



# An Analysis of Internet Content Delivery Systems

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# Outline

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- Introduction
- Overview
- Data Collection
- Data Summary
- Caching
- Conclusion



# Introduction

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- As complexity of data has increased, content delivery mechanisms have changed
- New systems, especially peer-to-peer, have affected the nature of Internet traffic
- Goal: Determine the effect of the latest types of content delivery systems on Internet traffic
  - Web vs. CDN vs. P2P



# Introduction

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## *What the authors did:*

- Traced all incoming and outgoing traffic at the University of Washington (pop. 60,000) over nine days.
- Totals were 500 million transactions and 20 terabytes of data.
- Separated traffic between Akamai, Kazaa, Gnutella, and regular Web traffic (all use HTTP).



# Overview of Content Delivery Systems

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## *World-Wide Web:*

- Most objects are 5-10KB
- Heavy-tailed distribution; some objects are very large while most of objects are small
- Increasing number of dynamic objects and total objects is reducing cache effectiveness
- The best existing or proposed caches have hit rates in the 40-50% range



# Overview of Content Delivery Systems

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## *Content Delivery Networks:*

- Servers distributed throughout the Internet with replicated content
- Largest CDNs have servers at every major point of traffic
- Invisible to users
- CDNs generally carry only **static** and **secondary content**



# Overview of Content Delivery Systems

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## *Peer-to-Peer Systems:*

- Collaborative content exchange systems
- Nodes act as both servers and clients
- Virtually all nodes are operated by end-users
- Most nodes are on slow, unreliable connections, unlike most content delivery systems



# Overview of Content Delivery Systems

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## *Peer-to-Peer Systems:*

- Systems use one of three methods for content searching
  - Search requests are made directly between peers
  - Central server maintains a complete index
  - Hybrid system: some peers are “supernodes” and index content to nearby peers
- After content has been located, transfer commences between peer and one or more peers that host the content





# Data Collection

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## *Methodology*

- Passive network monitoring used to collect traces of traffic data
- Switches connected to all campus backbones send data via a monitoring port to a monitoring host
- Kernel-level packet filter delivers packets to a user process for data collection



# Data Collection

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## *Distinguishing Traffic Types*

- HTTP traffic
  - WWW and Akamai on ports 80, 8080, 443 (SSL)
  - Gnutella on ports 6346 and 6347
  - Kazaa on port 1214
  - Mystery Traffic
- Non-HTTP traffic
  - Any other TCP traffic: NNTP, SMTP and etc.
  - Kazaa and Gnutella search traffic
- Did not include UDP traffic in the analysis.
  - 97% of UW traffic is TCP

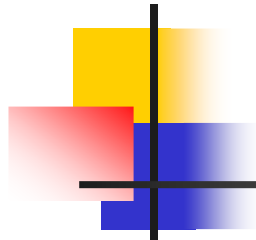


# Data Collection

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## *Traceability of P2P Traffic*

- Kazaa's architecture is proprietary, but some inferences about its operation are made by observing traffic
- Kazaa peers can download objects in fragments, which complicates tracing
- Download requests from external peers seen in their trace are often for fragments rather than entire object



# Data Summary

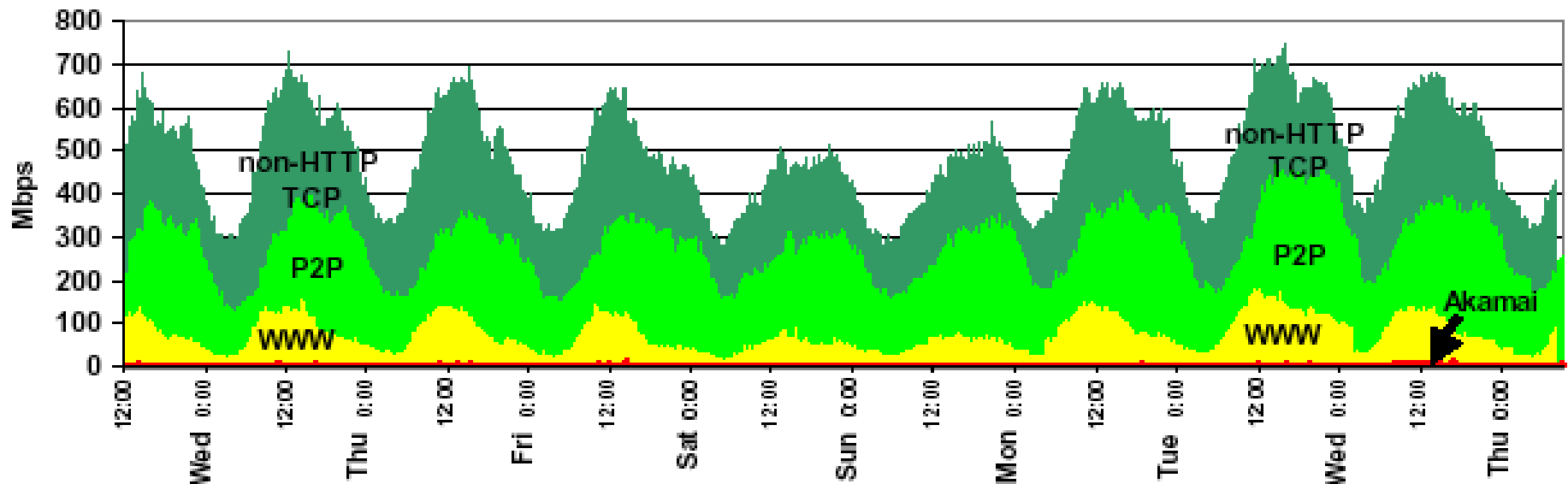
## *HTTP trace summary statistics*

	WWW		Akamai		Kazaa		Gnutella	
	inbound	outbound	inbound	outbound	inbound	outbound	inbound	outbound
HTTP transactions	329,072,253	73,001,891	33,486,508	N/A	11,140,861	19,190,902	1,576,048	1,321,999
unique objects	72,818,997	3,412,647	1,558,852	N/A	111,437	166,442	5,274	2,092
clients	39,285	1,231,308	34,801	N/A	4,644	611,005	2,151	25,336
servers	403,087	9,821	350	N/A	281,026	3,888	20,582	412
bytes transferred	1.51 TB	3.02 TB	64.79 GB	N/A	1.78 TB	13.57 TB	28.76 GB	60.38 GB
median object size	1,976 B	4,646 B	2,001 B	N/A	3.75 MB	3.67 MB	4.26 MB	4.08 MB
mean object size	24,687 B	82,385 B	12,936 B	N/A	27.78 MB	19.07 MB	19.16 MB	9.78 MB

- Compare outbound vs. inbound data
- Compare the volume of data transfer for each system
- Compare the object sizes

