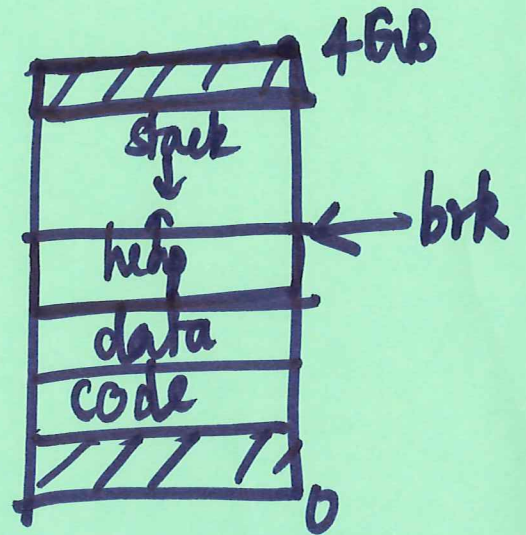


Lecture - 32

Review

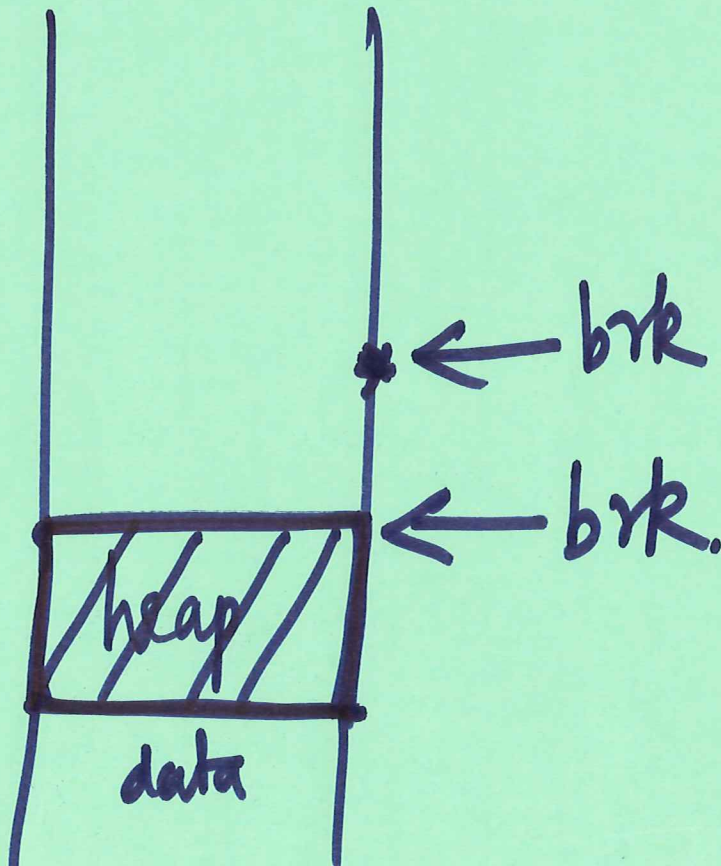
Virtual address space

malloc, free,
calloc, realloc.



int brk (void *addr);

void * sbrk (int incr); - most used.



