

Last Class

→ Jumps

→ Jb / For

→ Encoded

→ Relative

→ Absolute

① 804828 | : 74 05 je 8048296
 804829 |

(+) 05 ← 0000101

② 8048357 | 72 e7 jb 8048340
 8048359 |

(+) -19

11100111

16 | 25

1 - 9 ↑

③ 8048200 : 74 02 00 00 00 je 8048207

8048205 :

0x 00 00 00 02 \rightarrow 2

IA32 uses little endian (Little end comes first)

④ 80482bf : e9 e0 bb bb bb jmp

80482c4 :

\Rightarrow e9 e0 (equivalent)

0x bb bb bb e9

e9 e0 00 00 00

0x 00 00 00 e0

1111

e0

1110 0000

$$= -1 \times 2^8 + 1 \times 2^7 + 1 \times 2^6$$

$$= -128 + 64 + 32$$

$$= -32_{10} \quad 16 \mid 32$$

$$= -20_{16} = -0x20$$

2 -0 \uparrow

If / Else (Example the book)

goto statement

jmp Label

goto Label ;

Eg:

int a = 5 ;

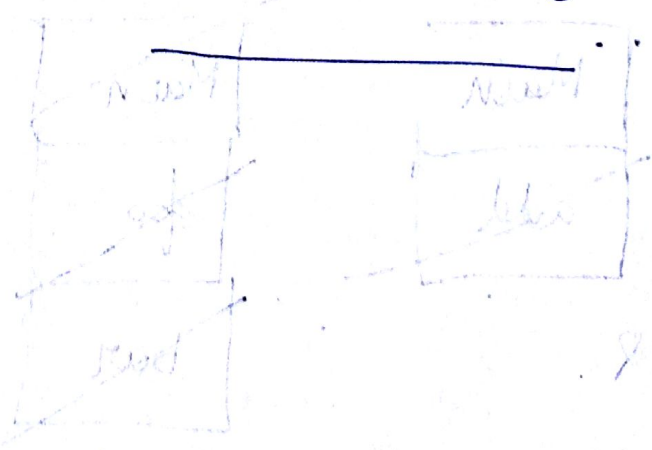
goto end_of_x ; //

end_of_x :

int b = 5 ;

;
;
;

Reading assignment (Monday)
3.6.5 // End of 3.6



3.7 Procedures

~~Proc~~

How does assembly handle / support functions?

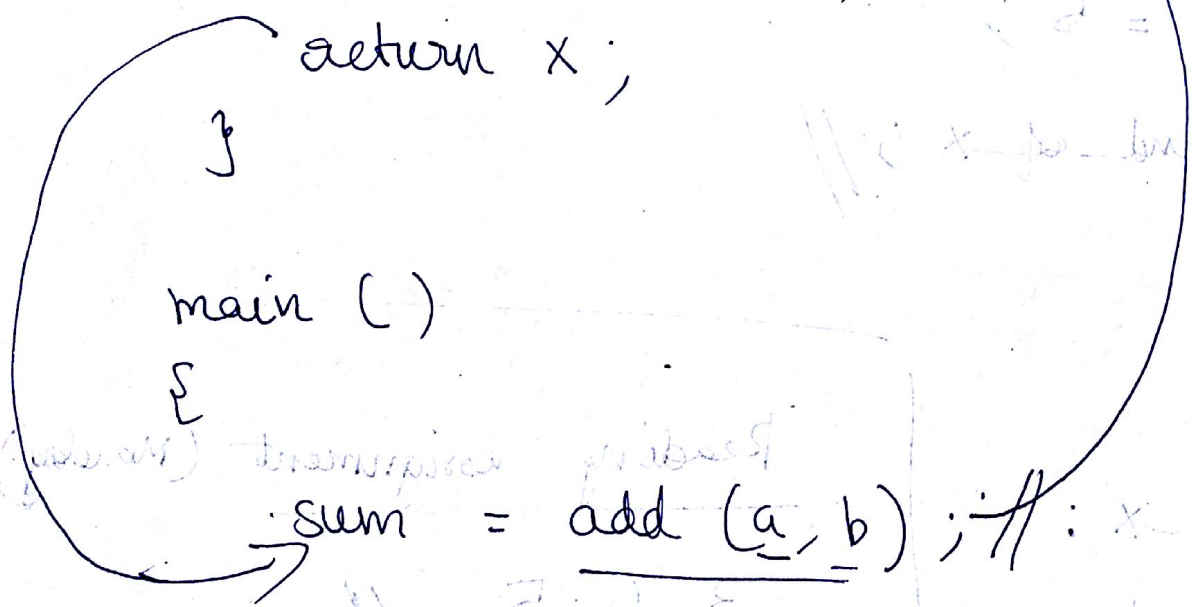
```

int add (int a, int b)
{
    int x = a + b;
    return x;
}
    
```

