

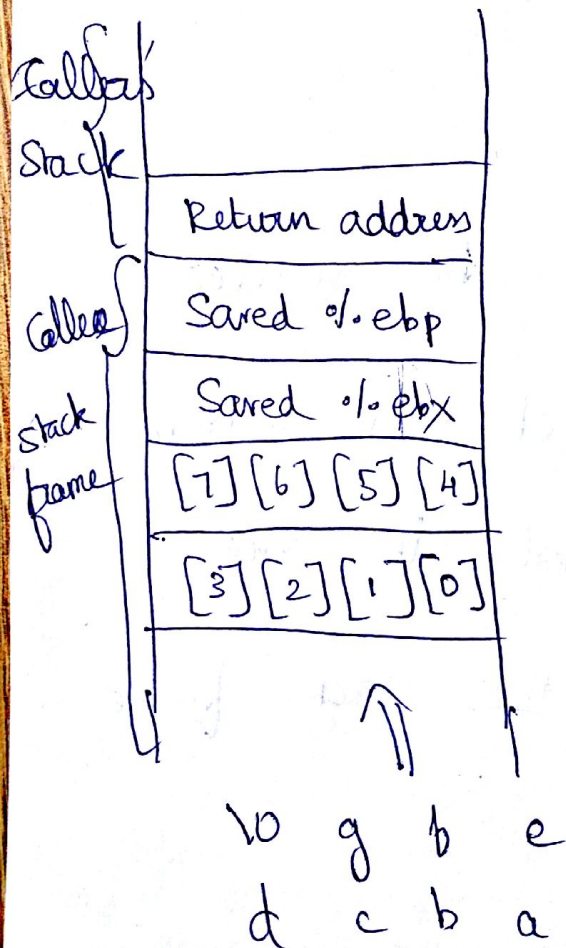
Today

- ① Buffer Overflow. ✓
- ② Assembly wrap-up.

→ No bounds-checking for array references.

```
int a[5];
```

```
printf("%d", a[10]); //
```



```
1: char buf[8];
```

```
2: gets(buf);
```

If I enter something that has 7 characters,

a b c d e f g \0

What will be corrupted

0-7	None
8-11	Saved value of 4e ebx
12-15	Saved value of ebp
16-19	Return address
20+	_____

Serious issue!

① Program will break.

② Security issue!

We know that → when the callee returns by using the ret instruction, it basically returns control to the address that's stored on the stack frame.

What if I am an attacker?

I just overwrite the return address.

- Load my own code (Exploit code)
- Modify the return address so that it returns to the exploit code.
- Now the exploit^{code} is executing ~~code~~
~~is~~ executing under the guise of your non-malicious code.
 - read materials (private)
 - get privileged access.
 - start another program that basically gives access to some remote computer.

1988

Morris Worm

(Robert Morris)

(Stack smashing)

How could we avoid this?

- ① Do not use vulnerable library functions.

gets X

fgets ✓

→

Discards the excess chars entered.

- ② We could make the OS better!

The stack structure looks exactly the same on different runs the program.

②a Stack Randomization!

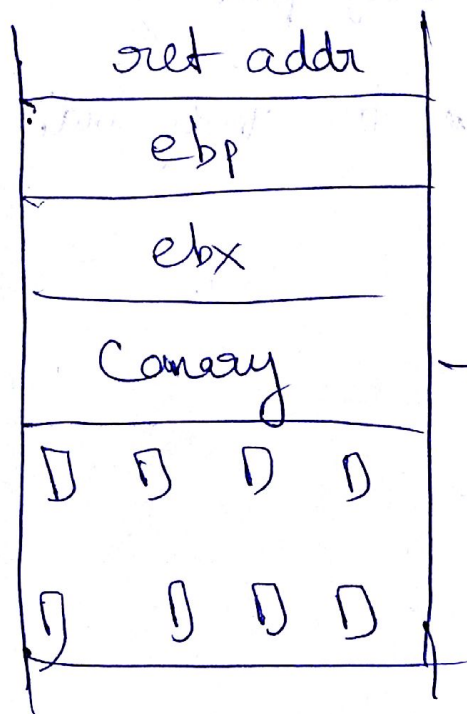
→ making the starting address of the stack vary between different

sums of the program.

(2b) Add a stack protector!



Canary value.



random number.

Before return, value of canary is checked.

