

THE abc-CONJECTURE AND THE LARGEST PRIME FACTOR OF $2^n + 1$

DENIS XAVIER CHARLES

ABSTRACT. We show that under the abc-conjecture the largest prime factor of $2^n + 1$ is $\Omega(n^{\frac{4}{3}})$ infinitely often.

1. MAIN RESULT

The abc-conjecture. (Oesterlé, Masser, Szpiro) For any $\epsilon > 0$, and $a, b, c \in \mathbb{Z}$ which are relatively prime satisfying $a + b = c$, we have

$$\max\{|a|, |b|, |c|\} \ll_{\epsilon} \left(\prod_{p \mid abc} p \right)^{1+\epsilon}.$$

This conjecture has some amazing consequences for example it implies the Fermat's Last Theorem for all large enough exponents. In this article we show that it yields some non-trivial lower bounds on the largest prime factor of the sequence $2^n + 1$.

Theorem 1.1. *The largest prime factor of $2^n + 1$ is $\Omega(n^{\frac{4}{3}})$ infinitely often.*

Proof : Let P_x be the largest prime factor of

$$\prod_{1 \leq n \leq x} (2^n + 1).$$

We will adopt the following notation:

$$\begin{aligned} \square_n &= \prod_{p \mid 2^n + 1} p \\ \blacksquare_n &= \prod_{p^{\nu} \mid 2^n + 1, \nu \geq 2} p^{\nu} \end{aligned}$$

The abc-conjecture can be used to lower bound the product

$$\prod_{1 \leq n \leq x} \square_n.$$

Since $2^n, 1$ and $2^n + 1$ are relatively prime in pairs and since $2^n + 1 = \square_n \blacksquare_n$ we have for any fixed $\epsilon > 0$,

$$\begin{aligned} 2^n + 1 &= \max\{2^n, 1, 2^n + 1\} \\ &\ll_{\epsilon} \left(2 \square_n \sqrt{\blacksquare_n} \right)^{1+\epsilon} \end{aligned}$$

since \blacksquare_n is the powerful part of $2^n + 1$. Since $\blacksquare_n \leq 2^n + 1$, we get

$$\square_n \gg 2^{\frac{n(1-\epsilon)}{2}}.$$

Thus we get

$$\log \left(\prod_{1 \leq n \leq x} \square_n \right) \gg_{\epsilon} x^2.$$

We also have

$$\prod_{1 \leq n \leq x} \square_n = \prod_{p \leq P_x} p^{N_x(p)}$$

where

$$N_x(p) = \sum_{1 \leq n \leq x, p \nmid 2^n + 1} 1.$$

Putting these together we get the lower bound

$$\sum_{p \leq P_x} N_x(p) \log p \gg_{\epsilon} x^2.$$

Now we will derive an upper bound for the sum $\sum_{p \leq P_x} N_x(p) \log p$ in terms of P_x to get the result.

We split up $\sum_{p \leq P_x} N_x(p) \log p$ into two parts $\Sigma_A = \sum_{p \leq x} N_x(p) \log p$ and $\Sigma_B = \sum_{x \leq p \leq P_x} N_x(p) \log p$.

Estimation of Σ_A : If $2^i + 1 \equiv 0 \pmod{p}$, we have $i \geq \log p$ this is in particular a lower bound on the order of 2 in the multiplicative subgroup of p . Thus we have $N_x(p) \leq \frac{x}{\log p} + 1$. Now

$$\begin{aligned} \Sigma_A &= \sum_{p \leq x} N_x(p) \log p \\ &\ll \sum_{p \leq x} \left(\frac{x \log p}{\log p} + \log p \right) \\ &= O\left(\frac{x^2}{\log x} \right). \end{aligned}$$

Estimation of Σ_B :

We can assume that

$$\forall p : x \leq p \leq P_x : \text{Ord}_p(2) \geq \sqrt{p}$$

since otherwise the theorem is already true. For $p \geq x$ we have $N_x(p) \leq \frac{\varphi(p)}{\sqrt{p}} \leq \sqrt{p}$. Thus

$$\begin{aligned} \Sigma_B &= \sum_{x \leq p \leq P_x} N_x(p) \log p \\ &\leq \sum_{x \leq p \leq P_x} \sqrt{p} \log p \\ &\leq \sqrt{P_x} \sum_{x \leq p \leq P_x} \log p \\ &\leq P_x^{1.5}. \end{aligned}$$

Together with the lower bound for $\Sigma_A + \Sigma_B$ we have

$$P_x^{1.5} \gg_{\epsilon} x^2.$$

This proves the theorem. \square

Remark 1.2. The main use of the abc-conjecture was to avoid the estimation of the sum

$$\sum_{p \leq P_x, p^a \nmid 2^n + 1, a > 1} N_x(p^a) \log p.$$

It seems like good upper bounds on the number of *Wieferich primes* (primes p such that $2^{p-1} \equiv 1 \pmod{p^2}$) are required to bound this sum. Though only two Wieferich primes are known below 3×10^9 its still an open problem whether there are infinitely many such primes (or even if there are infinitely many non-Wieferich

primes). Under the abc-conjecture we know that there are infinitely many non-Wieferich primes [Sil88] but the lower bound is only logarithmic.

REFERENCES

- [Sil88] Silverman, Joseph H.; *Wieferich's criterion and the abc-conjecture*, J. Number Theory, (30), no. 2, 226-237, (1988).