On error, sbrk return

(void *) -1



0x FF FF FF FF

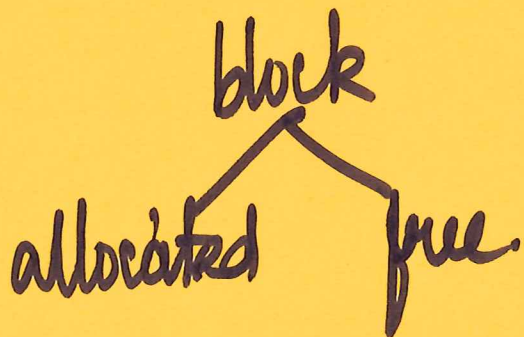
Allocators

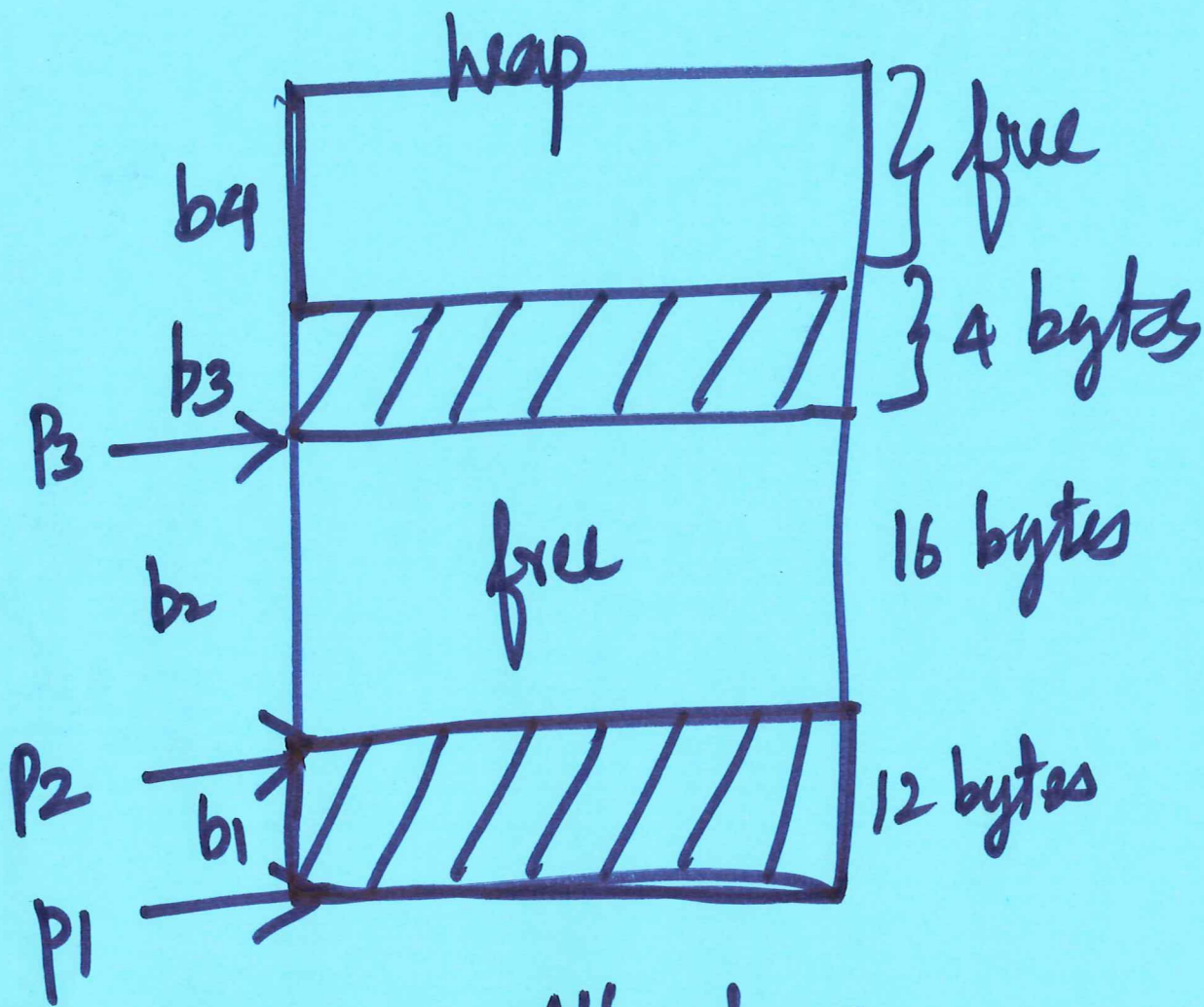


malloc

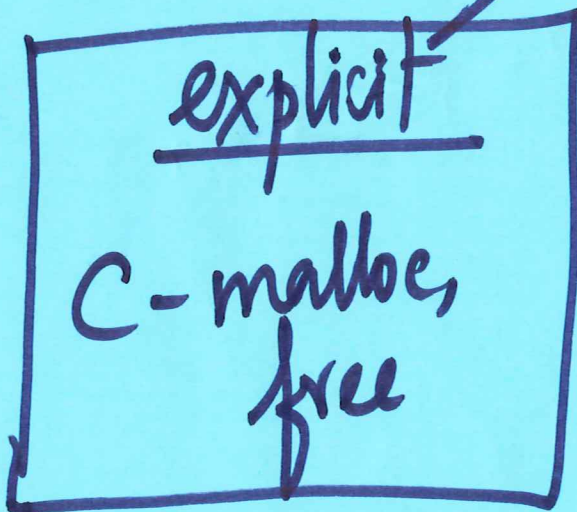
heap → collection of various-sized blocks

block → contiguous chunk of virtual mem.





