

# First Order Logic Part 1

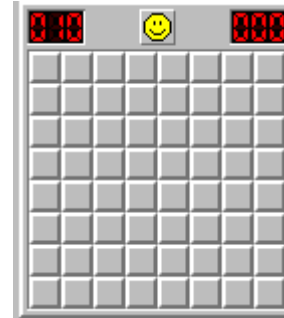
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# Problems with propositional logic

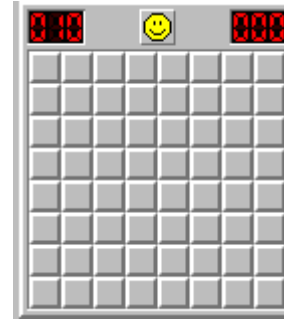
- Consider the game “minesweeper” on a 10x10 field with only one landmine.



- How do you express the knowledge, with propositional logic, that the squares adjacent to the landmine will display the number 1?

# Problems with propositional logic

- Consider the game “minesweeper” on a 10x10 field with only one landmine.



- How do you express the knowledge, with propositional logic, that the squares adjacent to the landmine will display the number 1?
- Intuitively with a rule like  
$$\text{landmine}(x,y) \Rightarrow \text{number1}(\text{neighbors}(x,y))$$
but propositional logic cannot do this...

# Problems with propositional logic

- Propositional logic has to say, e.g. for cell (3,4):
  - $\text{Landmine}_{3_4} \Rightarrow \text{number1}_{2_3}$
  - $\text{Landmine}_{3_4} \Rightarrow \text{number1}_{2_4}$
  - $\text{Landmine}_{3_4} \Rightarrow \text{number1}_{2_5}$
  - $\text{Landmine}_{3_4} \Rightarrow \text{number1}_{3_3}$
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  - $\text{Landmine}_{3_4} \Rightarrow \text{number1}_{4_3}$
  - $\text{Landmine}_{3_4} \Rightarrow \text{number1}_{4_4}$
  - $\text{Landmine}_{3_4} \Rightarrow \text{number1}_{4_5}$
  - And similarly for each of  $\text{Landmine}_{1_1}$ ,  
 $\text{Landmine}_{1_2}$ ,  $\text{Landmine}_{1_3}$ , ...,  $\text{Landmine}_{10_{10}}$ !
- Difficult to express large domains concisely
- Don't have objects and relations
- First Order Logic is a powerful upgrade

# Ontological commitment

- Logics are characterized by what they consider to be 'primitives'

<b>Logic</b>	<b>Primitives</b>	<b>Available Knowledge</b>
Propositional	facts	true/false/unknown
First-Order	facts, objects, relations	true/false/unknown
Temporal	facts, objects, relations, times	true/false/unknown
Probability Theory	facts	degree of belief 0...1
Fuzzy	degree of truth	degree of belief 0...1

# First Order Logic syntax

- **Term:** an object in the world
  - **Constant:** Jerry, 2, Madison, Green, ...
  - **Variables:** x, y, a, b, c, ...
  - **Function**(term<sub>1</sub>, ..., term<sub>n</sub>)
    - Sqrt(9), Distance(Madison, Chicago)
    - Maps one or more objects to another object
    - Can refer to an unnamed object: LeftLeg(John)
    - Represents a user defined functional relation
- A **ground term** is a term without variables.

# FOL syntax

- **Atom**: smallest T/F expression
  - **Predicate**(term<sub>1</sub>, ..., term<sub>n</sub>)
    - Teacher(Jerry, you), Bigger(sqrt(2), x)
    - Convention: read “Jerry (is)Teacher(of) you”
    - Maps one or more objects to a truth value
    - Represents a user defined relation
  - **term<sub>1</sub> = term<sub>2</sub>**
    - Radius(Earth)=6400km, 1=2
    - Represents the equality relation when two terms refer to the same object

# FOL syntax

- **Sentence:** T/F expression
  - Atom
  - Complex sentence using connectives:  $\wedge \vee \neg \Rightarrow \Leftrightarrow$ 
    - $\text{Spouse}(\text{Jerry}, \text{Jing}) \Rightarrow \text{Spouse}(\text{Jing}, \text{Jerry})$
    - $\text{Less}(11,22) \wedge \text{Less}(22,33)$
  - Complex sentence using quantifiers  $\forall, \exists$
- Sentences are evaluated under an interpretation
  - Which objects are referred to by constant symbols
  - Which objects are referred to by function symbols
  - What subsets defines the predicates



# FOL quantifiers

- Universal quantifier:  $\forall$
- Sentence is true **for all** values of  $x$  in the domain of variable  $x$ .
- Main connective typically is  $\Rightarrow$ 
  - Forms if-then rules
  - “all humans are mammals”  
$$\forall x \text{ human}(x) \Rightarrow \text{mammal}(x)$$
  - Means if  $x$  is a human, then  $x$  is a mammal

# FOL quantifiers

$$\forall x \text{ human}(x) \Rightarrow \text{mammal}(x)$$

- It's a big AND: Equivalent to the **conjunction** of all the instantiations of variable  $x$ :

$$\begin{aligned} &(\text{human}(\text{Jerry}) \Rightarrow \text{mammal}(\text{Jerry})) \wedge \\ &(\text{human}(\text{Jing}) \Rightarrow \text{mammal}(\text{Jing})) \wedge \\ &(\text{human}(\text{laptop}) \Rightarrow \text{mammal}(\text{laptop})) \wedge \dots \end{aligned}$$

- Common mistake is to use  $\wedge$  as main connective

$$\forall x \text{ human}(x) \wedge \text{mammal}(x)$$

- This means everything is human and a mammal!