# Data Summary

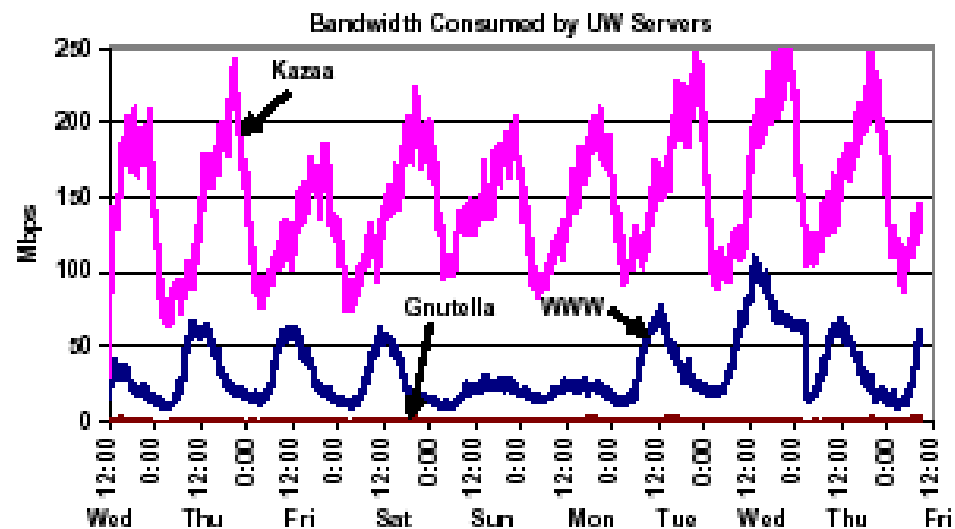
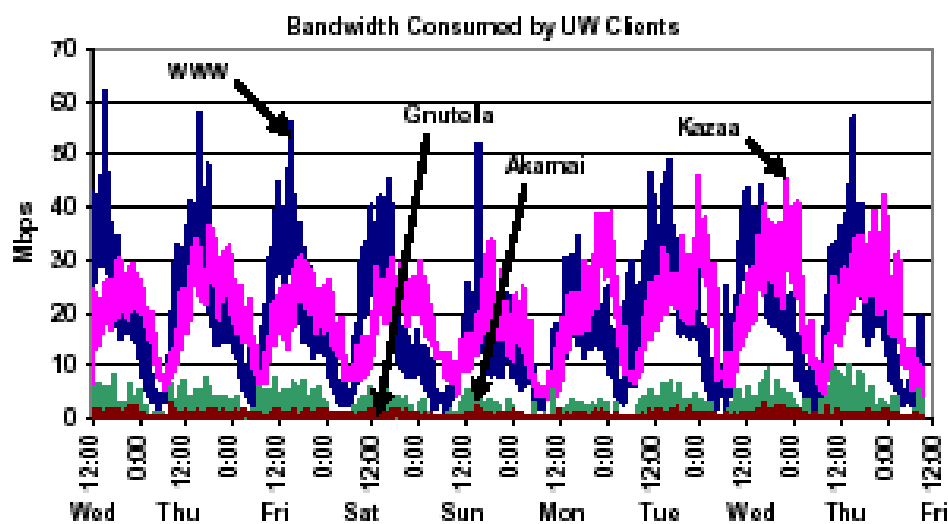
## *TCP bandwidth*



- Notice the pure volume of peer to peer traffic, compared to all TCP traffic (43% of all TCP)

# Data Summary

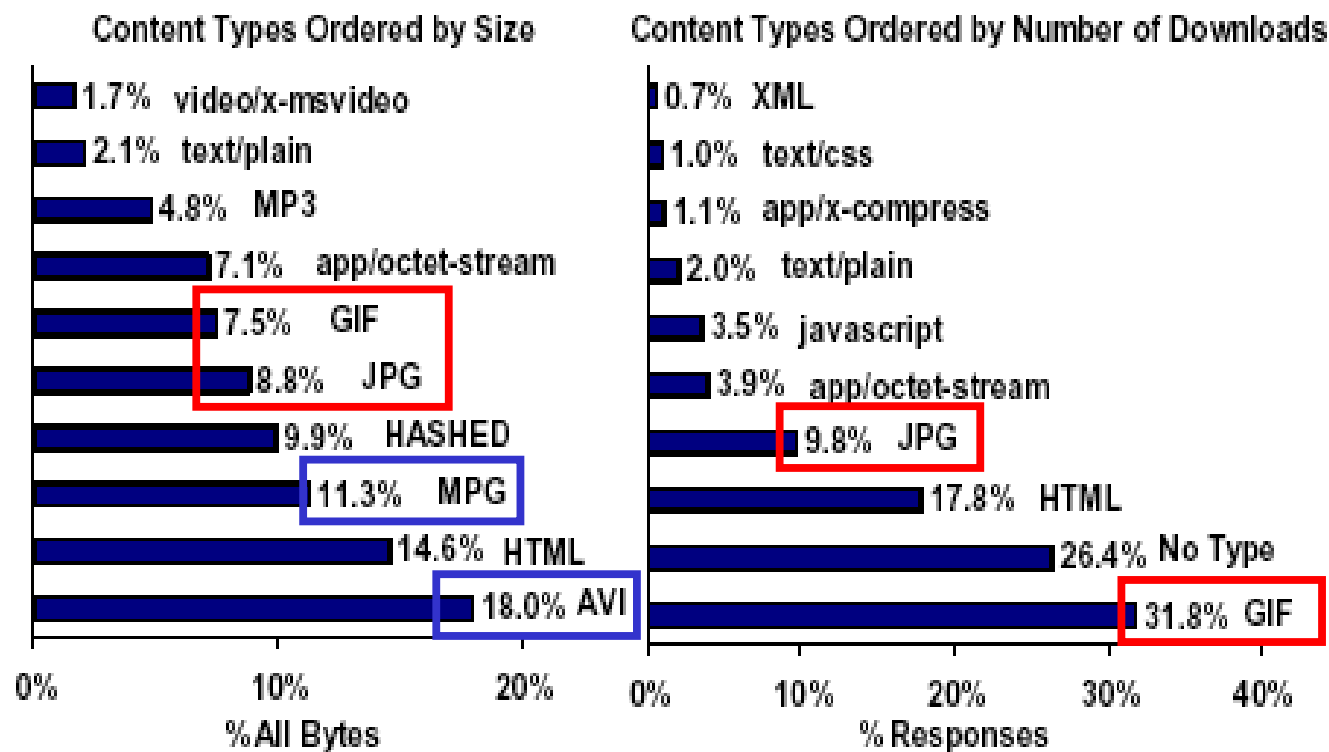
## *Client and server TCP bandwidth*



- Notice domination of Kazaa over outgoing bandwidth

# Data Summary

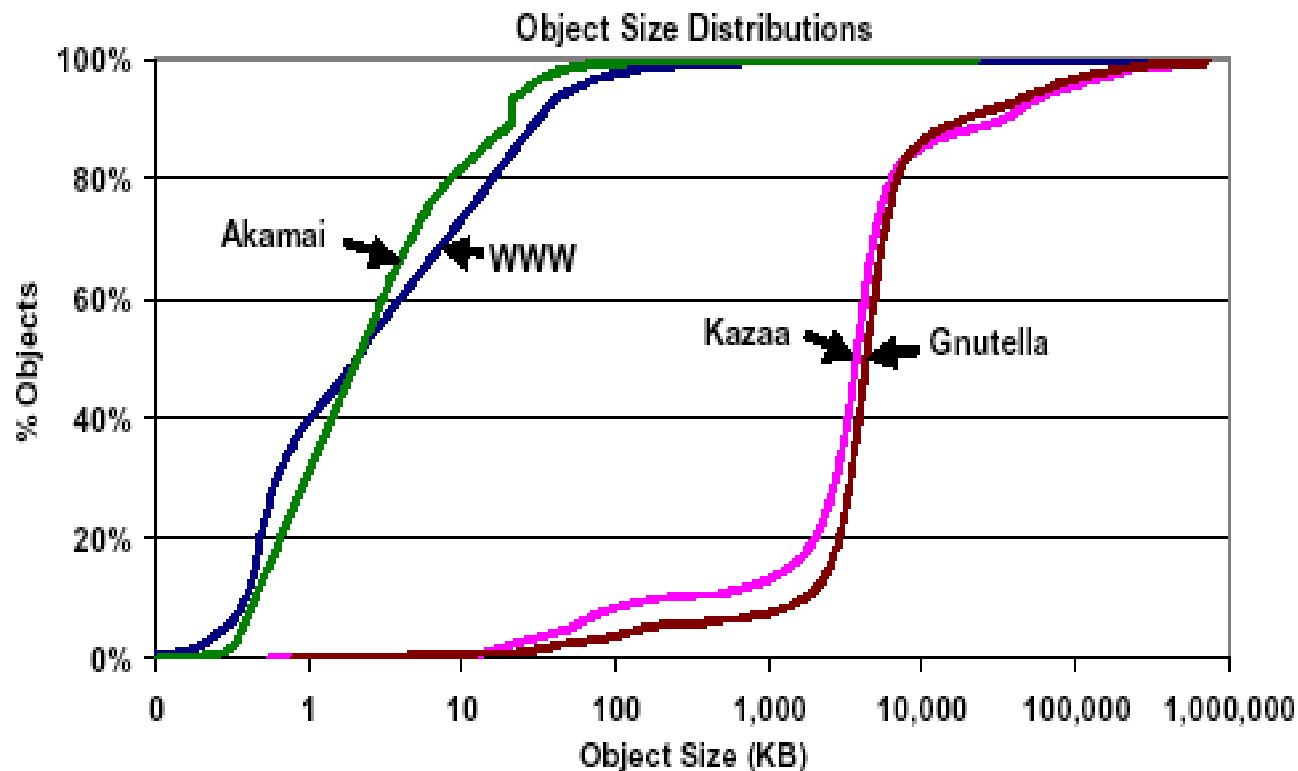
## *Content types downloaded*



- Hashed refers to unidentifiable Kazaa transfers

# Data Summary

## *Object size distribution*



Median Object size:

WWW = 2KB

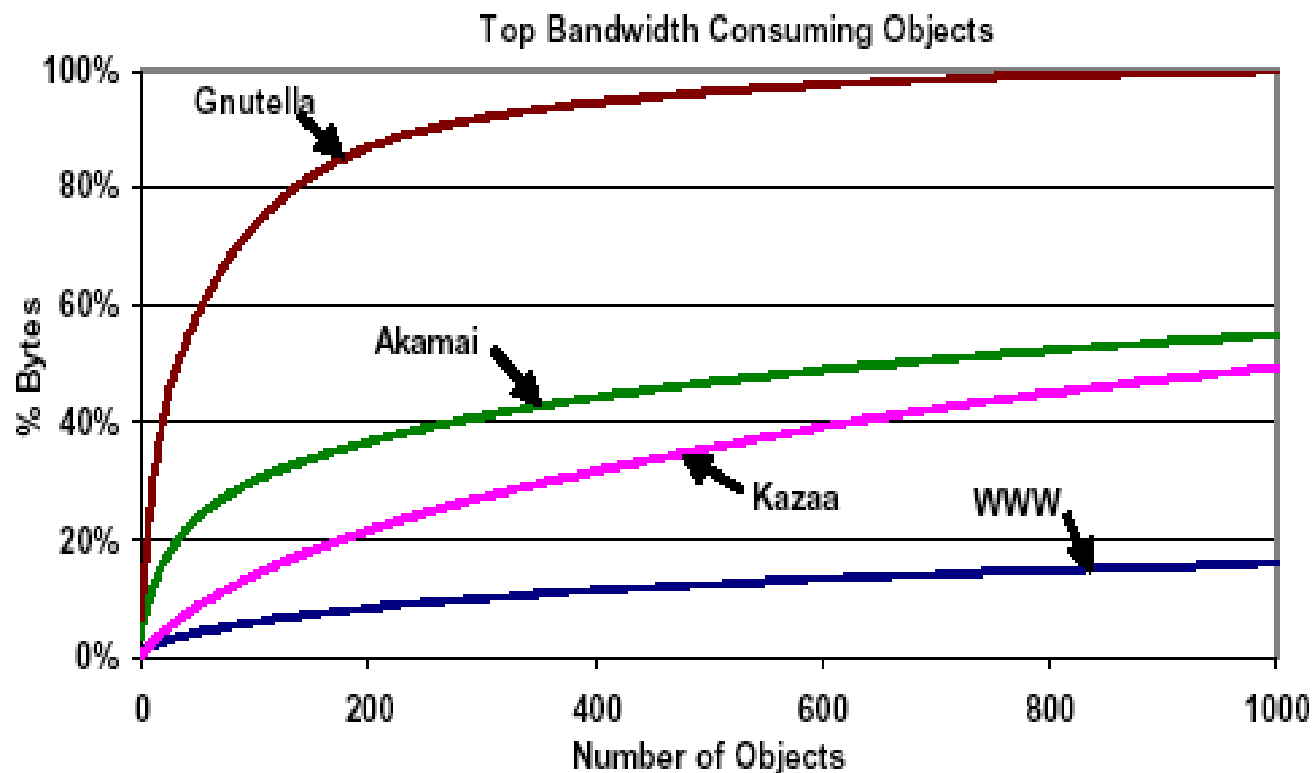
P2P = 4MB

- Notice x-axis is logarithmic

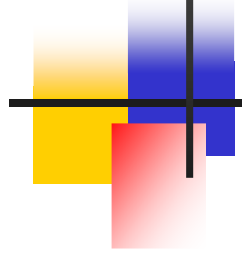


# Data Summary

## *Top bandwidth consuming objects*

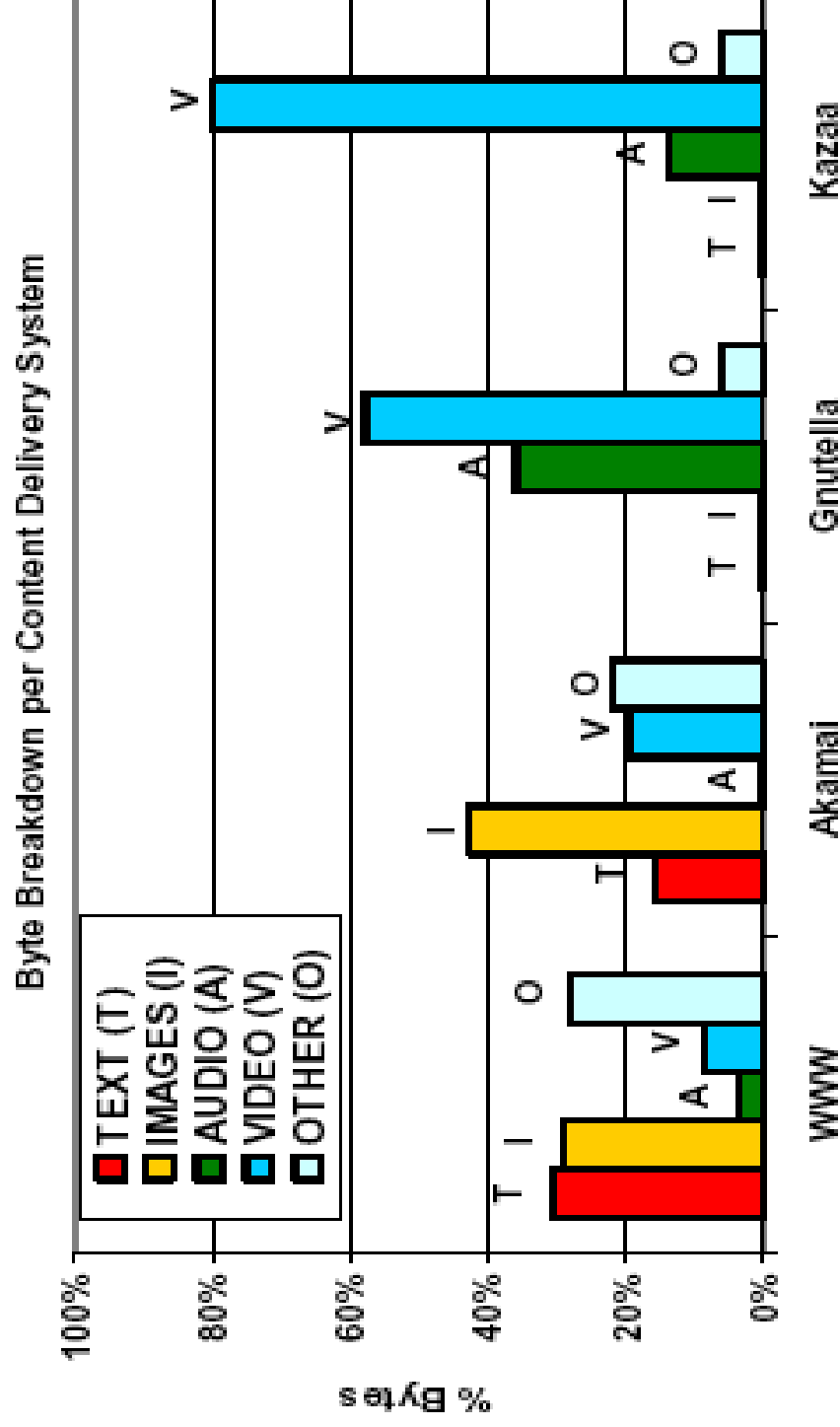


- Why are 200 objects 90% of Gnutella bandwidth?