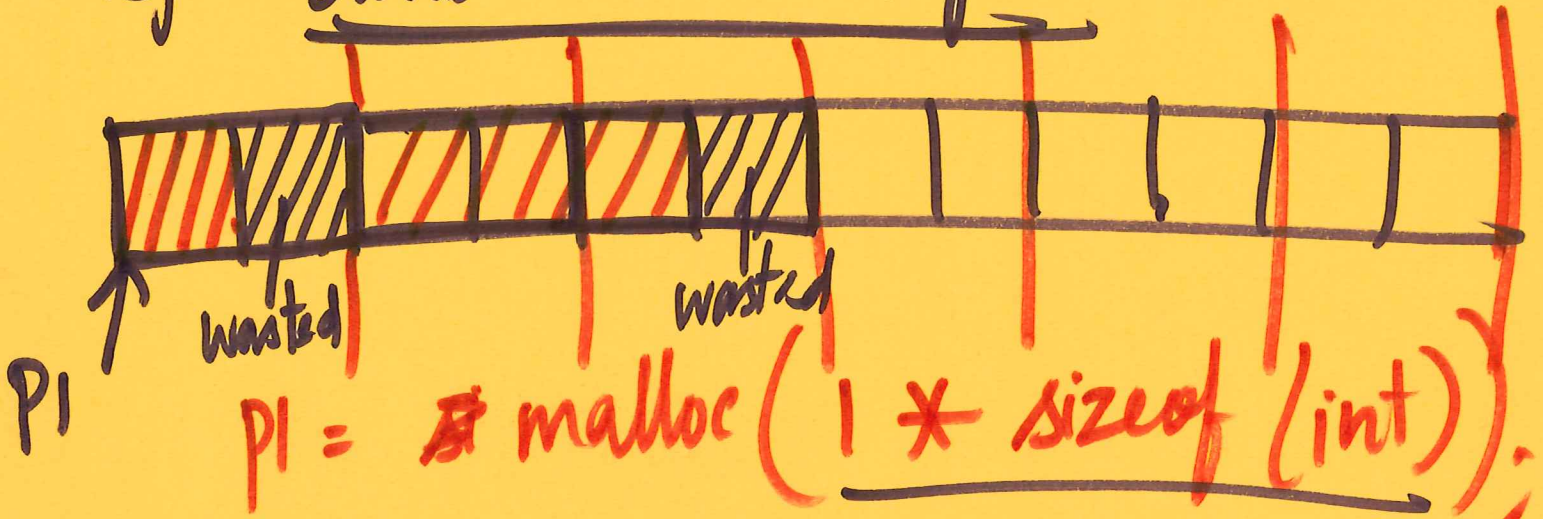
Allocators



implicit
Garbage collection
"new"
free - automatic.

Alignment requirements

eg. double - word aligned



In P5, align req. is word size.



$\text{malloc}(1) \rightarrow$ give you 4 bytes.

payload:

1 - used

padding:

3 - wasted

Implementation issues :

1. Free block organization.
 2. Placement
 3. Splitting into free blocks.
 4. Coalescing.
-

Requirements

1. handle arbitrary requests.
2. Respond immediately.
3. Metadata → stored only in the heap.
4. Aligns the blocks.
5. NOT modify allocated blocks.

