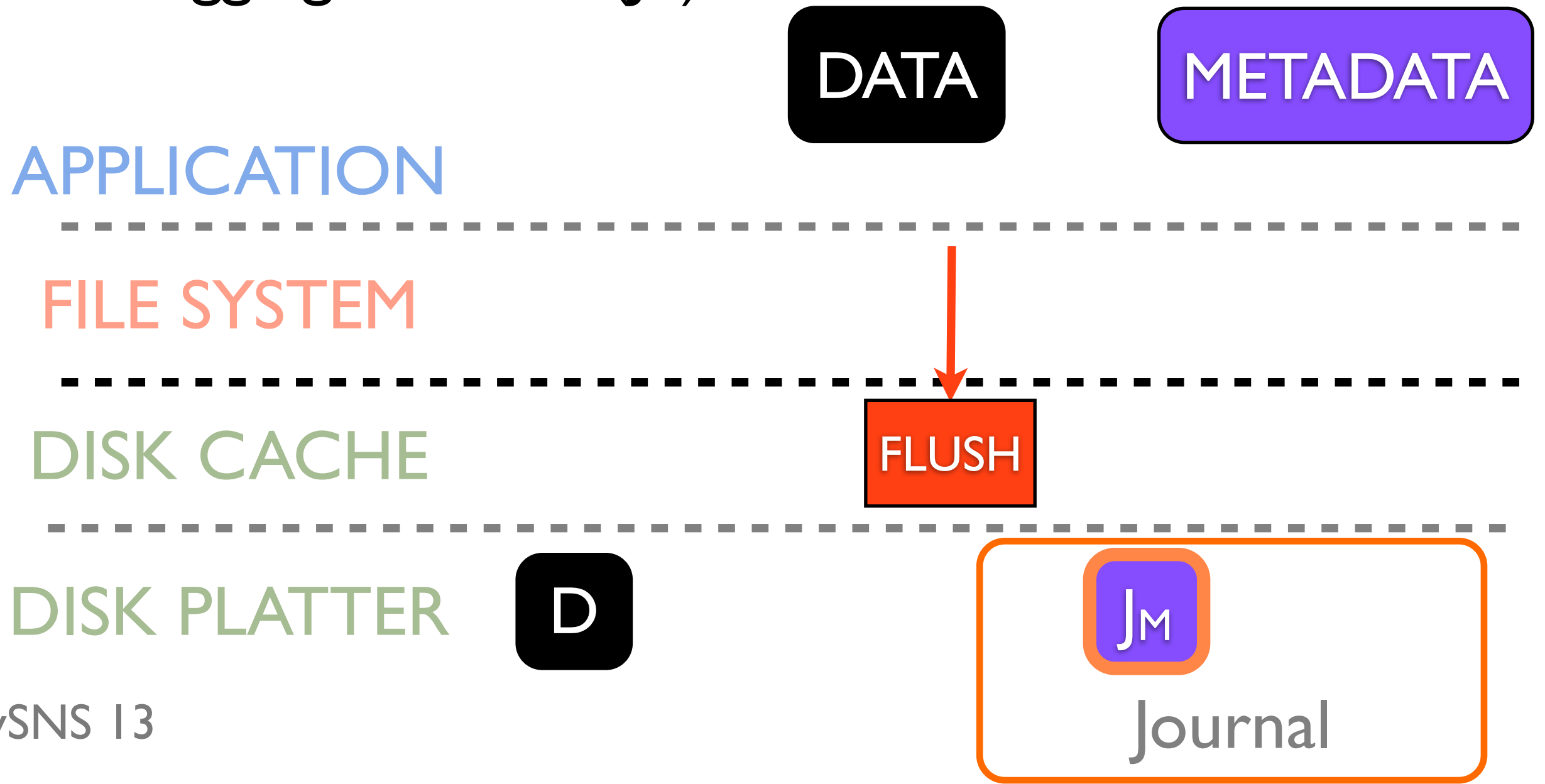
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Journaling with Flushes

Journaling protocol

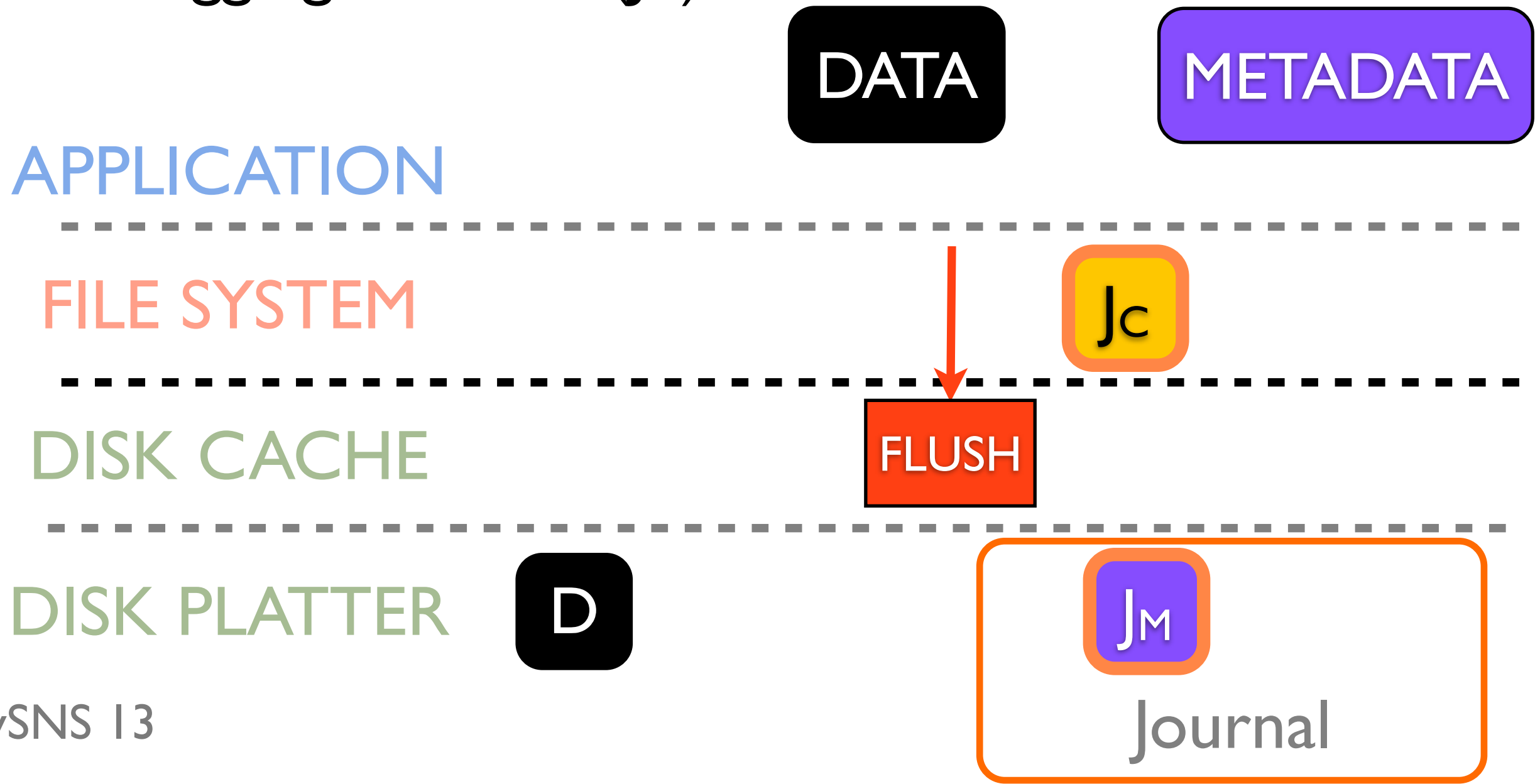
- Data write (D)
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- Logging Commit (Jc)



Journaling with Flushes

Journaling protocol

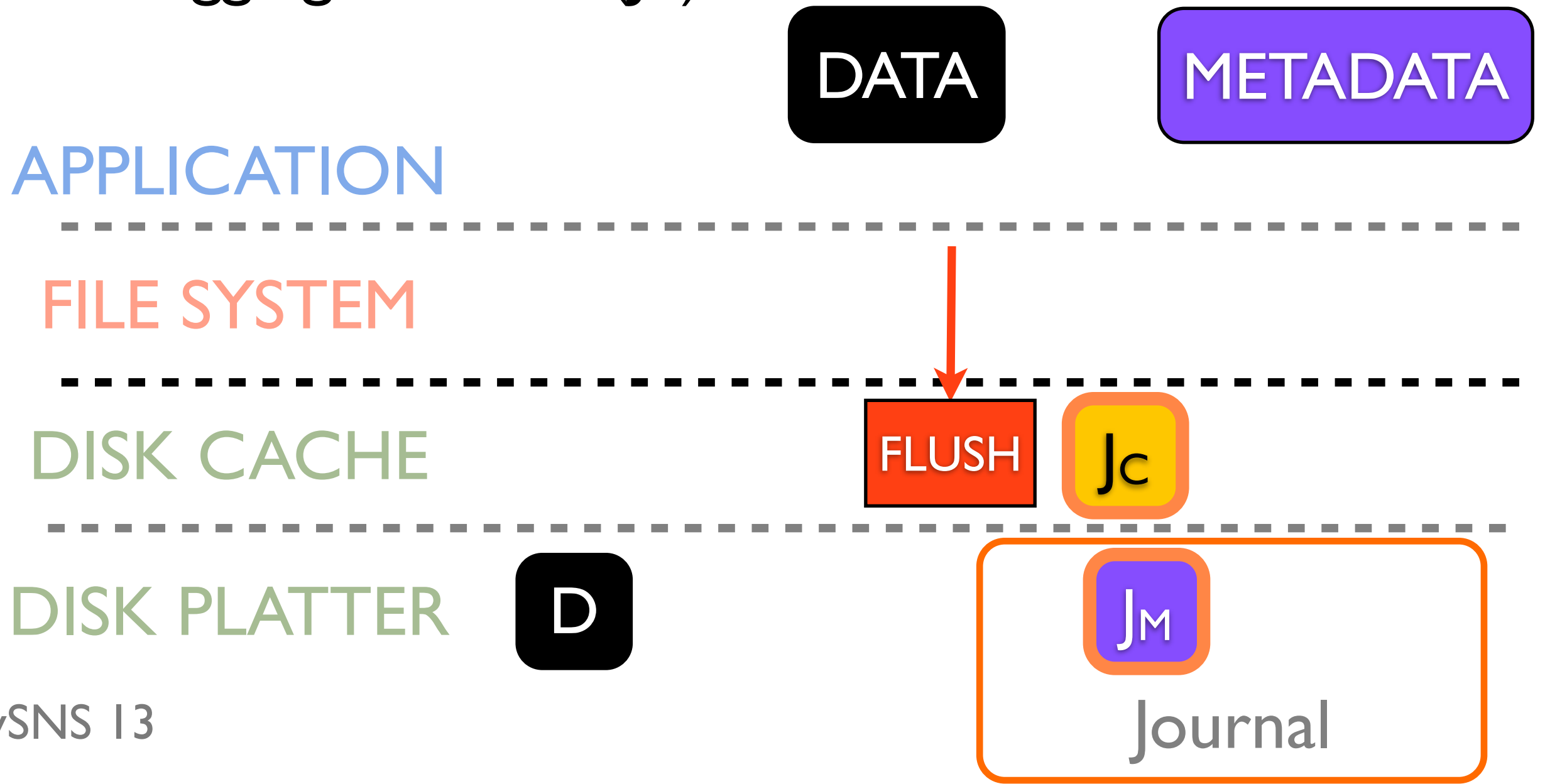
- Data write (D)
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- Logging Commit (J_c)



