

Recap Relative encoding

08048357

08048357 : 72 e7

08048359 :

00000019

(1) 10 0111

-256

19 16

08048359

bb b7 bb e7

340

804 82 00 :

02 00 00 00

je

804 82 05 :

804 82 07

Little Endian (Least significant byte comes first)

0x 00 00 00 02

02 87 92 00
 0x 00 92 87 02

③ 080482b6 : e9 e0 bb bb bb jmp 080482a4

080482c4 : _____

+ 0x bb bb bb e0

a4

3.7 Procedures

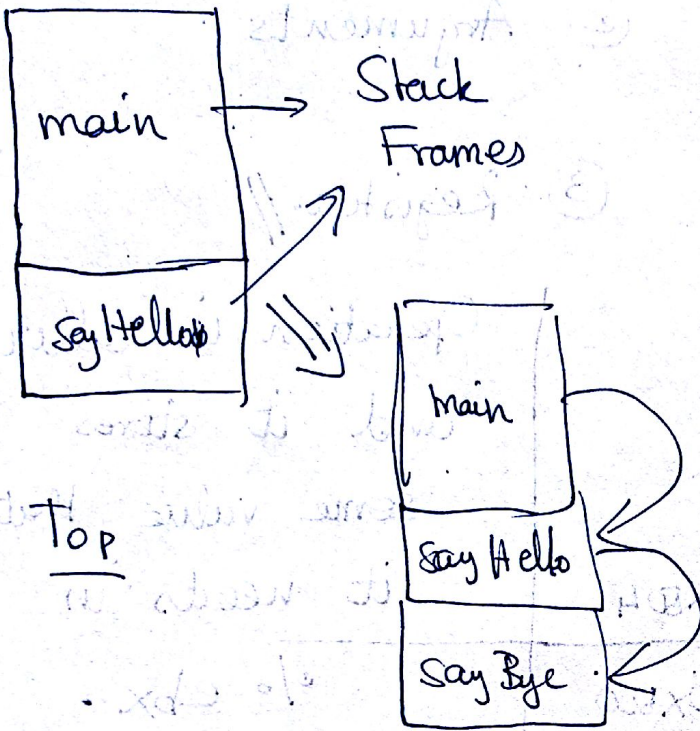
Transfer of control (and data)
from the caller to the callee and back.

We use stacks for this

- To pass arguments
 - To store return address & values
 - To save registers for later use
 - For local storage
- } all data

Stacks

Bottom



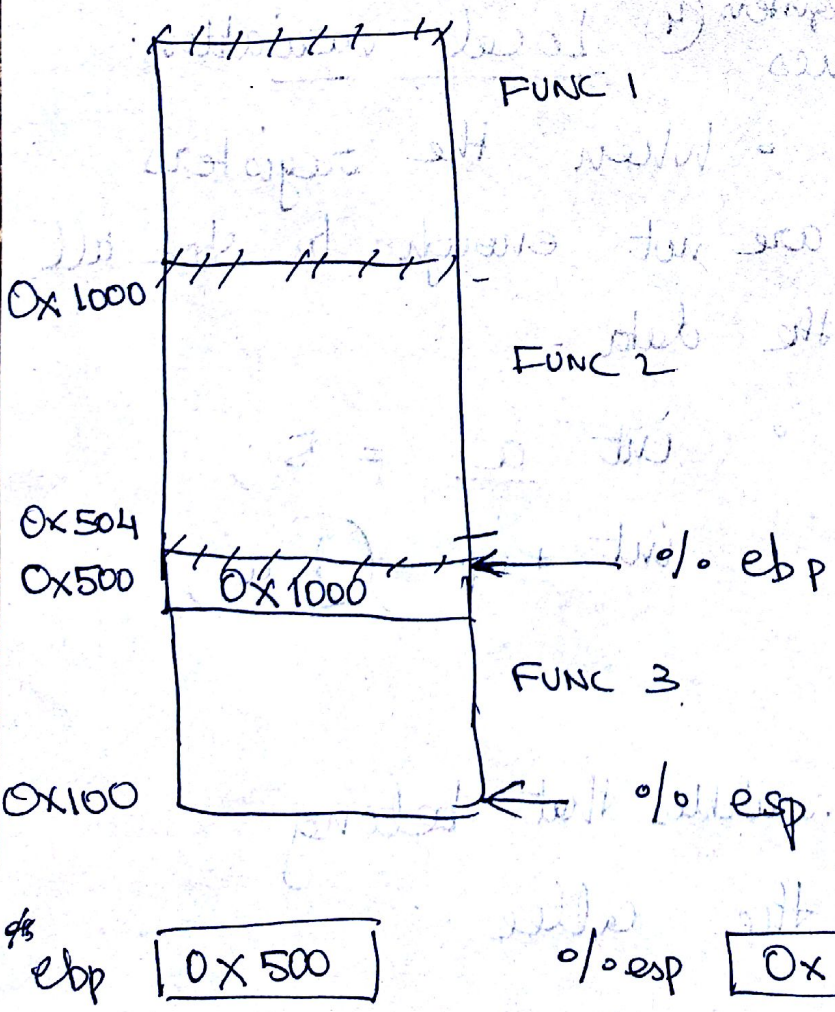
Top

Text Area

```
<main> :
movl . . .
incl . . .
call . . .
call . . .
```

```
<say-Hello> :
. . .
. . .
```

```
<say-Bye> :
. . .
. . .
```