# Data Summary

*Downloaded bytes by object type*





# Data Summary

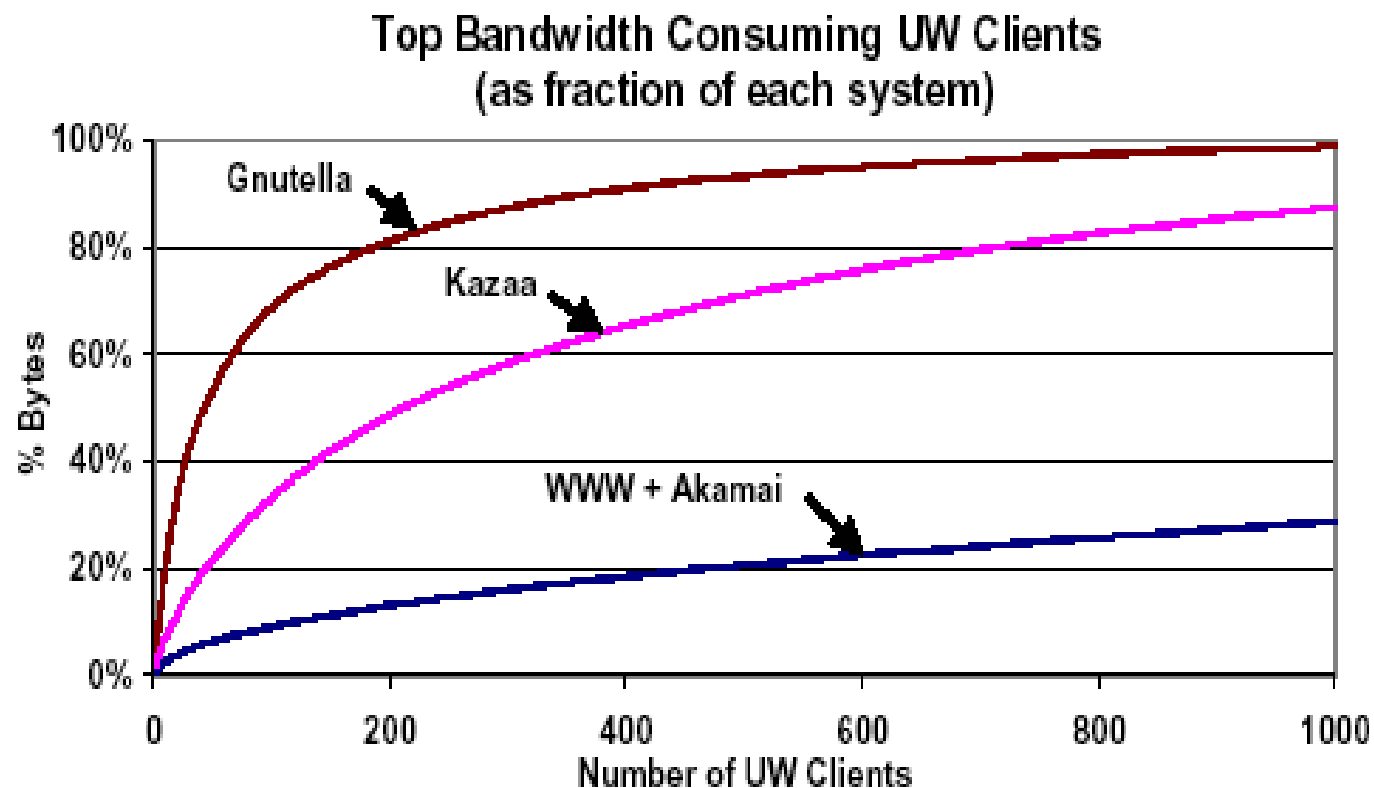
## *Top bandwidth consuming objects*

	WWW (inbound)			Akamai			Kazaa (inbound)				Kazaa (outbound)			
	object size (MB)	GB consumed	# requests	object size (MB)	GB consumed	# requests	object size (MB)	GB consumed	# clients	# servers	object size (MB)	GB consumed	# clients	# servers
1	0.009	12.29	1,412,104	22.37	4.72	218	694.39	8.14	20	164	696.92	119.01	397	1
2	0.002	6.88	3,007,720	0.07	2.37	45,399	702.17	6.44	14	91	699.28	110.56	1000	4
3	333	6.83	21	0.11	1.64	68,202	690.34	6.13	22	83	699.09	78.76	390	10
4	0.005	6.82	1,412,105	9.16	1.59	2,222	775.66	5.67	16	105	700.86	73.30	558	2
5	2.23	3.17	1,457	13.78	1.31	107	698.13	4.70	14	74	634.25	64.99	540	1
6	0.02	2.69	126,625	82.03	1.14	23	712.97	4.69	17	120	690.34	64.97	533	10
7	0.02	2.69	122,453	21.05	1.01	50	715.61	4.49	13	71	690.34	54.90	447	16
8	0.03	1.92	56,842	16.75	1.00	324	579.13	4.30	14	158	699.75	49.47	171	2
9	0.01	1.91	143,780	15.84	0.95	68	617.99	4.12	12	94	696.42	43.35	384	14
10	0.04	1.86	47,676	15.12	0.80	57	167.18	3.83	39	247	662.69	42.28	151	2

- Notice the 9KB object, it alone accounts for 1.9% of WWW traffic.

