Instructor Notes

Nooks
Improving the Reliability of Commodity Operating Systems
Swift, Bershad, and Levy
SOSP'03

1. What was the motivation for Nooks?

- Reliability is important - cost of failures is high
- Extensions (modules / device drivers) account for 70% of Linux.
  - written by less experienced programmers
  - yet reside in kernel space
- Extensions cause most OS failures (85% Win XP)

→ Treat extensions differently than rest of Kernel

2. What were the design principles and goals of Nooks?

Principles:

1) Fault resistance, not fault tolerance

(don't handle all)

2) Mistakes, not abuse

⇒ Better performance + better reliability

Goals:

1) Isolation: (detect ~~kernel~~ extension problem before infects rest of kernel)

2) Recovery: automatic recovery to permit applications to continue

3) Backward Compatibility: Existing, with minimal changes

Reviewer: Does Eval show they met their goals?

3. What are the components of the Nooks Isolation Manager (NIM)?

Layer ~~code~~ between OS Kernel + Extensions:

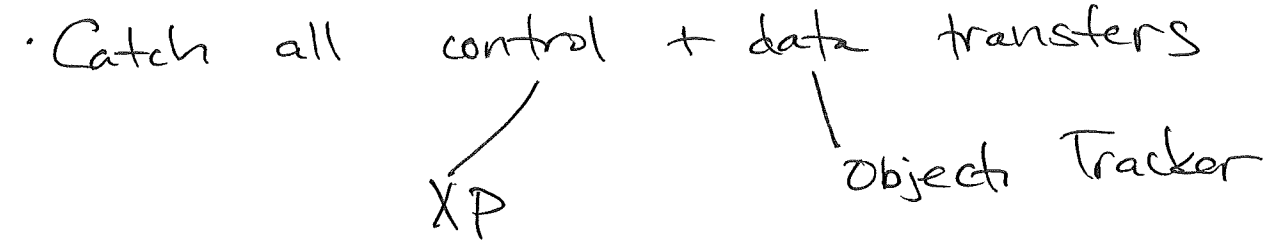
1) Isolation
2) Interposition
3) Object Tracking
4) Recovery

Should be transparent

4. Why is Isolation needed? At a high level, how is it provided?

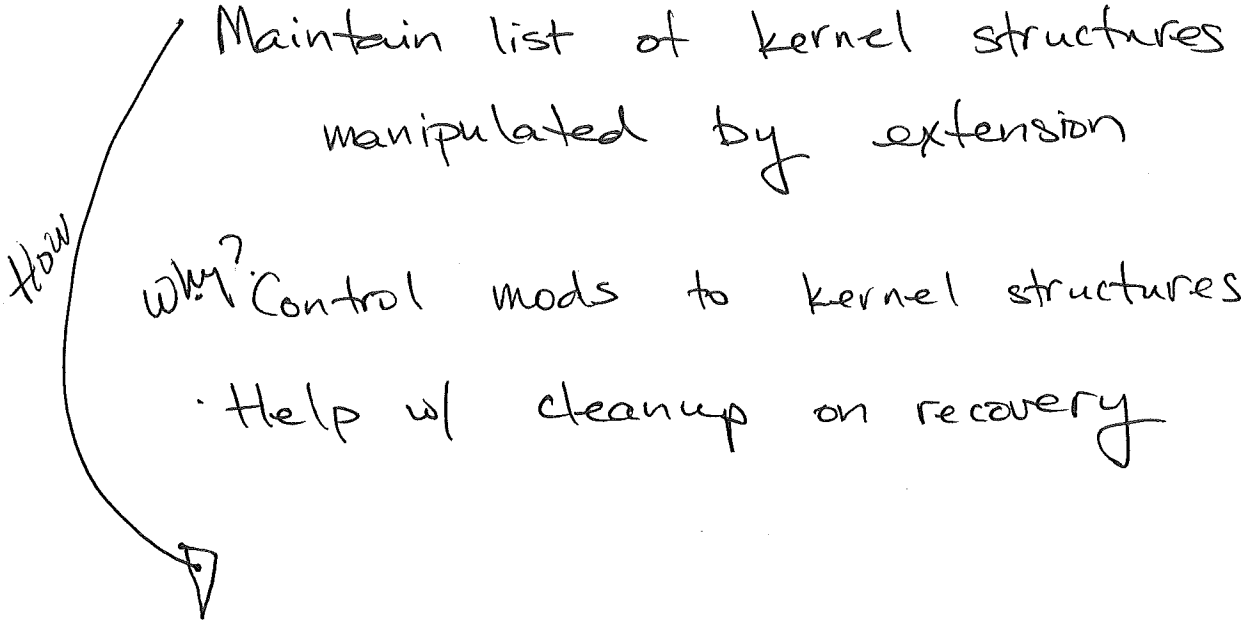
- Prevent damage to kernel or other extensions
- Lightweight protection domain:
  - run @ same processor privelege, but reduced memory access
- XPC - eXtension procedure call
  - to jmp between

5. Why is Interposition needed? At a high level, how is it provided?



How? Wrappers (stubs)
of kernel's extension API
+
extension entry points

6. Why is Object Tracking needed? At a high level, how is it provided?



7. Why is Recovery needed? At a high level, how is it provided?

- Need to be able to detect problems, restart extension

- S/W fault:
  - call routine w/ wrong args
  - too many resources

- H/W fault:
  - read/write pages w/o permission

8. How much work was it to implement Nooks in Linux 2.4.18?

Significant!

