

Flash-based SSDs

Solid State Storage ^① Device - SSD.

flash → NAND-based flash.

flash
page flash
block.

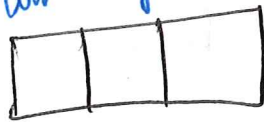
⇒ built only w/ transistors
(like DRAM).

1. storing a single bit

01 : SLC - Single Level Cell

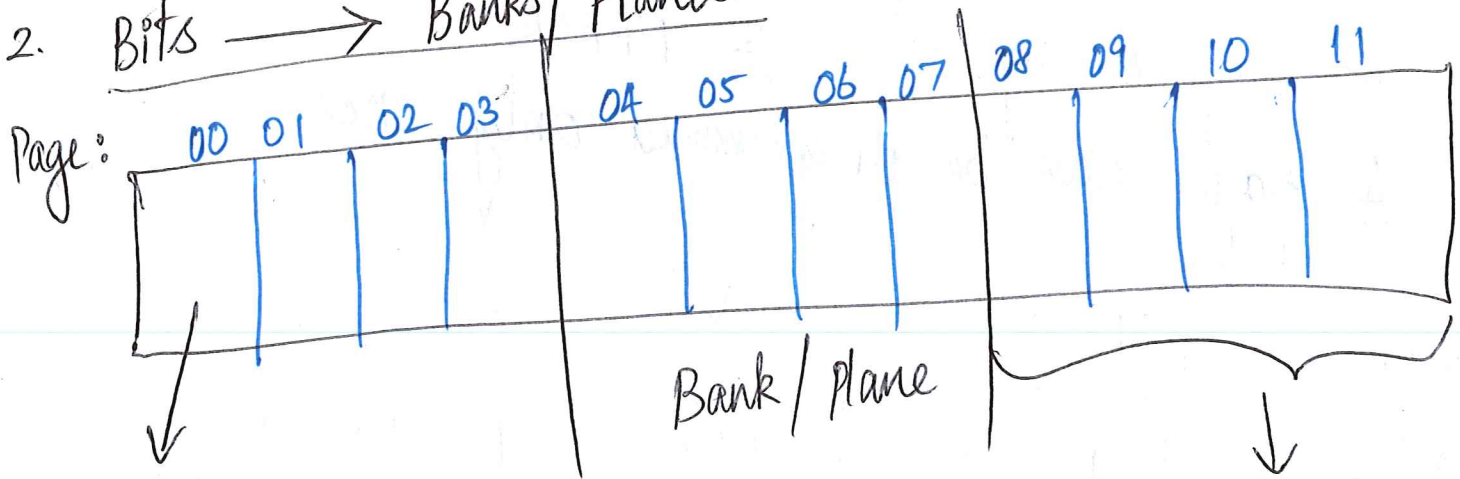
00, 01, 10, 11 : MLC - Multi " "

low somewhat low high high



: TLC - Triple-level cell.

2. Bits → Banks / Planes.



1 page
(eg, 4KB)

1 block.
(eg, 128KB or 256KB)