# Data Summary

## *Top bandwidth consuming Clients*

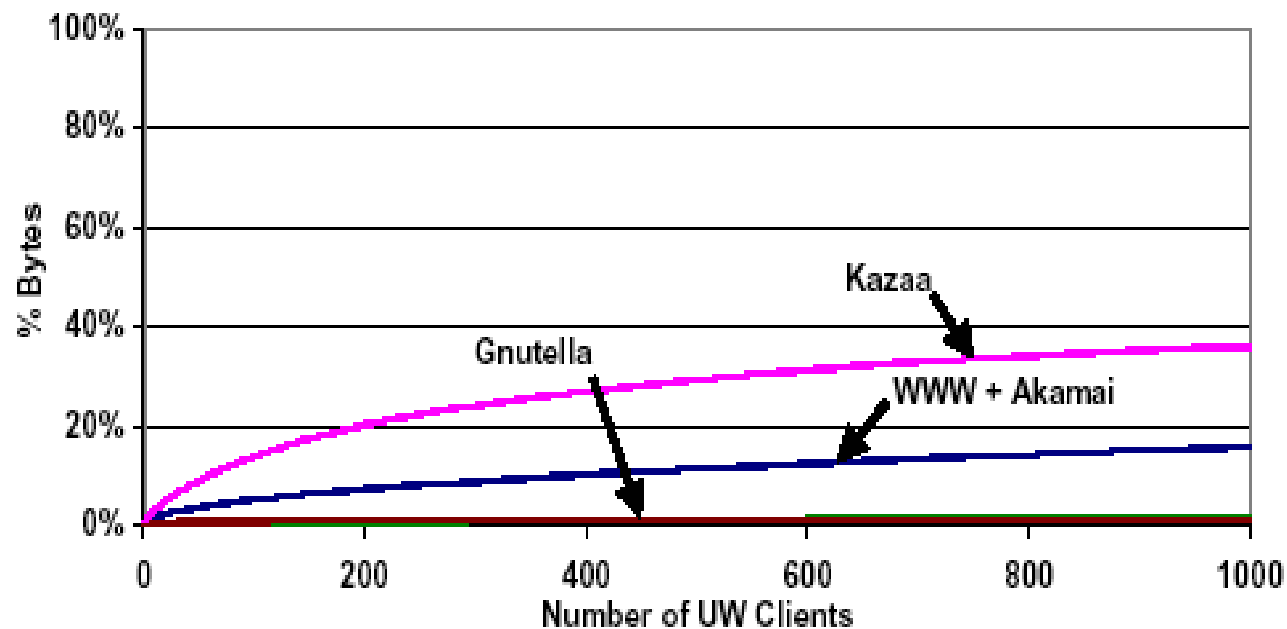


- 200 Clients consume 80% of Gnutella bandwidth

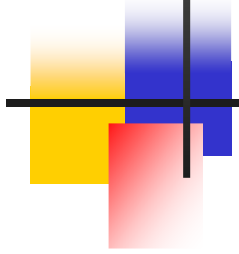
# Data Summary

## *Top bandwidth consuming Clients*

Top Bandwidth Consuming UW Clients  
(as fraction of total HTTP)