2 Goals

1. Maximize throughput

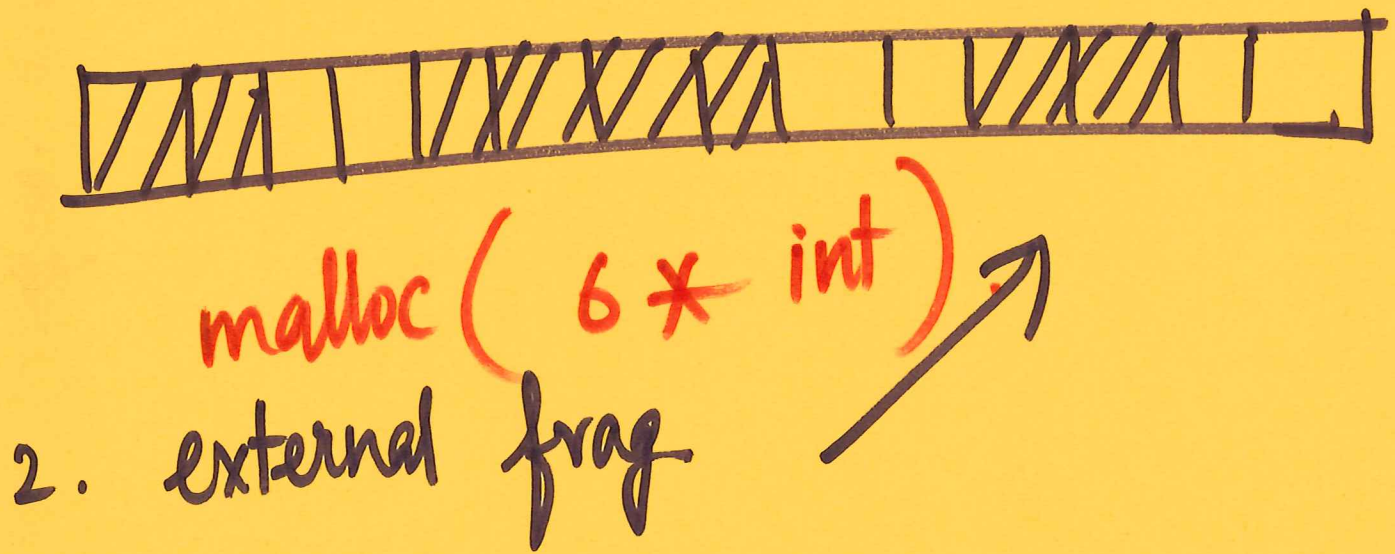
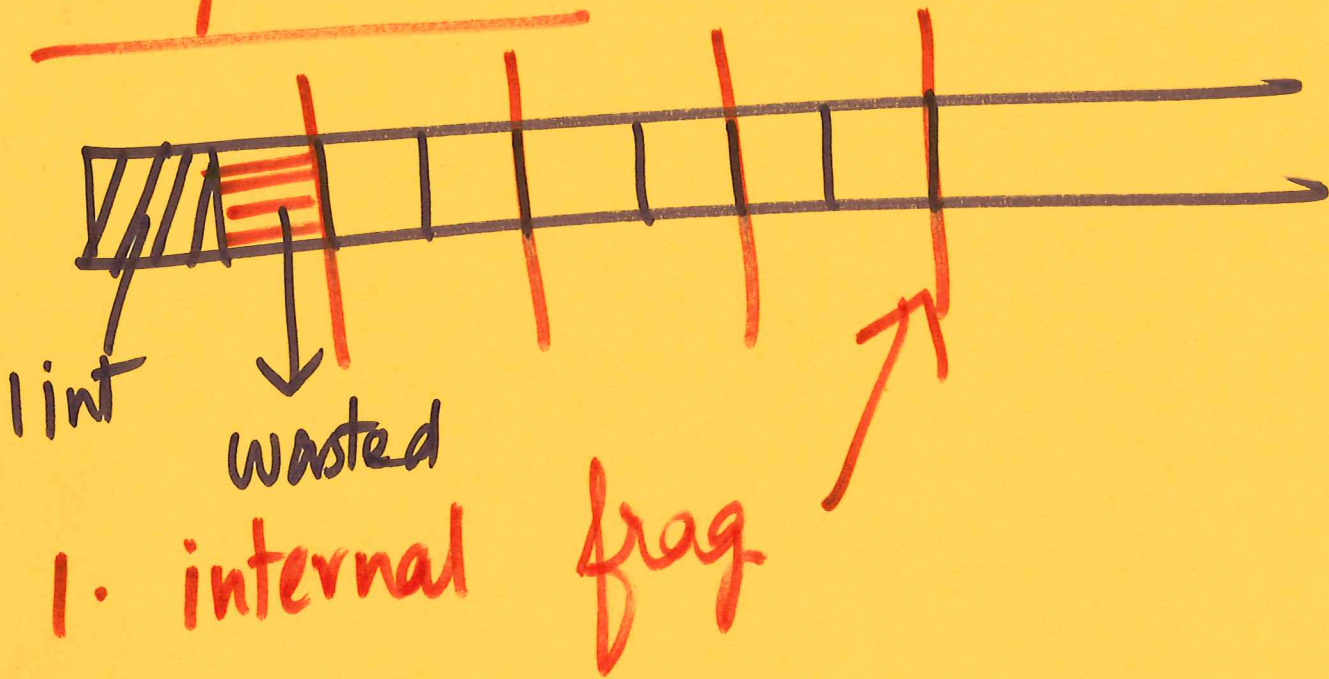
1000 mallocs + 1000 frees
in 1 sec.