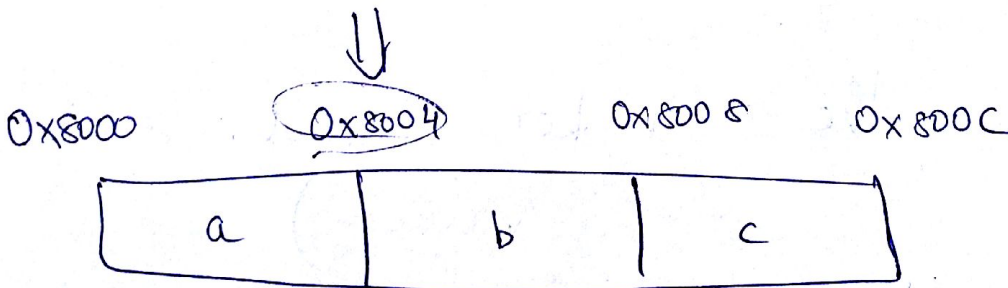
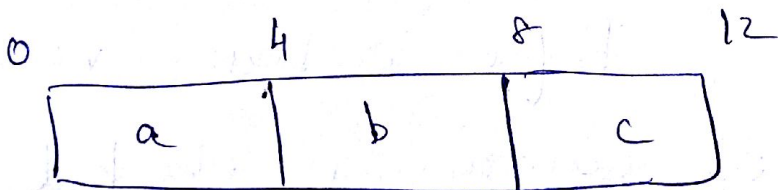
(Remember: the attacker cannot read values).

Review 1

Struct Addressing

struct s1

```
{  
    int a;  
    char int b;  
    int c;  
};
```



```
struct s1 x;
```

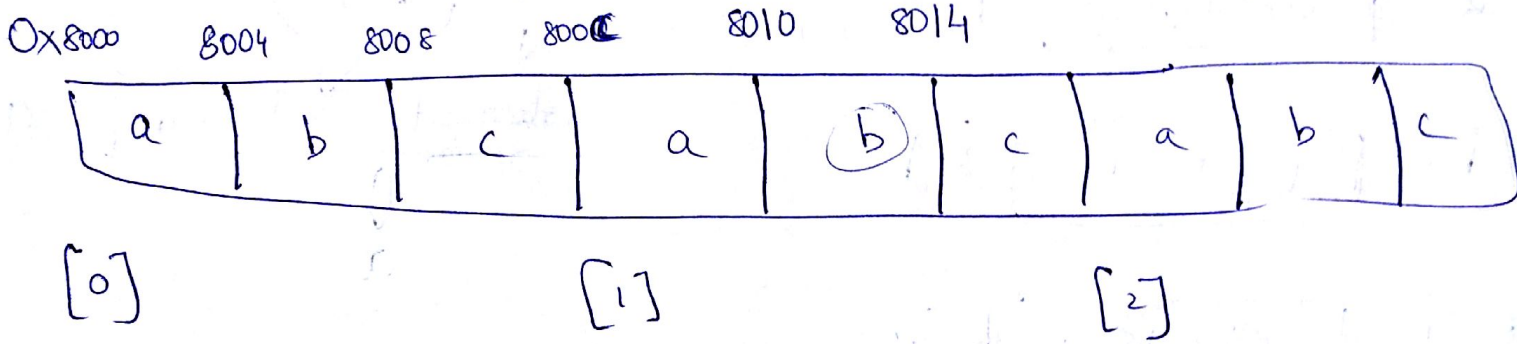
```
printf("%p", &x.b) => 0x8004
```

```
printf("%p", &x.a) => 0x8000
```

Things to Review

- ① Conversion (Hex to decimal)
(Decimal to binary)
- ② 2.4 (Floating point) Study only what we discussed in class
- ③ No need to study Gdb.
- ④

```
struct s1 arr [3];
```



```
printf ("%p", &arr [1].b); => 0x8010
```

② Array addressing

Assume var a starting at memory location

0x8049000

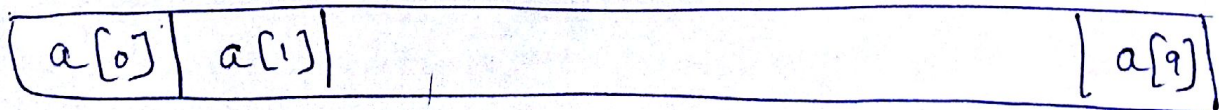
```
1: int a [10];
```

```
2: int *p = &a [7];
```

0x8049000 → 0x8049004

What is the value of p?