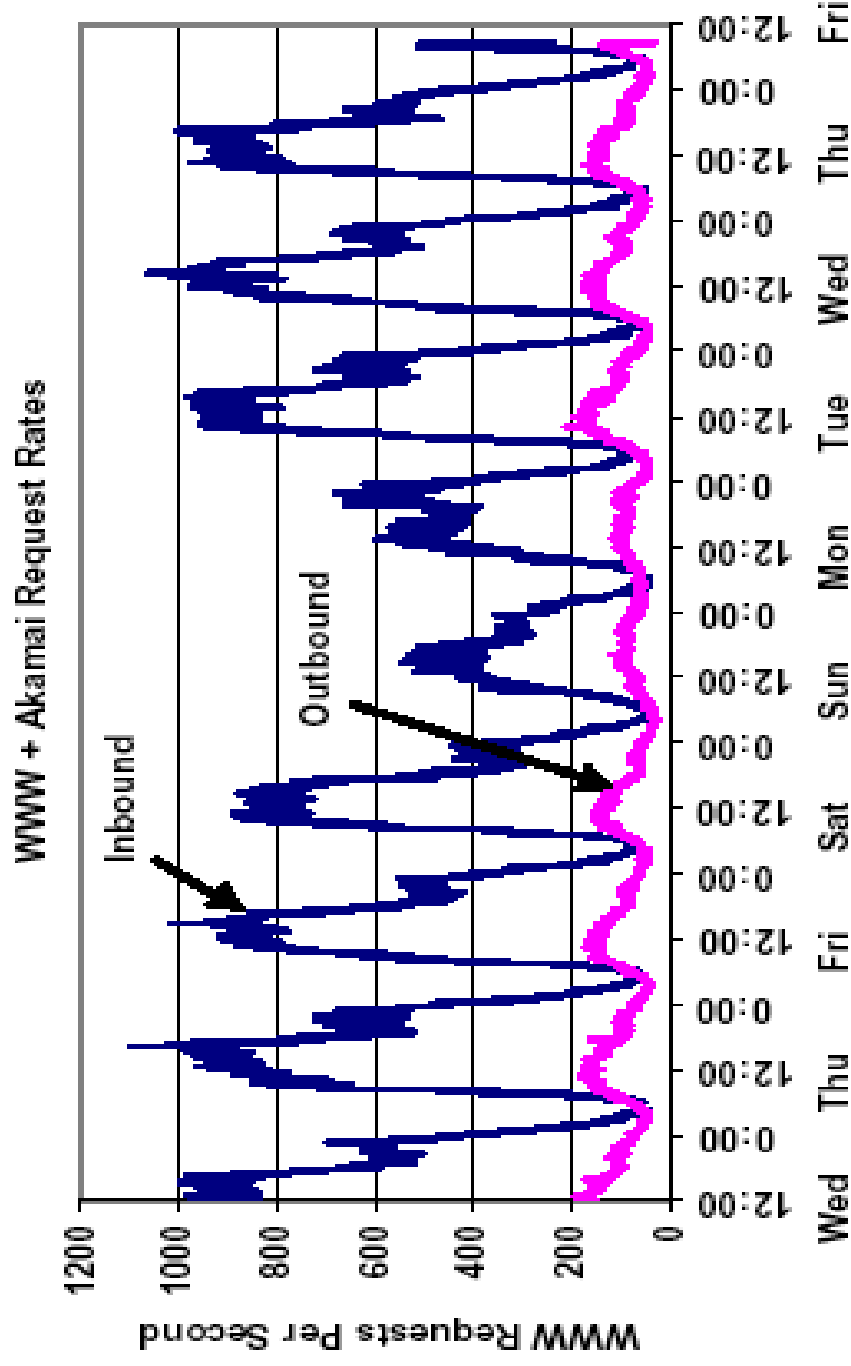


- 200 Kazaa clients consume 20% of all HTTP bandwidth

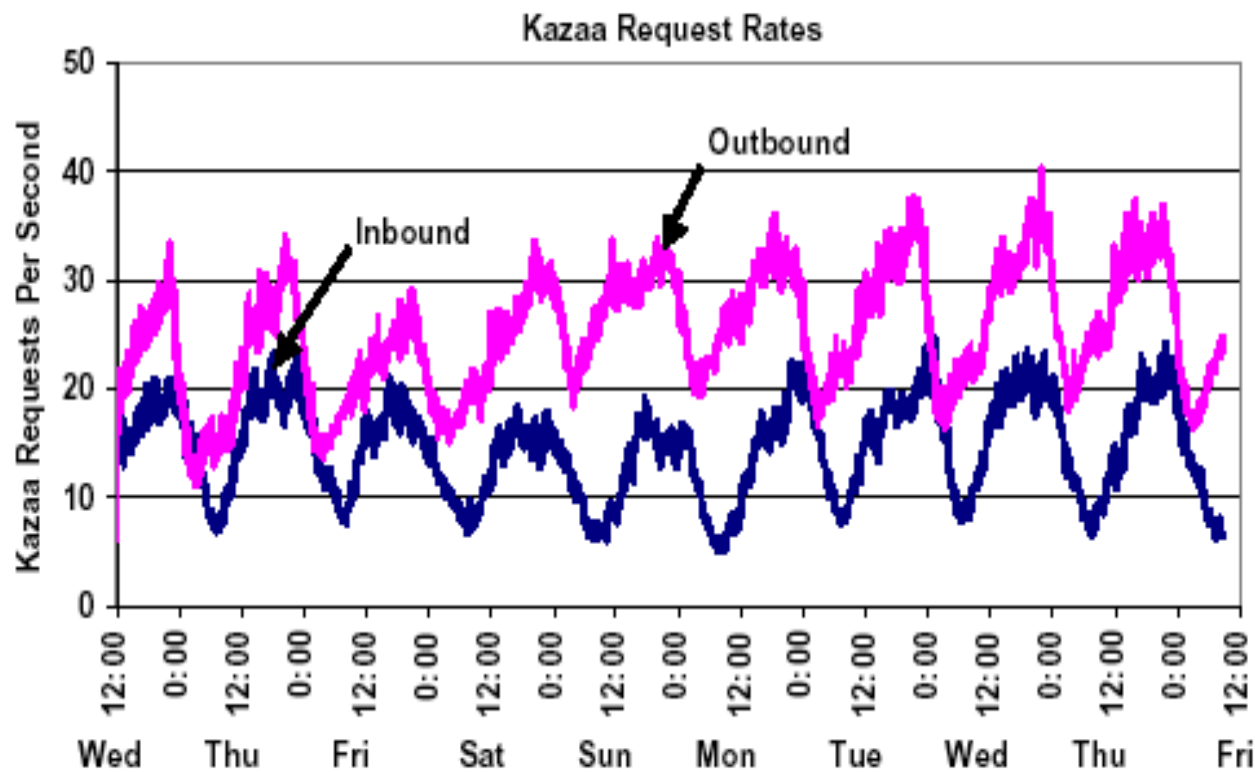


# Data Summary

*Request rates over time*



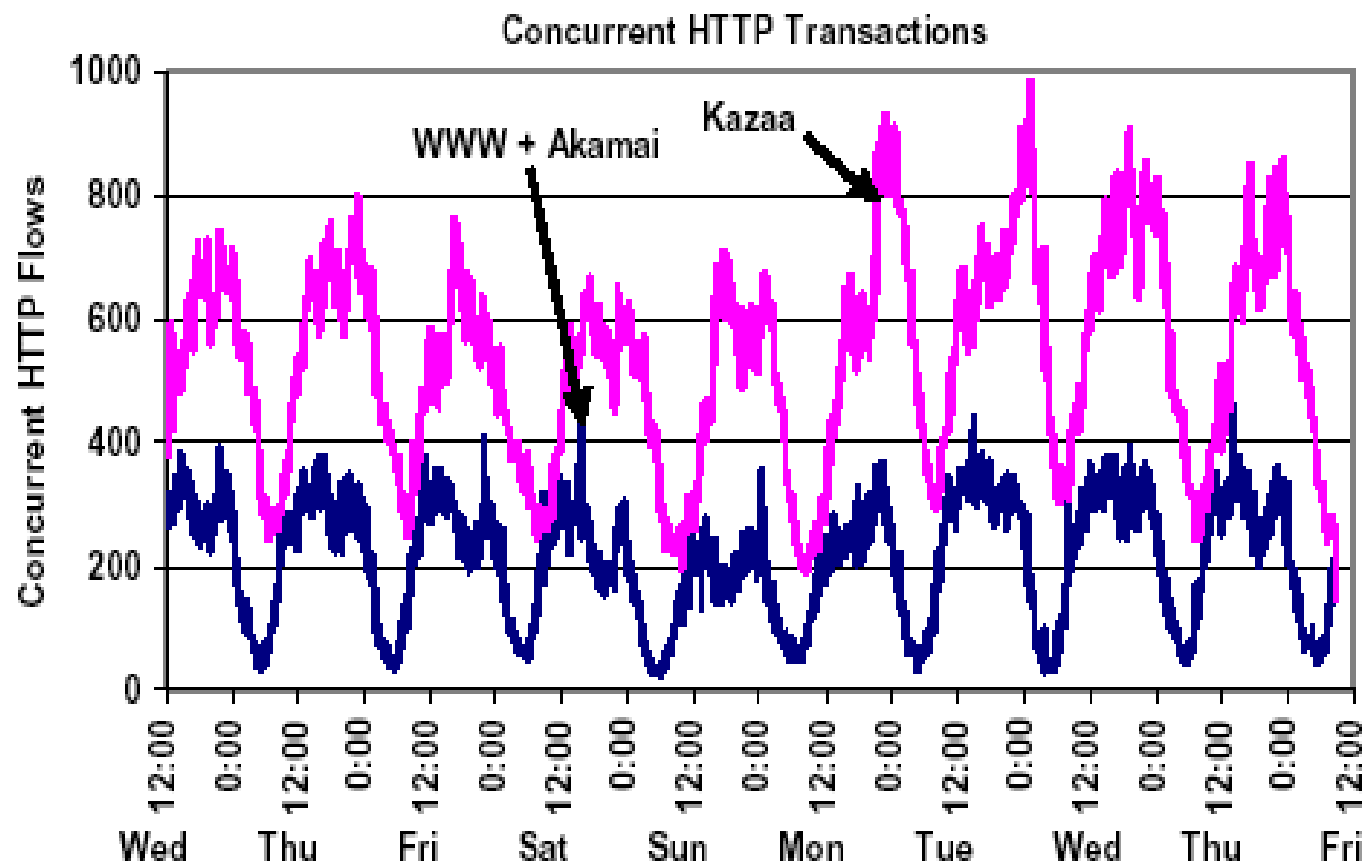
## Request rates over time



- Notice significantly fewer total requests per second than WWW and Akamai

# Data Summary

## *Concurrent HTTP transactions*



Medium duration  
of a request:

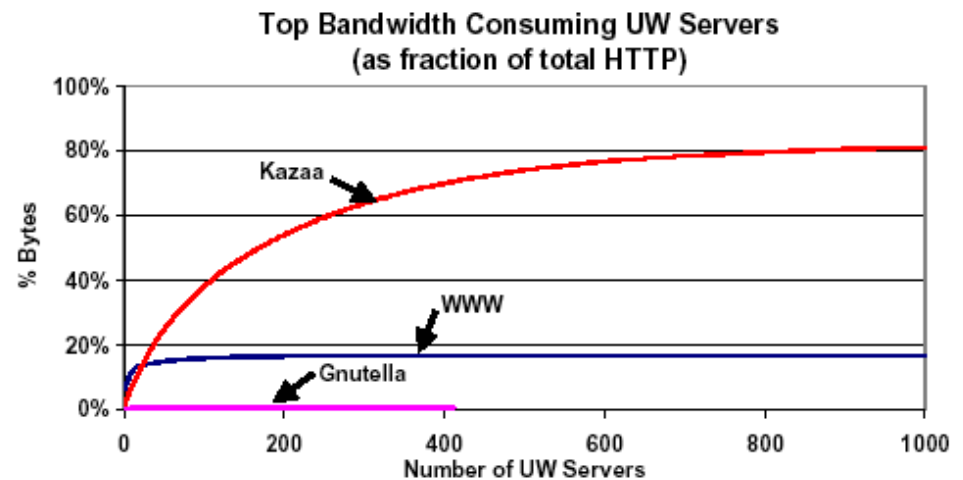
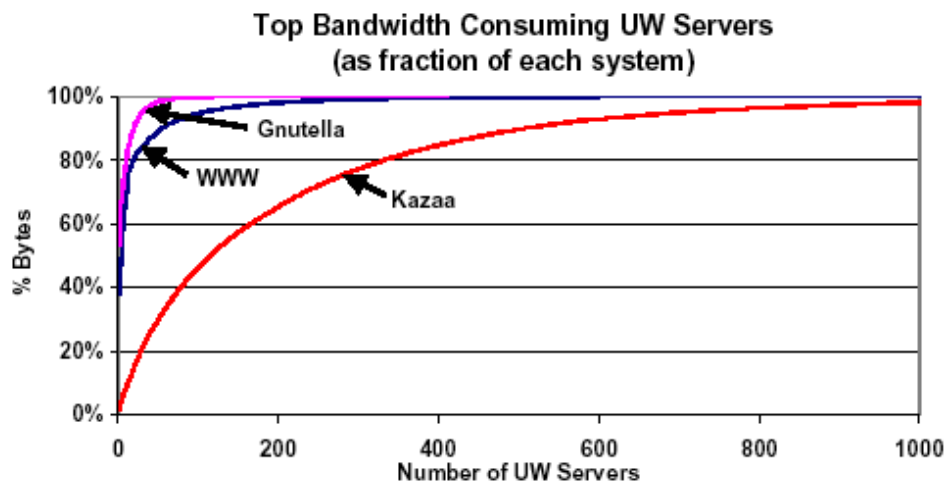
WWW = 120 ms