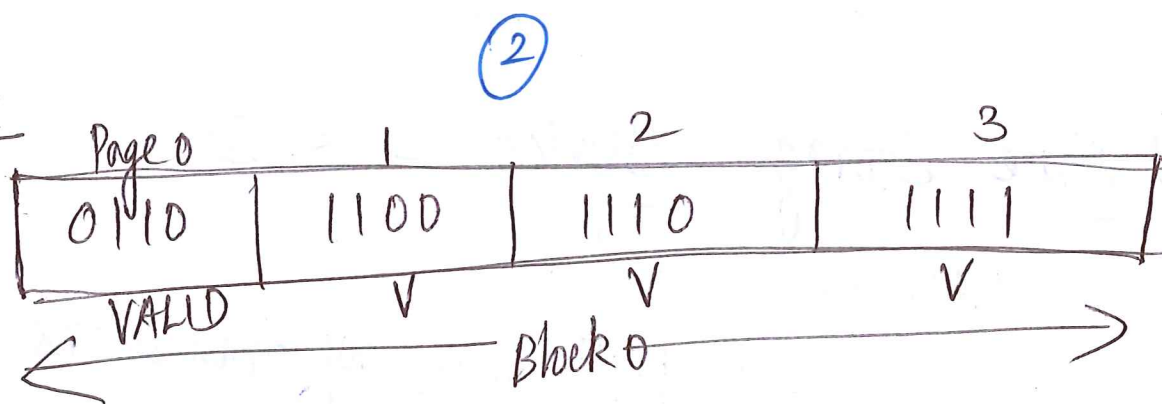
3. Basic Flash Operations

1. Read (a page) - 10s of μ s

2. Erase (a block) - few ms → COSTLY!

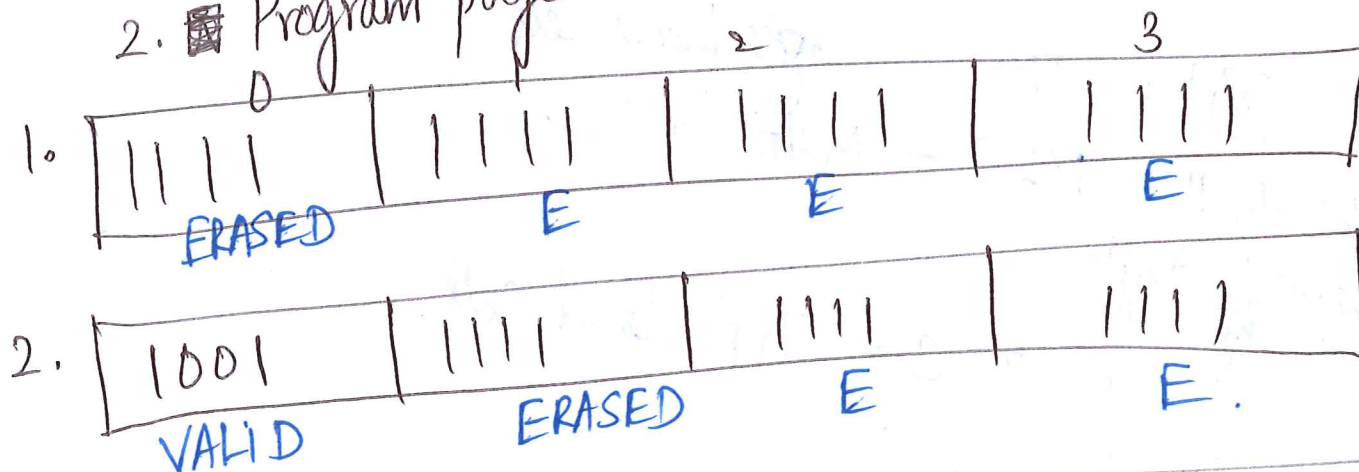
3. Program (a page) → 100s of μ s.

Example



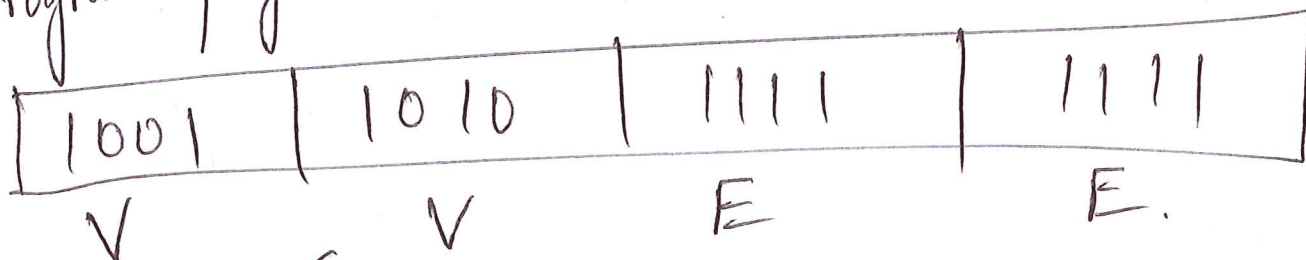
Write to page 0:

