

CS/ECE 252: INTRODUCTION TO COMPUTER ENGINEERING
UNIVERSITY OF WISCONSIN—MADISON

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Midterm Examination 4

In Class (50 minutes)

Wednesday, Dec 11, 2013

Weight: 17.5%

NO: BOOK(S), NOTE(S), OR CALCULATORS OF ANY SORT.

The exam has 10 pages. **Circle your final answers.** Plan your time carefully since some problems are longer than others. **You must turn in the pages 1-8.** Use the blank sides of the exam for scratch work.

Note: LC-3 instruction set is provided on Page 9. Trap Codes and Assembler Directives are provided on the last page.

LAST NAME: _____

FIRST NAME: _____

ID# _____

Problem	Maximum Points	Points Earned
1	4	
2	8	
3	8	
4	4	
5	6	
Total	30	

Problem 1: Assembly Errors

(4 Points)

Identify 4 errors in the following Assembly language program:

```
        .ORIG x3501

        LD R0, DEFAULT
LOOP    TRAP x20
        TRAP x21
        AND R0, R0, R0
        BRzP LOOP
        TRAP x20
        NOT R2, #3          ----- (a)
        ADD R0, R0, ONE    ----- (b)
        LD R1, R2, #3
        BRP LOOP

LOOP    HALT              ----- (c)
ONE     .FILL            x1
DEFAULT .ASCII          #2  ----- (d)
        .END
```

(a) **NOT** cannot have an immediate operand

(b) **ADD** cannot have a memory argument

(c) **Double declaration of Loop**

(d) **.ASCII** assembler directive does not exist in LC-3.

Problem 2: Two Pass Assembly Process

(8 Points)

(a) **(2 Points)** Consider the following assembly code. What will be the output on the console if you run this code on PennSim?

```
.ORIG x3800
LEA R3, INPUT
LD R1, SIZE
ADD R3, R3, R1

LOOP LDR R0, R3, 0
TRAP x21
ADD R3, R3, -1
ADD R1, R1, -1
BRp LOOP

HALT

INPUT .STRINGZ "RtbY"
STRING .BLKW 3
SIZE .FILL 3
.END
```

Answer: **Ybt**

(b) **(4 Points)** In the first pass, the assembler creates the symbol table. Fill in the symbol table created by the assembler for the program.

Label	Address
LOOP	3803
INPUT	3809
STRING	380E
SIZE	3811

(c) (2 Points) In the second pass, the assembler creates a binary (.obj) version of the program, using the entries from the symbol table. Assume that there exists another program at 0x3000, whose assembly instructions are as shown below. If the following symbol table entries were generated in the first pass of the assembly for this program, write the binary code generated by the assembler for the two instructions at 0x3000 and 0x3001.

Symbol Table:

Label	Address
INT	x3021
LOOP	x3011

Generated Binary code:

Address	Instruction	Binary Code
x3000	LD R0, INT	0010 0000 0010 0000
x3001	BRp LOOP	0000 0010 0000 1111

- 1) $PC' + \text{offset} = x3021$
 $x3000 + 1 + \text{offset} = x3021$
 $\Rightarrow \text{Offset} = x20$
- 2) $PC' + \text{offset} = x3011$
 $x3001 + 1 + \text{offset} = x3011$
 $\Rightarrow \text{Offset} = xF$

Problem 3: Subroutines and Traps**(8 Points)**

Suppose we want to write a new TRAP subroutine, TRAP x33, which takes a string input from the user. The trap subroutine starts from address x2200 and does the following:

- 1) It takes a character input from the user
- 2) It then displays this character (which the user inputs) on the console
- 3) After that, it stores the user input characters in consecutive memory locations starting from the address location present in register R1. It then repeats (1), (2) and (3) until user inputs 'Z'.
- 4) It uses a "callee-save" strategy and ensures that none of the register values are modified by it.
- 5) It uses R2 to store the ASCII value corresponding to -Z

(a) **(6 Points)** Fill in the missing parts of the trap subroutine.

```

.ORIG x2200
    ST R0, SAVEREG1
    ST R1, SAVEREG2
    ST R2, SAVEREG3
    ST R7, SAVEREG4
    LD R2, NEGZ

NEXT  TRAP x20
      TRAP x21
      STR R0, R1, #0
      ADD R1, R1, #1
      ADD R0, R0, R2
      BRnp NEXT

      LD R0, SAVEREG1
      LD R1, SAVEREG2
      LD R2, SAVEREG3
      LD R7, SAVEREG4
      RET

;Data Region
NEGZ      .FILL xFFA6      ;xFFA6 = FFFF - ASCII value of Z
SAVEREG1  .BLKW 1
SAVEREG2  .BLKW 1
SAVEREG3  .BLKW 1
SAVEREG4  .BLKW 1

```

(b) **(2 Points)** Given the following Trap vector table entry:

Address	Content
x44	X26

Give the assembly instruction that you would use to call the TRAP routine corresponding to this entry. Provide reasons to justify your answer.

Ans: **TRAP x44** because the starting location of trap routine is stored in x44.

Problem 4: I/O**(4 Points)**

(a) **(4 Points)** The following code segment should display the string specified at the “INPUT” label on to the console. Write the missing assembly instructions of the program (without using PUTS/PUTC/TRAP instructions).

- **Note:** The instructions which are missing should jump to halt if it is the end of the string. Else, it should print the character onto the console.

```
.ORIG    x3000

        LEA R2, INPUT

NEXT    LDR R0, R2, #0
        BRz END
POLL    LDI R1, DSR
        BRzp POLL
        STI R0, DDR

        ADD R2, R2, #1    ; Point to the next character
        BR NEXT

END     HALT

INPUT   .STRINGZ    "All the best!" ; String to display
DSR     .FILL      xFE04    ; Display status register location
DDR     .FILL      xFE06    ; Display data register location
        .END
```

Problem 5: Short Answer Questions

(6 Points)

- (a) **(1 Point)** Briefly state a scenario where you would prefer interrupt-driven I/O over polling based I/O?

