

## POPULAR SHOPPING CART APP PLUGS DOZENS OF XSS VULNERABILITIES

by **Tom Spring**

March 28, 2016 , 5:13 pm

Popular open source shopping cart app Zen Cart is warning its users of dozens of cross-site scripting vulnerabilities found in its software. Affected websites, security experts say, risk exposing customers to malware, theft of cookies data and site defacement.

Zen Cart, with an estimated 113,000 active users (according to BuiltWith.com), has told its users they will have to pro-actively install the patch. Affected customers told Threatpost that Zen Cart has notified them of the vulnerability offering

For its part, Trustwave told Threatpost that 50 XSS vulnerabilities were found in the admin section of the Zen Cart software along with one issue in the non-authentication portion of the application.

“We discovered the XSS on Zen Cart in a completely random manner,” said Alex Rothacker, senior security researcher at Trustwave. “We have a monthly Hack Friday event where the team goes ahead and picks something and tries to find vulnerabilities. In this case, one team member picked Zen Cart, because it was a popular solution.”

# network security

CS642

adam everspough computer security

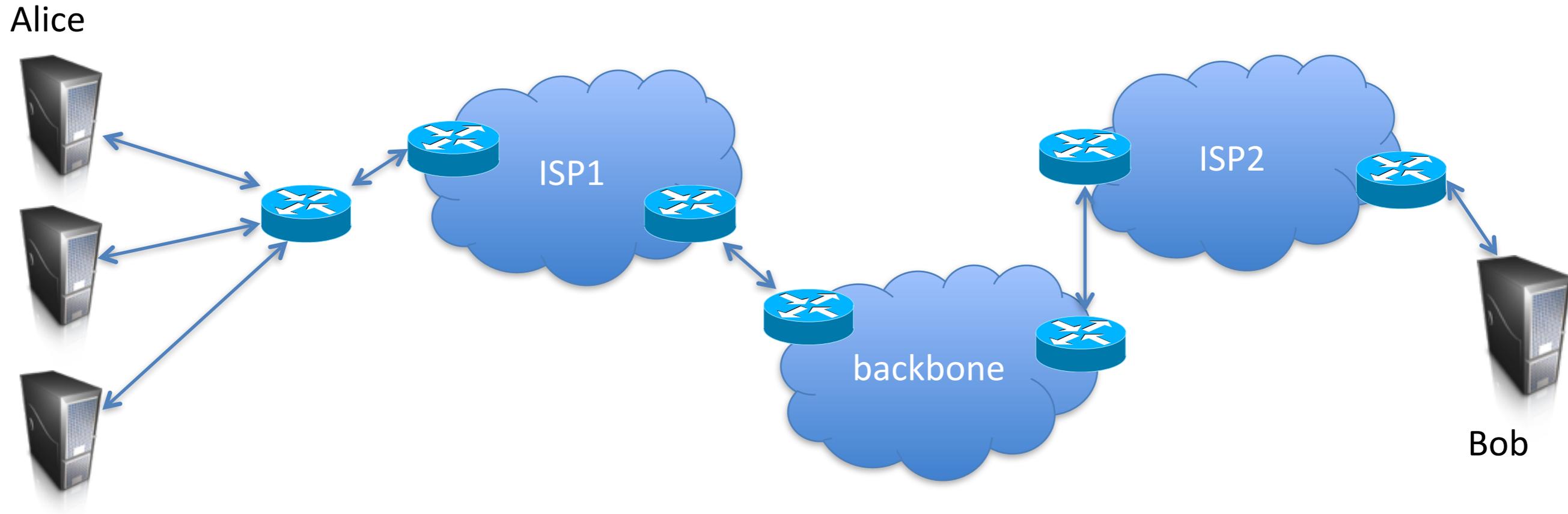
ace@cs.wisc.edu

# today

- \* Crypto exercise in-class
- \* Link layer (in-)security
- \* IP, TCP (in-)security

*crypto  
exercise*

# Internet



Local area network  
(LAN)

Ethernet

802.11

Internet

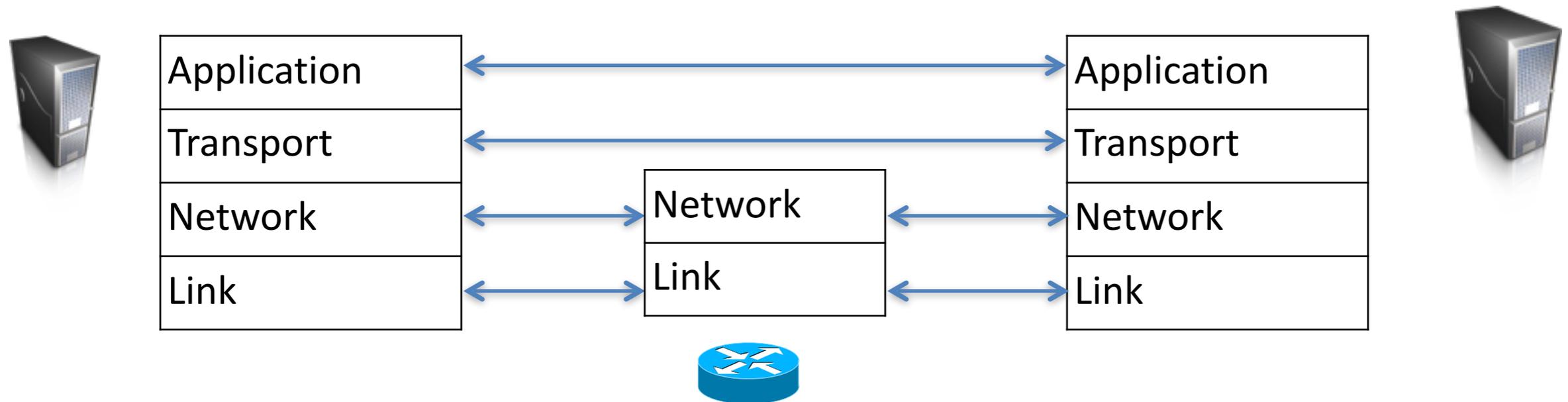
TCP/IP

BGP (border gateway protocol)

DNS (domain name system)

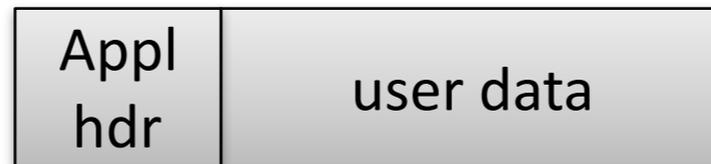
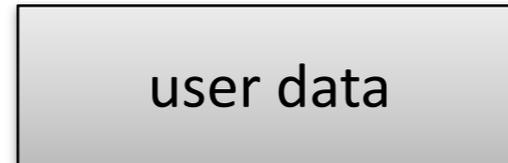
# Internet protocol stack

Application	HTTP, FTP, SMTP, SSH, etc.
Transport	TCP, UDP
Network	IP, ICMP, IGMP
Link	802x (802.11, Ethernet)