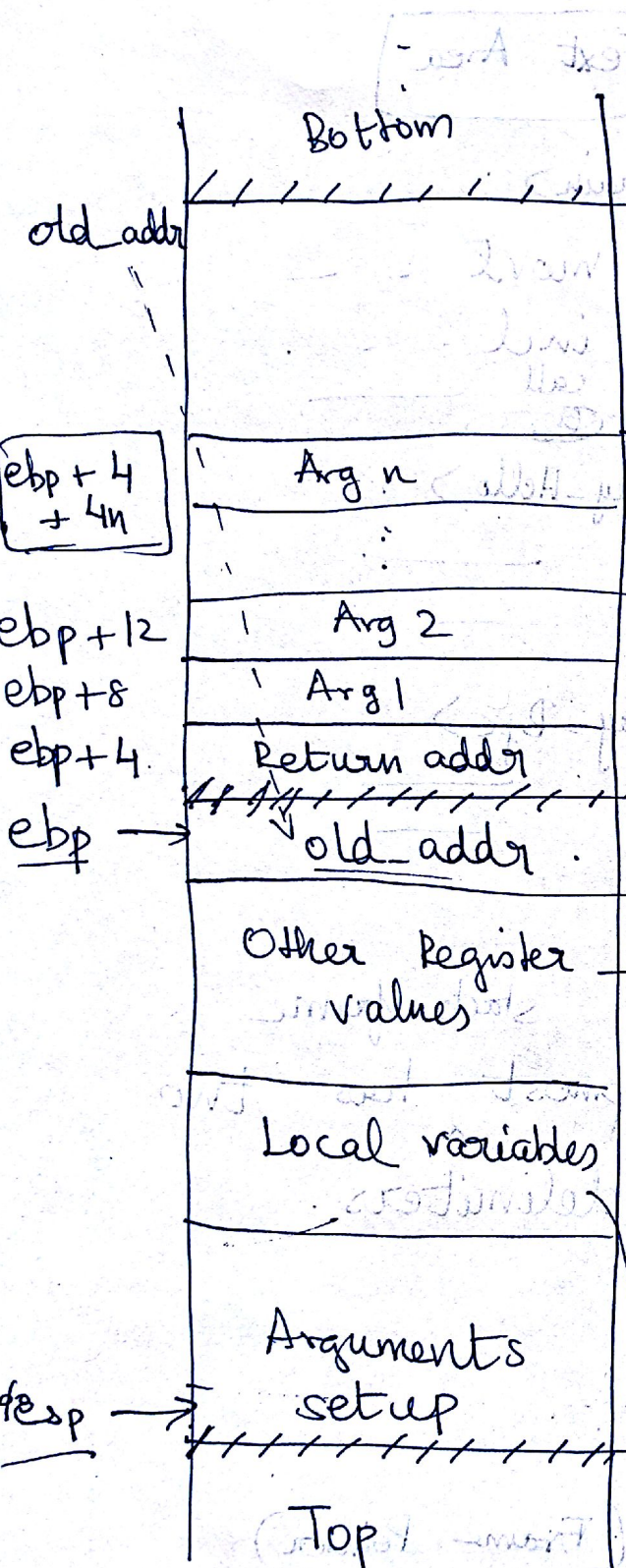
stack frame

Topmost has two delimiters.

(Frame Pointer)

(or)
(Base Pointer)

(Stack Pointer)



- ① Return Address
- ② Arguments
- ③ Registers //

Operation in caller and it stores some value that it needs in %ebx.

④ Local variables.

When the registers are not enough to store all the data.

```
int a = 5;
```

```
int *ap = (k)a;
```

The variables that belong to the callee.

Transfer Control

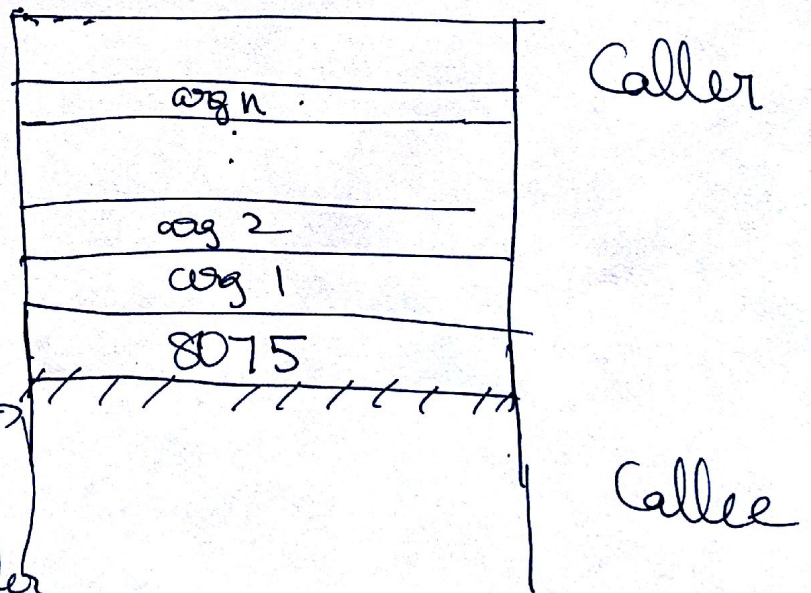
Call

Syntax → call Label

→ call *operand

It jumps to the target specified and it continues execution there.

Also pushes the return address onto the stack before jumping.



<callee>

;

ret

caller
<return> :

8071 : call <callee>

8075 : add

8078 : movl

!

Return

It pops the

Syntax \rightarrow ret

return address off the stack and it jumps to that location.

Before we do a return,

make sure that esp points to that location.