Kazaa = 130 sec



# Data Summary

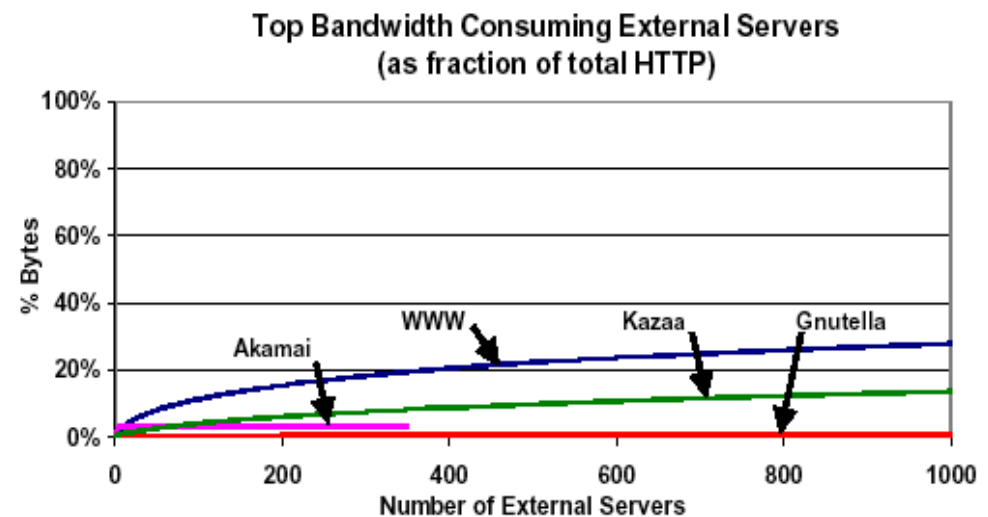
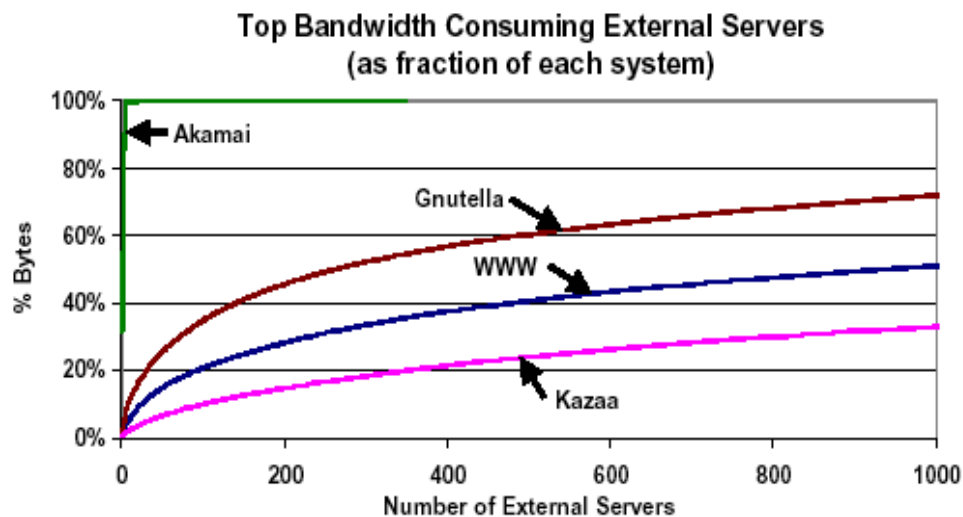
## *Top UW-internal bandwidth producing servers*



- Notice there are only 410 Gnutella servers at UW
- A very small amount of WWW servers are responsible for 40% of outgoing bytes.

# Data Summary

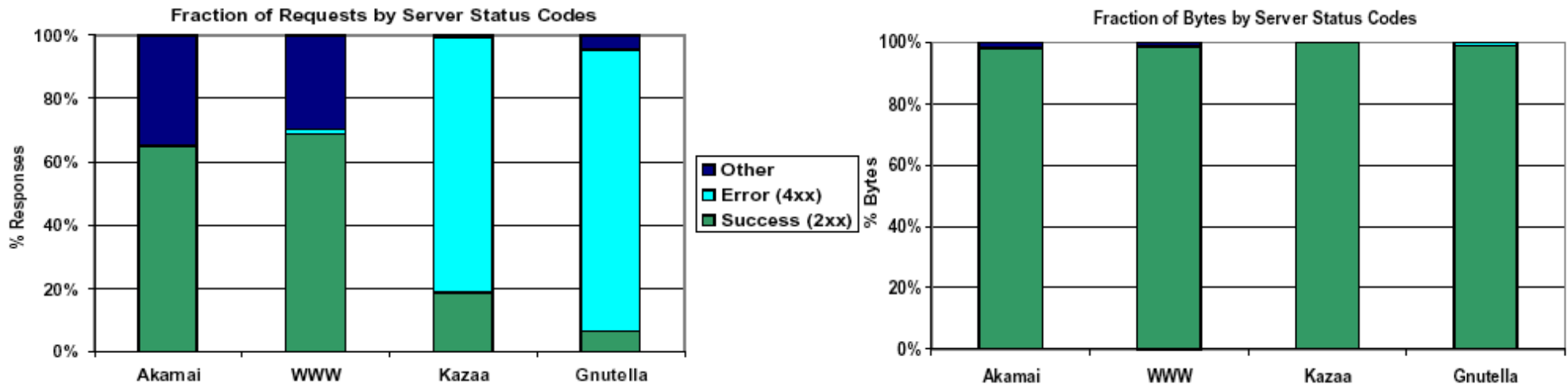
## *Top UW-external bandwidth producing servers*



- Of about 350 Akamai servers, only about 10 account for the majority of the bandwidth

# Data Summary

## *Server status codes*



- Why does Akamai have no apparent error rate?
- Why are Kazaa and Gnutella error rates high?



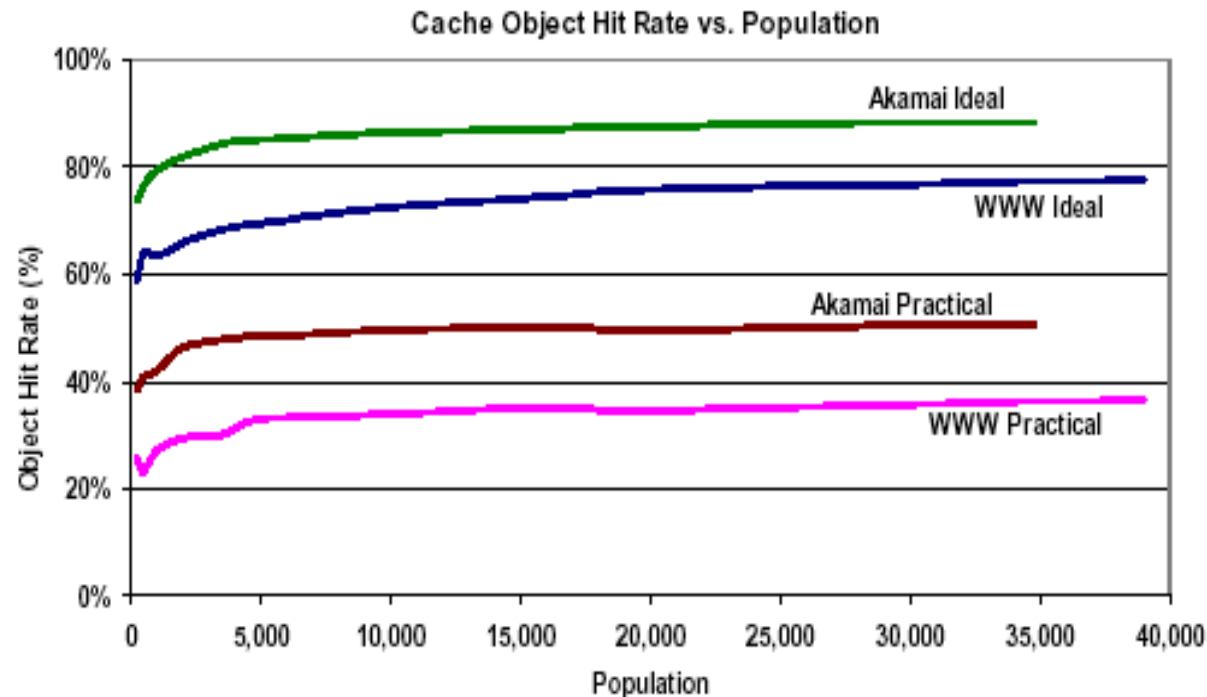
# Caching

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- The paper addresses the following questions:
  - Is the use of Akamai as an external CDN justified for content providers, given the existence of local caches?
  - Can caching be applied to Peer to Peer networks effectively?

