C5 354. Assembly Worksheet -1 Prepared by: Remzi Arpaci-Dusseau

x86 general-purpose registers							
(most	significant) (lead [] [] []	eax ax ah	32 bits 16 bits 8 bits 8 bits				
	[] [] []	bx bh					
	[] [] []	cx ch					
	[] [] []	dx dh					
I	[] []						
Referred to as %eax, %ebx, %ecx, %edx, %esi, %edi, etc.							

INSTRUCTION: addl SOURCE, DESTINATION

definition: adds SOURCE and DESTINATION, puts result into DESTINATION
 i.e., DESTINATION = DESTINATION + SOURCE

limited usage (for now):
 source=number ("immediate") destination=register
 source=register destination=register

INSTRUCTION: imull SOURCE, DESTINATION

definition: DESTINATION = DESTINATION * SOURCE

alternate:
 imull AUX, SOURCE, DESTINATION
 definition: DESTINATION = AUX * SOURCE

limited usage (for now):
 - source=number ("immediate") destination=register
 - source=register destination=register
 - (aux=immediate)

INSTRUCTION: idivl DIVISOR

definition: contents of %edx:%eax (64 bit number) divided by DIVISOR
 quotient -> %eax
 remainder -> %edx

limited usage (for now):
 divisor=register

Notes: A bit weird in its usage of VERY SPECIFIC registers!

Problem #1

```
Write assembly to:
```

- move value 1 into %eax
- add 10 to it and put result into %eax

Problem #2

```
Expression: 3 + 6 * 2
```

Use one register (%eax), and 3 instructions to compute this piece-by-piece

Problem #3

movl \$0, %edx

movl \$7, %eax

mov1 \$3, %ebx

idivl %ebx

movl %eax, %ecx

movl \$0, %edx

movl \$9, %eax

movl \$2, %ebx

idivl %ebx

movl %edx, %eax

addl %ecx, %eax

Write simple ${\tt C}$ expression that is equivalent to these instructions

Many x86 instructions can refer to memory addresses; these addresses take on many different forms. ABSOLUTE/DIRECT addressing definition: just use a number as an address movl 1000, %eax gets contents (4 bytes) of memory at address 1000, puts into %eax NOTE: DIFFERENT than movl \$1000, %eax (which just moves the VALUE 1000 into %eax) INDIRECT addressing definition: address is in register movl (%eax), %ebx treat contents of %eax as address, get contents from that address, put into %ebx BASE + DISPLACEMENT addressing definition: address in register PLUS displacement value (an offset) mov1 8(%eax), %ebx address = 8 + contents of eaxget contents from that address, put into %ebx INDEXED addressing definition: use one register as base, other as index movl 4(%eax, %ecx), %ebx address = 4 + contents[eax] + contents[ecx] get contents from that address, put into %ebx SCALED INDEXED addressing (most general form) definition: use one register as base, other as index, scale index by constant (e.g., 1, 2, 4, 8) movl 4(%eax, %ecx, 8), %ebx

address = 4 + contents[eax] + 8*contents[ecx]
get contents from that address, put into %ebx

Problem #4 (from	CSAPP 3.1)		
Memory			
Address	Value		
0x100	0xFF		
0x104	0xAB		
0x108	0x13		
0x10C	0×11		
Registers			
%eax	0x100		
%ecx	0x1		
%edx	0x3		
Value of:			
%eax			
ocax			
0×104			
\$0x108			
(%eax)			
4(%eax)			
9(%eax, %edx)			
260(%ecx, %edx)			
0xFC(,%ecx, 4)			
(%eax, %edx, 4)			

New register to help with stack: esp (extended stack pointer)							
Referred to as %esp							
[] []	eax ax ah al	32 bits 16 bits 8 bits 8 bits					
[] [] []	bx bh						
. [] [] []	ecx cx ch cl						
[] [] []	edx dx dh dl						
[]							
[]	esp	32 bits					
[]	eip	32 bits					
Points to "top of stack" when program is running Changes often (room for local variables, function call/return, etc.)							
Can use normal instructions to interact with it, e.g., addl, subl Can also use special instructions (we'll see this later)							

Problem #5

Use instructions to:

- Increase size of stack by 4 bytes
- Store an integer value 10 into the top of the stack
- Retrieve that value and put it into %ecx
- Add 5 to it
- Put final value into seax
 Decrease size of stack by 4 bytes.