700 kernel functions, 650 extension-entry functions

18 developer months

22,266 lines of code

14,396 wrapper lines!

924 linux kernel changes

9. For memory management, what memory rights does an extension have? What memory rights does the main Linux kernel have? How is this protection provided?

Extension: r/w own domain
r of kernel

Kernel: r/w of all

Each extension has own page tables
(changing protection domains → changing page tables)

10. Why is a synchronized copy of the kernel page table needed for each domain? Are there any implications of this? Why does the Nooks design prevent bugs but not malicious extensions? What is the performance cost of switching between lightweight protection domains?

Each extension
- Needs to know about pages in kernel address space

- Required changes to Linux when modify kernel page changes
(& costly w/ many extensions??)

- Nothing to prevent extension from modifying hw page table base register

- TLB flush on every switch between domains

11. How is control between an extension and kernel domain handled with XPC? What is the purpose of a deferred call?

nooks_driver_call: function ptr, args, protection domain _kernel_

- save context, find stack

~~*~~change page tables to target domain

call function

(wrappers around XPC provide transparency and do the checking of parameters)

Deferred: Used for batching multiple XPC calls; expensive to switch btwn domains frequently

12. Did the Linux kernel need to be modified to support isolation?

Yes

* 1) maintain coherency of kernel read p.t.

2) handle exceptions in Nooks domains

3) global variable for task ptr

13. In Linux, extensions sometimes directly access global data structures. How is this handled? When is XPC used? When is it not? How should one determine which approach to use?

- If just read, okay
- Write → Replace macros + inline functions w/ wrapped calls
- If direct in extension — have to find all of these!

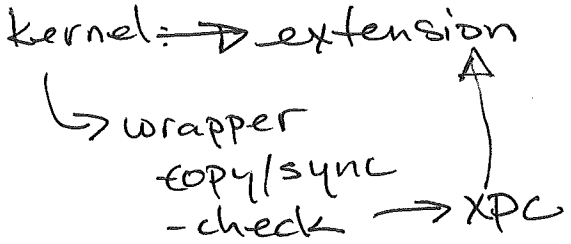
XPC: Not perf critical

Shadow copy of kernel object: perf critical in extension domain

— sync before + after multiple calls

Work, work, work...

14. What does a wrapper need to do? What is the difference between call-by-value-result and call-by-reference?



- checking of parameters
  - very context specific!
  (interact w/ object tracker + mem management)

* extension → kernel
↳ XPC → wrapper
. copy/syn
(runs in kernel mode)

- call-by-value-result
  arg V → copy to L during call
  → copy back to V when done

- single-threaded - same semantics as call by ref
  - can do checks on values/results when copy in + out

15. How much work is it to write a wrapper?

- Lots?
- By hand
- Requires knowledge of parameter use
+ Reusable by all extensions using same interface

16. What must the Object Tracker track?

- all kernel objects manipulated by extensions
- 43 different types - inspected every ~~[illegible]~~ interface to find set

1) Records addresses of objects used by extension
   - record in single XPC call table or per-protection-domain hash table if long-lived

2) ~~For~~ objects modified, track association between extension + kernel version

- Must know lifetime of objects

17. What does the Nooks recovery manager do?

1) Detects failure

2). Unwind executing tasks

order?
- Unload extension
- Release all resources (including refs to objs)
- Reload + restart

Does not ensure that apps can continue to use driver! (Swift's next paper)

18. What are some limitations of Nooks?

- can't prevent priv. instr.
- can't prevent infinite loops
- Check parameters, but not completely
- just kill + restart

19. Does Nooks meet its goals of Isolation, Recovery, and Backward Compatibility?

- Isolation: Prevent errors from crashing system?
  Fig 6 - positive

  Fig 7: Not so good w/ non-fatal errors
  - can't detect problem

- Recovery? can an app run??
  - If data is damaged before detection, problem
  - See Table 3 of Shadow Driver paper

- Backward Compatibility
  - some Linux changes
  - would need more extensions to know if have all wrappers...

20. Is the performance of Nooks acceptable?

Table 4.

ok if low RPC rate or low CPU util

not good for complex extension (kHTTPd) or VFAT

21. Conclusions?

+ Good motivation ← driver reliability important
different structure needed

+ Lightweight protection domain

- Lots of work by hand
- Performance impact of crossings
- Doesn't automatically recover for applications (shadow drivers)