[537] RAID

Chapter 38
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10/29/14

Review Disks/Devices

Device Protocol Variants

Status checks: polling vs. interrupts

Data: PIO vs. DMA

Control: special instructions vs. memory-mapped I/O

Disks

Doing an I/O requires:

- seek
- rotate
- transfer

What is expensive?

Schedulers

Strategy: reorder requests to meet some goal

- performance (by making I/O sequential)
- fairness
- consistent latency

Usually in both OS and H/W.

CFQ (Linux Default)

Completely Fair Queueing.

Queue for each process.

Do weighted round-robin between queues, with slice time proportional to priority.

Optimize order within queue.

Yield slice only if idle for a given time (anticipation).

RAID

Only One Disk?

Sometimes we want many disks — why?

Only One Disk?

Sometimes we want many disks — why?

- capacity
- performance
- reliability

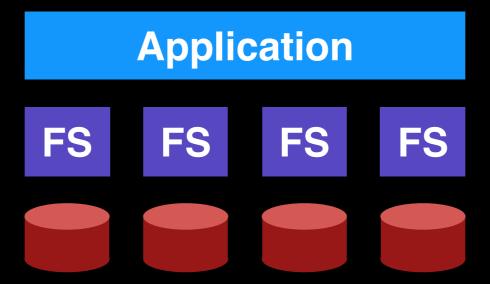
Only One Disk?

Sometimes we want many disks — why?

- capacity
- performance
- reliability

Challenge: most file systems work on only one disk.

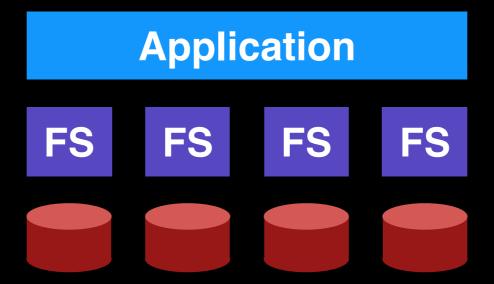
Solution 1: JBOD



Application is smart, stores different files on different file systems.

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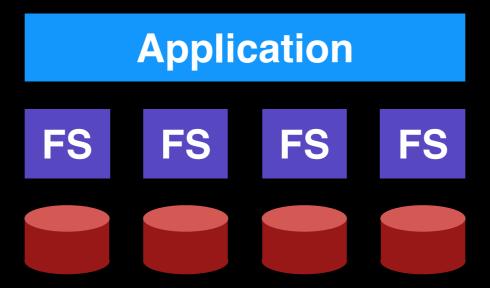
JBOD: Just a Bunch Of Disks (seriously!)



Application is smart, stores different files on different file systems.

Solution 1: JBOD

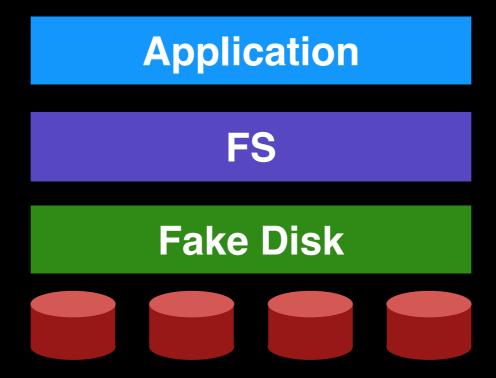
JBOD: Just a Bunch Of Disks (seriously!)



alternatives?

Application is smart, stores different files on different file systems.

RAID: Redundant Array of Inexpensive Disks

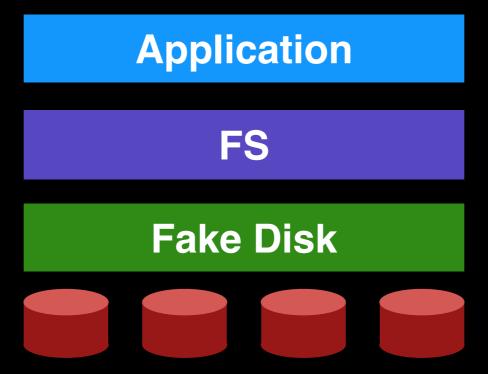


Build logical disk from many physical disks.

