

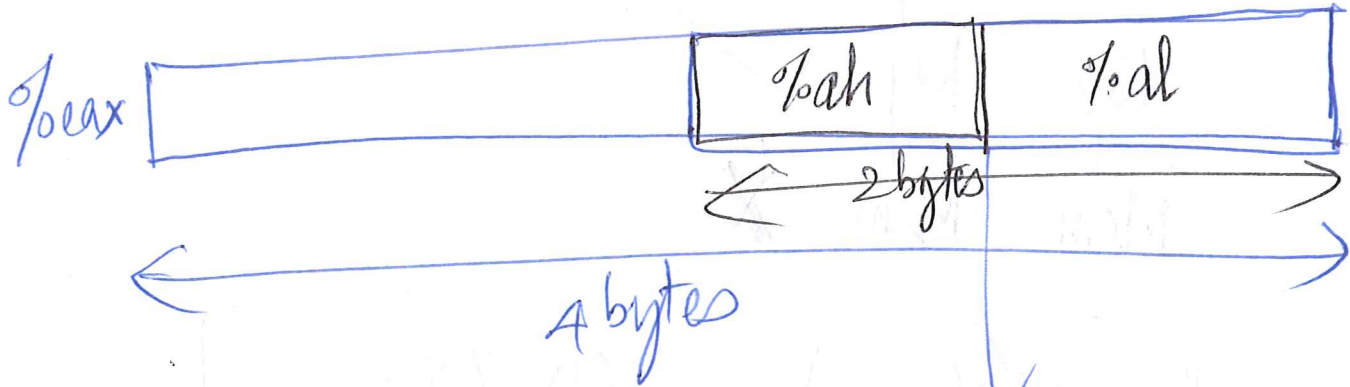
CS 354 - Lecture 7

Review

C $\xrightarrow{\textcircled{1}}$ Assembly

CPU registers. - 32-bits

(32-bits machines)



- %eax
- %ebx
- %ecx
- %edx
- %esi
- %edi

general purpose registers

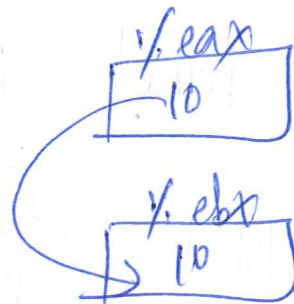
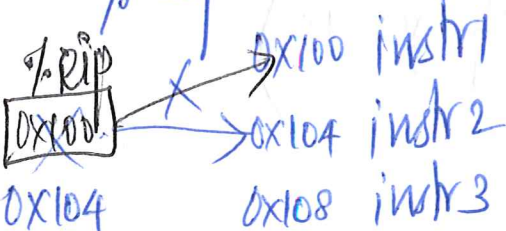
X86

AT&T

movl %eax, %ebx
S D

Special purpose registers

%eip - instruction pointer



(2)

mov S, D

Reg.	Mem.	✓
Mem	Reg	✓
Reg	Reg	✓
Mem	Mem	✗

movl 8(%eax), 0x7008

%eax
 0x3000

0x3000
 + 8
0x3008

Memory

movl 0x1000, %eax

movl N(R₁, R₂, K), %eax

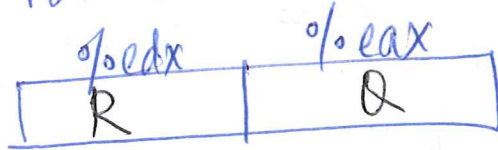
$$\text{addr} = N + \text{contents of } R_1 + (\text{contents of } R_2 * K)$$

(3)

Arithmetic

addl s, D $D \leftarrow D + S$

subl, imull, idivl



Bitwise

andl s, D $D \leftarrow D \& S$

orl, notl, xorl

Unary

incl D $D \leftarrow D + 1$

decl, negl, notl

shift

sarl K, D $D \leftarrow D \gg_A K$

shrl

sall / shll

Today

1. Control flow

- if/else

- loops

- functions.