0x804901c



$$\begin{array}{rcl} \text{Element at} & & \text{Start} \\ \text{index at} & = & \text{Address} \\ 7 & & \downarrow \\ & & \text{X} \end{array} \quad + \quad \begin{array}{r} 4 \\ \downarrow \\ L \end{array} \quad \times \quad \begin{array}{r} 7 \\ \downarrow \\ i \end{array}$$

③ Assembly to C.

- %eax holds &a (address of a)
- %ebx holds &b (address of b).

1: `movl (%eax), %esi` $\Rightarrow *(&a) \Rightarrow a$

2: `movl (%esi), %edi` $\Rightarrow *(a)$

~~3: `movl (%edi), %eax`~~

3: `movl %edi, (%ebx)` $\Rightarrow (*a) \Rightarrow *(&b)$
 $*a \Rightarrow b$

What is the equivalent C statement (s) for this set of instructions?

1: move the value at the address stored in %eax to %esi

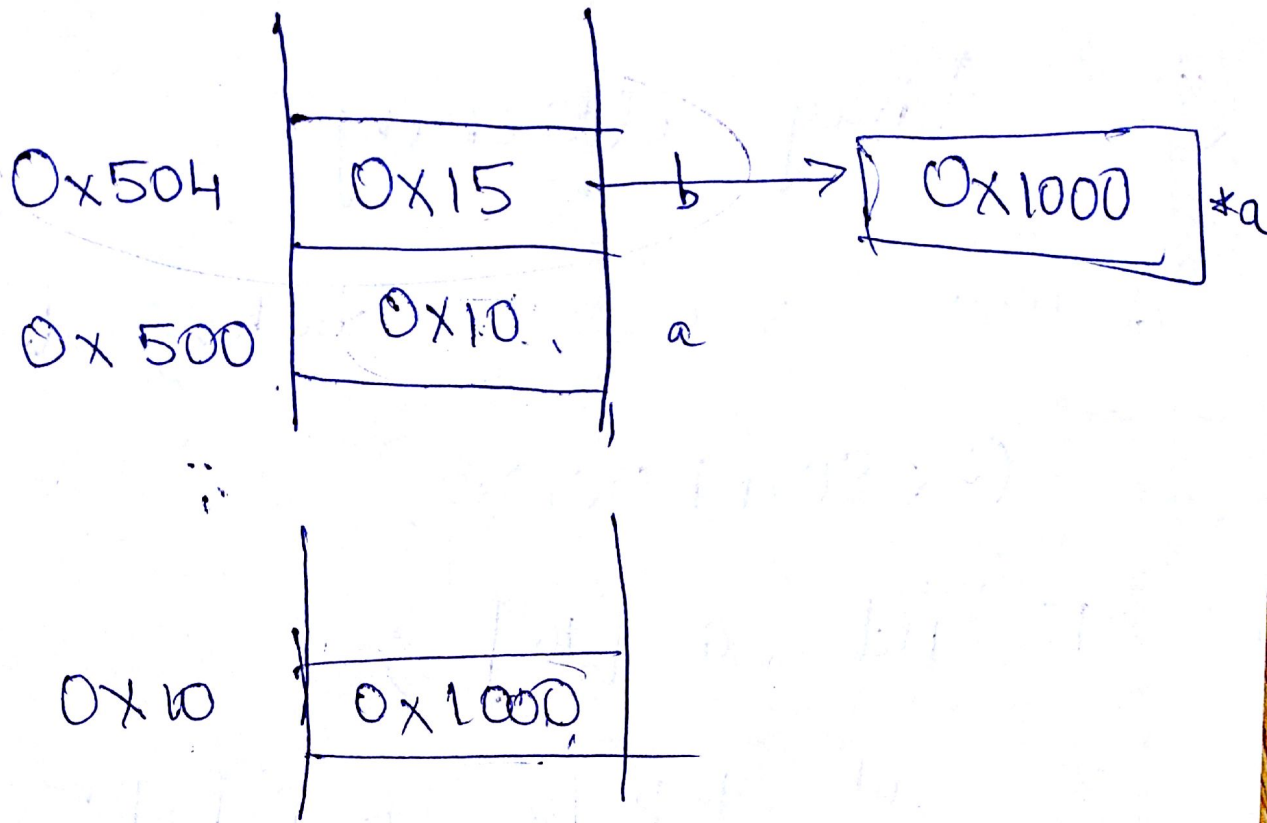
Check email for additional info about this problem!

eax 0x50A (k a)

ebx 0x504 (k b)

esi 0x10 (a)

edi 0x1000 (~~*~~ a)



$b = *a ;$