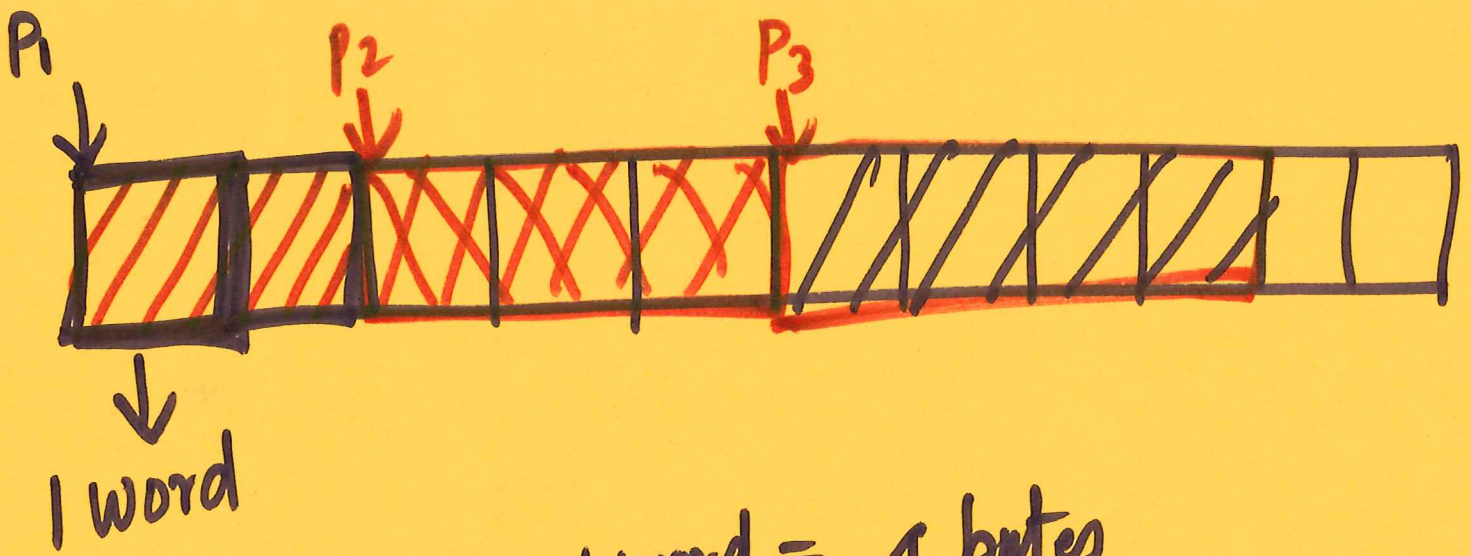
⇒ throughput of 2000 op/s.

2. Max. memory utilization.



Fragmentation





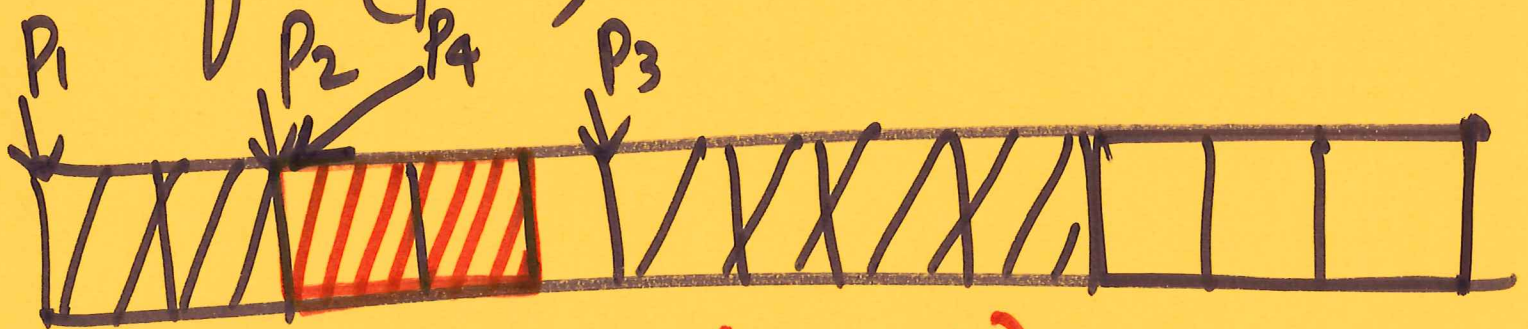
1 word = 4 bytes

```
p1 = malloc(2 * sizeof(int));
```

```
p2 = malloc(3 * sizeof(int));
```

```
p3 = malloc(4 * u);
```

```
free(p2);
```