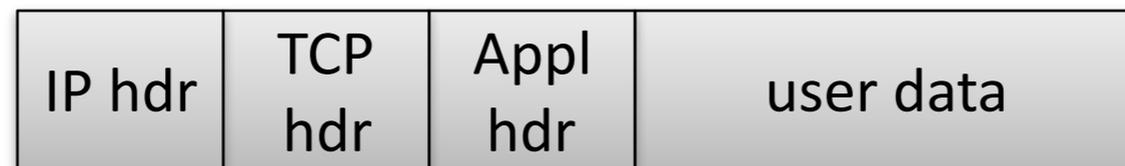


# Internet protocol stack

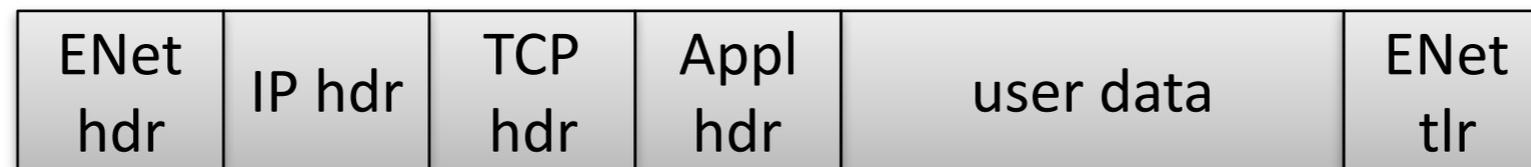
Application
TCP
IP
Ethernet



TCP segment



IP datagram



Ethernet frame

14

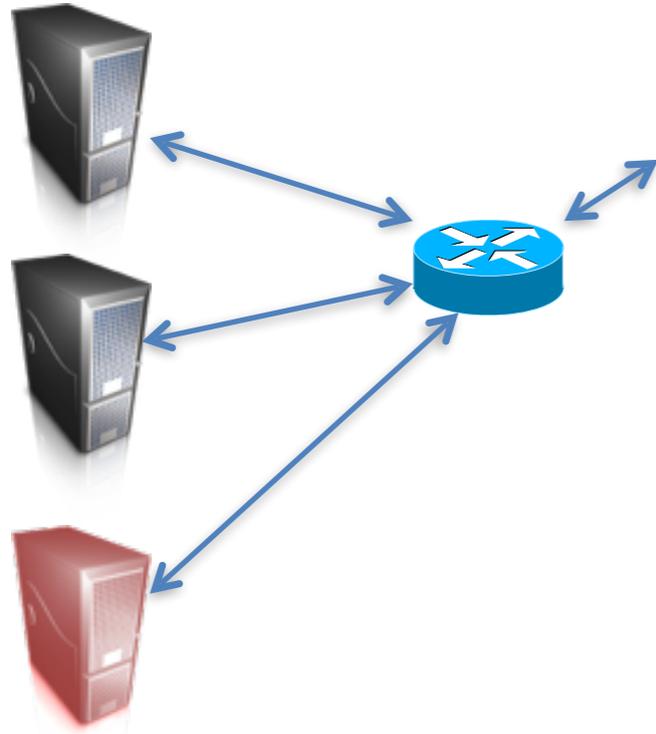
20

20



46 to 1500 bytes

# Address resolution protocol



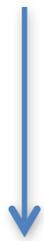
IP routing:

Figure out where to send an IP packet based on destination address.

Link layer and IP must cooperate to route packets

32-bit IP address

ARP



48-bit MAC address

ARP enables this cooperation by mapping IPs to MACs

# ARP caches

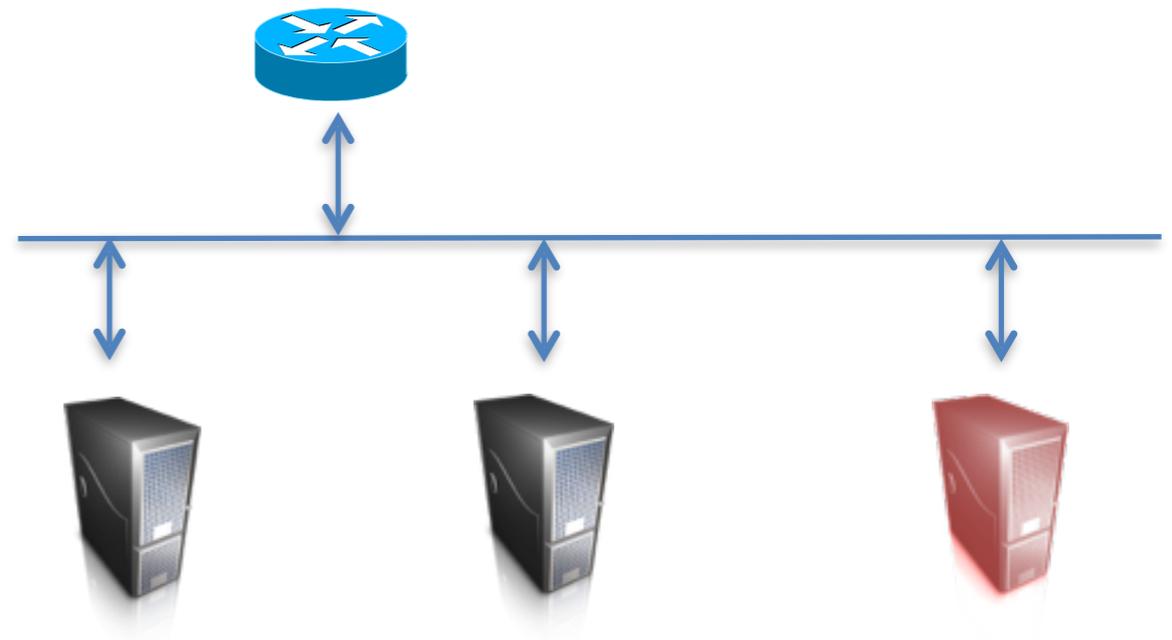
- Hosts maintain cache of ARP data
  - just a table mapping between IPs and MACs

```
rist@wifi-212:~/work/teaching/642-fall-2011/slides$ arp -a
? (172.16.219.1) at 0:50:56:c0:0:1 on vmnet1 ifscope permanent [ethernet]
? (172.16.219.255) at (incomplete) on vmnet1 ifscope [ethernet]
? (192.168.1.1) at 98:fc:11:91:73:92 on en1 ifscope [ethernet]
? (192.168.1.255) at (incomplete) on en1 ifscope [ethernet]
? (192.168.38.255) at (incomplete) on vmnet8 ifscope [ethernet]
rist@wifi-212:~/work/teaching/642-fall-2011/slides$
```

# ARP has no authentication

- Easy to sniff packets on (non-switched) ethernet
- What else can we do?

Easy Denial of Service (DoS):  
Send ARP reply associating gateway 192.168.1.1 with a non-used MAC address



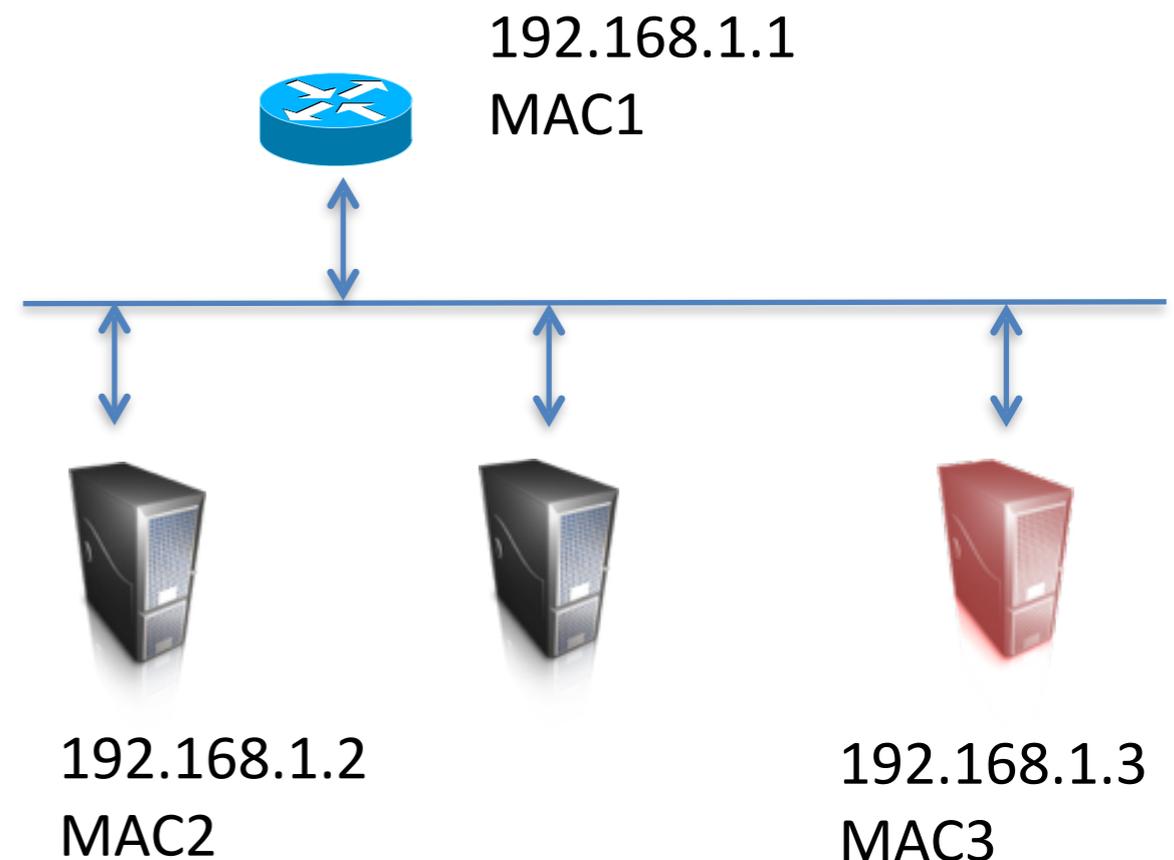
# ARP has no authentication

- Easy to sniff packets on (non-switched) ethernet
- What else can we do?