# Caching

## *Akamai vs. WWW*



- Notice in both ideal and practical cases, Akamai objects can be more effectively cached than normal web traffic



# Caching

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- Conclusion drawn:
  - Given the cache-ability of Akamai objects, the need for specialized content delivery networks like Akamai is greatly reduced if proxy caches are widely deployed enough
- Akamai objects are frequently requested after they expire from the cache. If properly marked with a longer time-to-live, the number of cache hits would increase significantly



# Caching

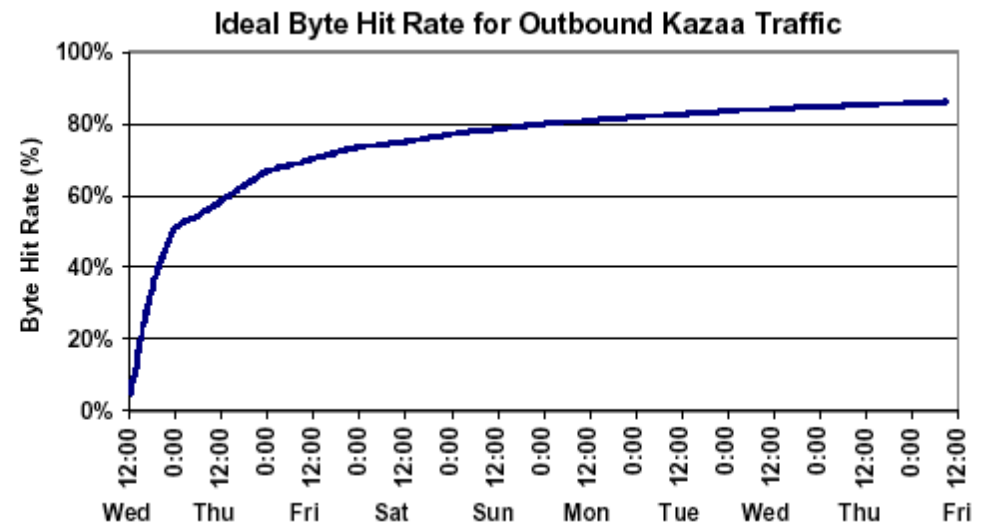
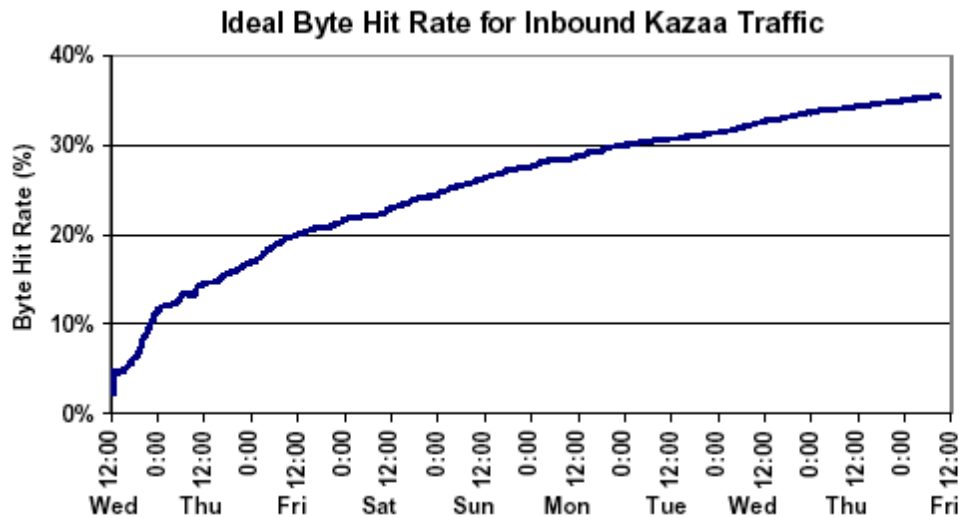
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## *Peer-to-Peer Caching*

- Authors designed a cache simulator for Kazaa
- The cache stores objects as 32K blocks
  - This differs from normal caches because it is possible that only parts of an object are cached, not necessarily a complete object.

# Caching

## *Peer-to-Peer Caching*



- After 2 days, three quarters (3/4) of outbound Kazaa traffic could come directly from the cache
- Why is inbound much lower?





# Caching

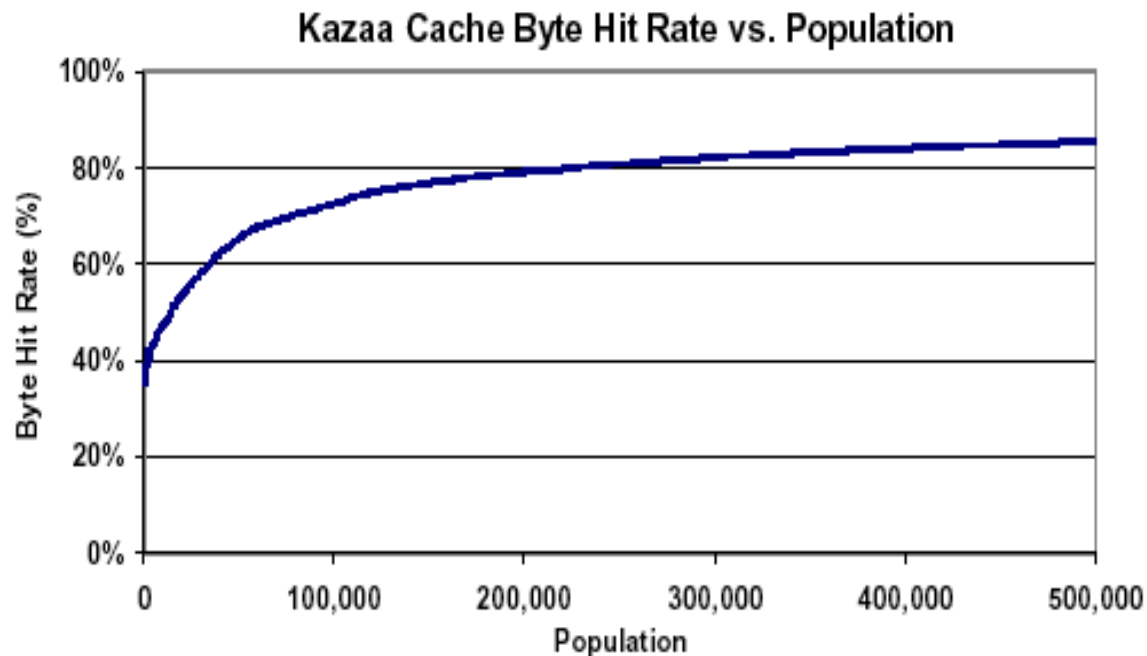
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## *Peer-to-Peer Caching*

- The top 300 bandwidth-consuming objects on Kazaa have an average size of 614MB. Why?
- A cache that stores the top 300 objects would have a hit rate of 38%

# Caching

## *Kazaa cache byte hit rate vs. population*



- Even for a small population, the hit rate is very good, starting at 40%.



# Caching

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## *Peer to Peer Conclusion*

- A cache for a system like Kazaa would be overwhelmingly beneficial, especially by cache standards
- This begs the question: Why are there no Kazaa caches?



## Conclusion

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- Peer to peer traffic accounts for the majority of HTTP traffic, despite relatively small amount of peer to peer clients vs. web clients.
- Unlike web traffic, peer to peer traffic creates a heavy bandwidth flow in both directions
- The majority of peer to peer bandwidth is taken up by a small number of very large objects
- Peer to peer caching would be greatly beneficial