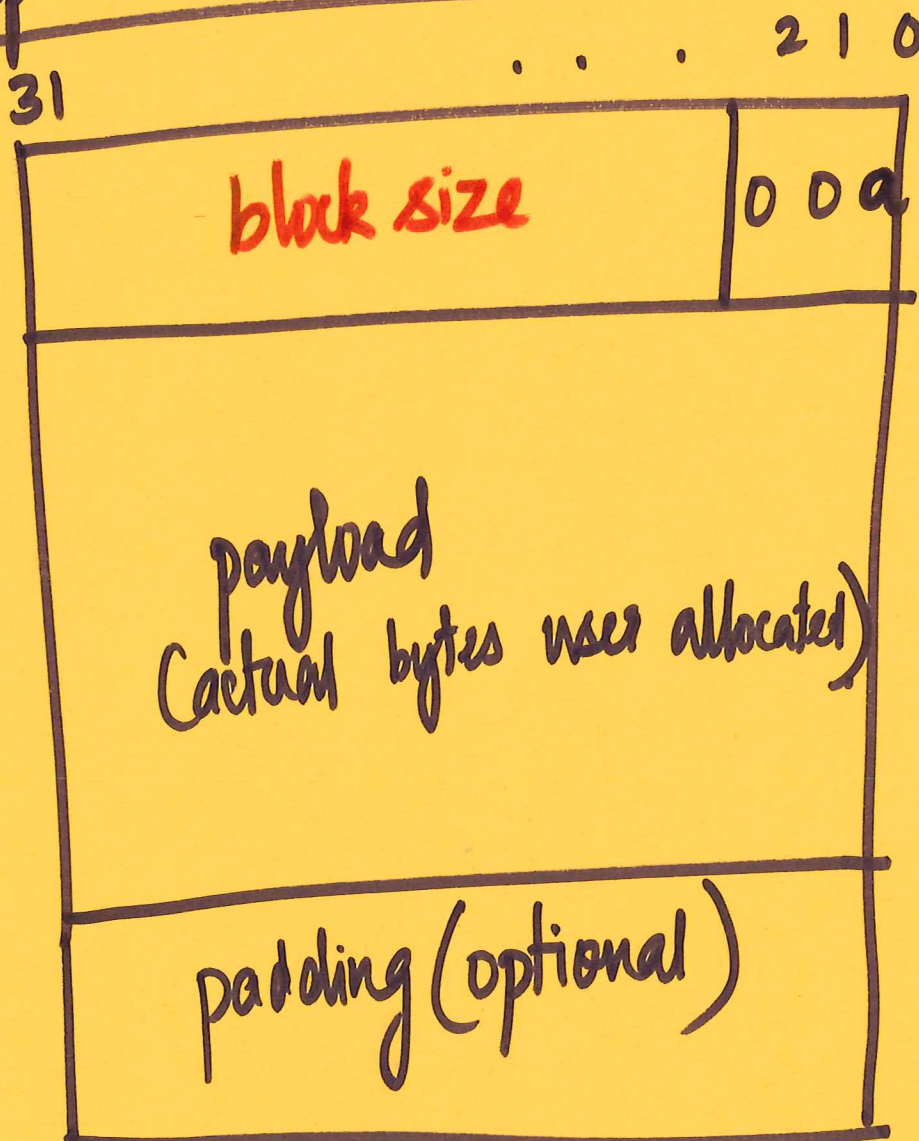
```
P4 = malloc(2 ints);
```


Free block organization

implicit free lists

double word aligned.

Heap block structure



$a=1 \Rightarrow$ allocated
 $a=0 \Rightarrow$ free.