## Active Man-in-the-Middle:

ARP reply to MAC2  
192.168.1.1 -> MAC3

ARP reply to MAC1  
192.168.1.2 -> MAC3



Now traffic “routed” through malicious box

# 802.11 (wifi)

**STA** = station

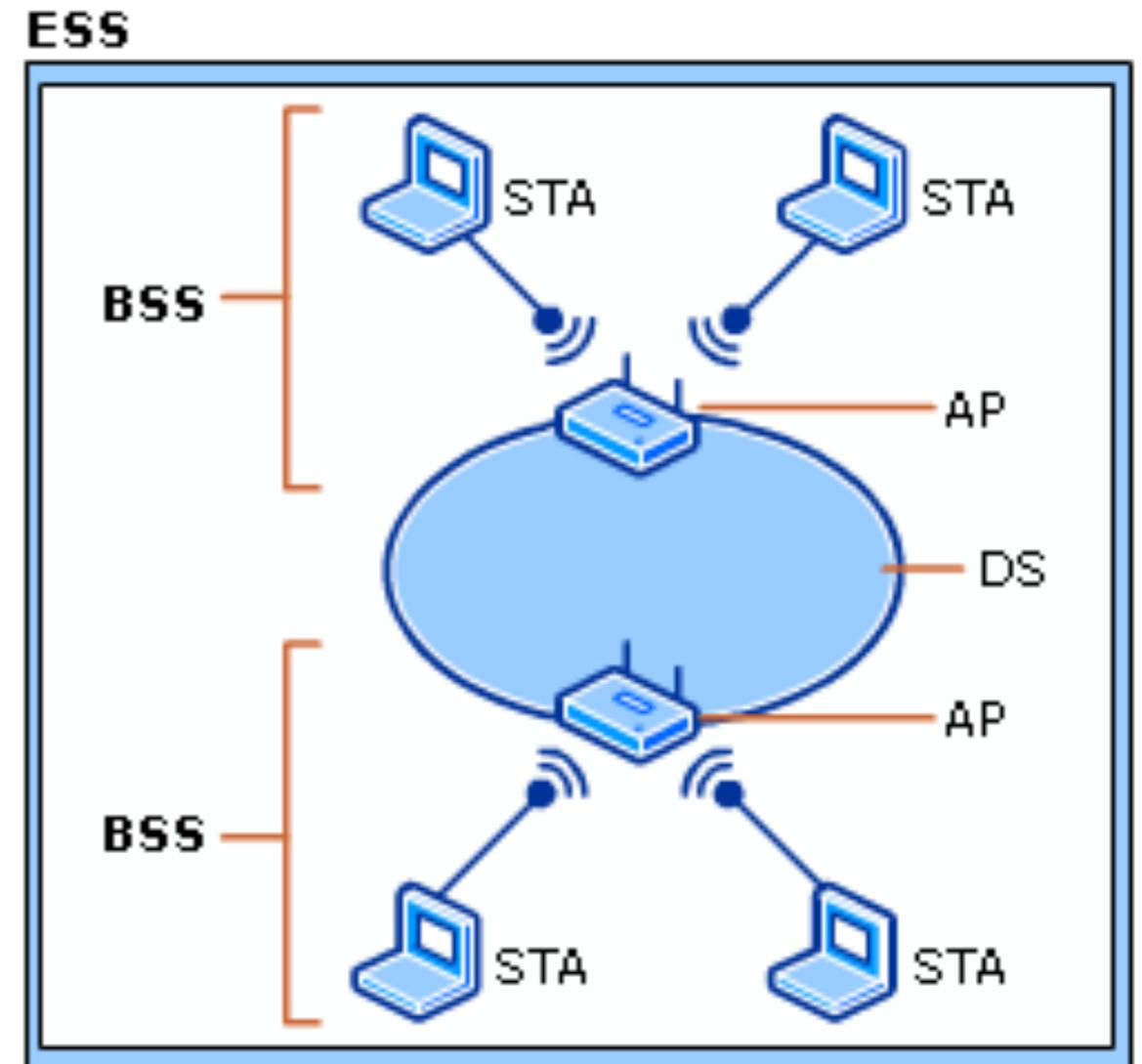
**AP** = access point

BSS = basic service set

DS = distribution service

ESS = extended service set

**SSID** (service set identifier)  
identifies the 802.11 network



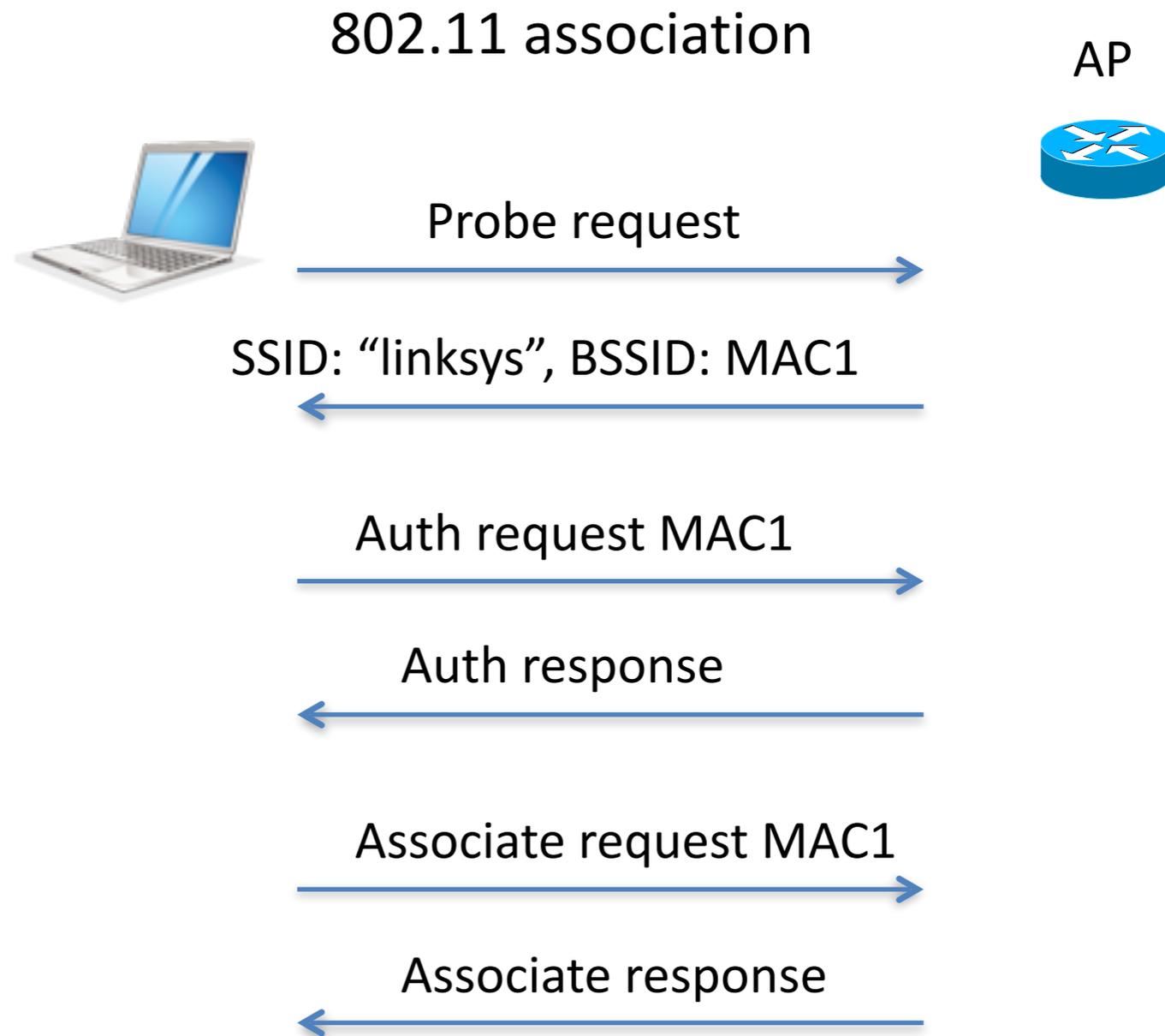
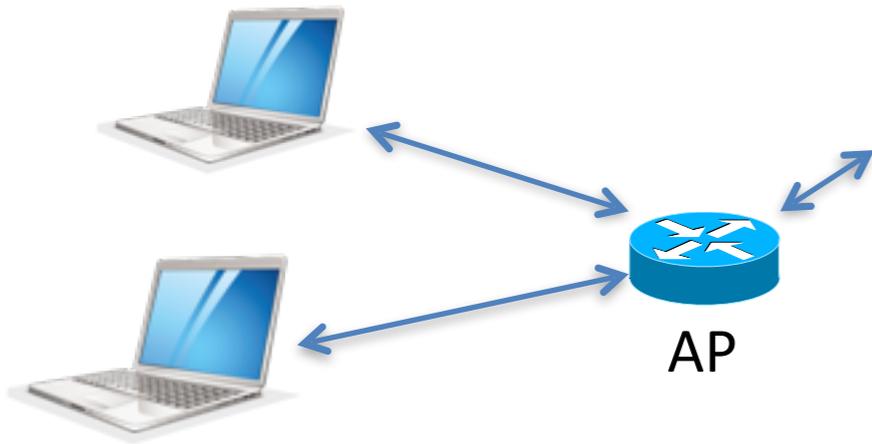
[http://technet.microsoft.com/en-us/library/cc757419\(WS.10\).aspx](http://technet.microsoft.com/en-us/library/cc757419(WS.10).aspx)