RAID: Redundant Array of Inexpensive Disks

RAID is:

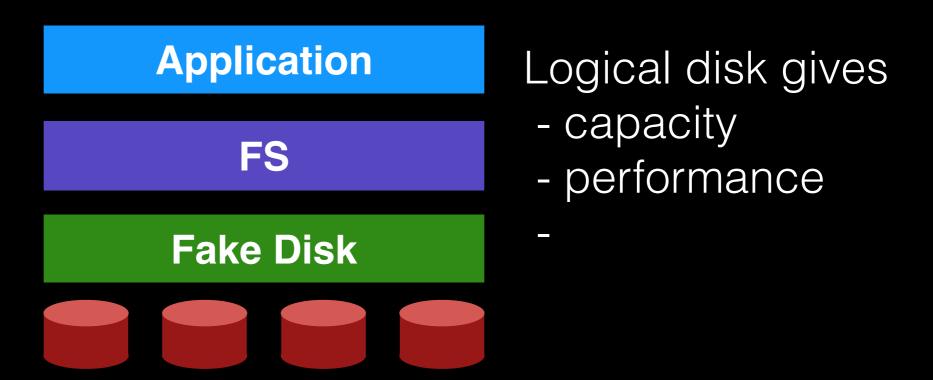
- transparent
- deployable



RAID: Redundant Array of Inexpensive Disks

RAID is:

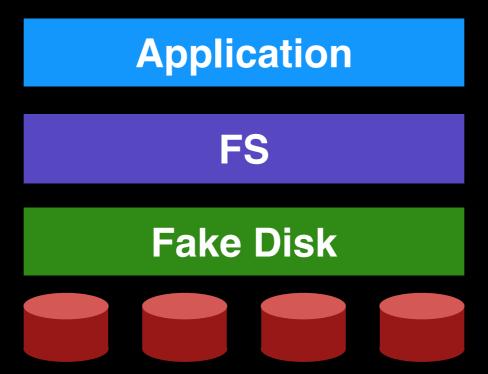
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RAID: Redundant Array of Inexpensive Disks

RAID is:

- transparent
- deployable



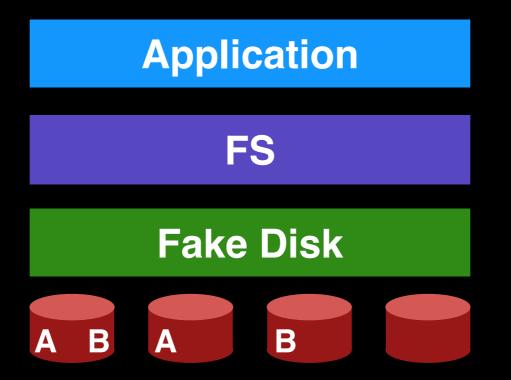
Logical disk gives

- capacity
- performance
- reliability?

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- transparent
- deployable



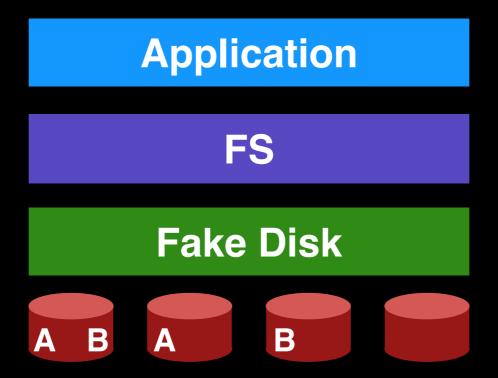
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- capacity
- performance
- reliability?

RAID: Redundant Array of Inexpensive Disks

RAID is:

- transparent
- deployable



Logical disk gives

- capacity
- performance
- reliability?

Why Inexpensive Disks?

Economies of scale! Cheap disks are popular.

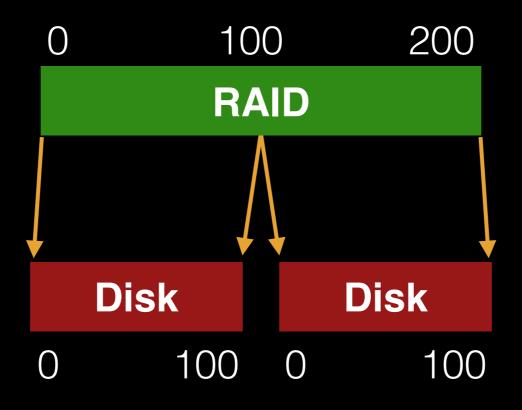
You can often get many commodity H/W components for the same price as a few expensive components.

Strategy: write S/W to build high-quality logical devices from many cheap devices.

Alternative to RAID: buy an expensive, high-end disk.