$$\begin{aligned} &(\text{human}(\text{Jerry}) \wedge \text{mammal}(\text{Jerry})) \wedge \\ &(\text{human}(\text{Jing}) \wedge \text{mammal}(\text{Jing})) \wedge \\ &(\text{human}(\text{laptop}) \wedge \text{mammal}(\text{laptop})) \wedge \dots \end{aligned}$$

# FOL quantifiers

- Existential quantifier:  $\exists$
- Sentence is true **for some** value of  $x$  in the domain of variable  $x$ .
- Main connective typically is  $\wedge$ 
  - “some humans are male”  
$$\exists x \text{ human}(x) \wedge \text{male}(x)$$
  - Means there is an  $x$  who is a human and is a male

# FOL quantifiers

$$\exists x \text{ human}(x) \wedge \text{male}(x)$$

- It's a big OR: Equivalent to the **disjunction** of all the instantiations of variable  $x$ :

$$(\text{human}(\text{Jerry}) \wedge \text{male}(\text{Jerry})) \vee$$

$$(\text{human}(\text{Jing}) \wedge \text{male}(\text{Jing})) \vee$$

$$(\text{human}(\text{laptop}) \wedge \text{male}(\text{laptop})) \vee \dots$$

- Common mistake is to use  $\Rightarrow$  as main connective
  - “Some pig can fly”

$$\exists x \text{ pig}(x) \Rightarrow \text{fly}(x) \quad \text{(wrong)}$$

# FOL quantifiers

$$\exists x \text{ human}(x) \wedge \text{male}(x)$$

- It's a big OR: Equivalent to the **disjunction** of all the instantiations of variable  $x$ :

$$(\text{human}(\text{Jerry}) \wedge \text{male}(\text{Jerry})) \vee$$

$$(\text{human}(\text{Jing}) \wedge \text{male}(\text{Jing})) \vee$$

$$(\text{human}(\text{laptop}) \wedge \text{male}(\text{laptop})) \vee \dots$$

- Common mistake is to use  $\Rightarrow$  as main connective
  - “Some pig can fly”

$$\exists x \text{ pig}(x) \Rightarrow \text{fly}(x) \quad (\text{wrong})$$

- This is true if there is something not a pig!

$$(\text{pig}(\text{Jerry}) \Rightarrow \text{fly}(\text{Jerry})) \vee$$

$$(\text{pig}(\text{laptop}) \Rightarrow \text{fly}(\text{laptop})) \vee \dots$$

# FOL quantifiers

- Properties of quantifiers:
  - $\forall \mathbf{x} \forall \mathbf{y}$  is the same as  $\forall \mathbf{y} \forall \mathbf{x}$
  - $\exists \mathbf{x} \exists \mathbf{y}$  is the same as  $\exists \mathbf{y} \exists \mathbf{x}$
- Example:
  - $\forall \mathbf{x} \forall \mathbf{y} \text{ likes } (\mathbf{x}, \mathbf{y})$   
Everyone likes everyone.
  - $\forall \mathbf{y} \forall \mathbf{x} \text{ likes } (\mathbf{x}, \mathbf{y})$   
Everyone is liked by everyone.

# FOL quantifiers

- Properties of quantifiers:
  - $\forall x \exists y$  is **not** the same as  $\exists y \forall x$
  - $\exists x \forall y$  is **not** the same as  $\forall y \exists x$
- Example:
  - $\forall x \exists y \text{ likes}(x, y)$   
Everyone likes someone (can be different).
  - $\exists y \forall x \text{ likes}(x, y)$   
There is someone who is liked by everyone.

# FOL quantifiers

- Properties of quantifiers:
  - $\forall \mathbf{x} P(\mathbf{x})$  when negated becomes  $\exists \mathbf{x} \neg P(\mathbf{x})$
  - $\exists \mathbf{x} P(\mathbf{x})$  when negated becomes  $\forall \mathbf{x} \neg P(\mathbf{x})$
- Example:
  - $\forall \mathbf{x} \text{sleep}(\mathbf{x})$   
Everybody sleeps.
  - $\exists \mathbf{x} \neg \text{sleep}(\mathbf{x})$   
Somebody does not sleep.



# FOL quantifiers

- Properties of quantifiers:
  - $\forall \mathbf{x} P(\mathbf{x})$  is the same as  $\neg \exists \mathbf{x} \neg P(\mathbf{x})$
  - $\exists \mathbf{x} P(\mathbf{x})$  is the same as  $\neg \forall \mathbf{x} \neg P(\mathbf{x})$
- Example:
  - $\forall \mathbf{x} \text{sleep}(\mathbf{x})$   
Everybody sleeps.
  - $\neg \exists \mathbf{x} \neg \text{sleep}(\mathbf{x})$   
There does not exist someone who does not sleep.

# FOL syntax

- A **free variable** is a variable that is not bound by an quantifier, e.g.  $\exists y \text{ Likes}(x, y)$ :  $x$  is free,  $y$  is bound
- A **well-formed formula** (wff) is a sentence in which all variables are quantified (no free variable)
- Short summary so far:
  - **Constants:** Bob, 2, Madison, ...
  - **Variables:**  $x, y, a, b, c, \dots$
  - **Functions:** Income, Address, Sqrt, ...
  - **Predicates:** Teacher, Sisters, Even, Prime...
  - **Connectives:**  $\wedge \vee \neg \Rightarrow \Leftrightarrow$
  - **Equality:**  $=$
  - **Quantifiers:**  $\forall \exists$

## More summary

- **Term:** constant, variable, function. Denotes an object. (A ground term has no variables)
- **Atom:** the smallest expression assigned a truth value. Predicate and =
- **Sentence:** an atom, sentence with connectives, sentence with quantifiers. Assigned a truth value
- **Well-formed formula (wff):** a sentence in which all variables are quantified