## Typical WiFi modes:

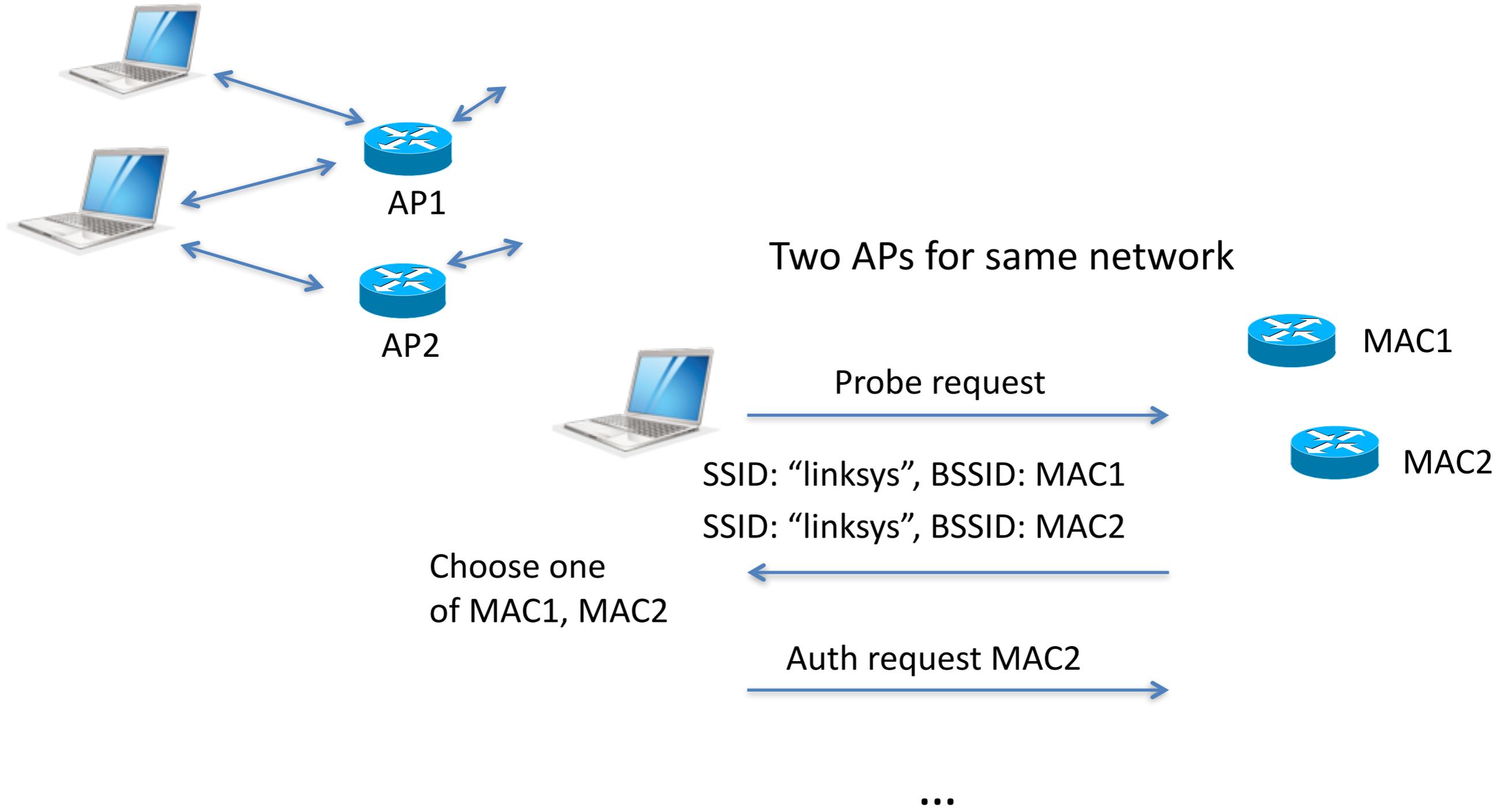
Unsecured

Wireless Protected Access (WPA2) - password authenticated, encrypted

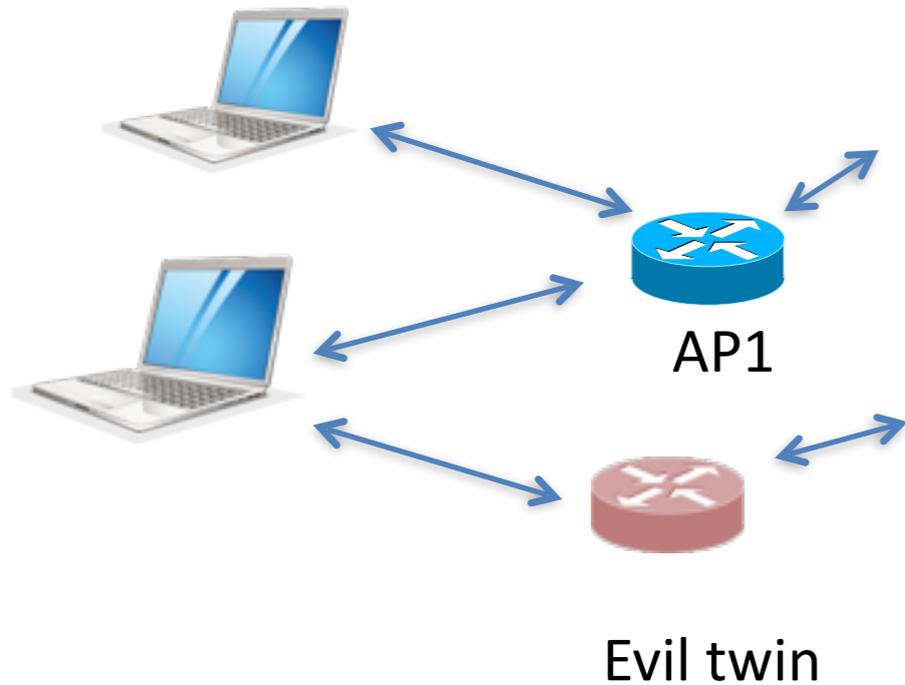
# 802.11 association



# 802.11 association



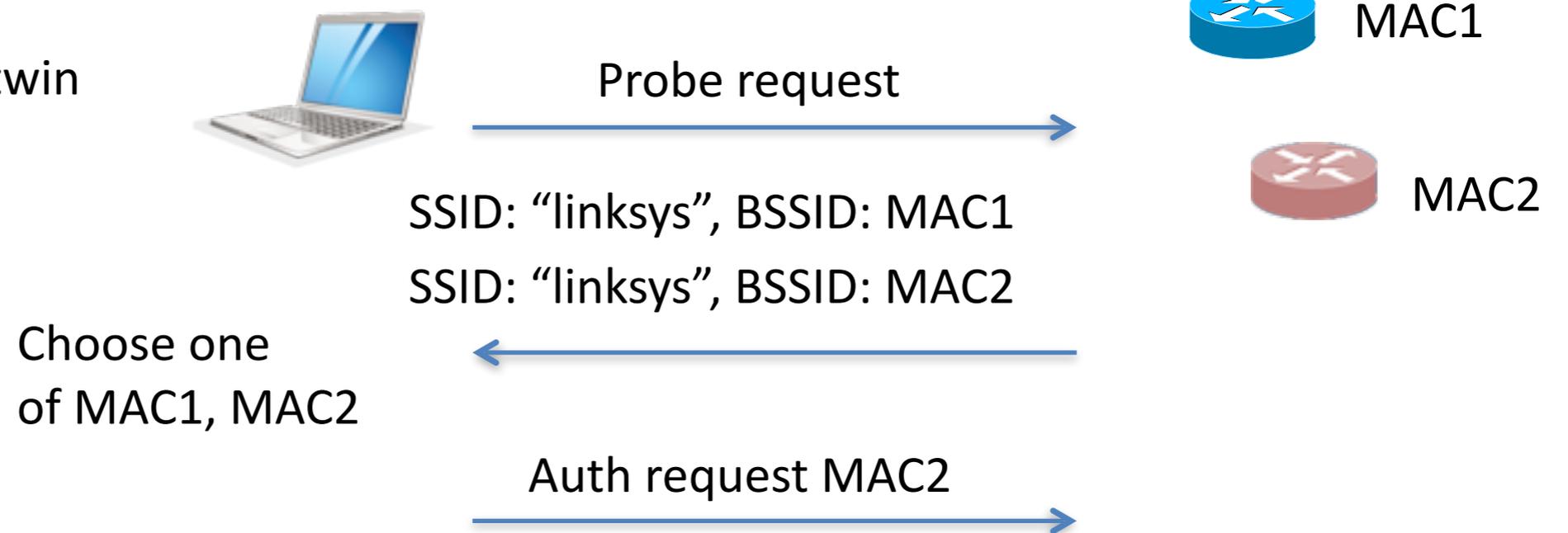
# 802.11 evil twins



Basic idea:

- Attacker pretends to be an AP to intercept traffic or collect data

Basic attack: rogue AP



...



## Parrot ARdrone

Drone is a WiFi access point

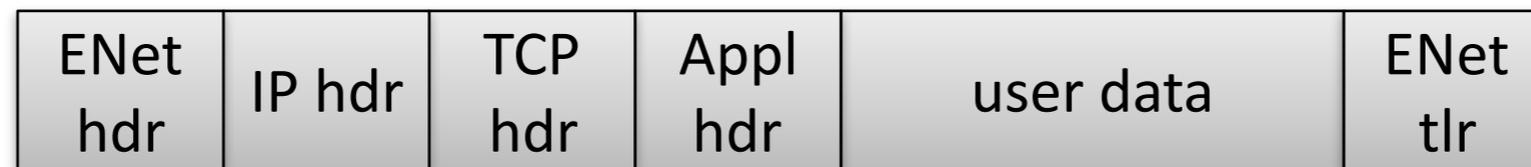
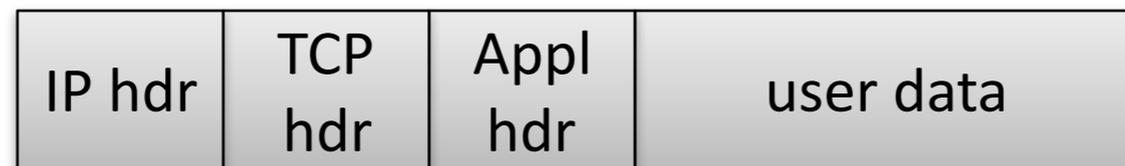
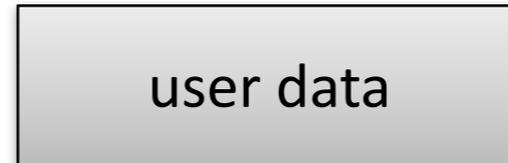
Uses unsecured 802.11 connection (WiFi)

Controlled from iPad or iPhone with an app

Uses MAC address for security

# Internet protocol stack

Application
TCP
IP
Ethernet



14

20

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46 to 1500 bytes

TCP segment

IP datagram

Ethernet frame

# IP protocol (IPv4)

- Connectionless
  - no state
- Unreliable
  - no guarantees
- ICMP (Internet Control Message Protocol)
  - error messages, etc.
  - often used by tools such as ping, traceroute