Journaling with Flushes

Journaling protocol

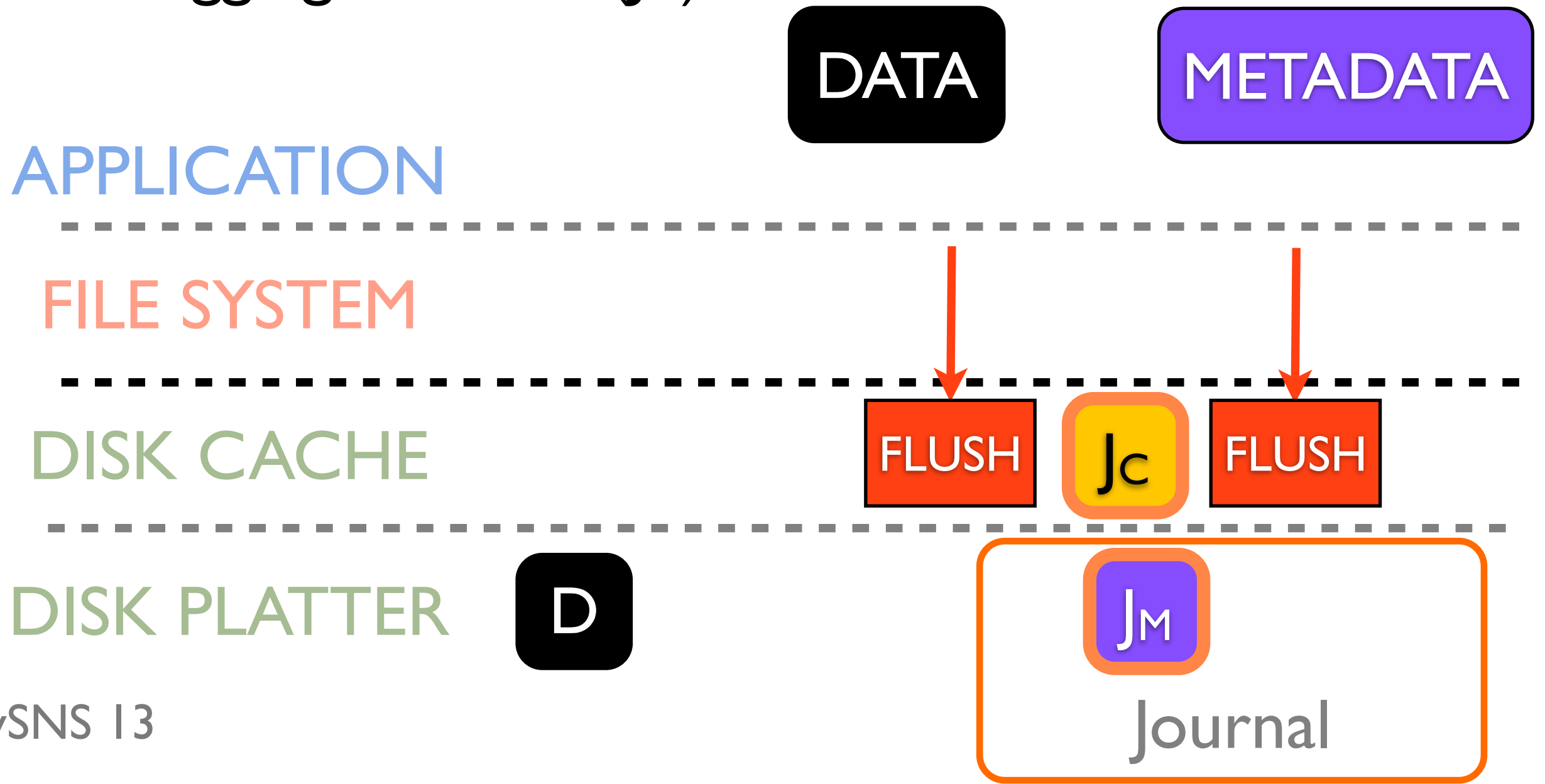
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Journaling with Flushes

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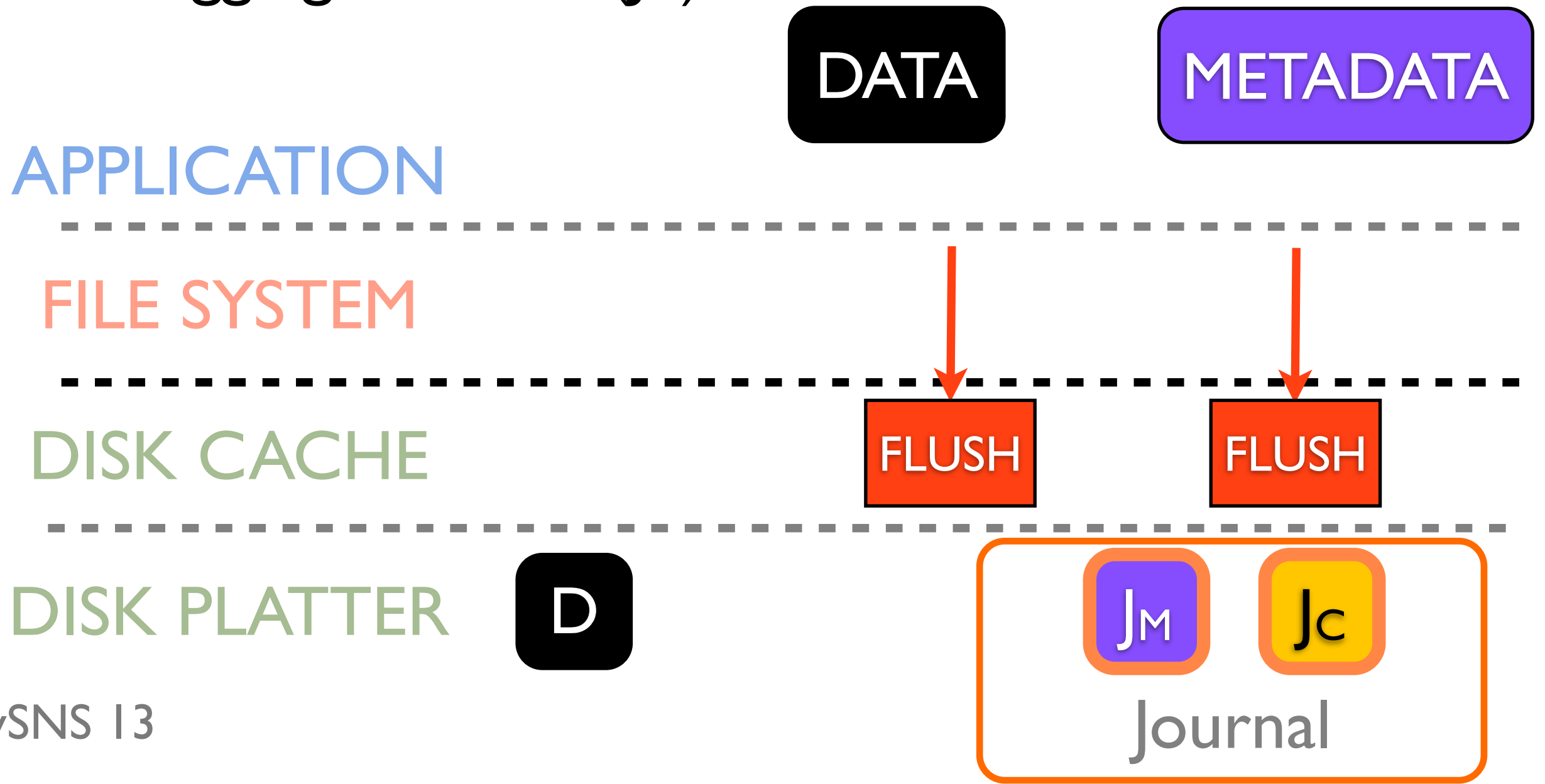
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Journaling with Flushes

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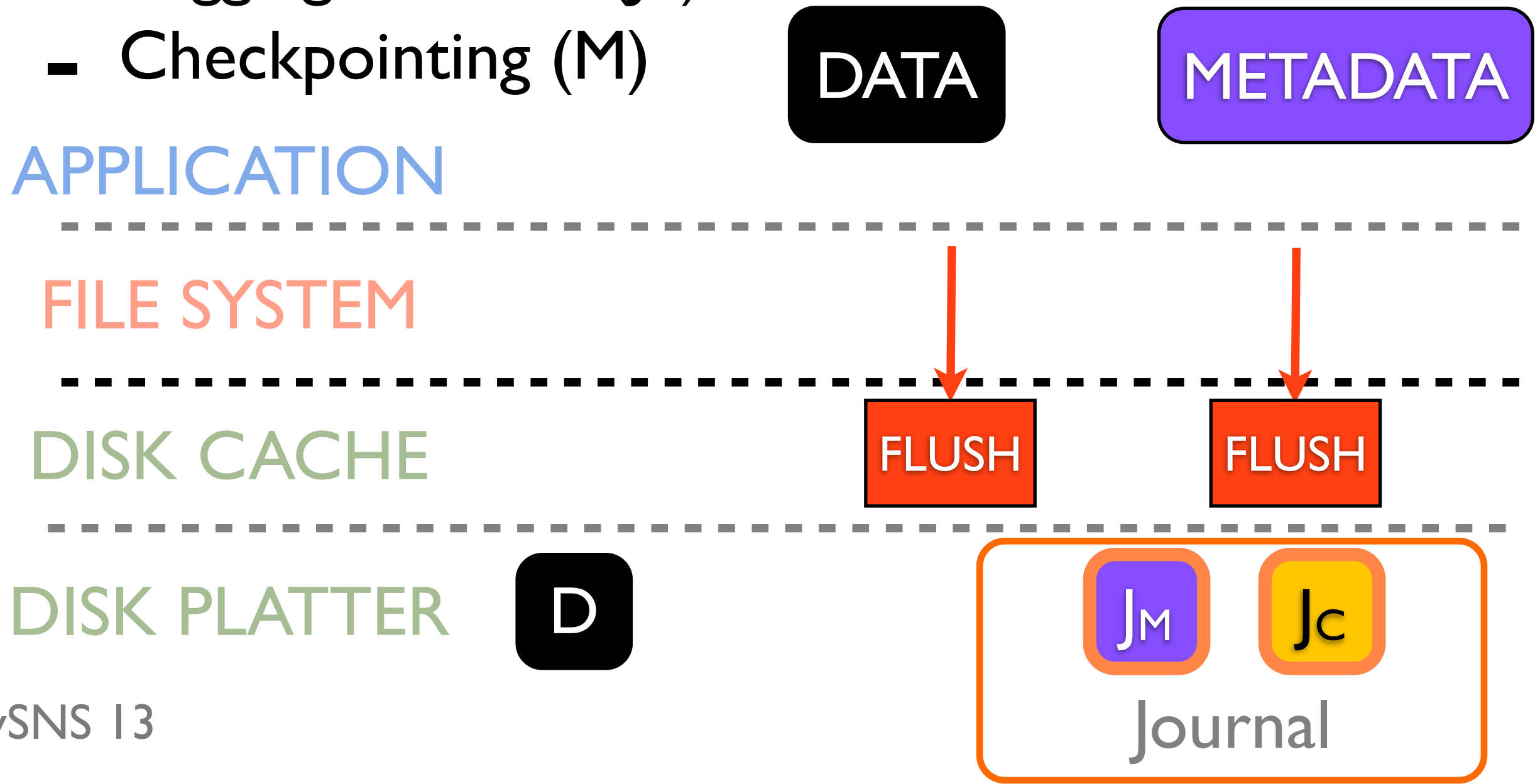
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Journaling with Flushes

Journaling protocol

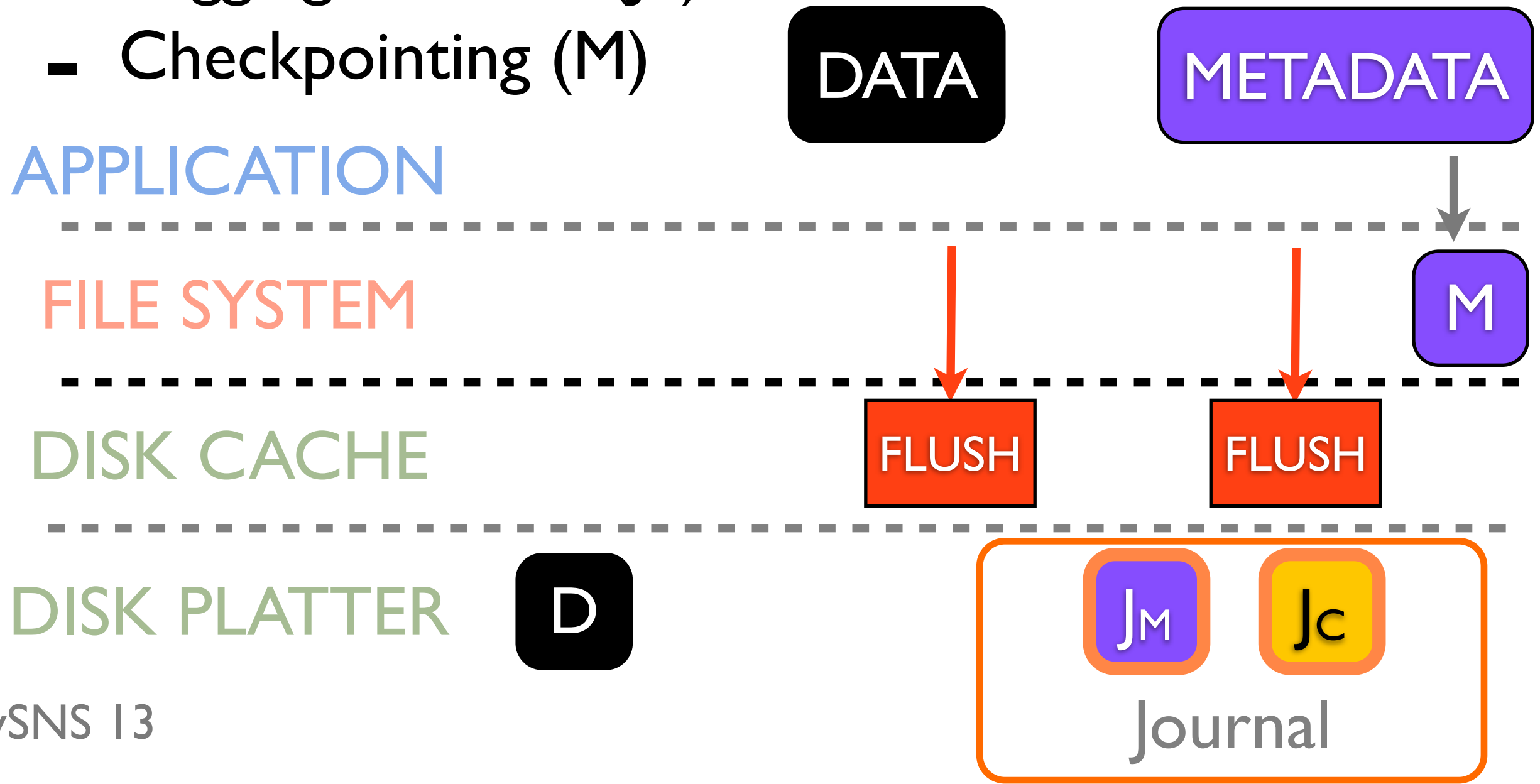
- Data write (D)
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Journaling with Flushes

