

## Previous Lecture

Two intuitive, but slow sorting algorithms

## Selection sort:

- Repeat for each key in list
- Find minimum key in unsorted portion
- Move to next position of sorted portion


## Insertion sort:

- Repeat for each key in unsorted list
- Insert into its correct position in sorted portion

Both algorithms $O\left(N^{2}\right)$ where $N$ is length of lis $\dagger$


## Recursive Algorithms

Algorithm is recursive if can be defined by:

- Simple base case
- Set of rules reducing other cases toward base case


Recursion: If you still don't get it, see: "Recursion".

## Recursive Definition of <br> Factorial

Example: Fact(5) $=5!=5 * 4 * 3 * 2 * 1$
Recursive definition:

- Fact $(1)=1$ [base case]
- For all integers $n>1:$ Fact $(n)=n$ * Fact ( $n-1$ )

Fact(5) $=$ ? ?
$=5$ * Fact (4)
$=5 * 4^{*}$ Fact (3)
$=5 * 4^{*} 3^{*}$ Fact(2)
$=5$ * 4* $^{*}{ }^{*} 2^{*}$ Fact (1)
$=5 * 4^{*} 3^{*} 2^{*} 1$ Recursion ends!

## Merging Two Sorted Runs

| $\longrightarrow$ | 2 |
| :---: | :---: |
| 5 | 4 |
| 6 | 8 |
| 10 | 13 |
|  |  |
|  |  |
|  |  |
|  |  |
|  |  |
|  | End |

Algorithm: Compare $1^{\text {st }}$ element of each list, remove the smaller as next element of sorted run

Very efficient! Very few comparisons needed for merge How many comparisons needed to create list of size $N$ ? $\mathrm{O}(\mathrm{N})$ comparisons

## Merge Sort Algorithm: <br> Uses Recursion

## Base case:

- If list of length 0 or 1, done (sorted)


## Otherwise:

- Divide unsorted list of size $M$ into two sublists of size M/2
- Sort each sublist recursively using mergesort
- Merge two sublists back into one sorted list

How to merge two lists into one?



## Merge Sort:

How many comparisons?
$2,5,8,10,11,13,14,14,15,16,19,35,46,66,72,89$


- $\log _{2} \mathrm{~N}$

How many comparisons to create next level? (last run? 2nd-to-last two runs?) - last run: $\mathrm{N}, 2^{\text {nd }}$ to last two runs: $2^{*} \mathrm{~N} / 2$, next: $4^{\text {* }} \mathrm{N} / 4$... always N ! Total comparison

- $N \log _{2} N$




## Quicksort (Qsort) Algorithm: <br> Recursive

Base case: list of size one is sorted by definition

## Otherwise:

Pick an element (pivot) from lis $\dagger$
Reorder:

- All keys < pivo $\rightarrow$ move key before pivo $\dagger$
- All keys > pivot $\rightarrow$ move key after pivo $\dagger$
- Equal values can go either way
- Pivot is now in its final sorted position

Recursively sort (w/ quick sort!) two sub-lists



| Sorting Algorithm | Comparison |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
|  | Selection sort | Insertion sort | Merge Sort | Quick sort |
| Worst case? | $O\left(N^{2}\right)$ | $O\left(N^{2}\right)$ | $O(N \log N)$ | $O\left(N^{2}\right)$ <br> If pick bad <br> pivot |
| Best case? | $O\left(N^{2}\right)$ | $O(N)$ <br> If sorted <br> already | $O(N \log N)$ | $O(N \log N)$ |
| Average case? | $O\left(N^{2}\right)$ | $O\left(N^{2}\right)$ | $O(N \log N)$ | $O(N \log N)$ |




## NOW-Sort: World Record Holder



Sorted 1 million keys (1997)

- Disk-to-disk
- < 2.5 seconds
- 100 machines on network


## Merge sort works well here

- Each machine starts with $1 / 100$ of keys (and data!) on local disk
- Sorts its own keys
- Each sends sorted run of keys (and data!) to destination machine
- After receive all keys, each machine:
- Merge 100 sorted runs



## Announcements

Sorting algorithms

- $O\left(N^{2}\right)$ sorting algorithms
- Selection sort: Find minimum and make next
- Insertion sort: Take next and insert in correct place
- $O(N \log N)$ sorting algorithms (expected, not worst-case)
- Merge sort: Recursively combine sub-lists into larger lists
- Quicksort: Recursively partition list into sub-lists around pivot


## Announcements

- Homework 5: Congrats to Fong Lor, Jake Hilborn, and Kameko Blair
- Homework 6: No Extra Credit
- email me if you think your project showed creativity, may re-evaluate
- Homework 7: Due Friday - No Programming

Explore Google Trends, Understand basic sorting algorithms, Reflect on Technology and Education (make sure you can watch Friday video)