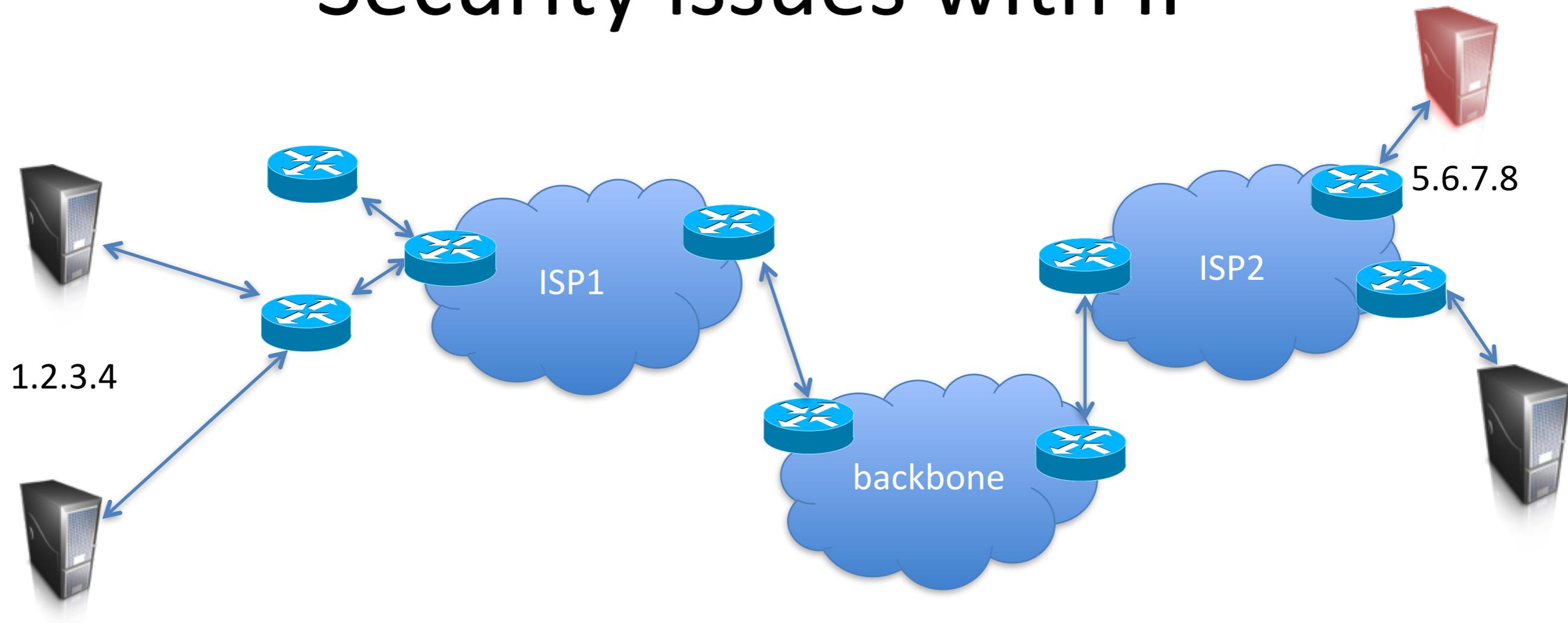
# IPv4



Ethernet frame  
containing  
IP datagram

4-bit version	4-bit hdr len	8-bit type of service	16-bit total length (in bytes)	
16-bit identification			3-bit flags	13-bit fragmentation offset
8-bit time to live (TTL)		8-bit protocol	16-bit header checksum	
32-bit source IP address				
32-bit destination IP address				
options (optional)				

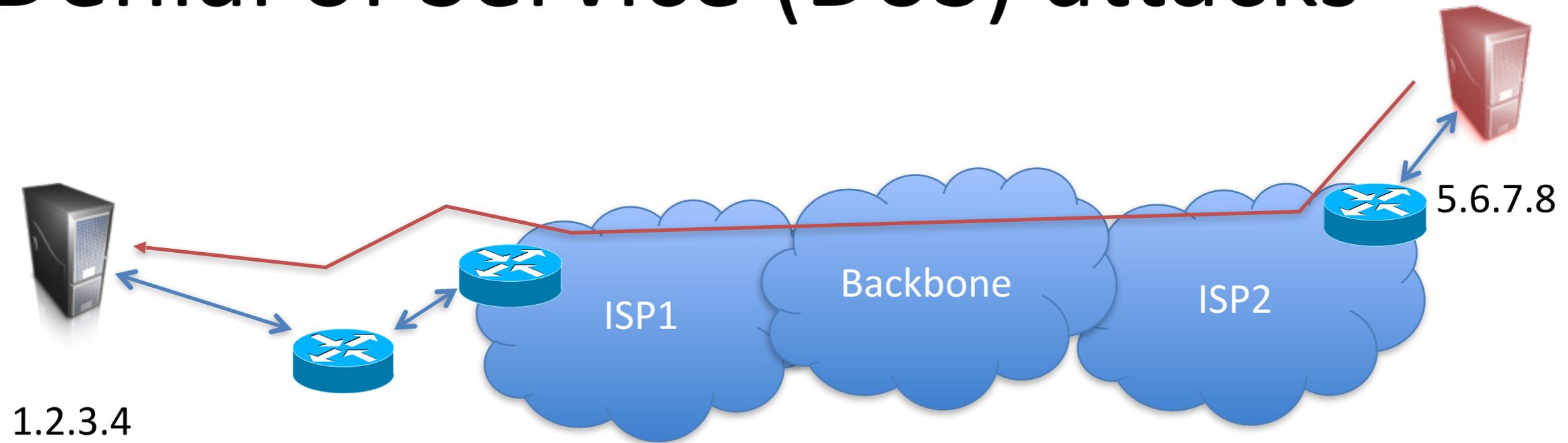
# Security issues with IP



Routing has issues, we'll get to that later  
What else?

- No source address authentication in general

# Denial of Service (DoS) attacks

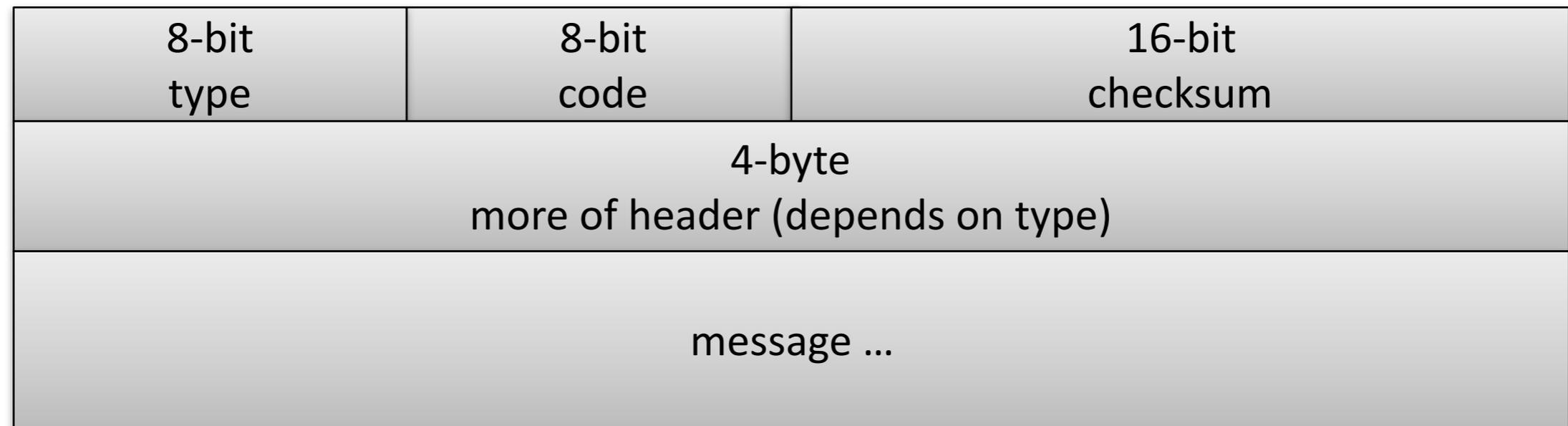


Goal: prevent legitimate users from accessing victim (1.2.3.4)