1. Erase Block 0
2. Program page 0.

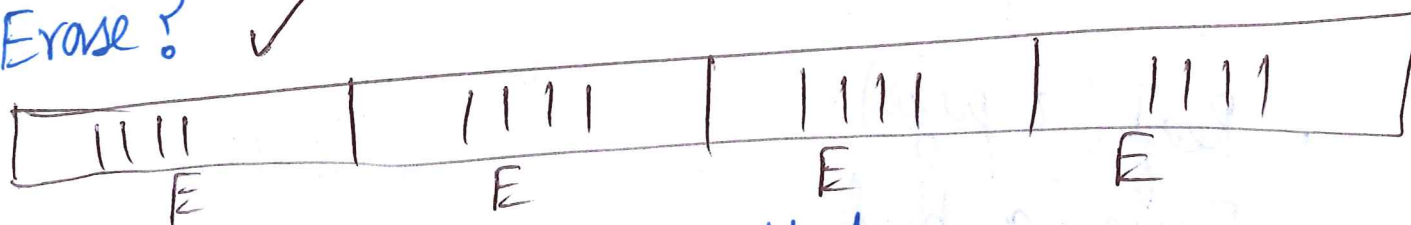


Program page 0 again? ⇒ ERROR
 A page can be programmed only once.

Program page 1 now? ⇒ ✓



Erase? ✓



Now: ALL pages programmable!

4. Flash - Performance & ^③ Reliability.

SHOW TABLE IN BOOK.

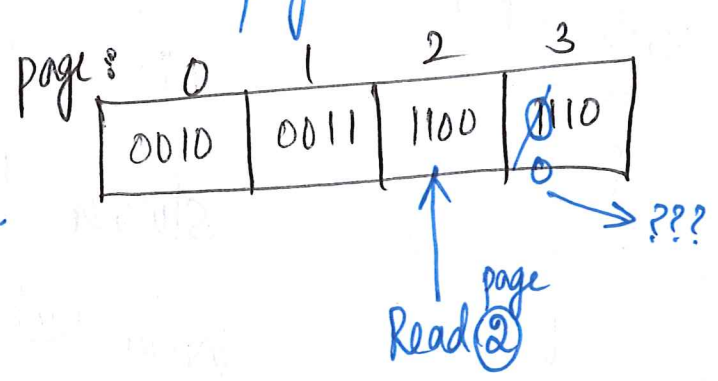
→ Head crash in HDD. - NOT in flash (:: just silicon)

→ wear out.