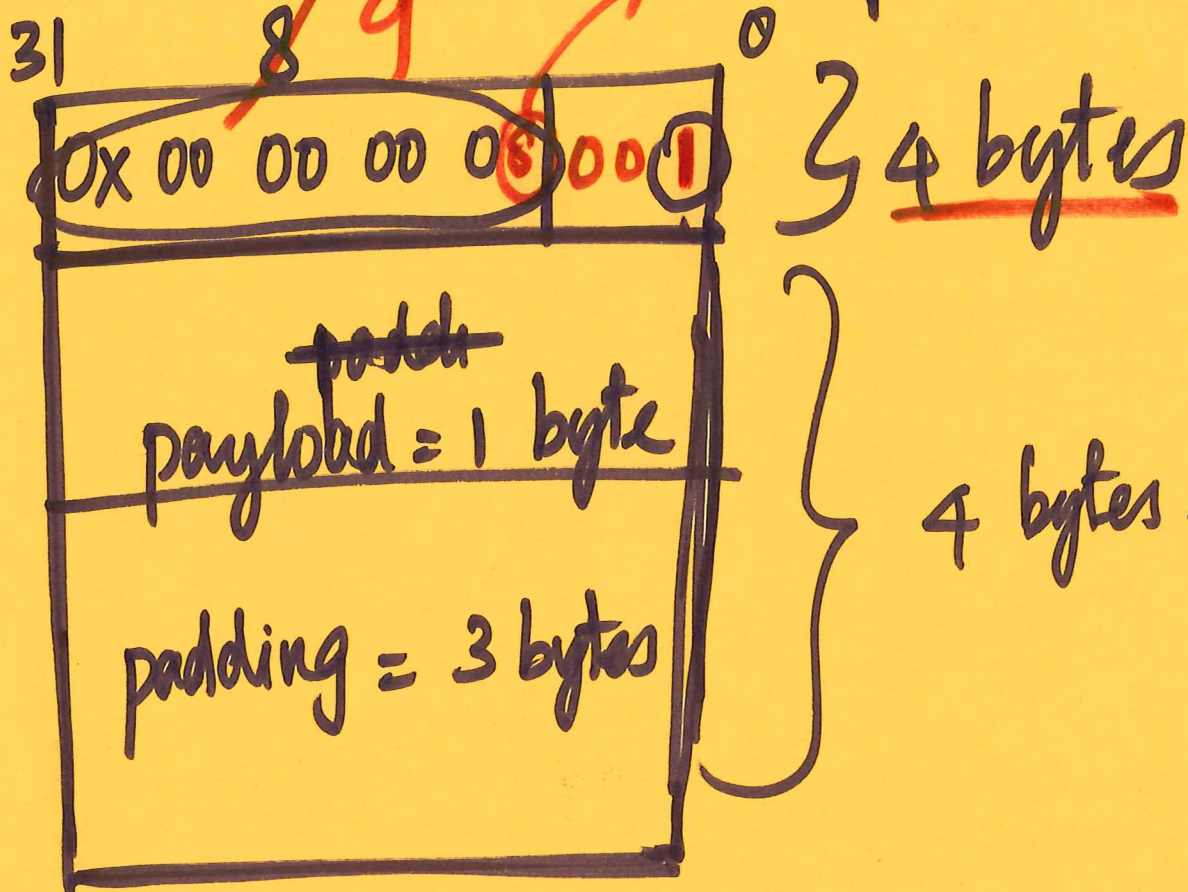
0x F1234578

0x F1234570

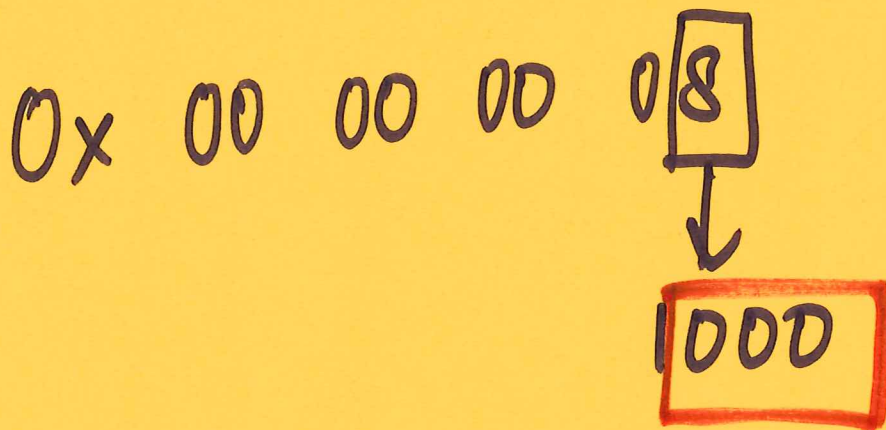


blocksize = header + payload + padding