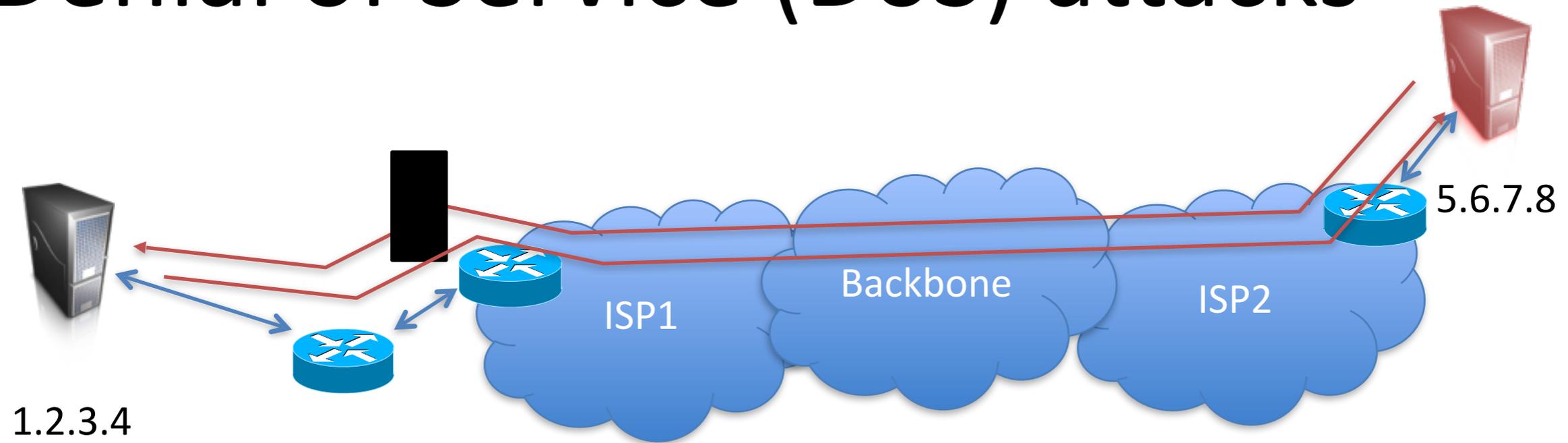
ICMP ping flood

# ICMP

## (Internet Control Message Protocol)



# Denial of Service (DoS) attacks

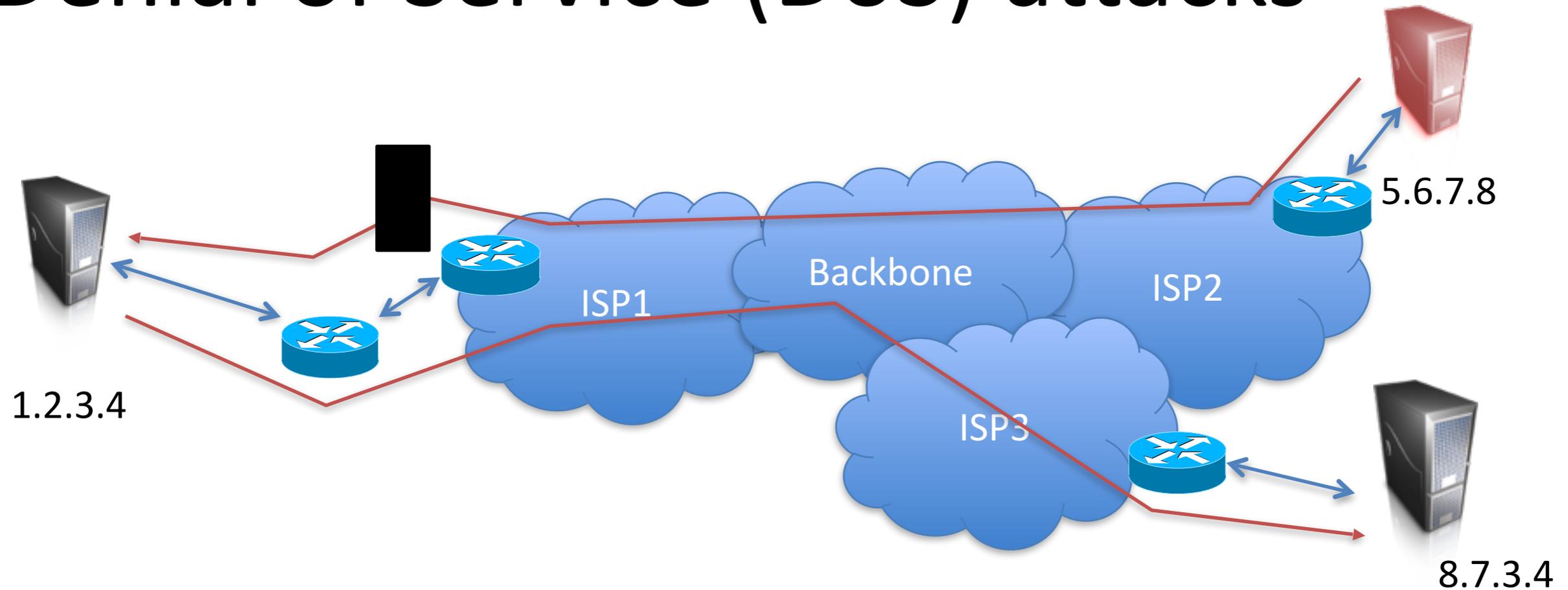


Goal is to prevent legitimate users from accessing victim (1.2.3.4)

## ICMP ping flood

- Attacker sends ICMP pings as fast as possible to victim
- When will this work as a DoS? **Attacker resources > victim's**
- How can this be prevented? **Ingress filtering near victim**

# Denial of Service (DoS) attacks



How can attacker avoid ingress filtering?