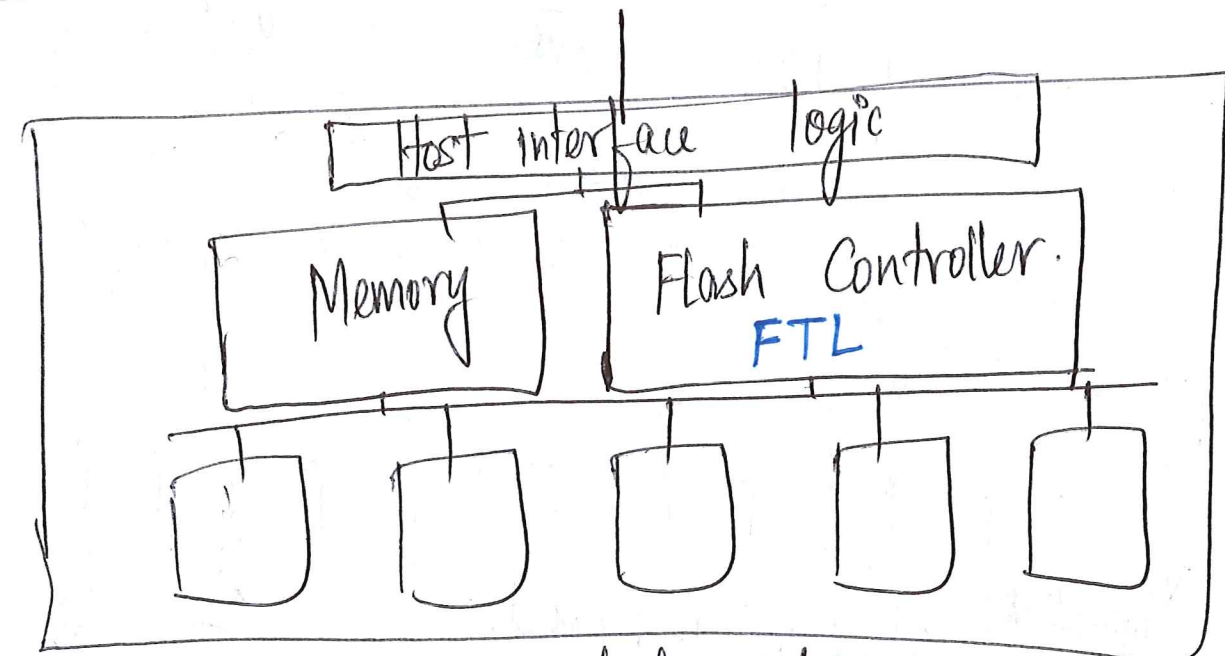
SLC - 100,000 P/E cycles

program → erase.

→ disturbance
read disturbs
program disturbs.



5. Raw Flash → Flash based SSDs



FTL → Flash Translation Layer.

Read/Write Logical block 0



Read/Erased/Program Physical blocks/pages

6. A Bad FTL organization ⁽⁴⁾

"direct mapped"

e.g. Logical ~~addr~~ page N

→ physical ~~addr~~ page N.

1. Read logical page N → Read phy. page N

2. WRITE.

Write Amplification

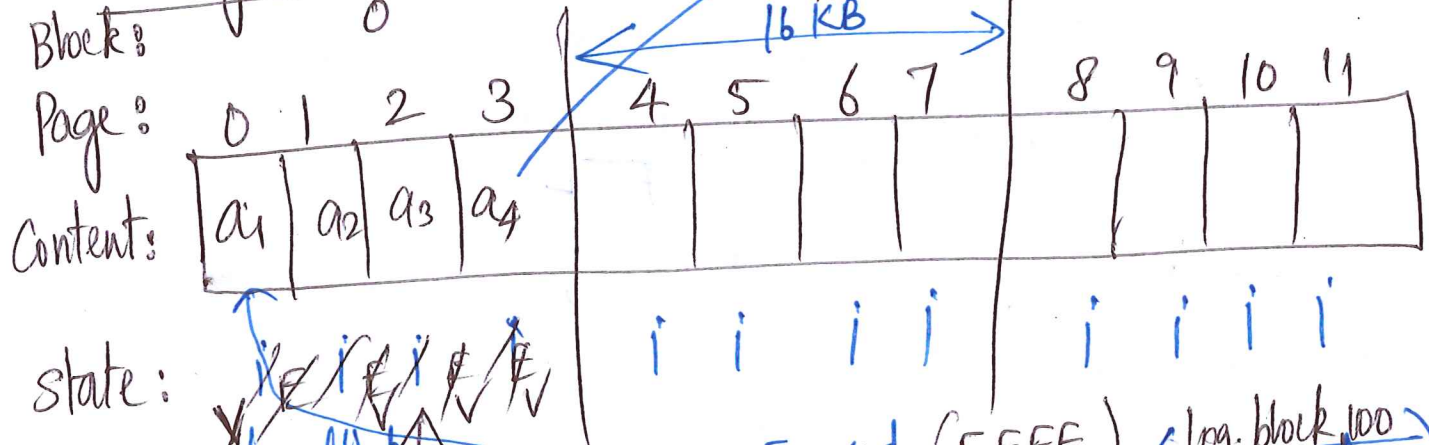
1. Read entire block with phy. page N.
2. ERASE the block
3. PROGRAM the old + new pages.

Problems?

→ Performance! - slower than HDDs.

→ Reliability? - Wear out of blocks
(metadata) freq. written.