Journaling protocol

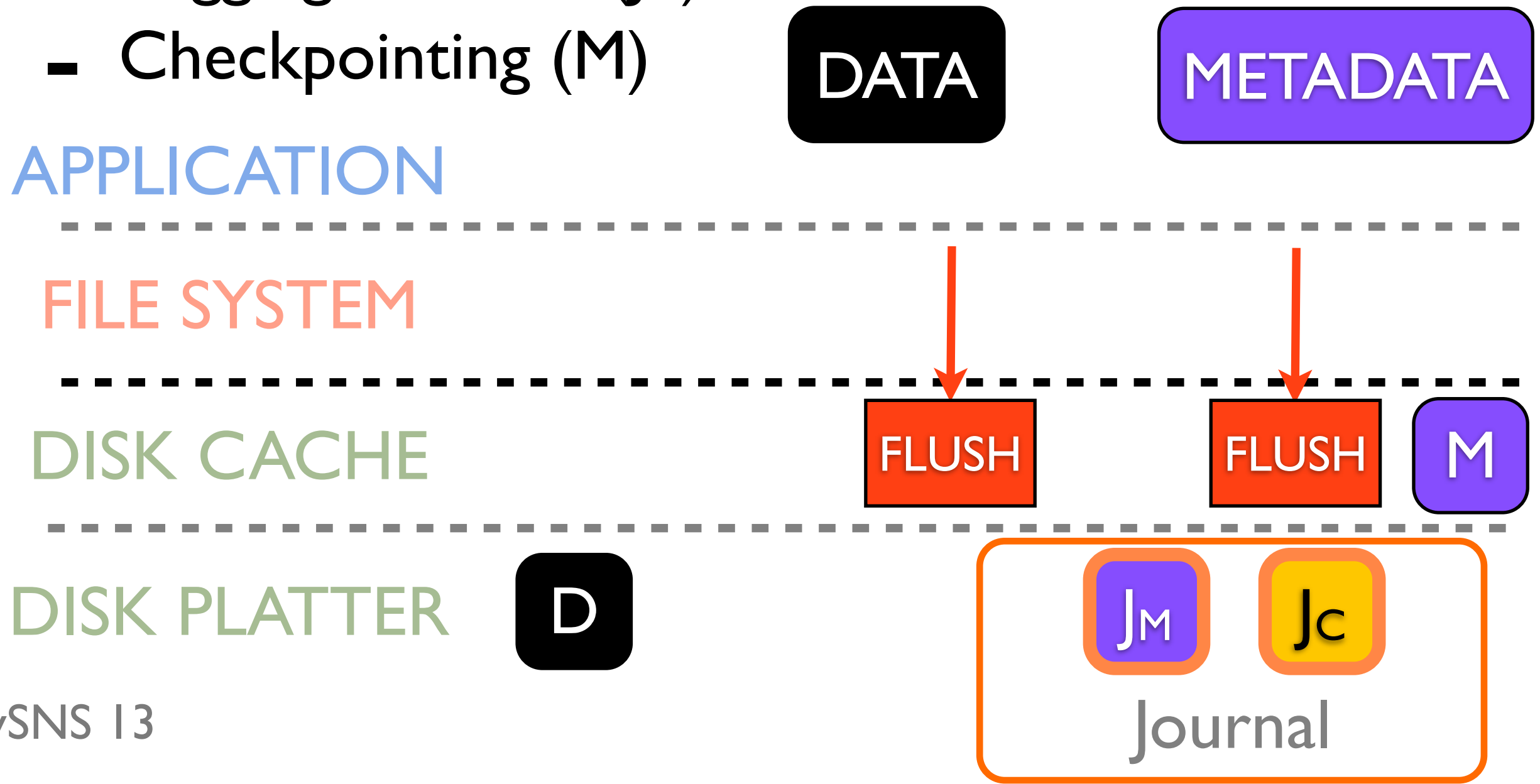
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Journaling with Flushes

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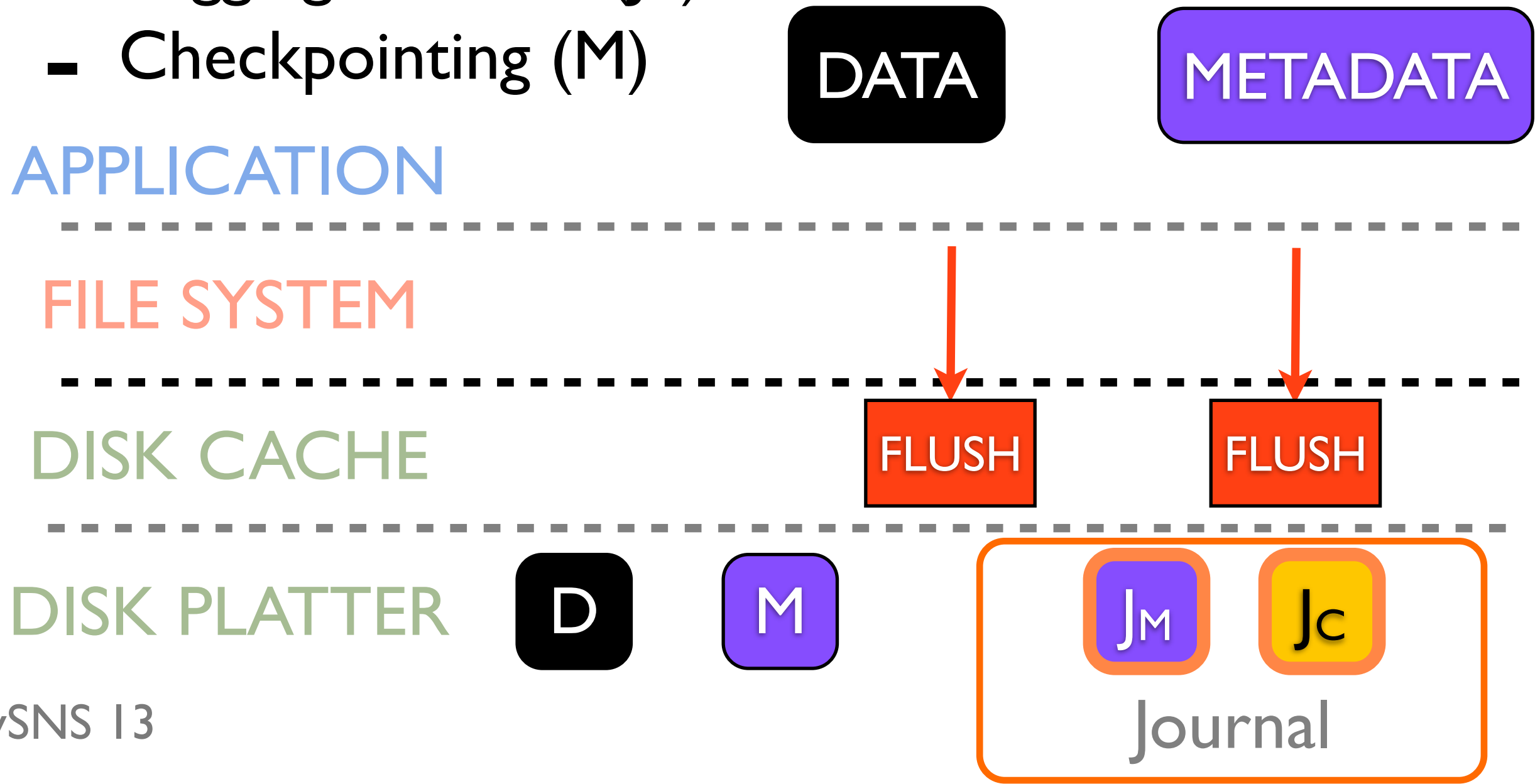
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Journaling with Flushes

Journaling protocol

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- Logging Metadata (JM)
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Optimistic File System

Results

Conclusion

Journaling without Ordering

Practitioners **turn off** flushes due to performance degradation

- E.g., ext3 by default did not enable flushes for many years

Observe crashes do not cause inconsistency for **some** workloads

We term this **probabilistic** crash consistency

- Studied in detail

Probabilistic Crash Consistency

p-inconsistency for different workloads

- Read-heavy workloads have low p-inconsistency
- Database workloads have high p-inconsistency

See paper for detailed study

- Factors that affect p-inconsistency

Turning off flushing provides performance, but **does not** ensure consistency

Additional techniques required to obtain **both** performance and consistency

Outline

Introduction

Ordering and Durability in Journaling

Optimistic File System

- **Overview**
- Handling Re-Ordering
- New File-system Primitives

Results

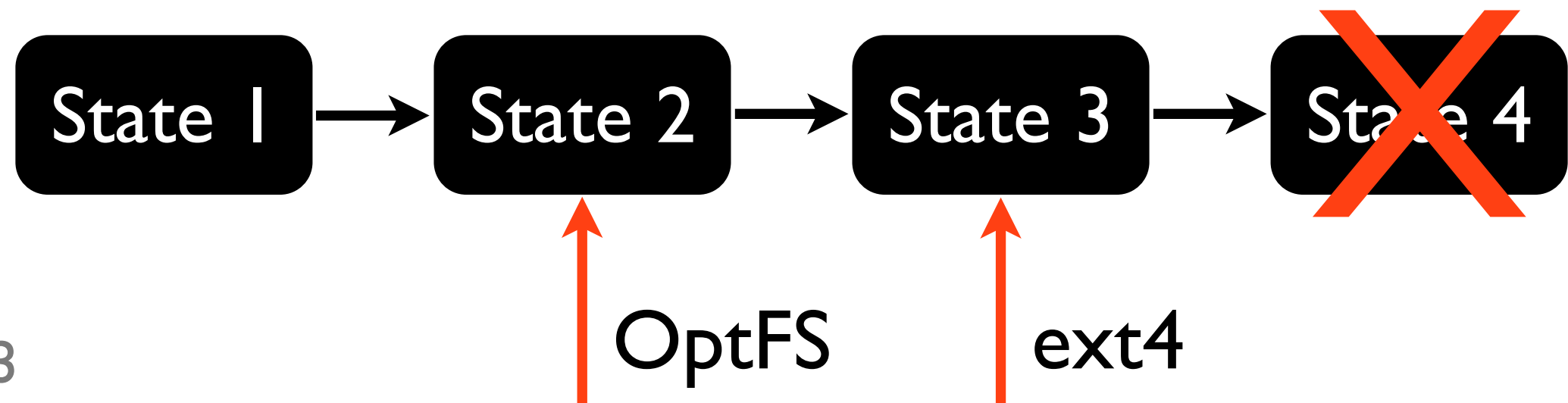
Conclusion

Optimistic File System

Achieves **both** performance and consistency by trading on **new** axis

Freshness indicates how up-to-date state is after a crash

OptFS provides strong consistency while **trading freshness** for increased performance



Optimistic File System

Eliminates flushes in the common case

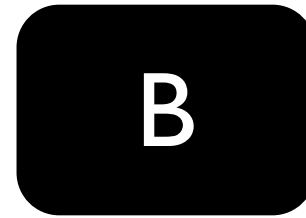
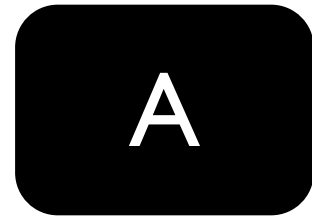
Blocks may be re-ordered without flushes

Optimistic Crash Consistency handles re-orderings with different techniques

- Some re-orderings are **detected** after crash
- Some re-orderings are **prevented** from occurring