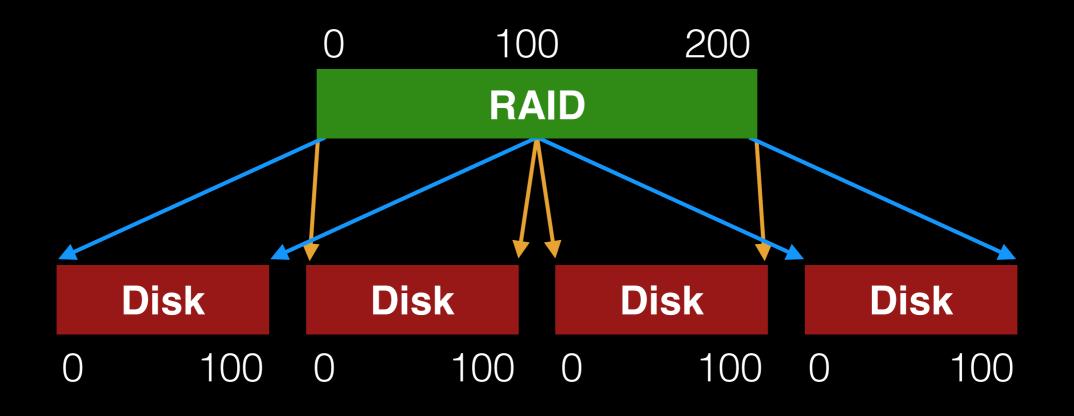
General Strategy

Build fast, large disk from smaller ones.



General Strategy

Add even more disks for reliability.



Mapping

How should we map logical to physical addresses?

How is this problem similar to virtual memory?

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Dynamic mapping: use data structure (hash table, tree)

- paging

Static mapping: use math

- RAID

Mapping

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Static mapping: use math

- RAID

RAID volume is fixed-sized, dense

Redundancy: how many copies?

System engineers are always trying to increase or decrease redundancy.

Increase: replication (e.g., RAID)

Decrease: deduplication (e.g., code sharing)

Redundancy: how many copies?

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Increase: replication (e.g., RAID)

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Why are we so fickle?

Increase: improves reliability

Decrease: improves space efficiency

One strategy: reduce redundancy as much is possible. Then add back just the right amount.

Increase: improves <u>reliability</u>

Decrease: improves space efficiency

One strategy: reduce redundancy as much is possible. Then add back just the right amount.

Take course on channel/source coding. (e.g., ECE 729)

RAID Analysis

Reasoning About RAID

Workload: types of reads/writes issued by app

RAID: system for mapping logical to physical addrs

Metric: capacity, reliability, performance

RAID "algebra", given 2 variables, find the 3rd:

$$f(W, R) = M$$

Reasoning About RAID

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RAID Decisions

Which logical blocks map to which physical blocks?

How do we use extra physical blocks (if any)?

Reasoning About RAID

Workload: types of reads/writes issued by app

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RAID "algebra", given 2 variables, find the 3rd:

$$f(W, R) = M$$

Workloads

Reads Writes

Workloads

```
Reads
One operation
Steady I/O
Writes
One operation
Steady I/O
```

Workloads

```
Reads
  One operation
  Steady I/O
     Sequential
     Random
Writes
  One operation
  Steady I/O
     Sequential
     Random
```

Reasoning About RAID

Workload: types of reads/writes issued by app

RAID: system for mapping logical to physical addrs

Metric: capacity, reliability, performance

RAID "algebra", given 2 variables, find the 3rd:

$$f(W, R) = M$$

Metrics

Capacity: how much space can apps use?

Reliability: how many disks can we safely lose?

Performance: how long does each workload take?

Metrics

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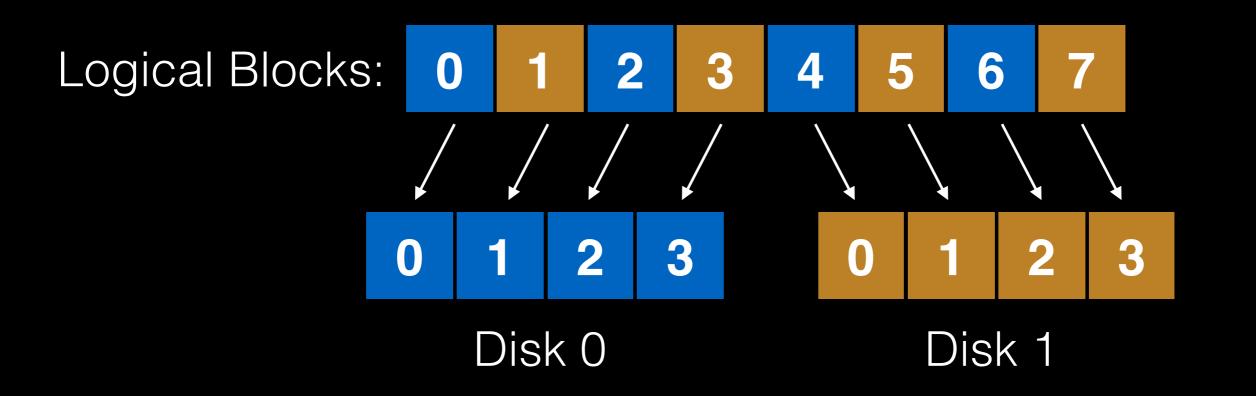
Performance: how long does each workload take?