7. A Log-Structured FTL



Write (100)

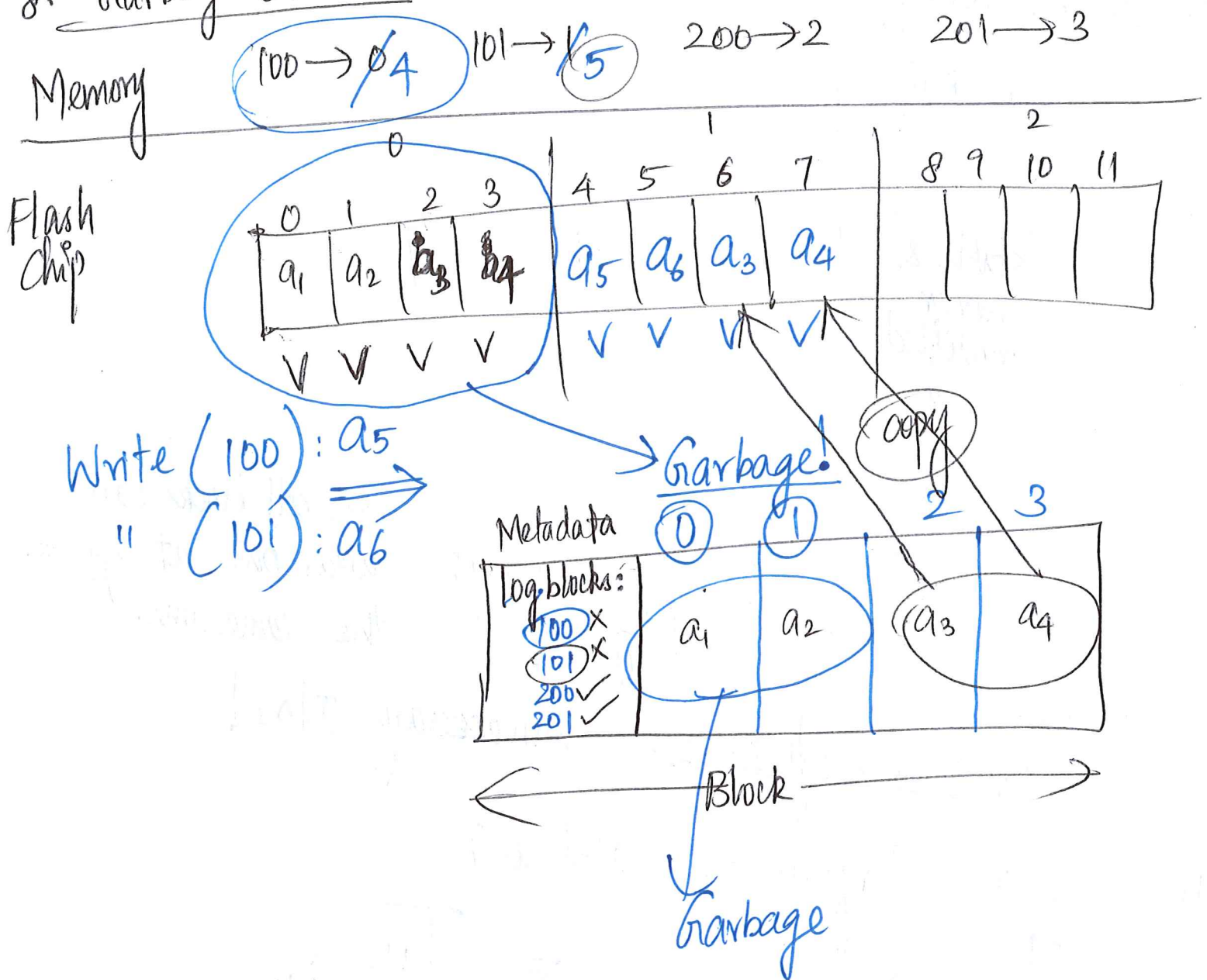
Block 0:

1. Erased. (EEEE)
2. Write page 0; (log. block 100 → phy. page 0)
3. update the mapping table in memory.

(X) Wear-leveling!