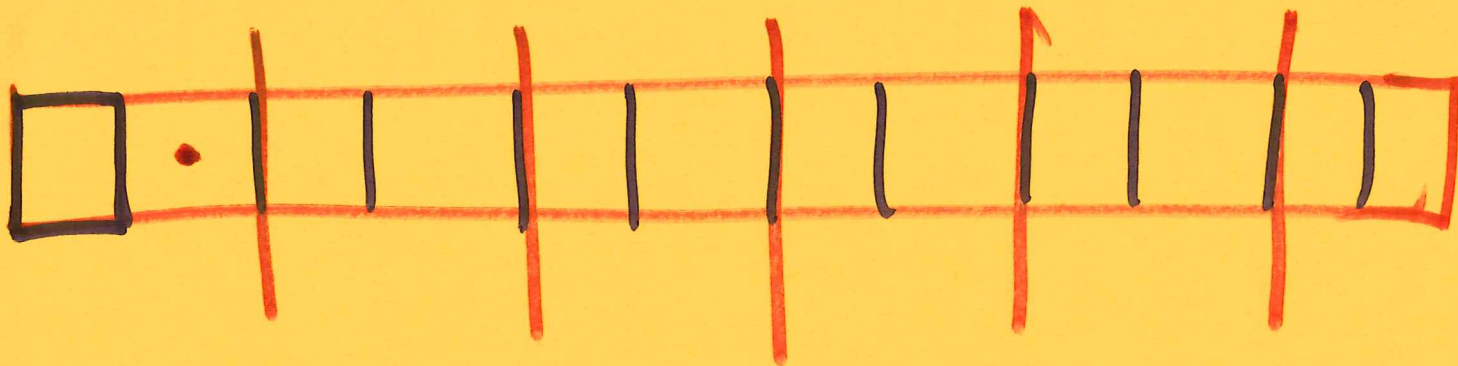
malloc(1)



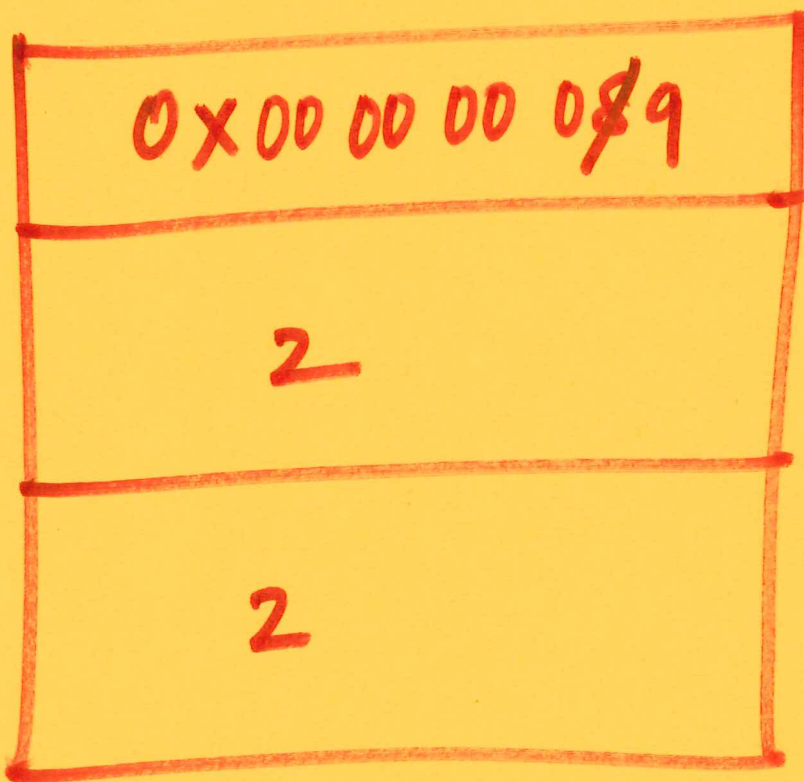
$$4 + \textcircled{1} + 3 = 8 \text{ bytes}$$



0



malloc(2)



} 4 bytes