Normalize each to characteristics of one disk (see definitions on worksheet).

RAID-0

RAID-0: Striping

Optimize for capacity. No redundancy (weird name).



Another View

Disk 0	Disk 1
O	1
2	3
4	5
6	7

Disk 0	Disk 1	Disk 2	Disk 4
O	1	2	3
4	5	6	7
8	9	10	11
12	13	14	15

	Disk 0	Disk 1	Disk 2	Disk 4
	O	1	2	3
stripe:	4	5	6	7
	8	9	10	11
	12	13	14	15

How to Map

Given logical address A, find:

$$Disk = ...$$

Offset = ...

Disk 0	Disk 1	Disk 2	Disk 4
O	1	2	3
4	5	6	7
8	9	10	11
12	13	14	15

How to Map

Given logical address A, find:

Disk = A % disk_count

Offset = A / disk_count

Disk 0	Disk 1	Disk 2	Disk 4
O	1	2	3
4	5	6	7
8	9	10	11
12	13	14	15

Disk 0	Disk 1	Disk 2	Disk 4
O	1	2	3
4	5	6	7
8	9	10	11
12	13	14	15

Disk 0	Disk 1	Disk 2	Disk 4
O	2	4	6
1	3	5	7
8	10	12	14
9	11	13	15

Disk 0	Disk 1	Disk 2	Disk 4
0	2	4	6
1	3	5	7
8	(10)	12	14
9	11	13	15

	Disk 0	Disk 1	Disk 2	Disk 4
	O	2	4	6
	1	3	5	7
otrino o	8	10	12	14
stripe:	9	11	13	15

	Disk 0	Disk 1	Disk 2	Disk 4
	O	2	4	6
	1	3	5	7
otrino	8	10	12	14
stripe:	9	11	13	15

We'll assume chunk size of 1 for today. Sizes of 64KB are typical in deployment.

	Disk 0	Disk 1	Disk 2	Disk 4
	O	2	4	6
	1	3	5	7
otrino	8	10	12	14
stripe:	9	11	13	15

When are small chunks better? When are big ones better?

(Problem C)

0a) What is capacity?

0b) How many disks can fail?

Oc) Throughput?

0a) What is capacity? N * C

0b) How many disks can fail?

Oc) Throughput?

Oa) What is capacity? N * C

0b) How many disks can fail?

Oc) Throughput?

Oa) What is capacity? N * C

Ob) How many disks can fail? 0

Oc) Throughput? N*S (i,ii) and N*R (iii,iv)

0a) What is capacity? N * C

Ob) How many disks can fail?

Oc) Throughput? N*S (i,ii) and N*R (iii,iv)

0d) Latency? D (i,ii)

0a) What is capacity? N * C

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Oc) Throughput? N*S (i,ii) and N*R (iii,iv)

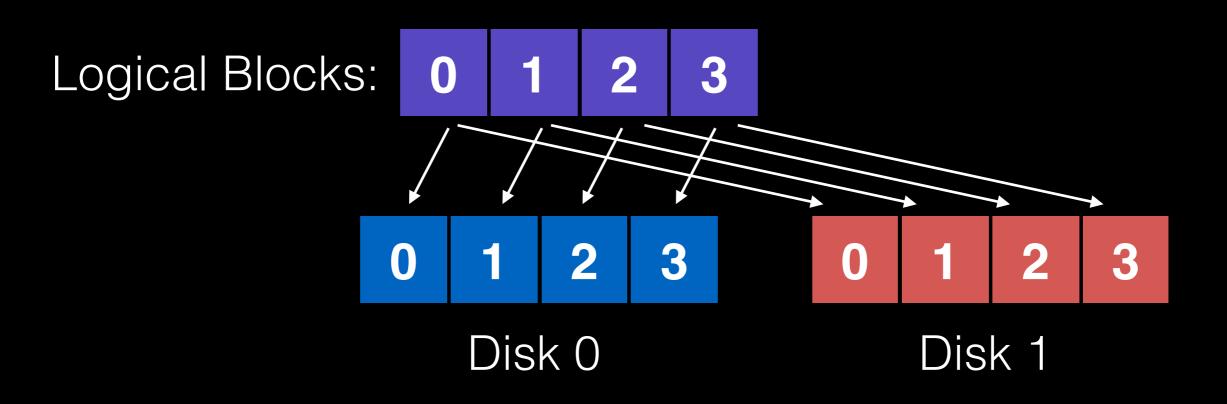
Od) Latency? D (i,ii)

Buying more disks improves throughput, but not latency!

RAID-1

RAID-1: Mirroring

Keep two copies of all data.



Assumptions

Assume disks are fail-stop.