2. Stack pointer (%esp)

3. Load Effective Address (leal).

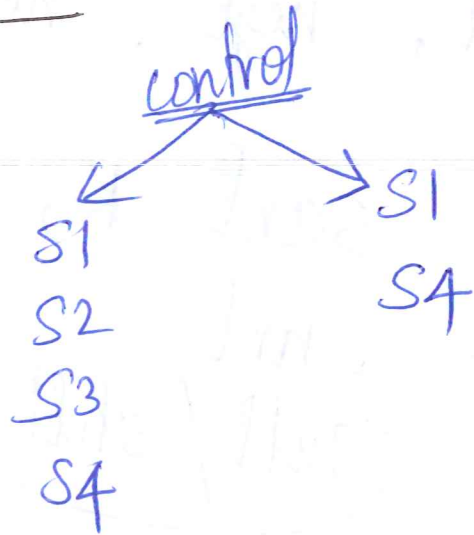
Control Flow

Decisions

```

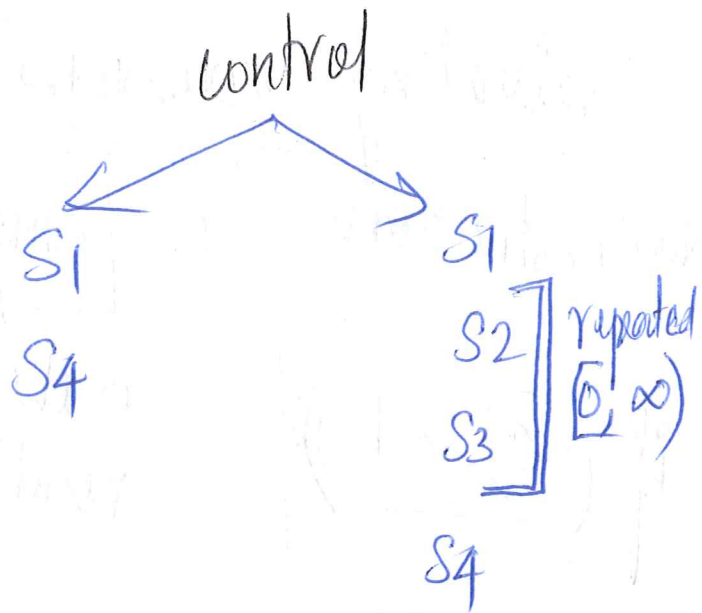
S1;
if (expr) {
    S2;
    S3;
}
S4;

```



5

```
S1;  
while (expr) {  
    S2;  
    S3;  
}  
S4;
```



What help do we need from Assembly to implement control flow?

- 1) need some mechanism to evaluate an expression. (eg boolean expr).
- 2). Ability to change the instruction pointer (program counter) to jump to some arbitrary instruction.