Attacker can send packet with fake source IP “spoofed” packet

Packet will get routed correctly

Replies will not

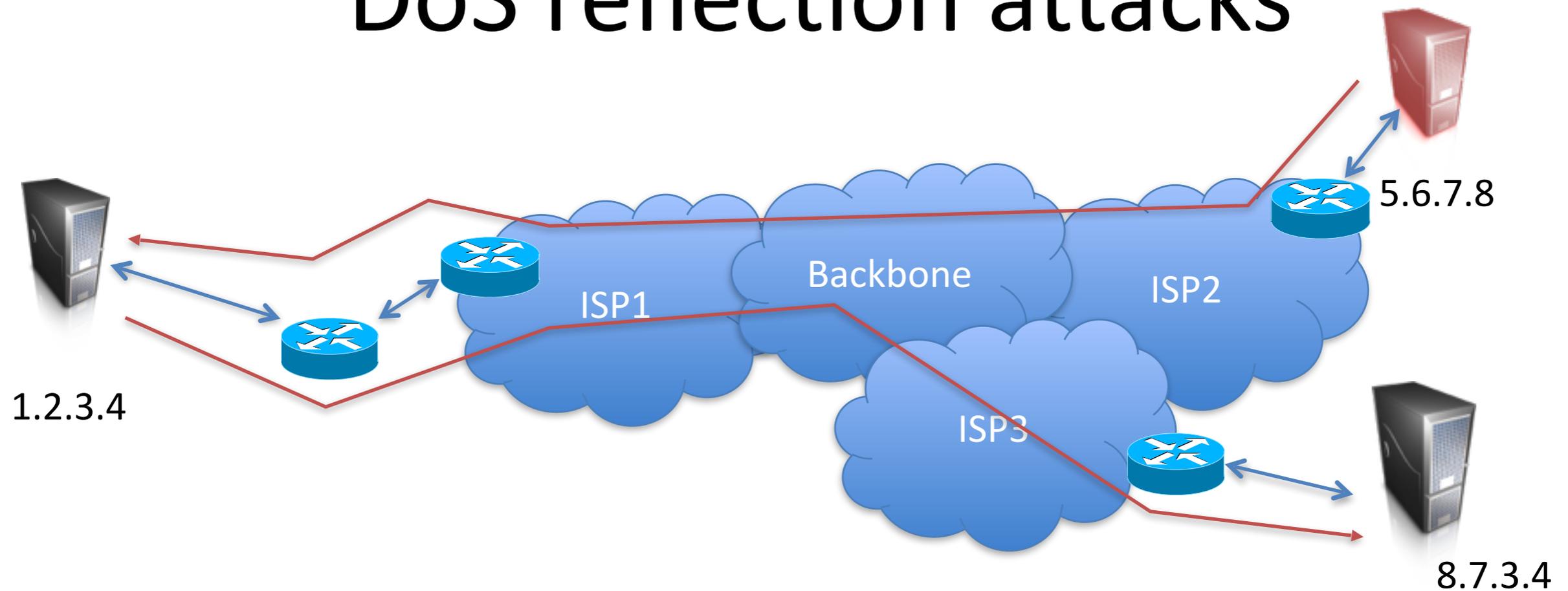
Send IP packet with

source: 8.7.3.4  
dest: 1.2.3.4

from 5.6.7.8

Filter based on source may be incorrect

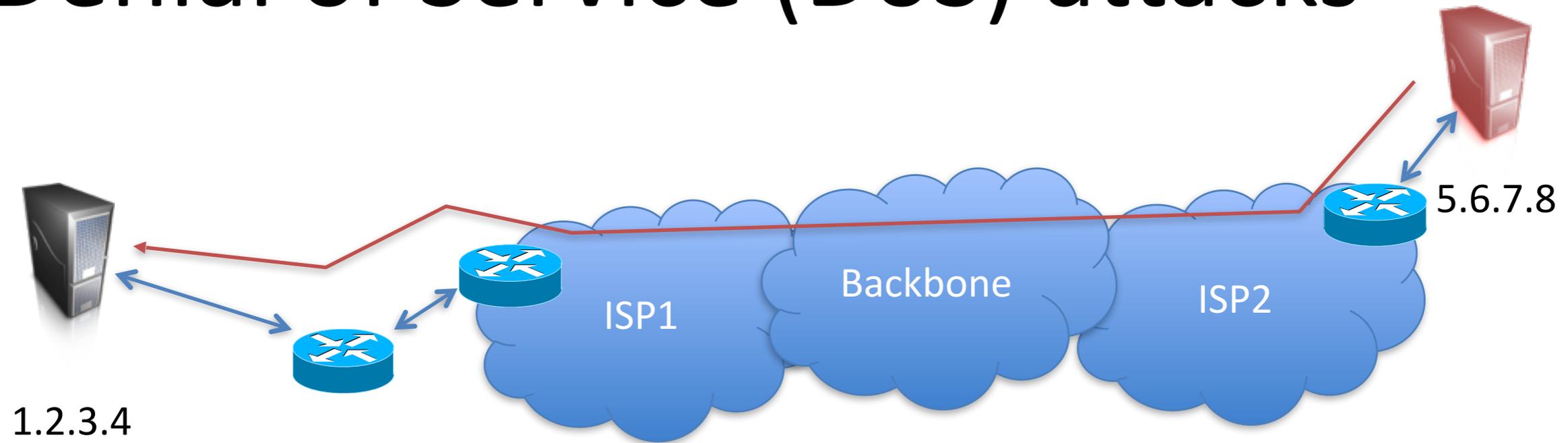
# DoS reflection attacks



Note a valid packet sends a reply to 8.7.3.4

- Attacker can bounce an attack against 8.7.3.4 off 1.2.3.4
- "Frame" 1.2.3.4
- Single-packet exploit (1.2.3.4 in foreign country)

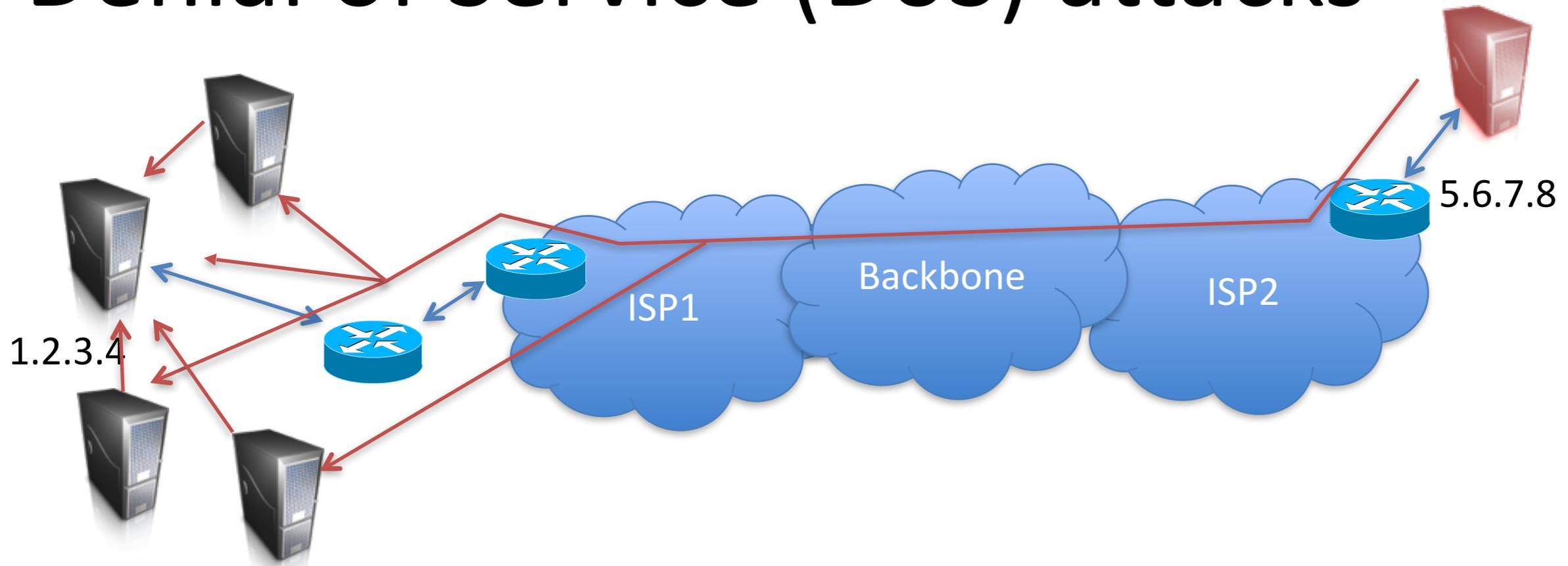
# Denial of Service (DoS) attacks



DoS works better when there is *asymmetry* between victim and attacker

- Attacker uses few resources to cause victim to consume lots of resources

# Denial of Service (DoS) attacks



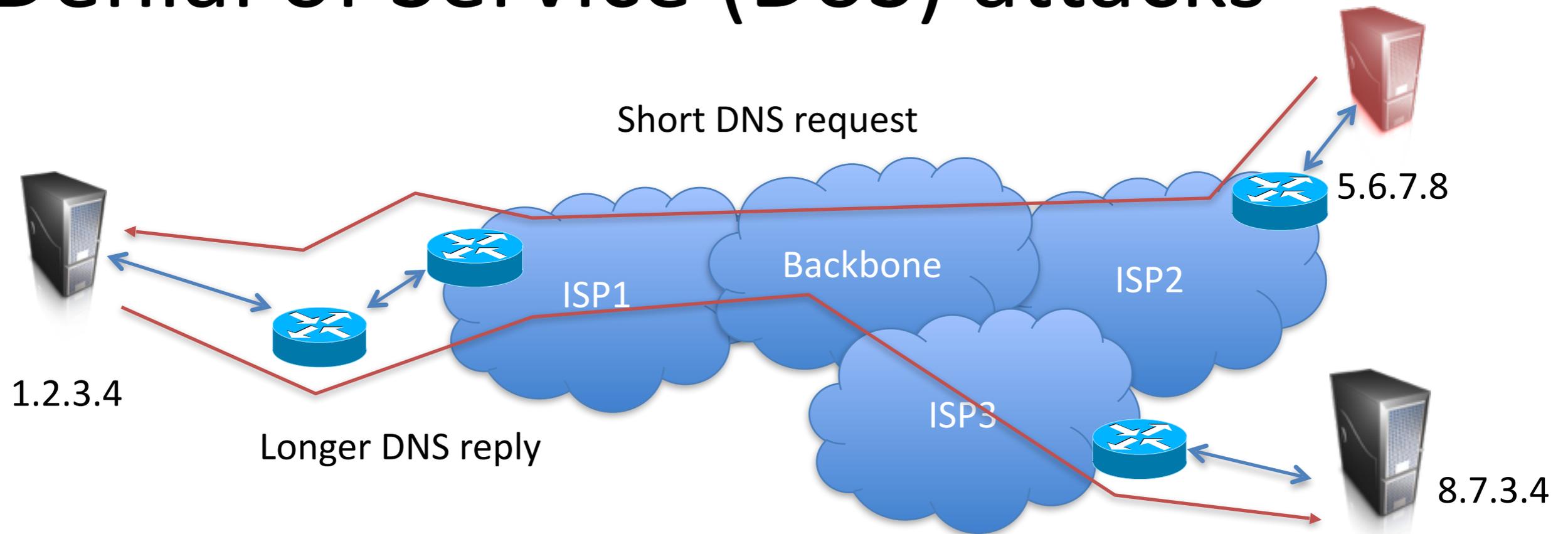
DoS works better when there is *asymmetry* between victim and attacker

- Attacker uses few resources to cause victim to consume lots of resources

Old example: Smurf attack

Router allows attacker to send broadcast ICMP ping on network. Attacker spoofs SRC address to be 1.2.3.4

# Denial of Service (DoS) attacks



DoS works better when there is *asymmetry* between victim and attacker

- Attacker uses few resources to cause victim to consume lots of resources

More recent: DNS reflection attacks

Send DNS request w/ spoofed target IP (~65 byte request)  
DNS replies sent to target (~512 byte response)

- \* In-class exercise
  - / Hybrid encryption, digital signatures, PBKDF
- \* Network Security
  - / ARP cache poisoning, MitM, DoS
  - / WiFi Evil Twins
  - / IP (in-security)
- \* Exit slips
  - / 1 thing you learned
  - / 1 thing you didn't understand

recap