- they work or they don't
- we know when they don't

Tougher Errors:

- latent sector errors
- silent data corruption

Disk 0	Disk 1
O	O
1	1
2	2
3	3

Disk 0	Disk 1	Disk 2	Disk 4
O	O	1	1
2	2	3	3
4	4	5	5
6	6	7	7

Disk 0	Disk 1	Disk 2	Disk 4
O	O	1	1
2	2	3	3
4	4	5	5
6	6	7	7

How many disks can fail?

Oa) What is capacity? N/2 * C

0b) How many disks can fail? 1 (or maybe N / 2)

Oc) Throughput? ???

0d) Latency? D (i,ii)

RAID-1: Throughput

What is steady-state throughput for

- sequential reads?
- sequential writes?
- random reads?
- random writes?

RAID-1: Throughput

What is steady-state throughput for

- sequential reads? N/2 * S
- sequential writes? N/2 * S
- random reads? N * R
- random writes? N/2 * R

Crashes

Disk0
 Disk1
 A
 A
 B
 B
 C
 C
 D

Crashes

Disk0
 Disk1
 A
 A
 B
 B
 C
 C
 D

write(A) to 2

Crashes

Disk0
 Disk1
 A
 A
 B
 B
 C
 D

write(A) to 2

Disk0
 Disk1
 A
 A
 B
 B
 A
 A
 A
 D

write(A) to 2

Disk0
 Disk1
 A
 A
 B
 B
 A
 A
 A
 D

Disk0
 Disk1
 A
 A
 B
 B
 A
 A
 A
 D

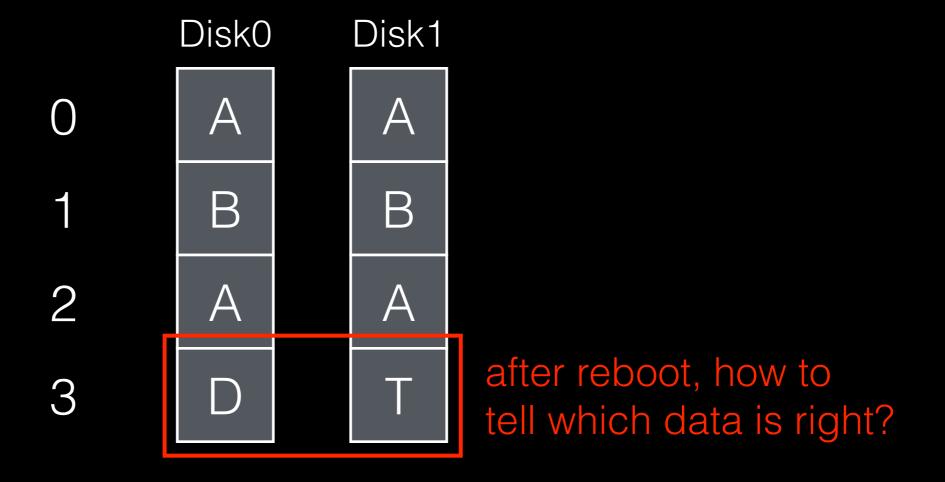
write(T) to 3

Disk0
 Disk1
 A
 A
 B
 B
 A
 A
 T

write(T) to 3

Disk0
 Disk1
 A
 A
 B
 B
 A
 A
 T

CRASH!!!