Evaluating expressions ^⑥

New Instructions

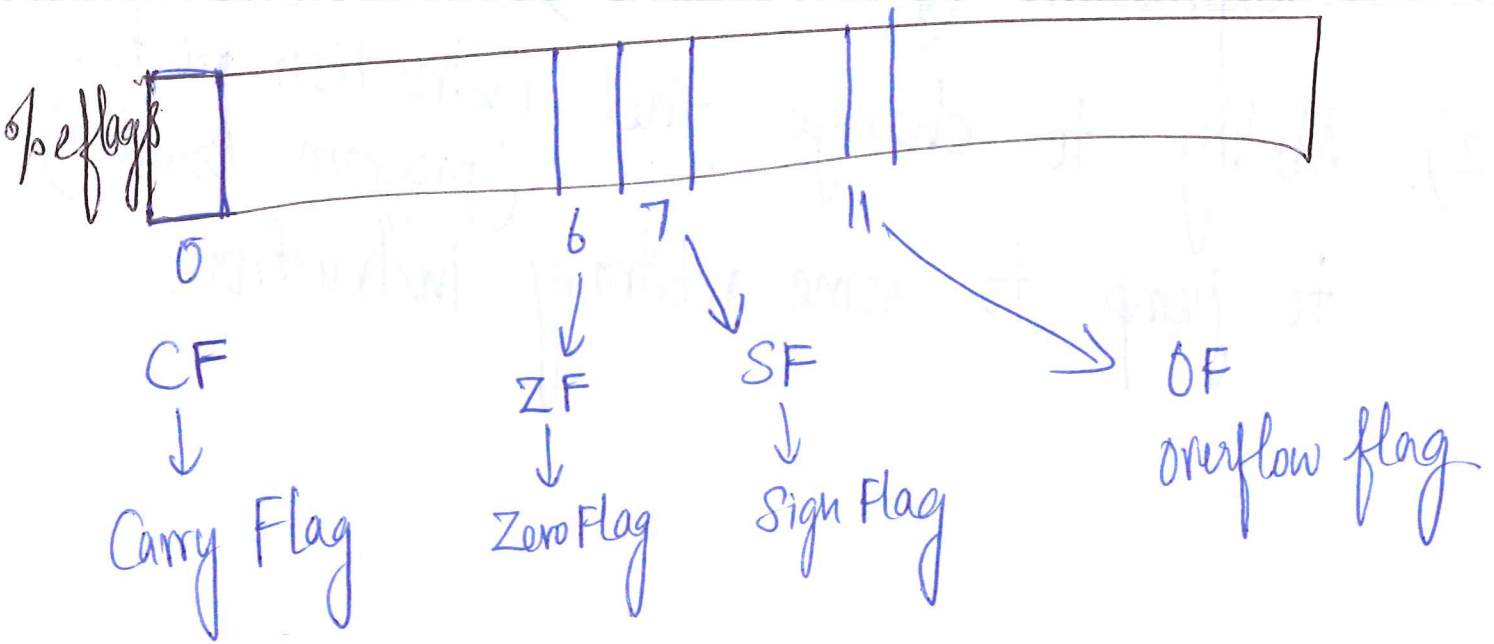
```
1. [ cmpl b, a ]
```

evaluate : $a - b$
result of this subtraction is discarded.

```
if (a > b)  
{  
}
```

→ update a special register

%eflags (hidden register)
(or)
Condition Code registers
(or)
PSW - Program Status Word.



(7)

Addition:

$$t = a + b$$

%eflags register are updated implicitly.

~~%eflags gets updated explicitly w/ compl instr.~~

int \rightarrow 3 bits

1) ZF - Zero Flag
if $(t == 0) \Rightarrow ZF = 1$ (set)
else $\Rightarrow ZF = 0$ (unset)

eg. $t = 3 + (-3) \Rightarrow$ ZF = 1

2) CF - Carry Flag. (unsigned numbers)

000 - 0
001
010
011
100
101
110
111 - 7

3 011
+ 4 100

7 111

No carry
CF = 0

3 011
+ 5 101

000

CF = 1

3) SF - Sign Flag (~~unsigned~~ signed numbers)

$$t = a + b$$

$$\text{if } (t < 0) \Rightarrow SF = 1$$

$$\text{else} \Rightarrow SF = 0.$$

eg $t = 3 + (-4) = -1 \Rightarrow \boxed{SF = 1}$

111

4) OF - Overflow flag

$$\boxed{t = a + b}$$

①

2	010
+1	001
+3	011

-4 -3 -2 -1 0 +1 +2 +3

Overflow

if $(a > 0 \ \&\& \ b > 0 \ \&\& \ t < 0)$

|| $(a < 0 \ \&\& \ b < 0 \ \&\& \ t > 0)$

②

3	011
+1	001
-4	100

$\Rightarrow \boxed{OF = 1}$

compl b, a

9

$$\bar{a} - b$$

1) ZF: $a - b == 0 \Rightarrow ZF = 1$

2) CF: $a - b < 0 \Rightarrow CF = 1$
(or) $a < b$

$a: 1 \Rightarrow \overset{0}{1} \overset{1}{0} \overset{1}{0} \overset{1}{1}$
 $-b: 2 \Rightarrow 010$