If the I/O device takes a lot of time to execute the command, then polling consumes a lot of cycles. In these cases, interrupt-driven I/O is preferred.

- (b) **(1 Point)** Suppose two I/O devices sends interrupts to the CPU at the same time. How does the CPU decide which interrupt to service first?

The one with a higher priority (ie, at a higher priority level) is executed first.

- (c) **(2 Points)** An LC-3 assembly program contains the following instruction:

```
FLOAT      LD R2, FLOAT
```

The symbol table entry for FLOAT is x3000. What will be the value of R2 after the execution of the above instruction?

Ans : x25FF (the binary code for this instruction).

- (d) **(2 Points)** Briefly state what happens in Linking and Loading phases for an assembly program?

During the Linking phase, the symbols between different object files which are linked together gets resolved.

During the loading phase, the executable image is copied onto the memory.

LC-3 Instruction Set (Entered by Mark D. Hill on 03/14/2007; last update 03/15/2007)

PC': incremented PC. setcc(): set condition codes N, Z, and P. mem[A]:memory contents at address A.

SEXT(immediate): sign-extend immediate to 16 bits. ZEXT(immediate): zero-extend immediate to 16 bits.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0						

0	0	0	1		DR		SR1		0	0	0		SR2								

ADD DR, SR1, SR2 ; Addition																					
DR ← SR1 + SR2 also setcc()																					

0	0	0	1		DR		SR1		1		imm5										

ADD DR, SR1, imm5 ; Addition with Immediate																					
DR ← SR1 + SEXT(imm5) also setcc()																					

0	1	0	1		DR		SR1		0	0	0		SR2								

AND DR, SR1, SR2 ; Bit-wise AND																					
DR ← SR1 AND SR2 also setcc()																					

0	1	0	1		DR		SR1		1		imm5										

AND DR,SR1,imm5 ; Bit-wise AND with Immediate																					
DR ← SR1 AND SEXT(imm5) also setcc()																					

0	0	0	0		n		z		p		PCoffset9										

BRx,label (where x={n,z,p,zp,np,nz,nzp}); Branch																					
GO ← ((n and N) OR (z AND Z) OR (p AND P))																					
if(GO is true) then PC←PC'+ SEXT(PCoffset9)																					

1	1	0	0		0	0	0		BaseR						0	0	0	0	0	0	

JMP BaseR ; Jump																					
PC ← BaseR																					

0	1	0	0		1		PCoffset11														

JSR label ; Jump to Subroutine																					
R7 ← PC', PC ← PC' + SEXT(PCoffset11)																					

0	1	0	0		0	0	0		BaseR						0	0	0	0	0	0	

JSRR BaseR ; Jump to Subroutine in Register																					
temp ← PC', PC ← BaseR, R7 ← temp																					

0	0	1	0		DR		PCoffset9														

LD DR, label ; Load PC-Relative																					
DR ← mem[PC' + SEXT(PCoffset9)] also setcc()																					

1	0	1	0		DR		PCoffset9														

LDI DR, label ; Load Indirect																					
DR←mem[mem[PC'+SEXT(PCoffset9)]] also setcc()																					

0	1	1	0		DR		BaseR						offset6								

LDR DR, BaseR, offset6 ; Load Base+Offset																					
DR ← mem[BaseR + SEXT(offset6)] also setcc()																					

1	1	1	0		DR		PCoffset9														

LEA, DR, label ; Load Effective Address																					
DR ← PC' + SEXT(PCoffset9) also setcc()																					

1	0	0	1		DR		SR		1	1	1	1	1	1	1						

NOT DR, SR ; Bit-wise Complement																					
DR ← NOT(SR) also setcc()																					

1	1	0	0		0	0	0		1	1	1		0	0	0	0	0	0			

RET ; Return from Subroutine																					
PC ← R7																					

1	0	0	0		0	0	0		0	0	0	0	0	0	0	0	0	0			

RTI ; Return from Interrupt																					
See textbook (2 nd Ed. page 537).																					

0	0	1	1		SR		PCoffset9														

ST SR, label ; Store PC-Relative																					
mem[PC' + SEXT(PCoffset9)] ← SR																					

1	0	1	1		SR		PCoffset9														

STI, SR, label ; Store Indirect																					
mem[mem[PC' + SEXT(PCoffset9)]] ← SR																					

0	1	1	1		SR		BaseR						offset6								

STR SR, BaseR, offset6 ; Store Base+Offset																					
mem[BaseR + SEXT(offset6)] ← SR																					

1	1	1	1		0	0	0		trapvect8												

TRAP ; System Call																					
R7 ← PC', PC ← mem[ZEXT(trapvect8)]																					

1	1	0	1																		

; Unused Opcode																					
Initiate illegal opcode exception																					
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0						

TRAP CODES

Code	Equivalent	Description
HALT	TRAP x25	Halt execution and print message to console.
IN	TRAP x23	Print prompt on console, read (and echo) one character from keybd. Character stored in R0[7:0].
OUT	TRAP x21	Write one character (in R0[7:0]) to console.
GETC	TRAP x20	Read one character from keyboard. Character stored in R0[7:0].
PUTS	TRAP x22	Write null-terminated string to console. Address of string is in R0.

ASSEMBLER DIRECTIVES

Opcode	Operand	Meaning
.ORIG	address	starting address of program
.END		end of program
.BLKW	n	allocate n words of storage
.FILL	n	allocate one word, initialize with value n
.STRINGZ	n-character string	allocate n+1 locations, initialize w/characters and null terminator