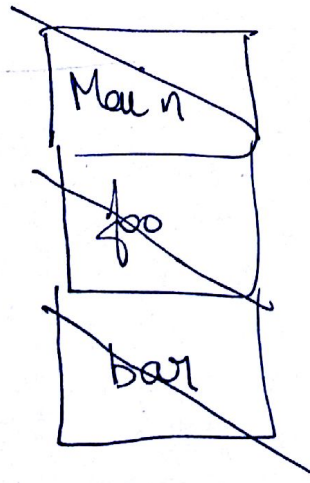
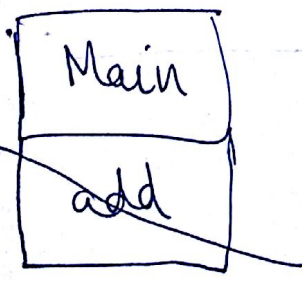
```

main ()
{
    sum = add (a, b);
}
    
```

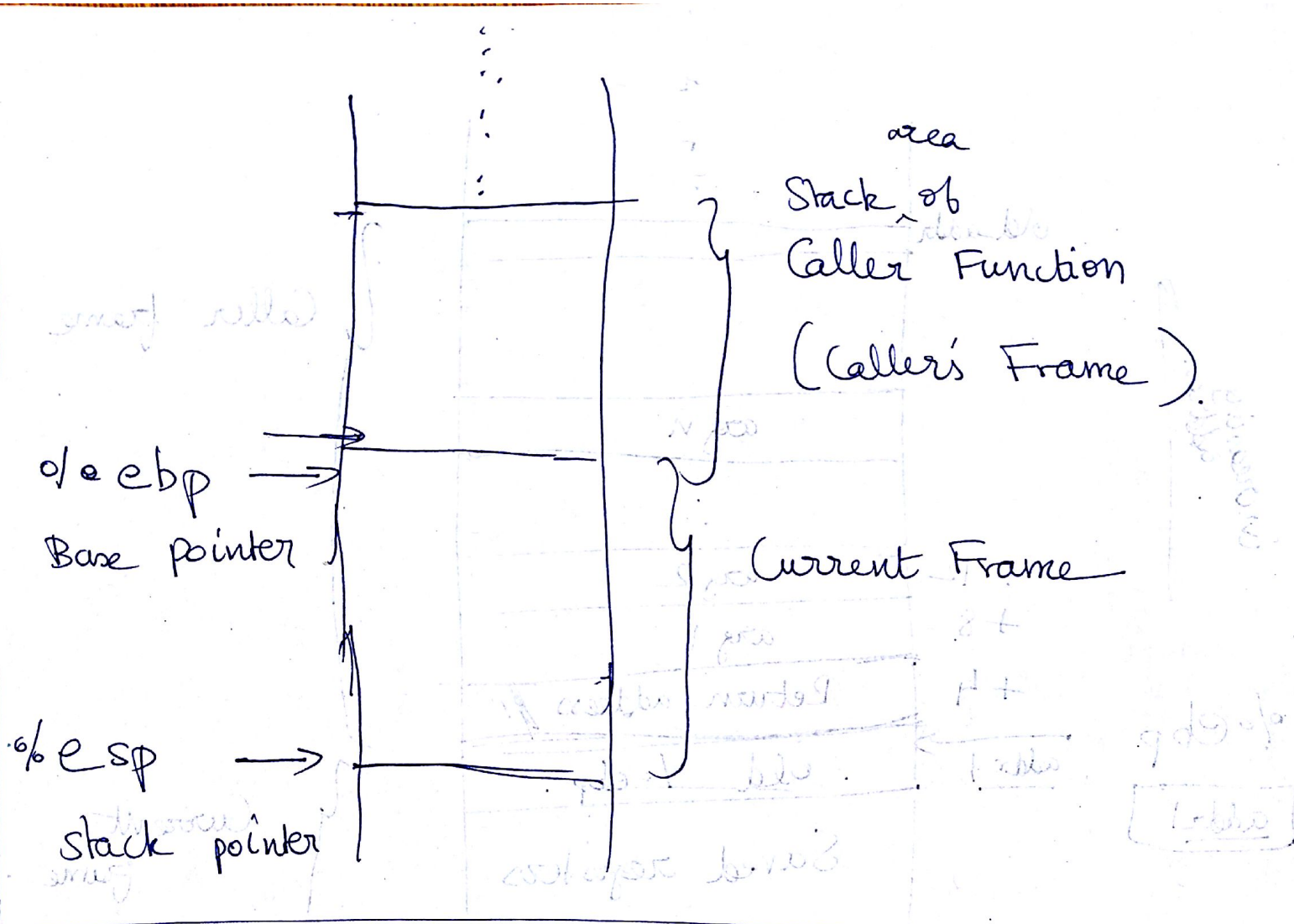
Transfer of data
 &
 Transfer of control