 $\textcircled{7} \Rightarrow \underline{\underline{111}}$

3) SF: $a - b < 0 \Rightarrow SF = 1$

4) OF: $\text{if } \left(\begin{aligned} &\cancel{a} > 0 \ \&\& \ b < 0 \ \&\& \ (a - b) < 0 \\ &\parallel \ (a < 0 \ \&\& \ b > 0 \ \&\& \ (a - b) > 0) \end{aligned} \right)$

$\Rightarrow \boxed{OF = 1}$

(10)

New instruction:

testl b, a

evaluate: a & b

→ explicitly sets the %eflags.

1) ZF: a & b == 0

2) SF: a & b < 0

testl %eax, %eax

1) if %eax = 0 : ZF = 1
SF = 0 } ⇒ value in %eax = 0

2) if %eax < 0 : ZF = 0
SF = 1 } ⇒ value in %eax < 0.
eg. 111
&111

011

3) if %eax > 0 : ZF = 0
SF = 0 } ⇒ value in %eax > 0
001
&001

001

Jump instructions

(11)

instr 1
jmp target

instr 2

instr 3

target:

instr 4

instr 5

instr 1

jz target

instr 2

instr 3

target:

instr 4

instr 5

cmpl b, a

Jumps

unconditional

jmp <target>

conditional

~~jz~~
~~jnz~~

jl
jle
jg
jge

signed numbers

jb
jbe
ja
jae

if (a < b) {
 ↓ ↓
 %eax %ebx
}

cmpl %ebx, %eax
jl label

a-b

(12)

Instruction

Jump condition

jnb jbe jae	cmpl b, a a-b CF == 1	CF
	CF == 1 OR ZF == 1	ZF == 1

CF == 0 and ZF == 0

CF == 0 ~~CF~~

jl

SF \wedge OF

cmpl b, a
a = 1; b = 2

$SF == 1 \wedge OF == 0$
= 1

$$\begin{array}{r} 1 \ 001 \\ - 2 \ 010 \\ \hline -1 \ 111 \end{array}$$

$SF == 0 \wedge OF == 1$
= 1

a = -4 b = 3

$$\begin{array}{r} -4 : 100 \\ - 3 : 011 \\ \hline 001 \end{array}$$

De Morgan's Law

$\sim (A \mid B) = \sim A \ \& \ \sim B$

(13)