Modified Disk Interface

Asynchronous Durability Notifications (ADN)
signal when block is made durable



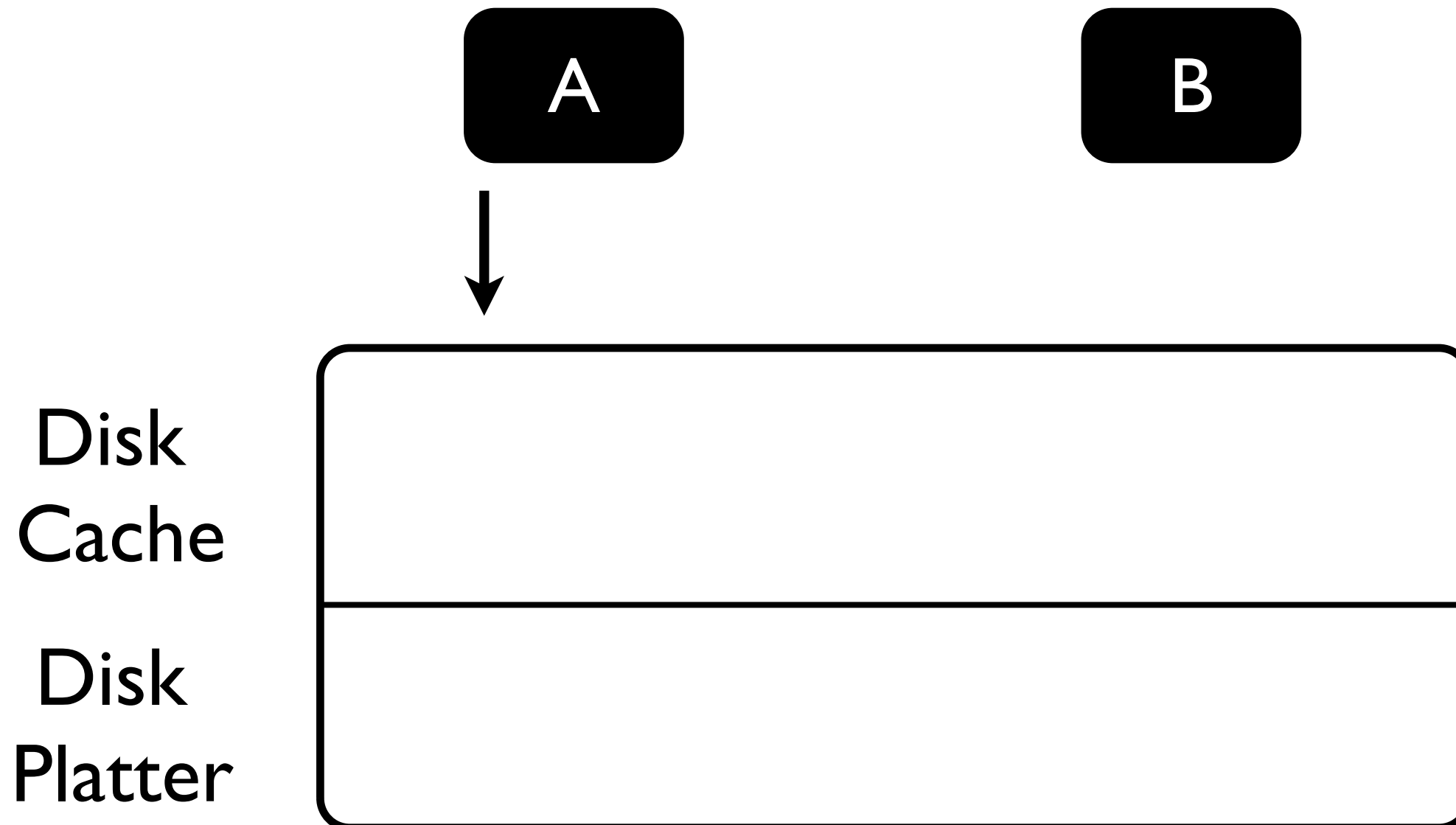
Disk
Cache

Disk
Platter



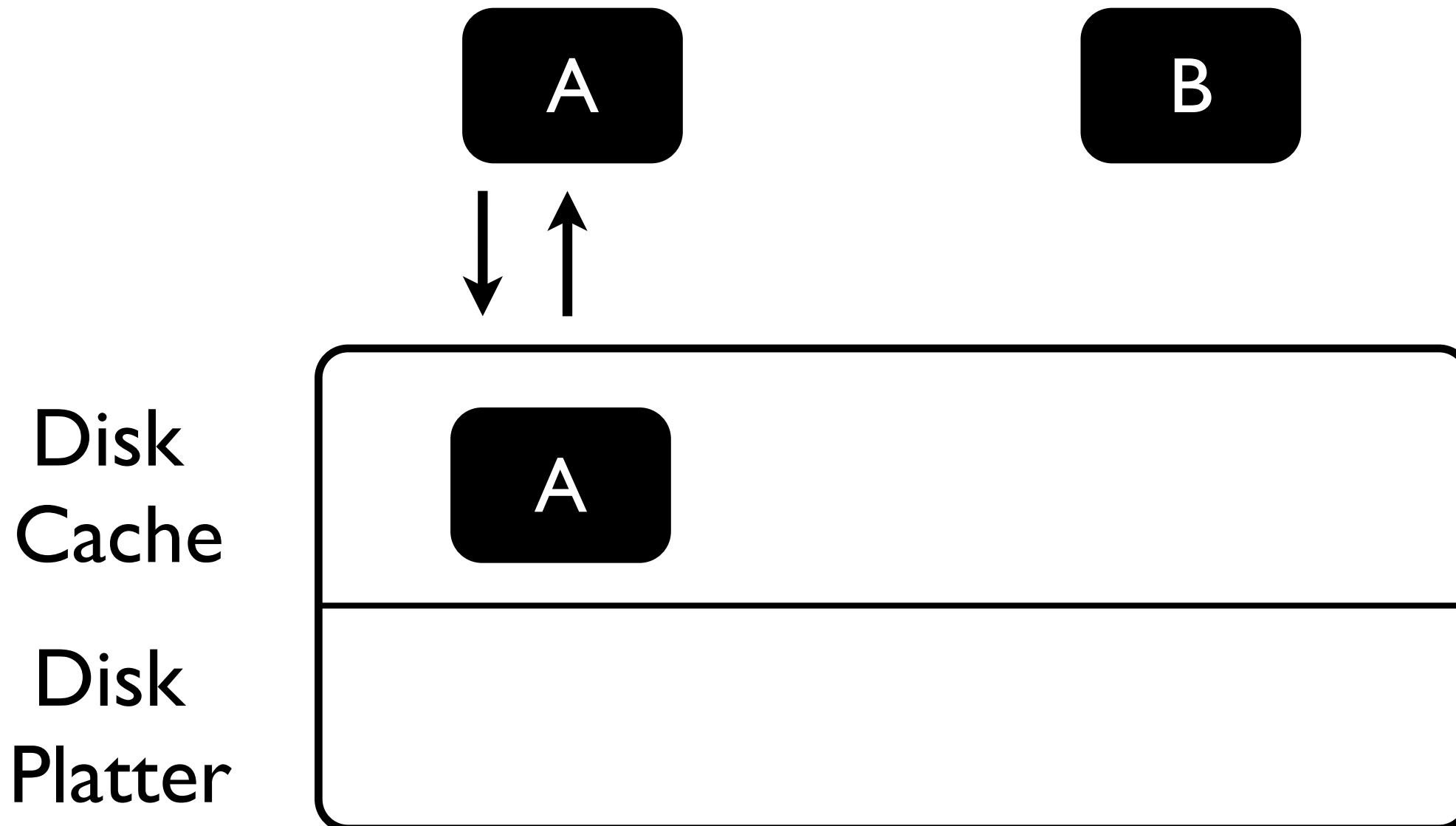
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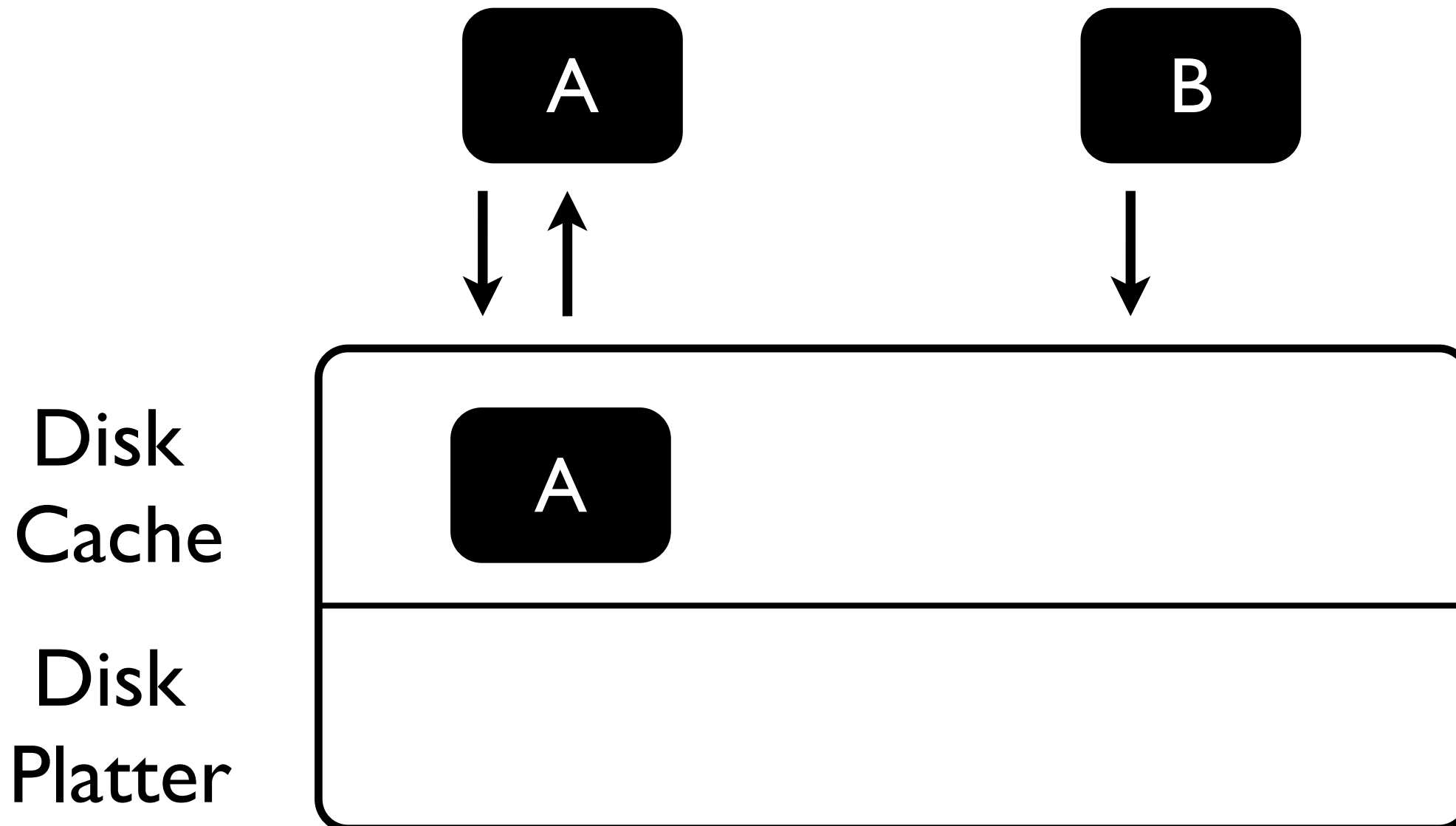
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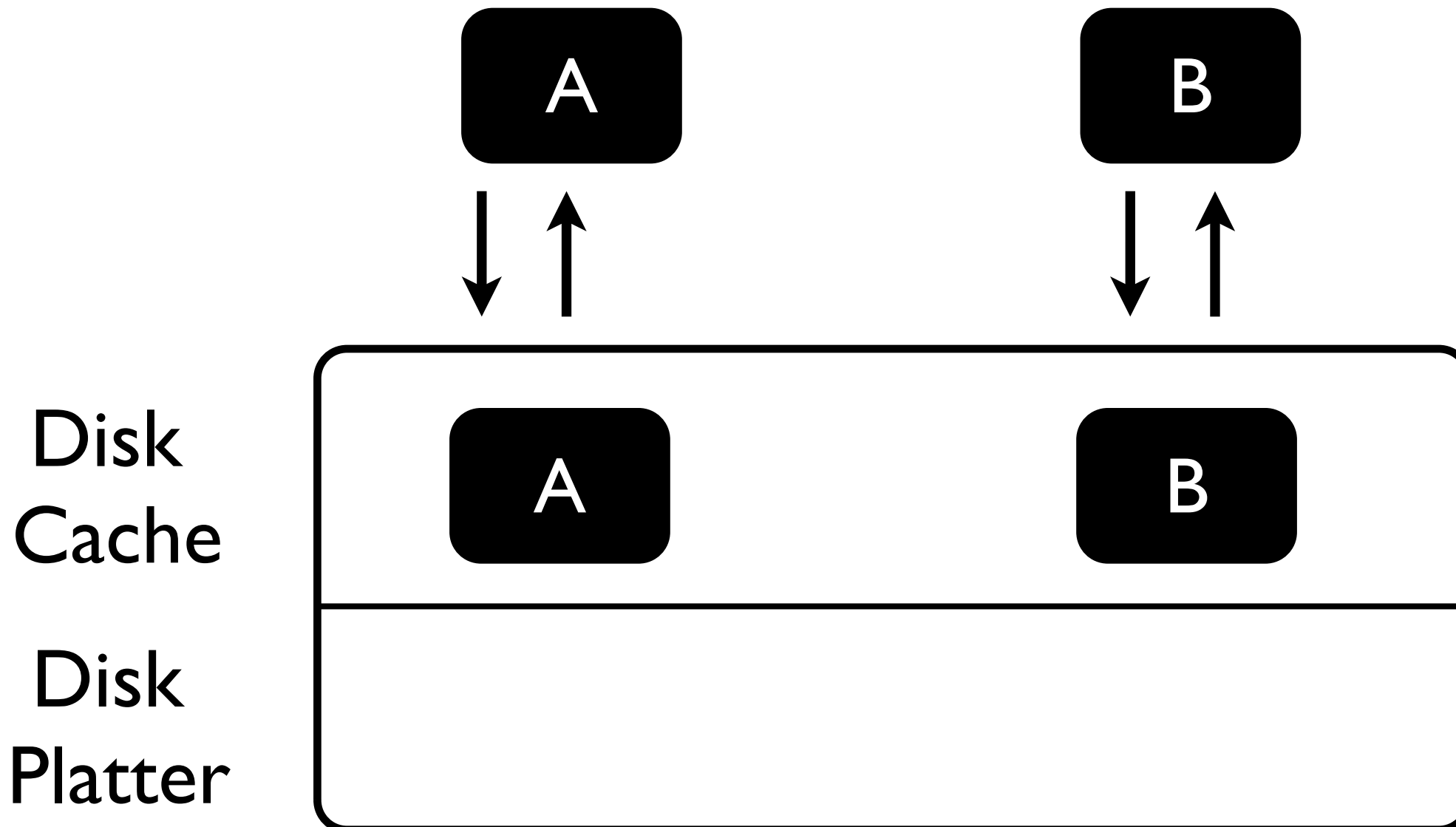
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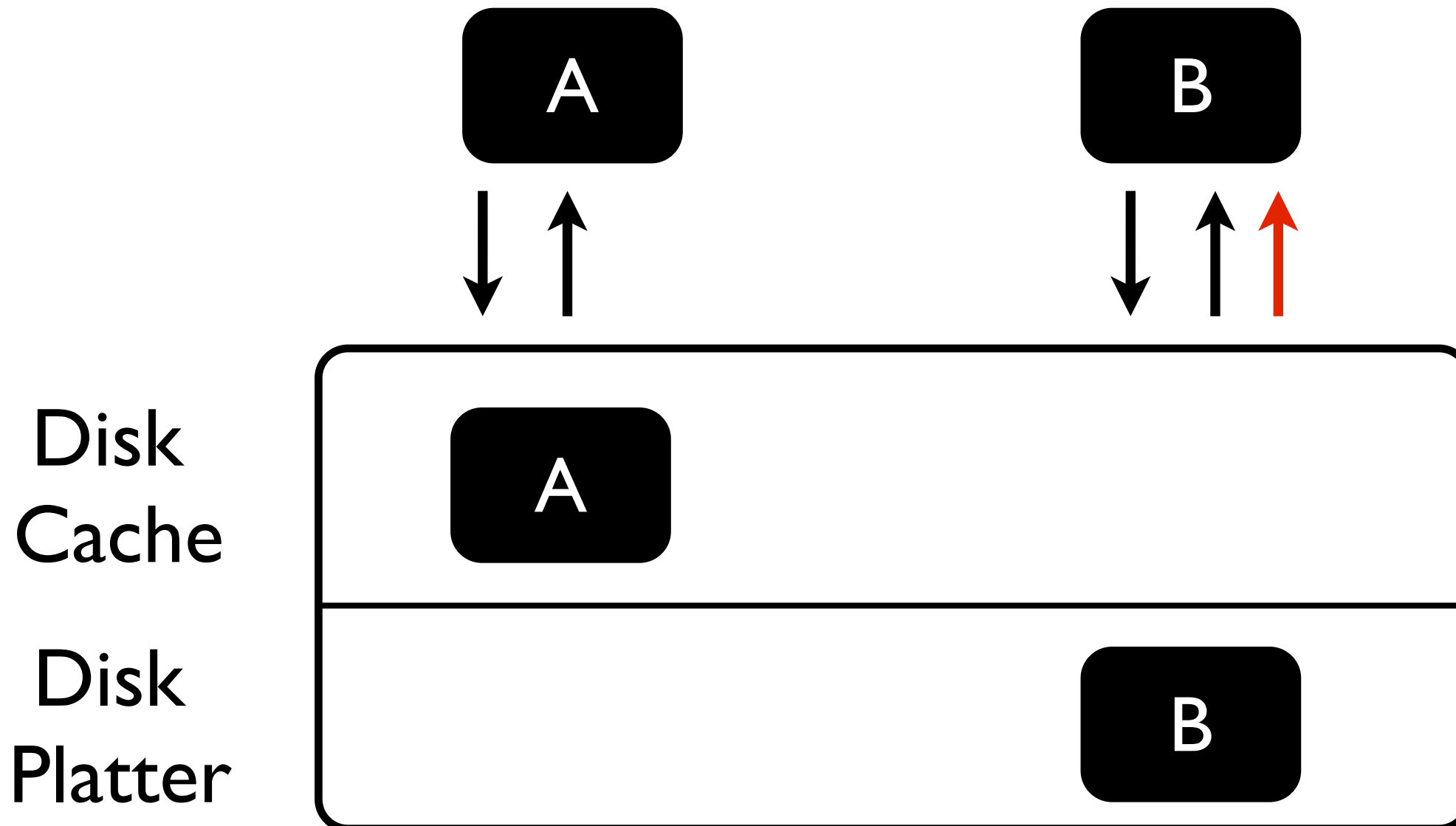
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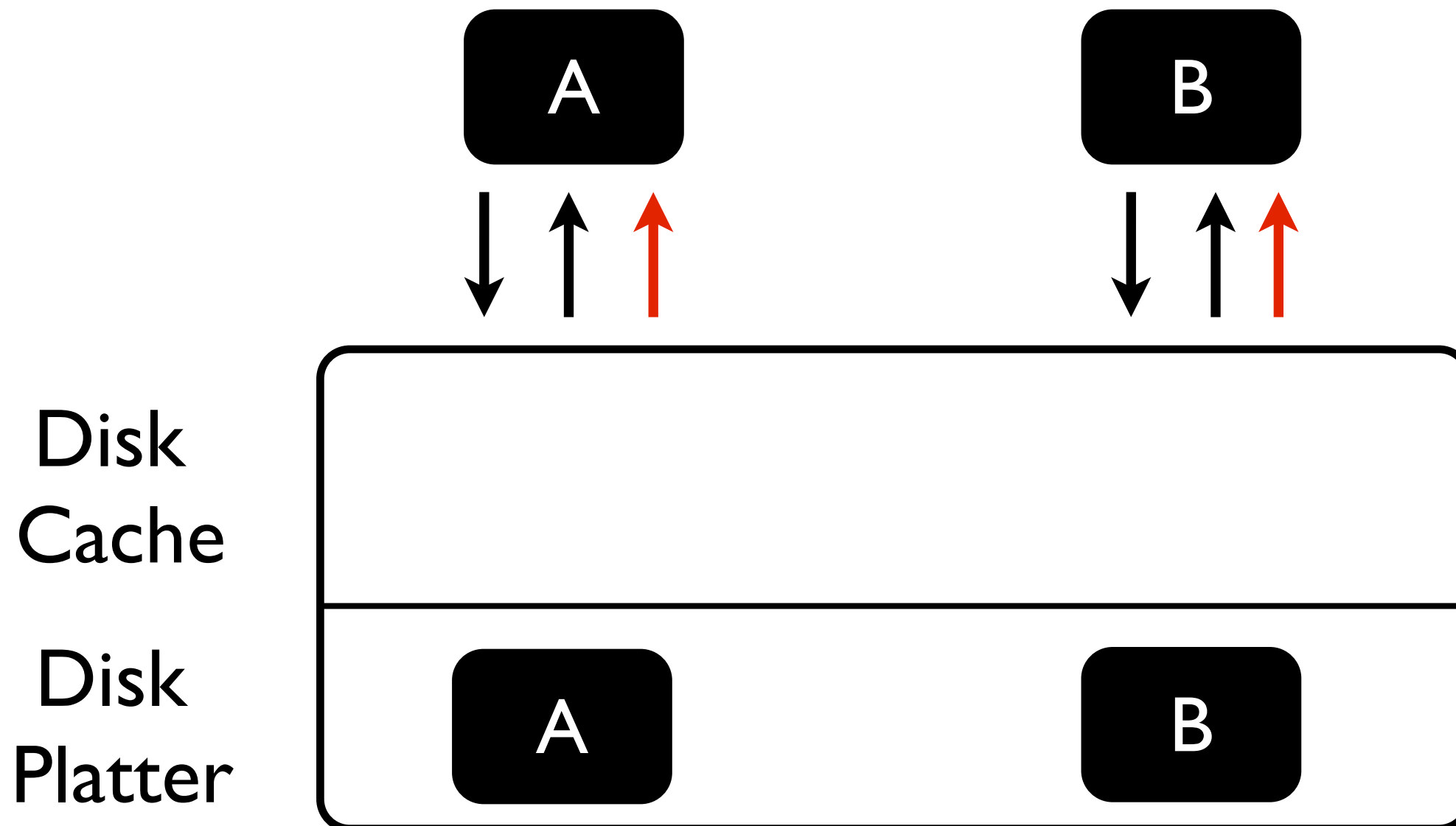
Modified Disk Interface