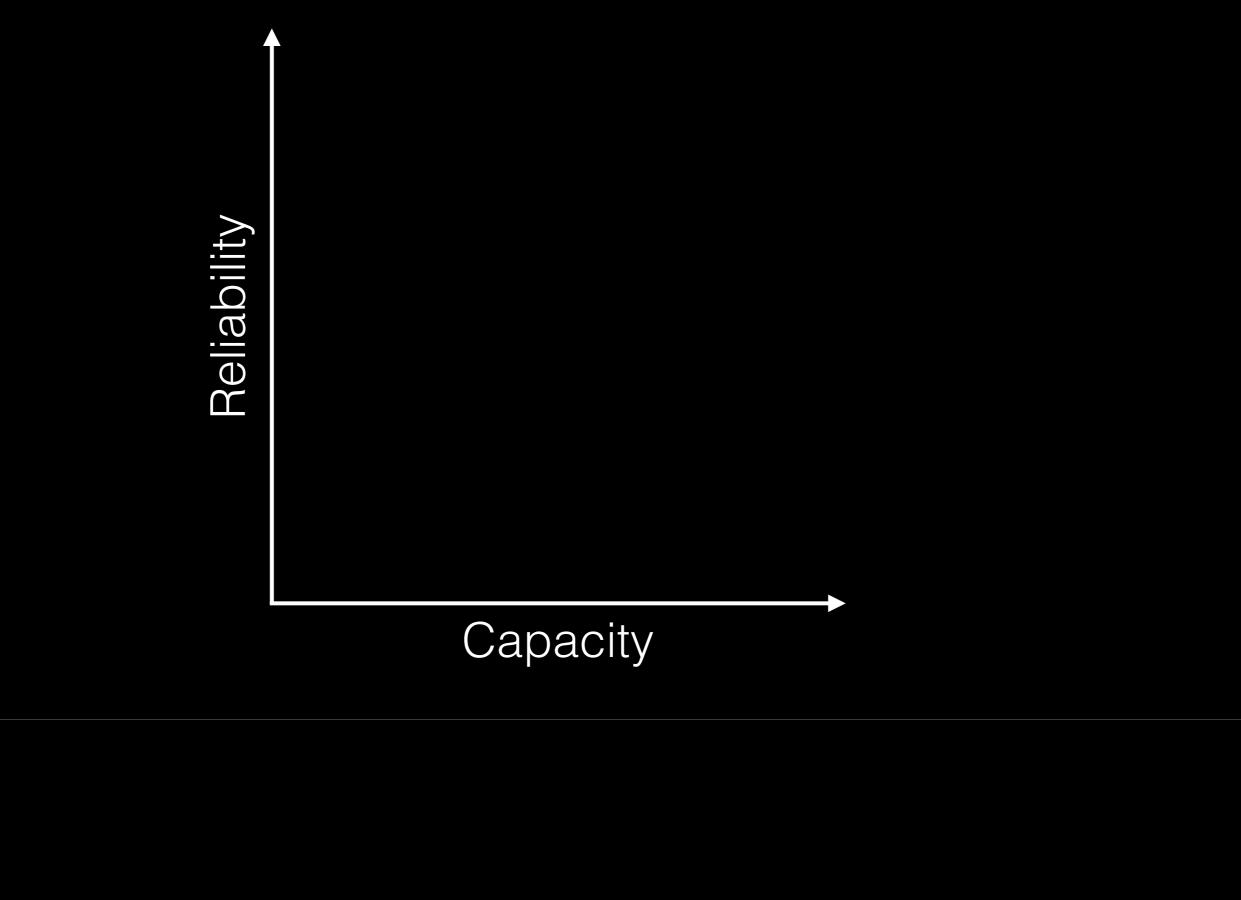
H/W Solution

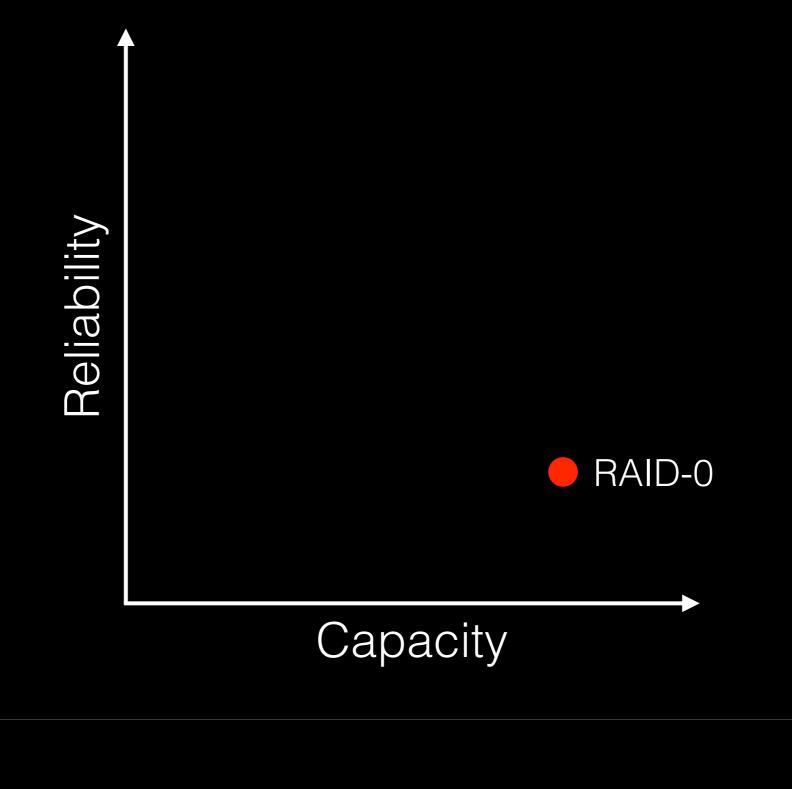
Problem: Consistent-Update Problem

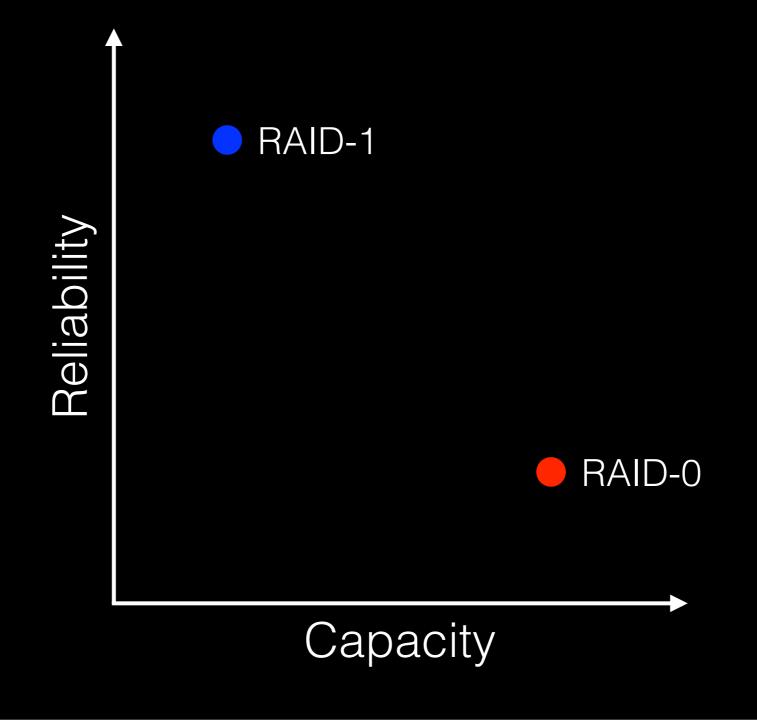
Use non-volatile RAM in RAID controller.

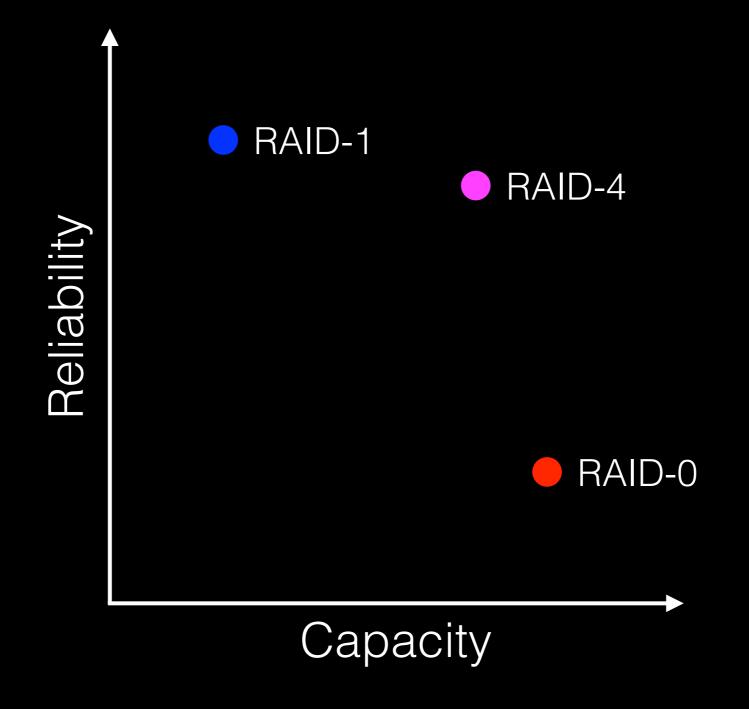
Software RAID controllers (e.g., Linux md) don't have this option.

RAID-4









Strategy

Use parity disk.

In algebra, if an equation has N variables, and N-1 are know, you can often solve for the unknown.

Treat the sectors across disks in a stripe as an equation.

A bad disk is like an unknown in the equation.

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:					

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:					
					(parity)

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:	5	3	O	1	
					(parity)

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:	5	3	O	1	9
					(parity)

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:	5	X	O	1	9
					(parity)

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:	5	3	O	1	9
					(parity)

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:	2	1	1	X	5
					(parity)

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:	2	1	1	1	5
					(parity)

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:	3	O	1	2	Χ
					(parity)

	Disk0	Disk1	Disk2	Disk3	Disk4
Stripe:	3	O	1	2	6
					(parity)

Parity Functions

Problem C.

Which functions could we use to compute parity?

Remember, disk blocks are finite (assume modulo arithmetic).

RAID-4: Analysis

0a) What is capacity?

0b) How many disks can fail?

Oc) Throughput?

0d) Latency?

Disk0	Disk1	Disk2	Disk3	Disk4
3	0	1	2	6

(parity)

RAID-4: Analysis

```
Oa) What is capacity? (N-1) * C
```

0b) How many disks can fail? 1

Oc) Throughput? ???

0d) Latency? D (i), 2*D (ii)

RAID-4: Throughput

- sequential reads?
- sequential writes?
- random reads?
- random writes?