Loops

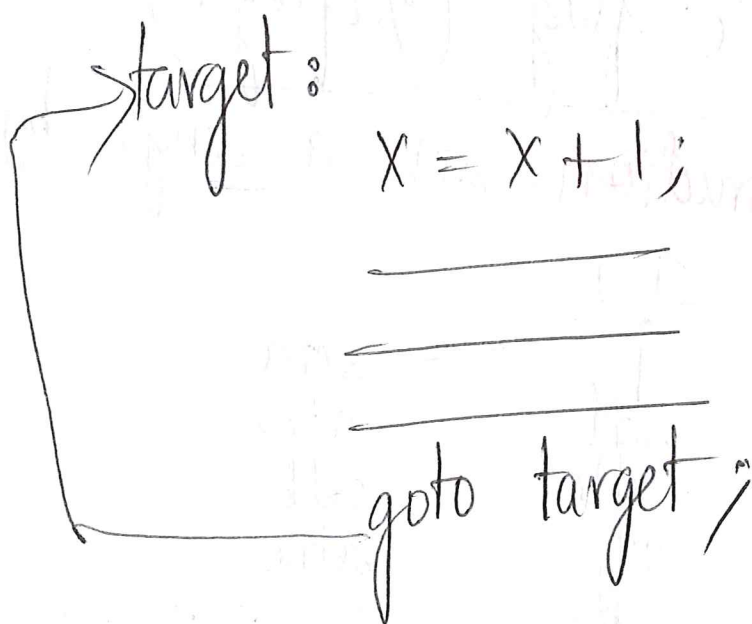
for

while

do-while

```
for (init-expr; cond; post-expr) {  
    loop body  
}
```

```
init-expr;  
while (cond) {  
    loop body;  
    post-expr;  
}
```

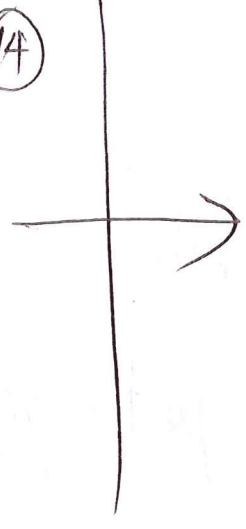


14

```

while (expr) {
  loop body;
}

```



```

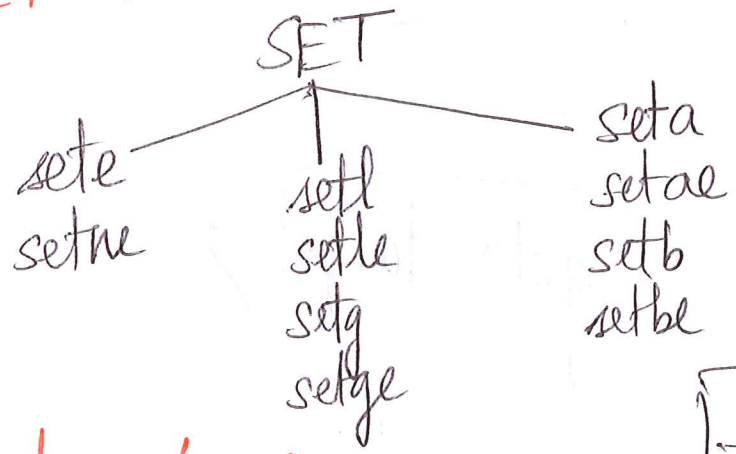
if (!expr)
  goto end;
do {
  body of loop;
} while (expr);

```

end: _____

Access the CC flags (%eflags)

SET instruction - sets a single byte to 0 or 1.



```

cmlt %ebx, %eax

```

```

setl %al

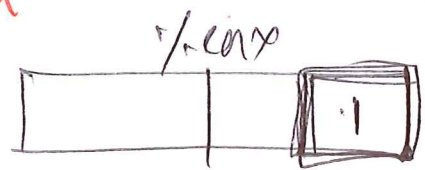
```

```

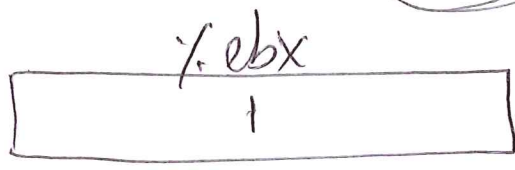
movzbl %al, %ebx