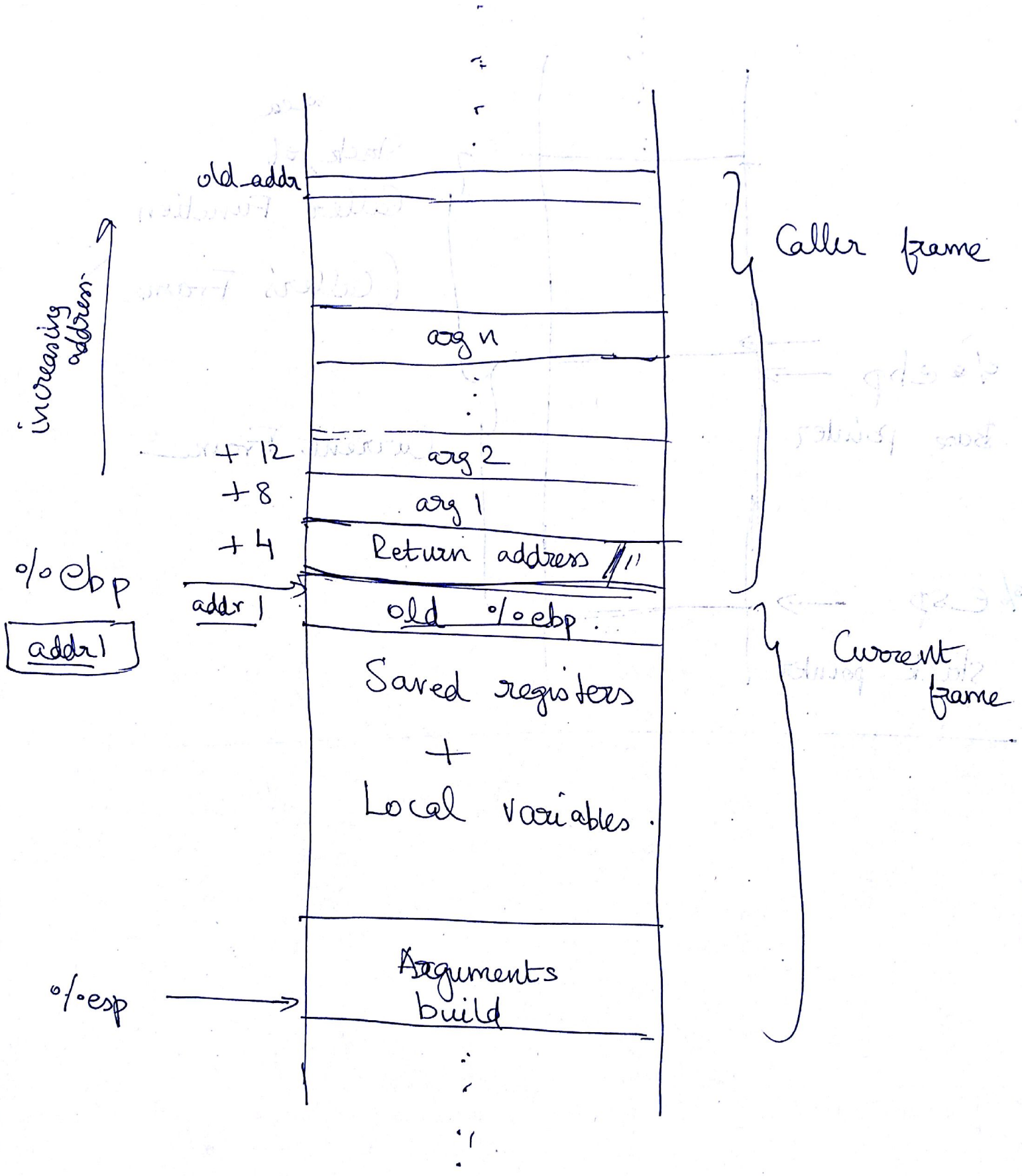
Bottom



Top



level
 level
 level



Return address

Where the program should resume execution when the called function returns.

Arguments

the values that need to be sent / passed to the callee function.

First arg \rightarrow $\%ebp + 8$

To get the n^{th} argument \rightarrow $\%ebp + 4 + 4n$

ebp stores the value of the old ebp so that it can point to it again when it returns.

registers

Local Variables

When? (Local Variables)

variables must be

→ not enough registers to hold all the ~~local~~ local data.

→ When there are references.

k a → address of the first value in a

2 + 7/2

← proc start

10 + 1 + 9/2

← to get the arguments

9/2 shows the value of the

if that it can be used to

scope

local variables

How do we transfer control?

call instruction

Syntax

call Label

call *operand.

} Jump to the target specified and continue execution there.

But before jumping, it pushes the return address on to the stack.

↓
Address of the instruction after the call instruction.

call <add>

⊙ mov %eax, %ebx

Return

(syntax → ret)

The return instruction pops this address off the stack and transfers execution to that address.