RAID-4: Throughput

```
- sequential reads? (N-1) * S
```

- sequential writes? (N-1) * S
- random reads? (N-1) * R
- random writes? R/2

RAID-4: Throughput

(R/2)

What is steady-state throughput for

```
- sequential reads? (N-1) * S
```

- sequential writes? (N-1) * S
- random reads? (N-1) * R
- random writes?

how to avoid parity bottleneck?

RAID-5

RAID-5

Disk0	Disk1	Disk2	Disk3	Disk4
_	_	-	-	Р
-	-	-	Р	-
		Р	_	-

RAID-5: Analysis

```
Oa) What is capacity? (N-1) * C
```

0b) How many disks can fail? 1

Oc) Throughput? ???

0d) Latency? D (i), 2*D (ii)

RAID-5: Throughput

- sequential reads?
- sequential writes?
- random reads?
- random writes?

RAID-5: Throughput

```
- sequential reads? (N-1) * S
```

- sequential writes? (N-1) * S
- random reads? N * R
- random writes? N/4 * R

RAID-5: Throughput

N/4 * R

- sequential reads? (N-1) * S
- sequential writes? (N-1) * S
- random reads? N * R
- random writes?

```
(N-1) * R
R/2
RAID-4
```

All RAID

	Reliability	Capacity
RAID-0	Ο	C*N
RAID-1	1	C*N/2
RAID-4	1	N-1
RAID-5	1	N-1

	Read Latency	Write Latency
RAID-0	D	D
RAID-1	D	D
RAID-4	D	2D
RAID-5	D	2D

	Read Latency	Write Latency
RAID-0	D	D
RAID-1	D	D
RAID-4	D	2D
RAID-5	D	2D

but RAID-5 can do more in parallel

	Seq Read	Seq Write	Rand Read	Rand Write
RAID-0	N * S	N * S	N * R	N * R
RAID-1	N/2 * S	N/2 * S	N * R	N/2 * R
RAID-4	(N-1)*S	(N-1)*S	(N-1)*R	R/2
RAID-5	(N-1)*S	(N-1)*S	N * R	N/4 * R

	Seq Read	Seq Write	Rand Read	Rand Write
RAID-0	N * S	N * S	N * R	N * R
RAID-1	N/2 * S	N/2 * S	N * R	N/2 * R
RAID-4	(N-1)*S	(N-1)*S	(N-1)*R	R/2
RAID-5	(N-1)*S	(N-1)*S	N * R	N/4 * R

RAID-5 is strictly better than RAID-4

	Seq Read	Seq Write	Rand Read	Rand Write
RAID-0	N * S	N * S	N * R	N * R
RAID-1	N/2 * S	N/2 * S	N * R	N/2 * R
RAID-5	(N-1)*S	(N-1)*S	N * R	N/4 * R

	Seq Read	Seq Write	Rand Read	Rand Write
RAID-0	N * S	N * S	N * R	N * R
RAID-1	N/2 * S	N/2 * S	N * R	N/2 * R
RAID-5	(N-1)*S	(N-1)*S	N * R	N/4 * R

RAID-0 is always fastest and has best capacity. (but at cost of reliability)

	Seq Read	Seq Write	Rand Read	Rand Write
RAID-0	N * S	N * S	N * R	N * R
RAID-1	N/2 * S	N/2 * S	N * R	N/2 * R
RAID-5	(N-1)*S	(N-1)*S	N * R	N/4 * R

	Seq Read	Seq Write	Rand Read	Rand Write
RAID-0	N * S	N * S	N * R	N * R
RAID-1	N/2 * S	N/2 * S	N * R	N/2 * R
RAID-5	(N-1)*S	(N-1)*S	N * R	N/4 * R

RAID-5 better than RAID-1 for sequential.

	Seq Read	Seq Write	Rand Read	Rand Write
RAID-0	N * S	N * S	N * R	N * R
RAID-1	N/2 * S	N/2 * S	N * R	N/2 * R
RAID-5	(N-1)*S	(N-1)*S	N * R	N/4 * R

	Seq Read	Seq Write	Rand Read	Rand Write
RAID-0	N * S	N * S	N * R	N * R
RAID-1	N/2 * S	N/2 * S	N * R	N/2 * R
RAID-5	(N-1)*S	(N-1)*S	N * R	N/4 * R

RAID-1 better than RAID-4 for random write.

Summary

Many engineering tradeoffs with RAID. (capacity, reliability, different types of performance).

H/W RAID controllers can handle crashes easier.

Transparent, deployable solutions are popular.

Announcements

Extra late days!

p3b due Friday.

I/O schedulers tomorrow (discussion).

Watch email for office hours today (but assume usual if you hear nothing).