Asynchronous Durability Notifications (ADN)
signal when block is made durable



Modified Disk Interface

Asynchronous Durability Notifications (ADN)
signal when block is made durable



Modified Disk Interface

ADNs increase disk **freedom**

- Blocks can be destaged in **any order**
- Blocks can be destaged at **any time**
- Only requirement is to inform upper layer

OptFS uses ADNs to control what blocks are **dirty** at the **same time** in disk cache

- Re-ordering can only happen among these blocks

Outline

Introduction

Ordering and Durability in Journaling

Optimistic File System

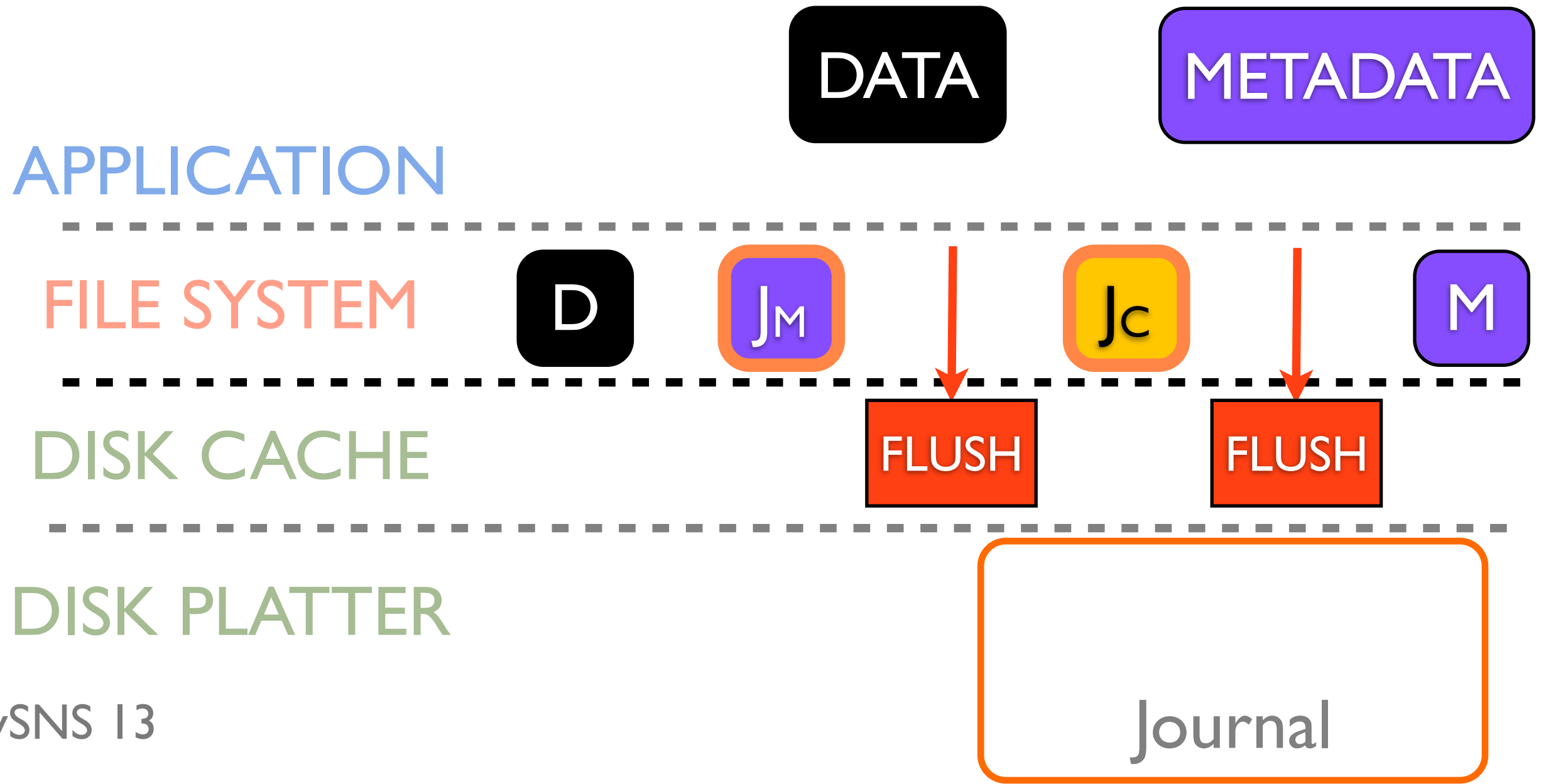
- Overview
- **Handling Re-Ordering**
- New File-system Primitives

