Type Equivalence in Classes

In C, C++ and Java, instances of the same struct or class are type-equivalent, and mutually assignable.

For example:

class MyClass { ... }
MyClass v1, v2;
v1 = v2; // Assignment is OK

We expect to be able to assign values of the same type, including class objects.

However, sometimes a class models a data object whose size or shape is set upon creation (in a constructor).
Then we may *not* want assignment to be allowed.

class Point {
    int dimensions;
    float coordinates[];
    Point () {
        dimensions = 2;
        coordinates = new float[2];
    }
    Point (int d) {
        dimensions = d;
        coordinates = new float[d];
    }
}
Point plane = new Point();
Point solid = new Point(3);
plane = solid; //OK in Java

This assignment is allowed, even though the two objects represent points in different dimensions.
Subtypes

In C++, C# and Java we can create subclasses—new classes derived from an existing class. We can use subclasses to create new data objects that are similar (since they are based on a common parent), but still type-inequivalent.

Example:

```java
class Point2 extends Point {
    Point2() {super(2); }
}
class Point3 extends Point {
    Point3() {super(3); }
}
Point2 plane = new Point2();
Point3 solid = new Point3();
plane = solid; //Illegal in Java
```
**Parametric Polymorphism**

We can create distinct subclasses based on the values passed to constructors. But sometimes we want to create subclasses based on distinct *types*, and types can’t be passed as parameters. (Types are not values, but rather a *property* of values.)

We see this problem in Java, which tries to create general purpose data structures by basing them on the class `Object`. Since any object can be assigned to `object` (all classes must be a subclass of `object`), this works—at least partially.
class LinkedList {
    Object value;
    LinkedList next;
    Object head() {return value;}
    LinkedList tail(){return next;}
    LinkedList(Object O) {
        value = O; next = null;
    }
    LinkedList(Object O,
               LinkedList L){
        value = O; next = L;
    }
}

Using this class, we can create a linked list of any subtype of Object.

But,

• We can’t guarantee that linked lists are type homogeneous (contain only a single type).

• We must cast Object types back into their “real” types when we extract list values.
• We must use wrapper classes like `Integer` rather than `int` (because primitive types like `int` aren’t objects, and aren’t subclass of `Object`).

For example, to use `LinkedList` to build a linked list of `ints` we do the following:

```java
LinkedList l =
    new LinkedList(new Integer(123));
int i =
    ((Integer) l.head()).intValue();
```

This is pretty clumsy code. We’d prefer a mechanism that allows us to create a “custom version” of `LinkedList`, based on the type we want the list to contain.
We can’t just call something like \( \text{LinkedList}(\text{int}) \) or \( \text{LinkedList}(\text{Integer}) \) because types can’t be passed as parameters.

*Parametric polymorphism* is the solution. Using this mechanism, we *can* use type parameters to build a “custom version” of a class from a general purpose class.

C++ allows this using its template mechanism. Tiger Java also allows type parameters.

In both languages, type parameters are enclosed in “angle brackets” (e.g., \( \text{LinkedList}<T> \) passes \( T \), a type, to the \( \text{LinkedList} \) class).
Thus we have

class LinkedList<T> {
    T value; LinkedList<T> next;
    T head() {return value;}
    LinkedList<T> tail() {
        return next;
    }
    LinkedList(T O) {
        value = O; next = null;
    }
    LinkedList(T O,LinkedList<T> L) {
        value = O; next = L;
    }
}
LinkedList<int> l =
    new LinkedList(123);
int i = l.head();
Overloading and Ad-hoc Polymorphism

Classes usually allow overloading of method names, if only to support multiple constructors.

That is, more than one method definition with the same name is allowed within a class, as long as the method definitions differ in the number and/or types of the parameters they take.

For example,

class MyClass {
    int f(int i) { ... }
    int f(float g) { ... }
    int f(int i, int j) { ... }
}

Overloading is sometimes called “ad hoc” polymorphism, because, to the programmer, it appears that one method can take a variety of different parameter types. This isn’t true polymorphism because the methods have different bodies; there is no sharing of one definition among different parameter types. There is no guarantee that the different definitions do the same thing, even though they share a common name.
Issues in Overloading

Though many languages allow overloading, few allow overloaded methods to differ only on their result types. (Neither C++ nor Java allow this kind of overloading, though Ada does). For example,

class MyClass {
    int f() { ... }
    float f() { ... }
}

is illegal. This is unfortunate; methods with the same name and parameters, but different result types, could be used to automatically convert result values to the type demanded by the context of call.
Why is this form of overloading usually disallowed?

It’s because overload resolution (deciding which definition to use) becomes much harder. Consider

class MyClass {
    int f(int i, int j) { ... }
    float f(float i, float j) { ... }
    float f(int i, int j) { ... }
}

each of the three calls?
Getting the correctly answer can be tricky, though solution algorithms do exist.
**Operator Overloading**

Some languages, like C++ and C#, allow operators to be overloaded. You may add new definitions to existing operators, and use them on your own types. For example,

```cpp
class MyClass {
    int i;
    public:
        int operator+(int j) {
            return i+j;
        }
}
MyClass c;
int i = c+10;
int j = c.operator+(10);
int k = 10+c; // Illegal!
```
The expression 10+c is illegal because there is no definition of + for the types int and MyClass&. We can create one by using C++’s friend mechanism to insert a definition into MyClass that will have access to MyClass’s private data:

```cpp
class MyClass {
    int i;
    public:
        int operator+(int j) {
            return i+j; }
        friend int operator+(int j, MyClass& v){
            return j+v.i; }
    }

MyClass c;
int k = 10+c; // Now OK!
```
C++ limits operator overloading to existing predefined operators. A few languages, like Algol 68 (a successor to Algol 60, developed in 1968), allow programmers to define brand new operators. In addition to defining the operator itself, it is also necessary to specify the operator’s precedence (which operator is to be applied first) and its associativity (does the operator associate from left to right, or right to left, or not at all). Given this extra detail, it is possible to specify something like

```cpp
op +++ prec = 8;
int op +++(int& i, int& j) {
    return (i++)+(j++);
}
```

(Why is `int&` used as the parameter type rather than `int`?)