```

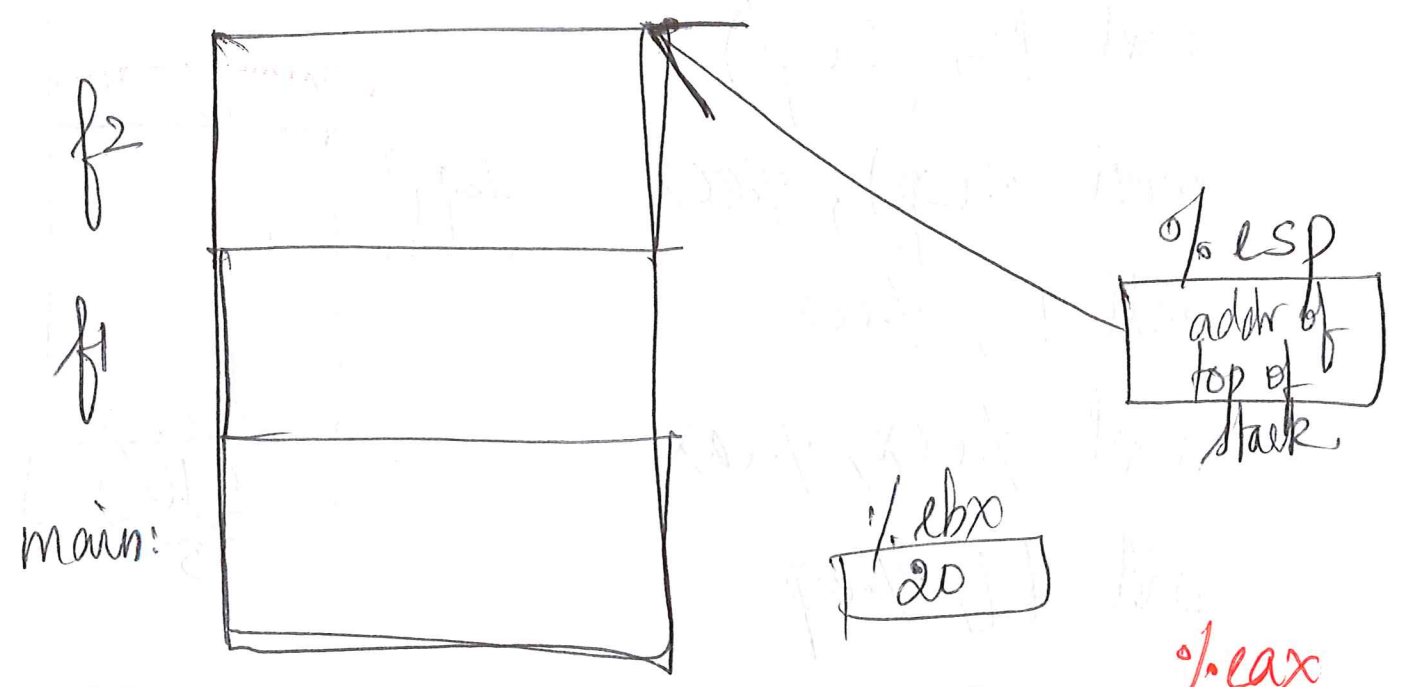
%eax	%ebx
a: 1	b: 2



(a - b)