Results

Conclusion

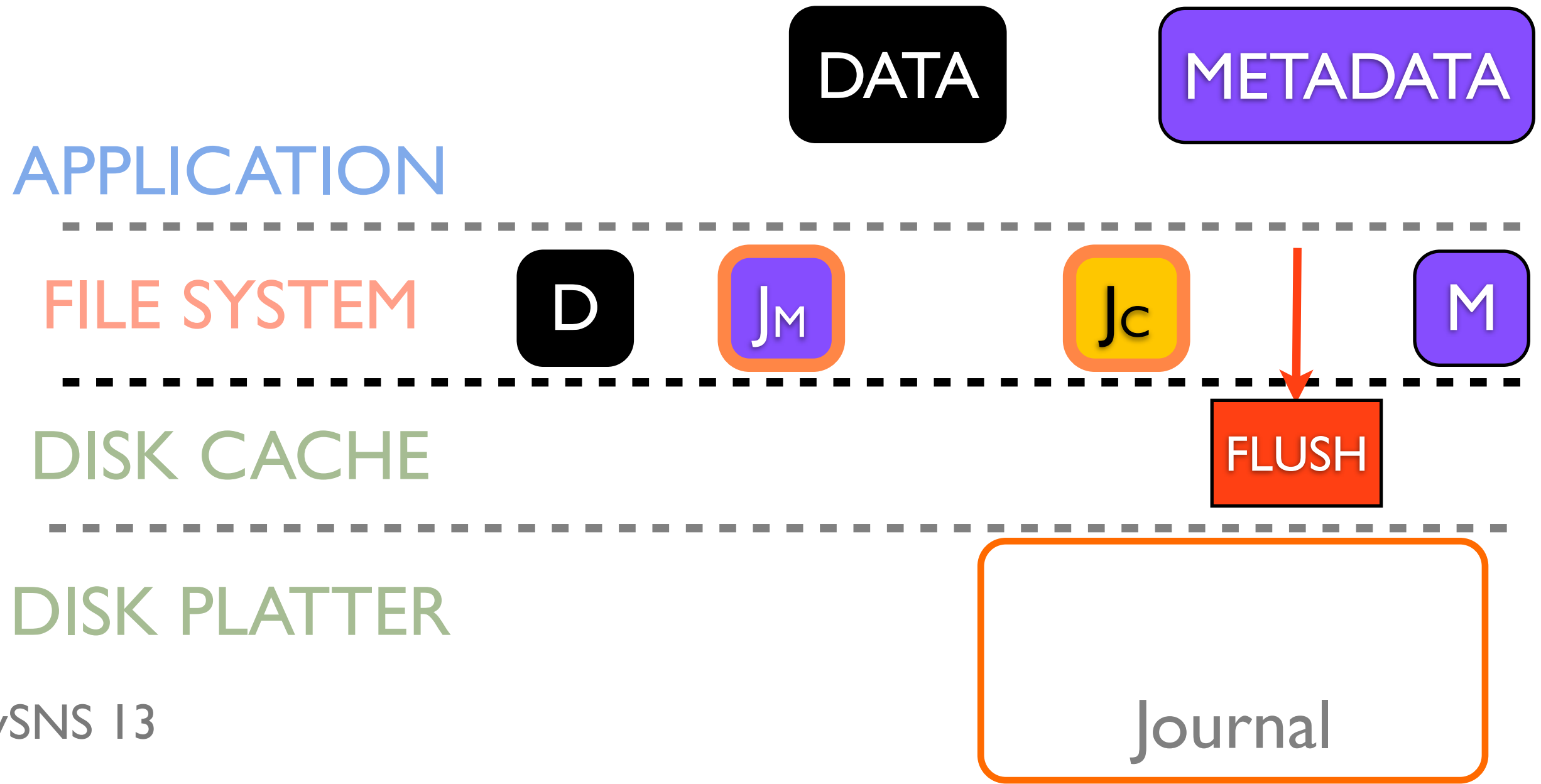
Optimistic Journaling

Checksums and Delayed Writes handle reordering from removing flushes



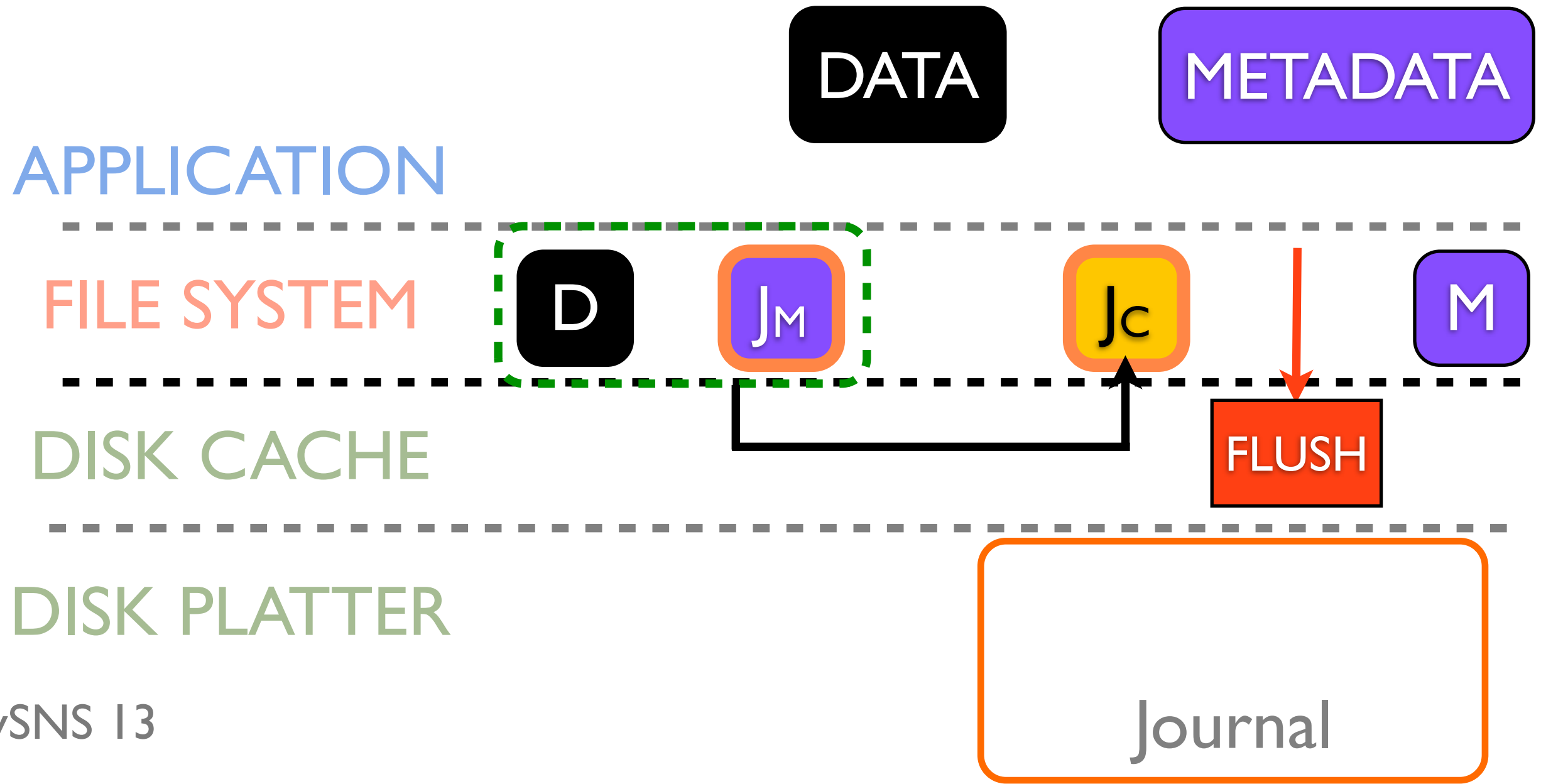
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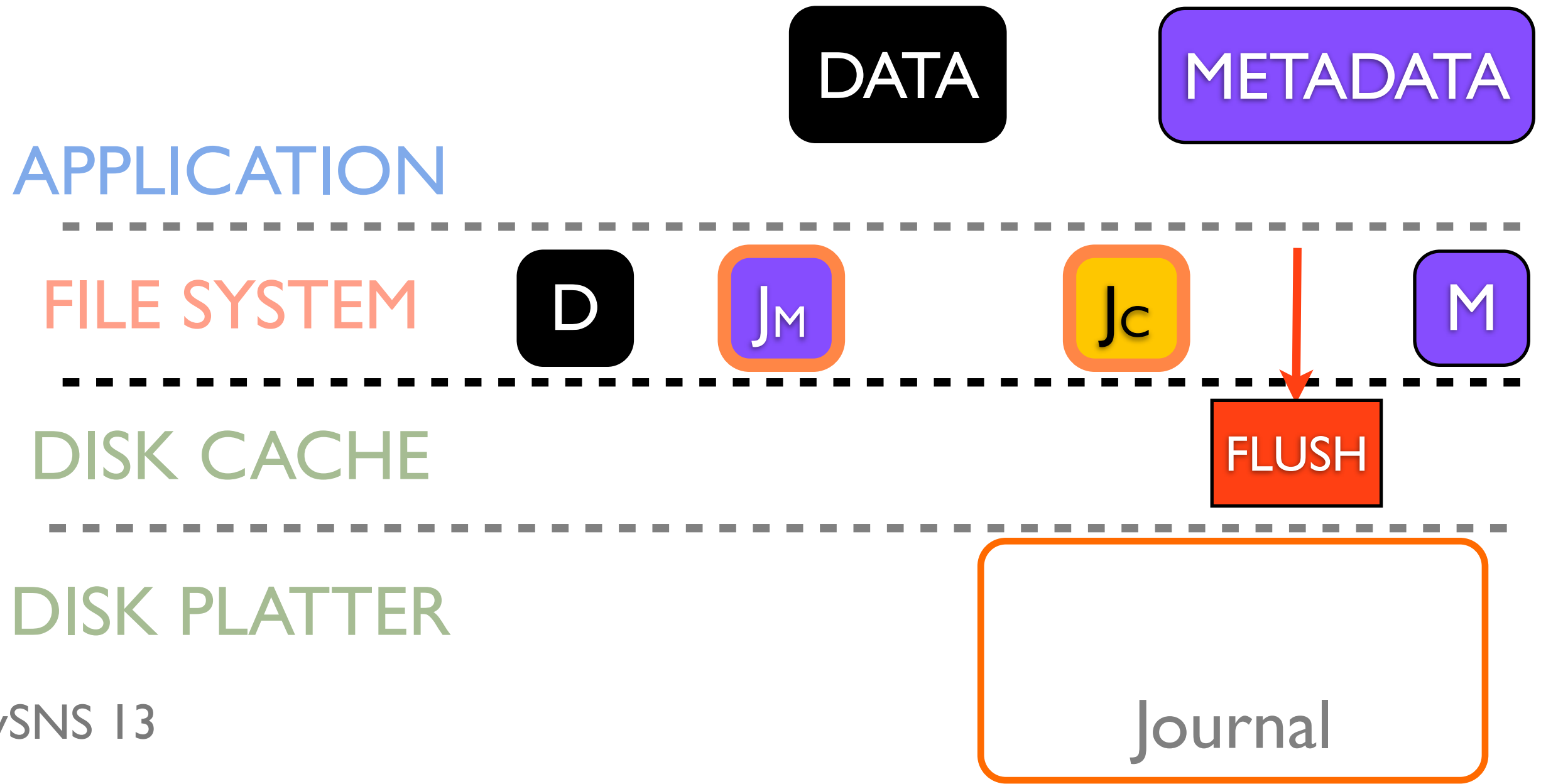
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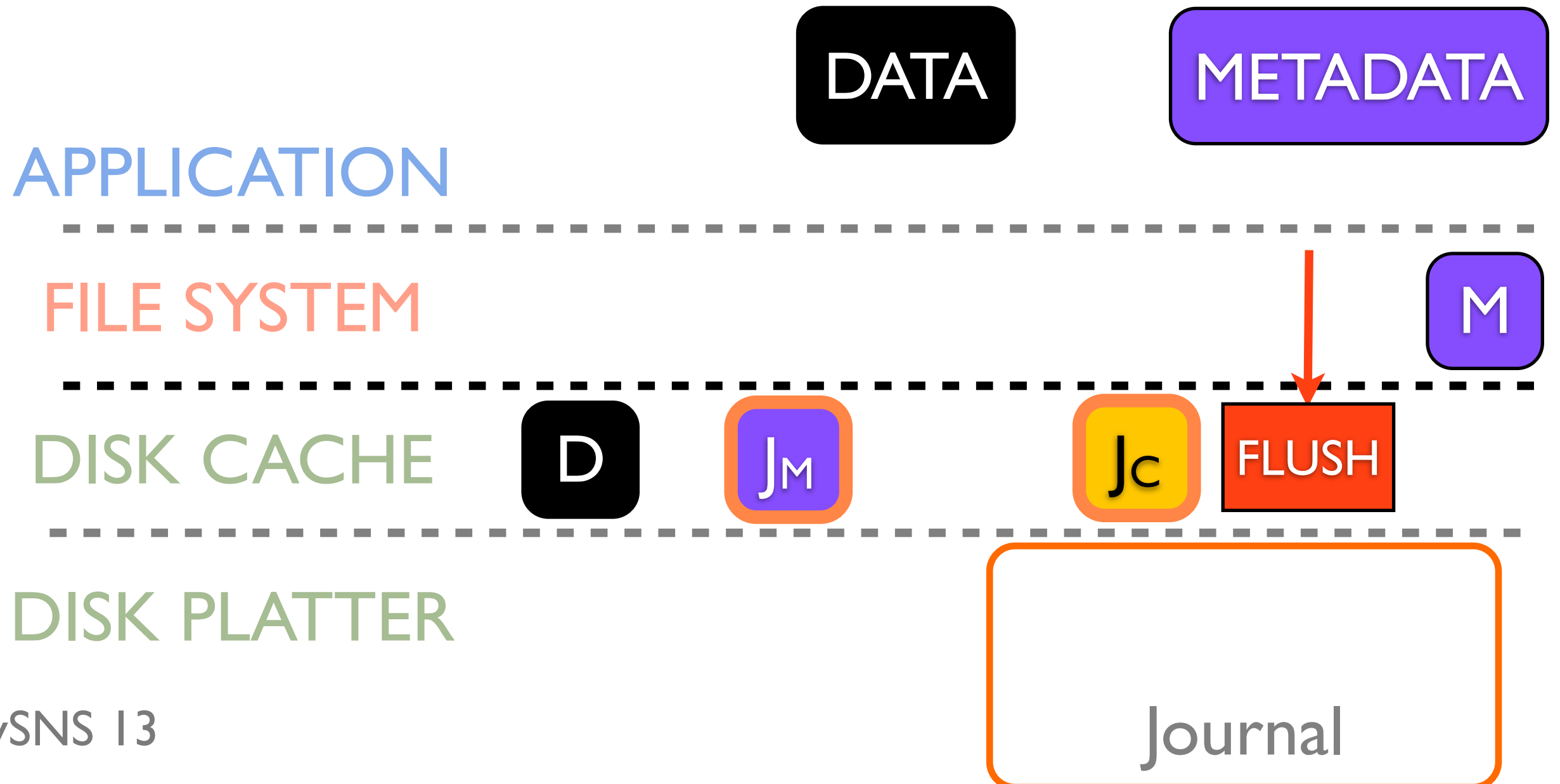
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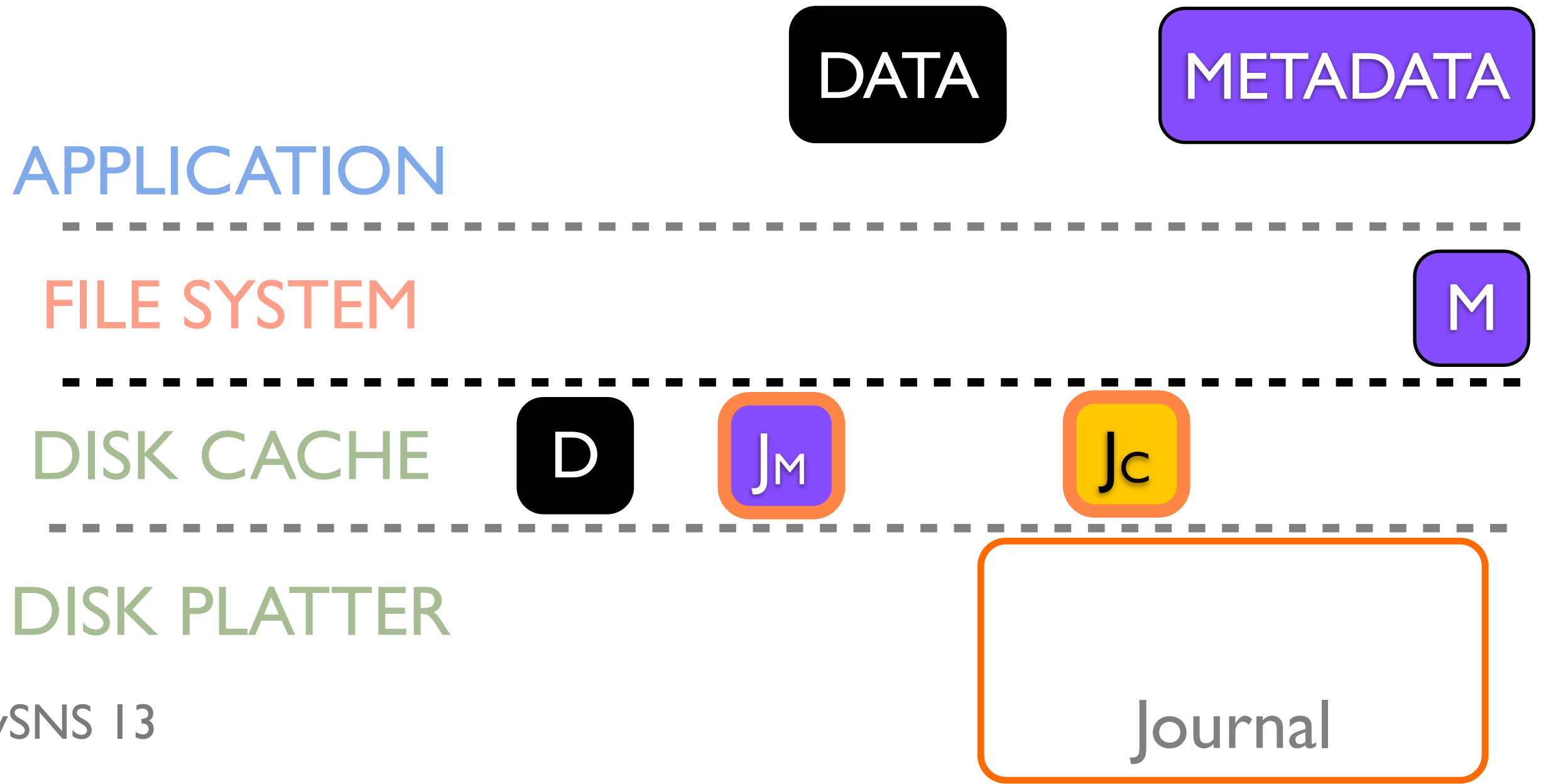
