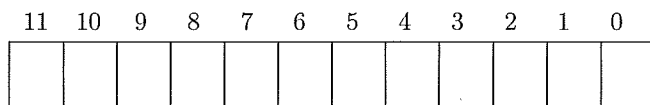


## Lecture 13 - Set Associative Caches

- The memory is byte addressable.
- Memory accesses are to **1-byte words** (not 4-byte words).
- Physical addresses are 12 bits wide.
- The cache is 4-way set associative, with a 2-byte block size and 32 total lines.

4-way Set Associative Cache																
Index	Tag	Valid	Byte 0	Byte 1	Tag	Valid	Byte 0	Byte 1	Tag	Valid	Byte 0	Byte 1	Tag	Valid	Byte 0	Byte 1
0	29	0	34	29	87	0	39	AE	7D	1	68	F2	8B	1	64	38
1	F3	1	0D	8F	3D	1	0C	3A	4A	1	A4	DB	D9	1	A5	3C
2	A7	1	E2	04	AB	1	D2	04	E3	0	3C	A4	01	0	EE	05
3	3B	0	AC	1F	E0	0	B5	70	3B	1	66	95	37	1	49	F3
4	80	1	60	35	2B	0	19	57	49	1	8D	0E	00	0	70	AB
5	EA	1	B4	17	CC	1	67	DB	8A	0	DE	AA	18	1	2C	D3
6	1C	0	3F	A4	01	0	3A	C1	F0	0	20	13	7F	1	DF	05
7	0F	0	00	FF	AF	1	B1	5F	99	0	AC	96	3A	1	22	79

<i>CO</i>	The block offset within the cache line
<i>CI</i>	The cache index
<i>CT</i>	The cache tag



## Part 2

For the given physical address, indicate the cache entry accessed and the cache byte value returned **in hex**. Indicate whether a cache miss occurs.

If there is a cache miss, enter “-” for “Cache Byte returned”.

**Physical address:** 0x3B6

Physical address format (one bit per box)

11	10	9	8	7	6	5	4	3	2	1	0
<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>	<input type="text"/>

Physical memory reference

Parameter	Value
Cache Offset (CO)	0x
Cache Index (CI)	0x
Cache Tag (CT)	0x
Cache Hit? (Y/N)	
Cache Byte returned	0x

## Part 3

In the 4-way Set Associative Cache given above, list all the hex memory addresses that will hit in **Set 7**.

## Cache Miss Rate Analysis

You are evaluating the cache performance of the following code on a machine with a **64-byte direct-mapped data cache (C = 64)** with **block size of 16-bytes (B = 16)**.

You are given the following definitions:

```
struct point {
    int x;
    int y;
};
struct point grid[4][4];
int total_x = 0, total_y = 0;
int i, j;
```

You should also assume the following:

- `sizeof(int) == 4`.
- `grid` begins at memory address 0.
- The cache is initially empty.
- The only memory accesses are to the entries of the array `grid`.
- Variables `i`, `j`, `total_x`, and `total_y` are stored in registers.

A. Determine the cache performance for the following **code snippet 1**:

**Code snippet 1:**

```
for (i = 0; i < 4; i++) {
    for (j = 0; j < 4; j++) {
        total_x += grid[i][j].x;
    }
}
for (i = 0; i < 4; i++) {
    for (j = 0; j < 4; j++) {
        total_y += grid[i][j].y;
    }
}
```

1. What is the total number of reads that miss in the cache? \_\_\_\_\_
2. What is the miss rate? \_\_\_\_\_

B. Determine the cache performance for the following **code snippet 2**:

**Code snippet 2:**

```
for (i = 0; i < 4; i++) {  
    for (j = 0; j < 4; j++) {  
        total_x += grid[i][j].x;  
        total_y += grid[i][j].y;  
    }  
}
```

1. What is the total number of reads that miss in the cache? \_\_\_\_\_
2. What is the miss rate? \_\_\_\_\_

C. Which of these 2 code snippets is better with respect to cache performance?

- a. Code snippet #1
- b. Code snippet #2

Why? (just a single line explanation is sufficient)

---

# Linking

Consider the three files **sum.h**, **sum.c** and **main.c** as shown below and answer the questions that follow.

**sum.h**

=====

```
1.  #ifndef SUM_H
2.  #define SUM_H
3.  extern int sum(int, int);
4.  #endif
```

**sum.c**

=====

```
1.  #include "sum.h"
2.  extern int num_ops;
3.  static int global_sum = 0;
4.
5.  int sum(int x, int y)
6.  {
7.      global_sum += (x+y);
8.      num_ops++;
9.      return x+y;
10. }
```

**main.c**

=====

```
1.  #include "sum.h"
2.  #define SUCCESS 0
3.
4.  int num_ops = 0;
5.
6.  int main()
7.  {
8.      int a = 10;
9.      int b = 3;
10.     int result = sum(a,b);
11.     return SUCCESS;
12. }
```

The final executable **a.out** is produced by running the following command:

```
% gcc sum.c main.c -m32
```

**Questions:**

- A. The variable `num_ops` is **ONLY declared but not defined** in the file \_\_\_\_\_
- B. The variable `num_ops` is **defined** in the file \_\_\_\_\_
- C. Do the following variables / functions need **relocation** when the executable (a.out) is formed? (Just answer: **Yes** or **No**).
  - a. Variable `global_sum` in `sum.c` - \_\_\_\_\_
  - b. Variable `result` in `main.c` - \_\_\_\_\_
  - c. Variable `num_ops` in `main.c` - \_\_\_\_\_
  - d. Function `sum()` in `sum.c` - \_\_\_\_\_
  - e. Function `main()` in `main.c` - \_\_\_\_\_
- D. Can the variable `global_sum` be accessed in the file `main.c` without changing the type of the variable `global_sum` in `sum.c`? Yes OR No.
- E. What type of object file is **sum.o** ?  
**sum.o** is generated using the command: `gcc -c sum.c -m32`
  - i. **Relocatable Object File**
  - ii. **Executable Object File**
- F. Which part of **program memory** (code, data, stack, heap) are the following variables / functions stored?
  - a. Variable `global_sum` in `sum.c` - \_\_\_\_\_
  - b. Variable `result` in `main.c` - \_\_\_\_\_
  - c. Variable `num_ops` in `main.c` - \_\_\_\_\_
  - d. Function `sum()` in `sum.c` - \_\_\_\_\_
  - e. Function `main()` in `main.c` - \_\_\_\_\_