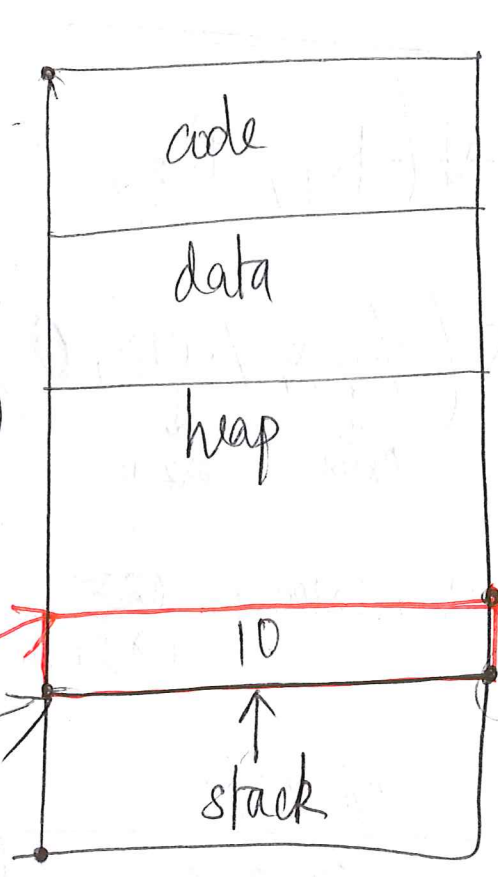


Stack pointer



```
push %eax
```

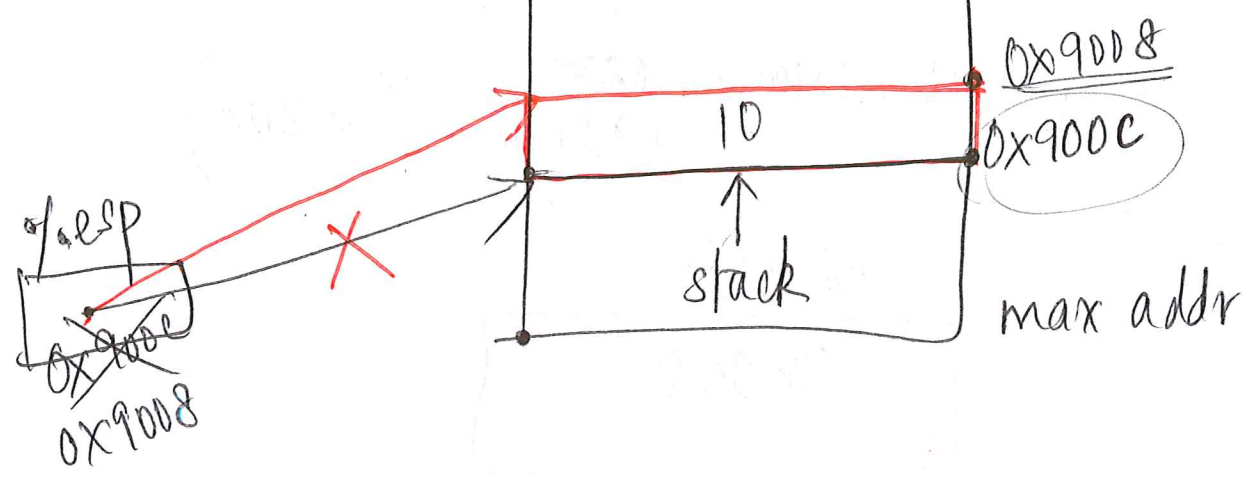
```
subl $4, %esp
movl %eax, (%esp)
```



```
popl %ebx
movl (%esp), %ebx
addl $4, %esp
```

%eax
10

```
popl %ebx
```



16

add \$-4, %esp

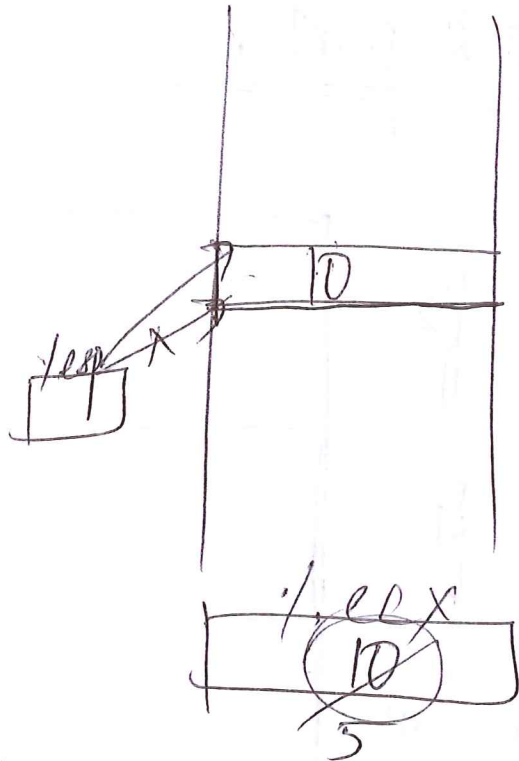
movl \$10, (%esp)