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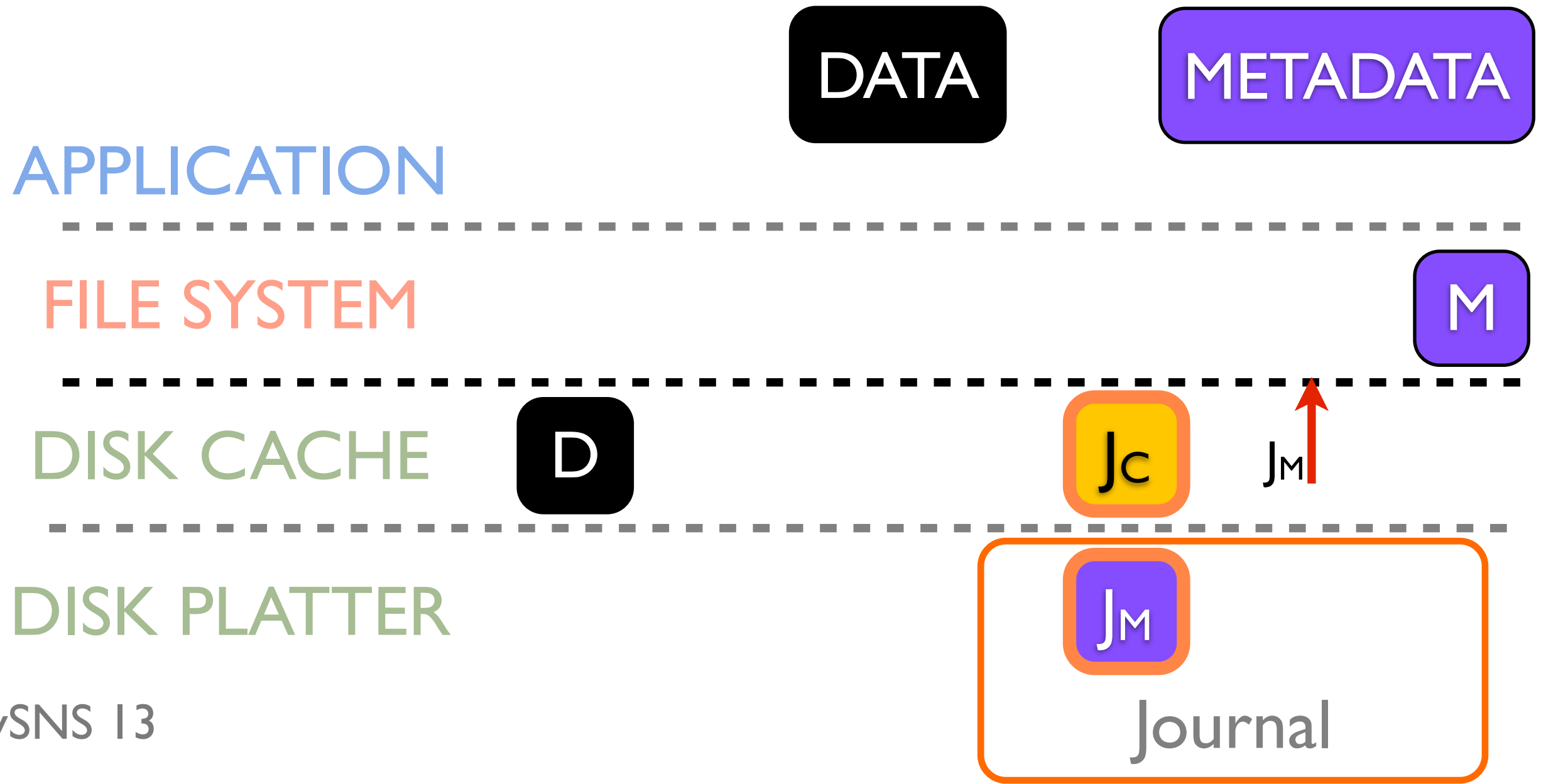
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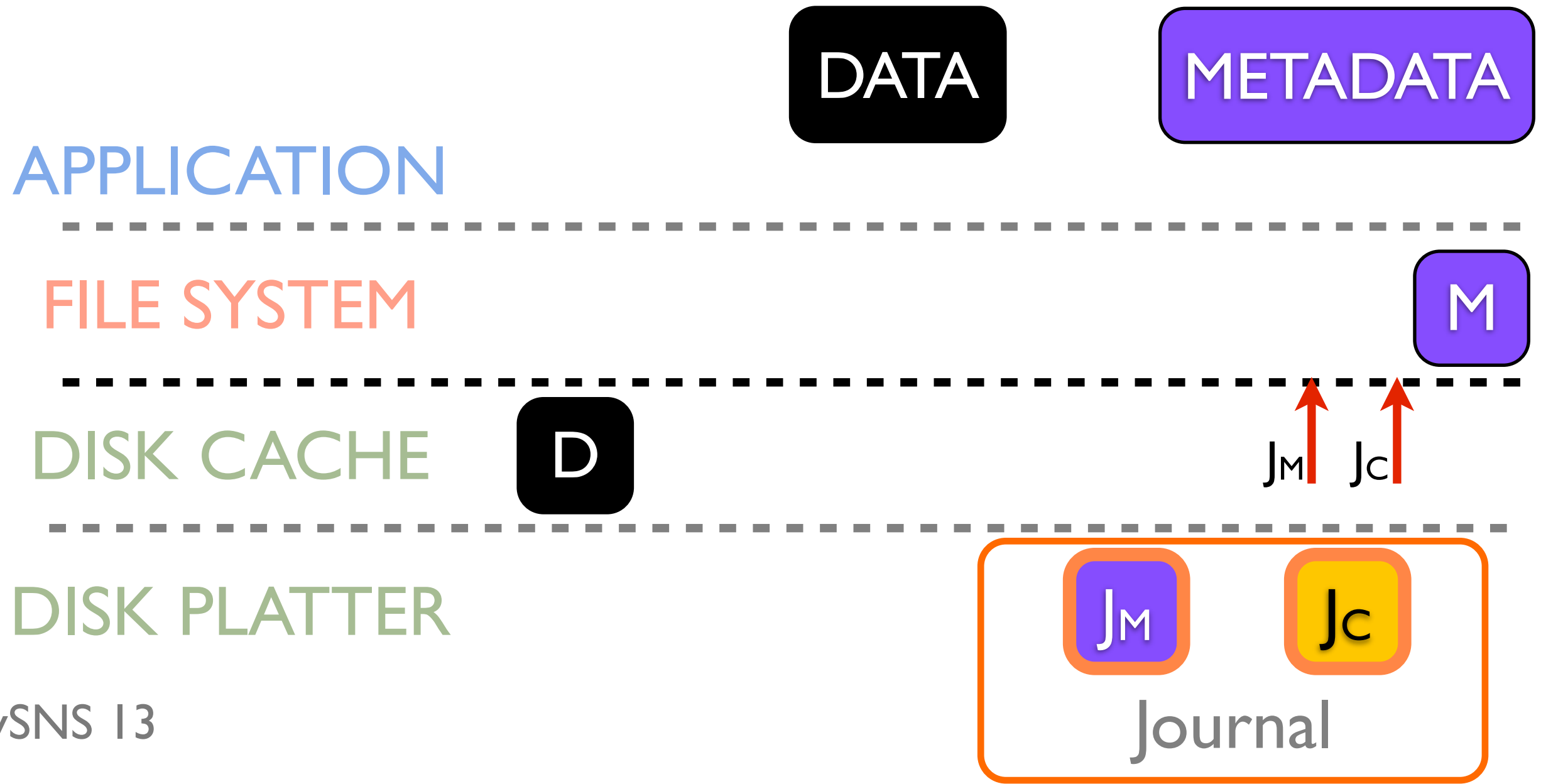
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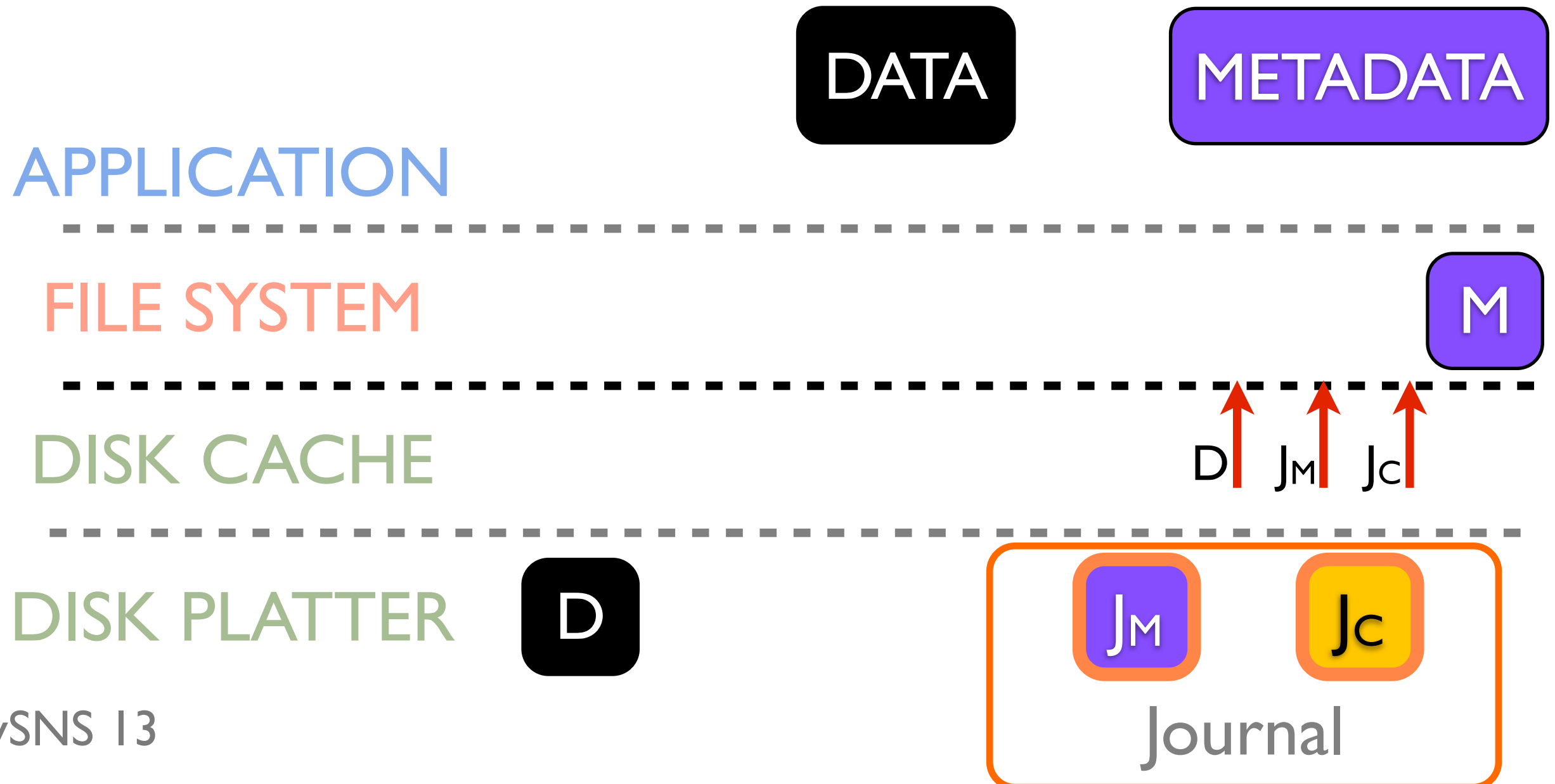
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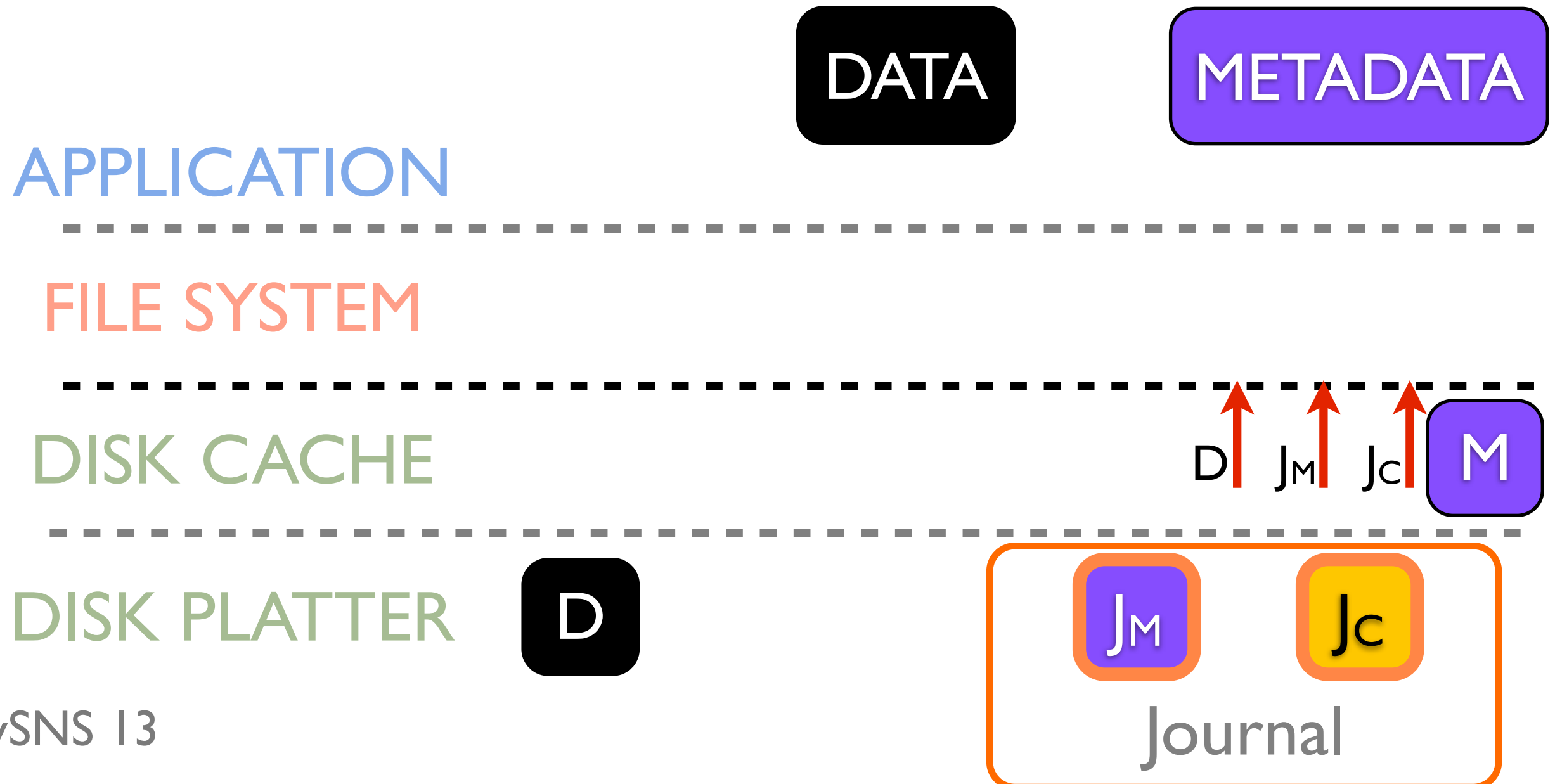
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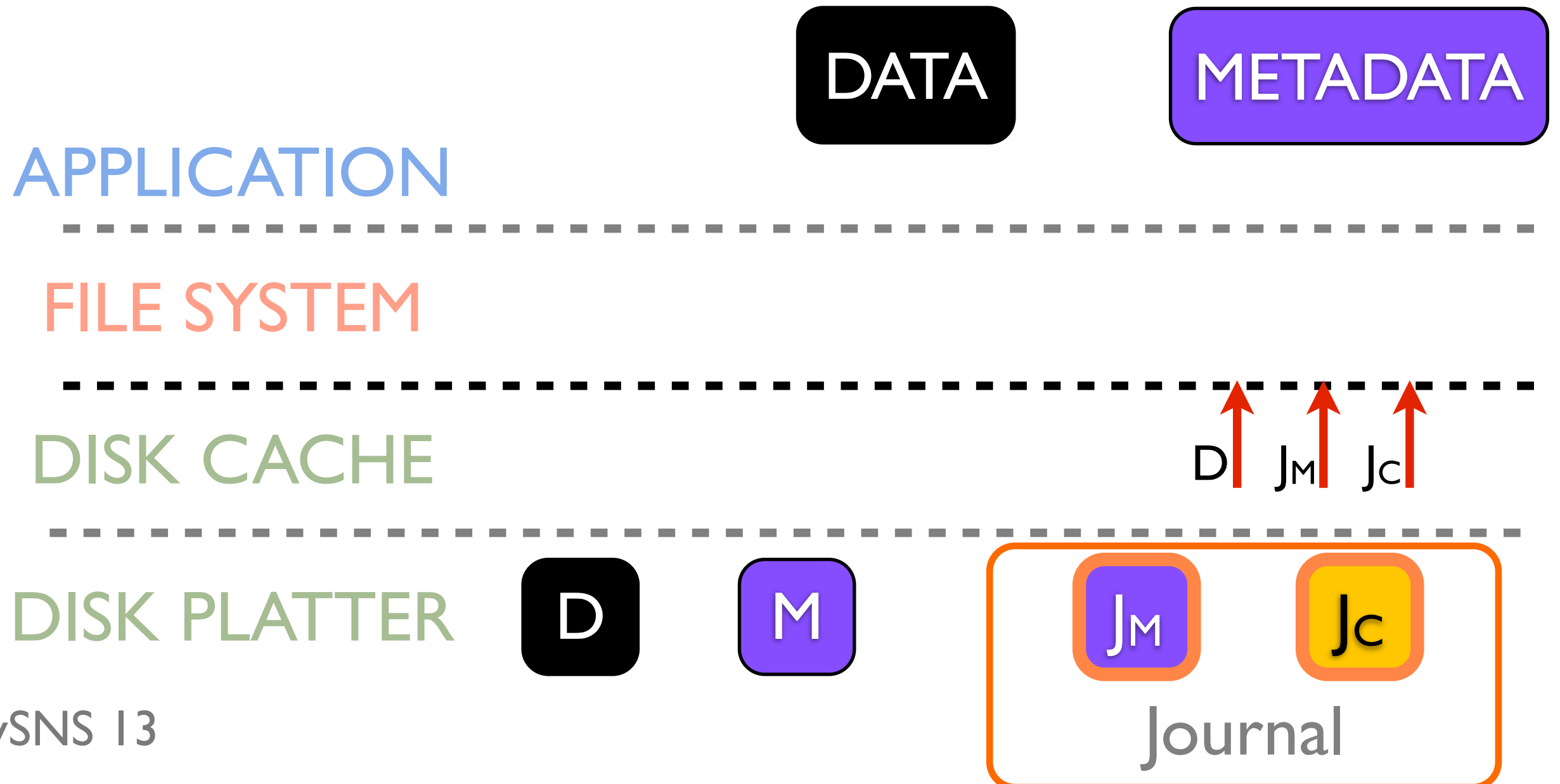
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Optimistic Techniques

