Reading Assignment

- MULTILISP: a language for concurrent symbolic computation, by Robert H. Halstead (linked from class web page)
Lazy Evaluation

Lazy evaluation is sometimes called “call by need.” We do an evaluation when a value is used; not when it is defined.

Scheme provides for lazy evaluation:

(delay expression)

Evaluation of expression is delayed. The call returns a “promise” that is essentially a lambda expression.

(force promise)

A promise, created by a call to delay, is evaluated. If the promise has already been evaluated, the value computed by the first call to force is reused.
Example:

Though \texttt{and} is predefined, writing a correct implementation for it is a bit tricky.

The obvious program

\begin{verbatim}
(define (and A B)
  (if A
    B
    #f)
)
\end{verbatim}

is incorrect since \texttt{B} is always evaluated whether it is needed or not. In a call like

\begin{verbatim}
(and (not (= i 0)) (> (/ j i) 10))
\end{verbatim}

unnecessary evaluation might be fatal.
An argument to a function is strict if it is always used. Non-strict arguments may cause failure if evaluated unnecessarily.

With lazy evaluation, we can define a more robust and function:

```scheme
(define (and A B)
  (if A
      (force B)
       #f
    )
)
```

This is called as:

```scheme
(and (not (= i 0))
    (delay (> (/ j i) 10)))
```

Note that making the programmer remember to add a call to delay is unappealing.
Delayed evaluation also allows us a neat implementation of suspensions. The following definition of an infinite list of integers clearly fails:

(define (inflist i)
  (cons i (inflist (+ i 1))))

But with use of delays we get the desired effect in finite time:

(define (inflist i)
  (cons i
    (delay (inflist (+ i 1))))))

Now a call like (inflist 1) creates

1 promise for (inflist 2)
We need to slightly modify how we explore suspended infinite lists. We can’t redefine \texttt{car} and \texttt{cdr} as these are far too fundamental to tamper with.

Instead we’ll define \texttt{head} and \texttt{tail} to do much the same job:

\begin{verbatim}
(define head car)
(define (tail L)
  (force (cdr L)))
\end{verbatim}

\texttt{head} looks at \texttt{car} values which are fully evaluated.

\texttt{tail} forces one level of evaluation of a delayed \texttt{cdr} and saves the evaluated value in place of the suspension (promise).
Given

\[(\text{define } IL \text{ (inflist 1)})\]

\[(\text{head (tail } IL))\text{ returns 2 and expands } IL\text{ into}\]

\[
\begin{array}{c}
\text{promise for} \\
(\text{inflist 3})
\end{array}
\]
Exploiting Parallelism

Conventional procedural programming languages are difficult to compile for multiprocessors. Frequent assignments make it difficult to find independent computations.

Consider (in Fortran):

```fortran
  do 10  I = 1,1000
    X(I) = 0
    A(I) = A(I+1)+1
    B(I) = B(I-1)-1
    C(I) = (C(I-2) + C(I+2))/2
  10 continue
```

This loop defines 1000 values for arrays $X$, $A$, $B$ and $C$. 
Which computations can be done in parallel, partitioning parts of an array to several processors, each operating independently?

- \( X(I) = 0 \)
  
  Assignments to \( X \) can be readily parallelized.

- \( A(I) = A(I+1) + 1 \)
  
  Note that each computation of \( A(I) \) uses an \( A(I+1) \) value that is yet to be changed. Thus a whole array of new \( A \) values can be computed from an array of “old” \( A \) values in parallel.

- \( B(I) = B(I-1) - 1 \)
  
  This is less obvious. Each \( B(I) \) uses \( B(I-1) \) which is defined in terms of \( B(I-2) \), etc. Ultimately all new \( B \) values depend only on \( B(0) \) and \( I \). That is, \( B(I) = B(0) - I \). So this
computation can be parallelized, but it takes a fair amount of insight to realize it.

- \( C(I) = \frac{C(I-2) + C(I+2)}{2} \)

It is clear that even and odd elements of \( C \) don’t interact. Hence two processors could compute even and odd elements of \( C \) in parallel. Beyond this, since both earlier and later \( C \) values are used in each computation of an element, no further means of parallel evaluation is evident. Serial evaluation will probably be needed for even or odd values.
Exploiting Parallelism in Scheme

Assume we have a shared-memory multiprocessor. We might be able to assign different processors to evaluate various independent subexpressions.

For example, consider

\[(\text{map (lambda(x) (* 2 x)) '}(1 2 3 4 5))\]

We might assign a processor to each list element and compute the lambda function on each concurrently:

```
1 2 3 4 5

Processor 1 ...

2 4 6 8 10

Processor 5
```
How is Parallelism Found?

There are two approaches:

- We can use a “smart” compiler that is able to find parallelism in existing programs written in standard serial programming languages.

- We can add features to an existing programming language that allows a programmer to show where parallel evaluation is desired.
Concurrentization

Concurrentization (often called parallelization) is the process of automatically finding potential concurrent execution in a serial program.

Automatically finding current execution is complicated by a number of factors:

- **Data Dependence**
  
  Not all expressions are independent. We may need to delay evaluation of an operator or subprogram until its operands are available.

  Thus in
  
  
  \( (+ (* x y) (* y z)) \)

  we can’t start the addition until both multiplications are done.
• Control Dependence

Not all expressions need be (or should be) evaluated.

In

\[(\text{if } (\text{=} \ a \ 0))\]

0
\n
\[(/ \ b \ a)\]  

we don’t want to do the division until we know \(a \neq 0\).

• Side Effects

If one expression can write a value that another expression might read, we probably will need to serialize their execution.

Consider

\[(\text{define rand!})\]

\[(\text{(let ((seed 99))})\]

\[(\text{(lambda ()})\]

\[(\text{(set! seed (mod (* seed 1001) 101101) seed))})\]
Now in

\((+ (f (\text{rand!})) (g (\text{rand!})))\)

we can't evaluate \((f (\text{rand!}))\) and \((g (\text{rand!}))\) in parallel, because of the side effect of \text{set!} in \text{rand!}. In fact if we did, \(f\) and \(g\) might see exactly the same “random” number! (Why?)

- **Granularity**

  Evaluating an expression concurrently has an overhead (to setup a concurrent computation). Evaluating very simple expressions (like \((\text{car x})\) or \((+ x 1)\)) in parallel isn’t worth the overhead cost.

  Estimating where the “break even” threshold is may be tricky.