CRASH → OOB And Area
in each page 100 → 0, 101 → 1, 200 → 2, 201 → 3

8. Garbage Collection



9. Mapping Table Size

page-level mapping.

block-level mapping

hybrid mapping.

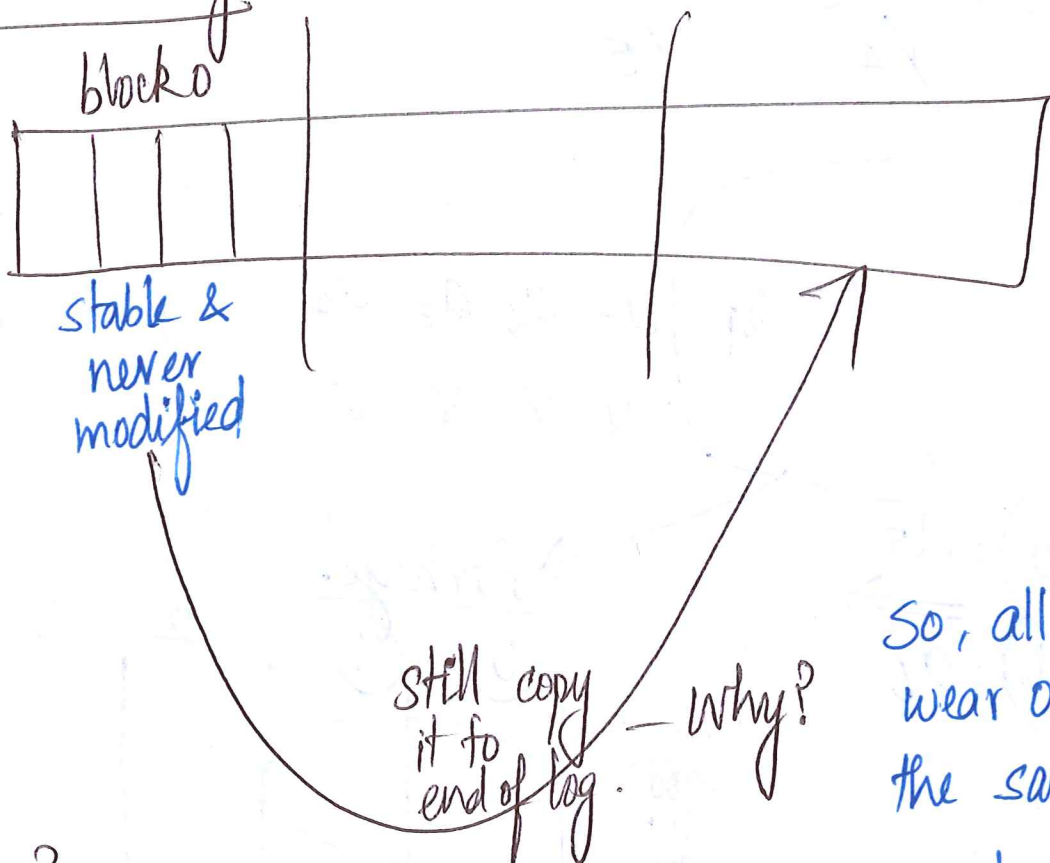
READ OSTEP FOR DETAILS!

$$\frac{1\text{TB (Flash mem)}}{4\text{KB}} = 2^{28}$$

if 1 mapping = 4 bytes $\Rightarrow$ (1GB) for mapping table.

10. Wear Leveling

(6)



So, all blocks can wear out at approx. the same time.

Con?

Write amplification — unnecessary I/Os!

11. SSD Performance and Cost

SSD Random I/O $\ll$ HDD Random I/O

$T_{\text{Random Reads}_{\text{SSD}}} > T_{\text{Random Writes}_{\text{SSD}}}$

why?

$\text{Cost}_{\text{SSD}} > \text{Cost}_{\text{HDD}}$
 60 cents/GB 5 cents/GB

	Cap.	Cost
SSD	250GB	\$150
HDD	1TB	\$50.