Other Techniques

- In-order journal recovery and release
- Reuse after notification
- Selective data journaling

See paper for more details

Outline

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File-system Primitives

`fsync()` provides ordering and durability

OptFS splits `fsync()`

- `osync()` for only **ordering** and high performance
- `dsync()` for durability

Primitives can increase performance

- Ex: SQLite

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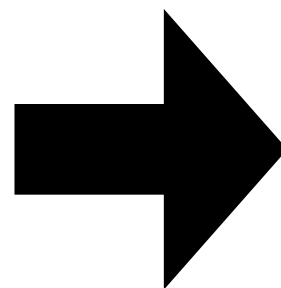
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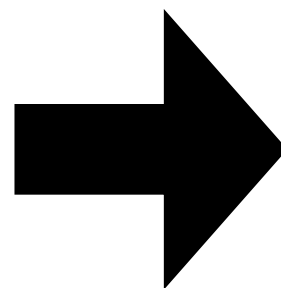
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Evaluation

Does OptFS preserve **file-system consistency** after crashes?

- OptFS consistent after **400** random crashes

How does OptFS **perform**?

- OptFS **4-10x** better than ext4 with flushes

Can meaningful **application-level consistency** be built on top of OptFS?

- SQLite provides ACI semantics at 10x performance

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Summary

Problem: providing **both** performance and consistency

Solution: **decoupling** ordering and durability in OptFS

OptFS maintains consistency while trading freshness for increased performance

osync() provides a cheap primitive to order application writes

Conclusion

Storage-stack layers are increasing

- 18 layers between application and storage [Thereska13]
- Interfaces that provide **freedom** to each layer are the way forward

First impulse: trade consistency for performance

- Trade-off not required in distributed systems [Escriva12]
- By trading freshness, we can obtain both consistency and high performance

Thank You

Source code

<http://research.cs.wisc.edu/adsl/Software/optfs/>

<http://github.com/vijay03/optfs>

Questions?



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