movl (%esp), %ecx

addl \$5, %ecx

movl %ecx, %eax

addl \$4, %esp



leal N(R₁, R₂, K), %eax

leal 8(%eax, %ebx, 0), %eax

addr: $8 + 0x100 + \cancel{0x200} = 0x308$

%eax
 0x308

Problem #6

Assume value of a is in %eax, and value of b is in %ebx

Write x86 assembly code for:

```
if (a > b) {
    a++;
}
```

```
cmpl %ebx, %eax (a-b)
jg do
jmp dont
do: addl $1, %eax
dont: instr.
```

```
cmpl %ebx, %eax
jle dont
addl $, %eax
dont: instr.
```

(a-b)

Problem #7

Assume value of a is in %eax, and value of b is in %ebx

Write x86 assembly code for:

```
if (a > b) {
    a++;
} else {
    b = a;
}
```

```
cmpl %ebx, %eax (a-b)
jle dont
addl $1, %eax
jmp end
dont: movl %eax, %ebx
end: instr after if/else
```

(a-b)

Problem #8

Assume value of a is in %eax, and value of b is in %ebx

Write x86 assembly code for:

```
while (b > 0) {
    a++;
    b--;
}
```

```
loop:
cmpl $0, %ebx (b-0)
jle end
incl %eax
decl %ebx
jmp loop
```

(b-0)

end: instr(s)

Condition codes: new bits in hidden %eflags register.

Some instructions set those bits based on comparisons:

cmp, test

Other instructions change control flow (%eip) based on results:

jmp family

INSTRUCTION: **cmpl B, A**

computes A-B (but doesn't put result anywhere)

condition codes (incomplete):

zero flag : ZF=1 if (A-B) == 0 otherwise ZF=0

signed flag : SF=1 if (A-B) < 0 otherwise SF=0

INSTRUCTION: **jmp TARGET** always changes %eip to TARGET

INSTRUCTION: **je TARGET** %eip=TARGET if ZF==1

INSTRUCTION: **jne TARGET** %eip=TARGET if ZF== _____

INSTRUCTION: **jg TARGET** %eip=TARGET if _____

INSTRUCTION: **jge TARGET** %eip=TARGET if _____

INSTRUCTION: **jl TARGET** %eip=TARGET if _____

INSTRUCTION: **jle TARGET** %eip=TARGET if _____

Problem #6

Assume value of a is in %eax, and value of b is in %ebx

Write x86 assembly code for:

```
if (a > b) {  
    a++;  
}
```

Problem #7

Assume value of a is in %eax, and value of b is in %ebx

Write x86 assembly code for:

```
if (a > b) {  
    a++;  
} else {  
    b = a;  
}
```

Problem #8

Assume value of a is in %eax, and value of b is in %ebx

Write x86 assembly code for:

```
while (b > 0) {  
    a++;  
    b--;  
}
```

Condition codes: new bits in hidden %eflags register.

Some instructions set those bits based on comparisons:

cmp, test

Other instructions change control flow (%eip) based on results:

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INSTRUCTION: `cmpl B, A`

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zero flag : ZF=1 if (A-B) == 0 otherwise ZF=0

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INSTRUCTION: `jmp TARGET` always changes %eip to TARGET

INSTRUCTION: `je TARGET` %eip=TARGET if ZF==1

INSTRUCTION: `jne TARGET` %eip=TARGET if ZF== 0

INSTRUCTION: `jg TARGET`

%eip=TARGET if

$\neg(SF \wedge OF)$ & $\neg ZF$

INSTRUCTION: `jge TARGET`

%eip=TARGET if

$\neg(SF \wedge OF)$

INSTRUCTION: `jle TARGET`

%eip=TARGET if

$SF \wedge OF$

INSTRUCTION: `jle TARGET`

%eip=TARGET if

$(SF \wedge OF) \vee ZF$