

1 S B A G2

2 ABCD

3 A. min B. left column, 2nd row

4

A. the 2nd node has candidate labels reduced to {-}

B. all nodes get their correct label

5

A. x converge to 0 the minimum

B. x diverges

6 A

7 (top down, left to right) 5.5, 5.5, -1, 2, 7

8

3 2

1 1

9 0.5

10 two remaining cells on the 3rd row: (8,6), (5,6)

11 7

12 Same as P :-) Q

13

A. {x/A, y/right(A)}

B. {x/y, y/A, z/right(A)}

C. fail

14 prune two branches: that to -2 and to 2. root=1.

15

KB=>CNF:

1. $\sim P \vee Q$

2. R

3. $\sim Q$

Negate query

4. $\sim R \vee P$

Add 4 to KB, resolve:

2, 4 => 5. P

5, 1 => 6. Q

6, 3 => false

16

1. $V(I, Duck, Tele) \wedge L(I, Park)$

2. $V(I, Duck, Tele) \wedge L(Duck, Park)$

3. $\exists x: V(I, Duck, x) \wedge L(I, Park) \wedge L(Tele, Park)$

4. $\exists x: V(I, Duck, x) \wedge L(Duck, Park) \wedge L(Tele, Park)$

5. $V(I, Duck, Tele) \wedge L(I, Park) \wedge L(Duck, Park)$

We can introduce another predicate $Cut(x, y, z)$: x cuts (using a sawing motion) y with tool z

Then one can replace all $V()$ with $Cut()$ for more interpretations.

17. This is DFS and it will loop forever.

18. Show that the KB entails empty.

19. C.

$o \xrightarrow{(1)} o \xrightarrow{(-1)} o \xrightarrow{(2)} o$ is OK

a loop with negative edge weight is not fine

20

A. strict dominating strategy: (work, work)

B. if A,B communicate, they can agree on (slack, slack)

C. you can change the outcome matrix, e.g. let (slack, slack)=-1 (cut off their communication is not a good idea :-)