Xv6 File System

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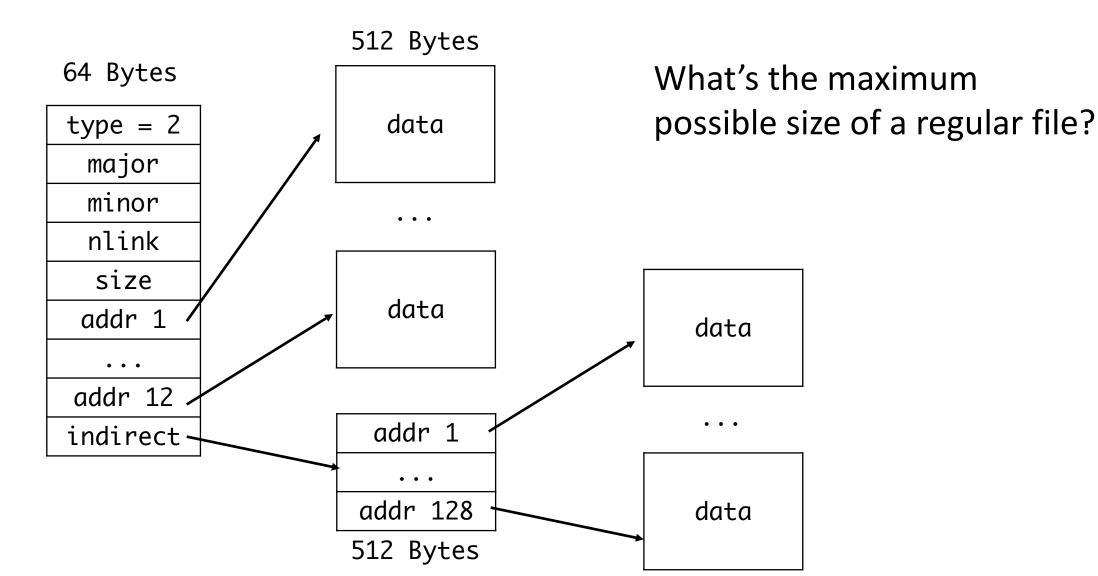
11/28/2017

File System Layout in xv6

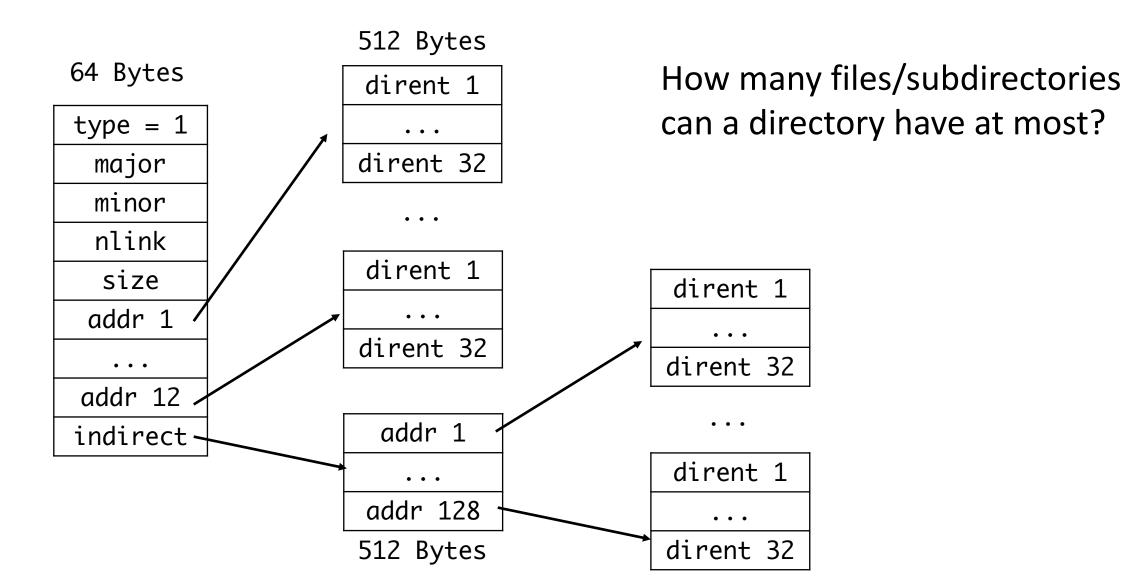
Unused | Superblock | Inodes ... | Unused | Bitmap | Data ...

```
50 struct dirent {
51  ushort inum;
52  char name[DIRSIZ];
53 };
```

Inode of a regular file in xv6



Inode of a directory in xv6



Bitmap in xv6

- Each bit in the bitmap is associated with a block, NOT an inode.
- Although the very first block is unused, it is always marked as 1 in the first bit of the bitmap, and so are all the blocks where the inodes and bitmap itself reside.
- Bitmap is grouped in byte.
- Intel x86 processors use little-endian.
- Example:
 - ff c2 => 1111 1111 1100 0010 7 6 5 4 3 2 1 0 | 15 14 13 12 11 10 9 8

Demos

- How to build your own file system image and reflect in xv6
 - Closer look at mkfs.c
- How xv6 files change the image
 - Closer look at fs.img with xxd
 - Some tricks editing file images with vim:
 - :%!xxd to open the image; :%!xxd -r to save changes
- How to read image? mmap()!

Demos

- Using GDB, go through the entire write system call.
 - Key methods to understand:
 - writei() in kernel/sysfile.c
 - bmap() in kernel/fs.c
 - bread(), bwrite() in kernel/bio.c
- How does xv6 normally handle large files?
- What do we do to handle small files here?

