

# Homework 1

Due on October 6

## A: QUERY CONTAINMENT [50%]

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1. [9%] For each of the following pair of queries  $q, q'$  decide whether  $q \subseteq q'$ ,  $q' \subseteq q$ , both or none. Explain your answer.
  - (a)  $q(x, y) : -R(x, z), S(z, z), R(z, y)$   
 $q'(y, w) : -R(y, z), S(z, t), R(t, w)$
  - (b)  $q(x) : -R(x, y), R(y, z), R(z, x)$   
 $q'(x) : -R(x, y), R(y, z), R(z, u), R(u, v), R(v, z)$
  - (c)  $q(x) : -R(x, y), R(x, z), R(z, y)$   
 $q'(x) : -R(x, y), R(y, z), R(z, u)$
2. [6%] Find the minimal equivalent CQ to  $S_k(x) : -R(x, z_1), R(x, z_2), \dots, R(x, z_k)$  for any  $k \geq 1$ . Explain your answer.
3. [10%] Consider the following two CQs:  $q_1(x) : -R(x, y), R(y, z)$  and  $q_2(x) : -R(x, y)$ . The notation  $q \subsetneq q'$  means that  $q \subseteq q'$ , but that  $q$  is not equivalent to  $q'$ . Does there exist a conjunctive query  $q$  such that  $q_1 \subsetneq q \subsetneq q_2$ ? Explain your answer.
4. [10%] Consider the following two queries in  $\text{CQ}^<$ :

$$q() : -R(x, 10), x < 10$$
$$q'() : -R(y, 10), R(9, y), y \leq 10$$

Decide whether  $q \subseteq q'$ ,  $q' \subseteq q$ , both or none. Explain your answer.

5. [15%] We say that a conjunctive query  $q$  has a *self join* if its body contains two atoms of the same relation. For example, the query  $q() : -R(x, y), R(y, z)$  has a self-join, but the query  $q'(x) : -R(x, y), S(y, z, w)$  does not. What is the complexity for query containment of CQs that have no self join? You should either provide a polynomial time algorithm or show that the problem is NP-complete.

## B: QUERY COMPLEXITY [50%]

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1. [15%] We have shown that the combined complexity of query evaluation for CQs is NP-complete. However, the reduction uses queries where the relations are binary (have arity 2).
  - (a) What is the combined complexity of CQ evaluation if we restrict the relations to be *unary* (i.e. have arity exactly 1)? For example, such a query would be  $q(x) = R(x), S(x), T(y)$ .
  - (b) What is the combined complexity of CQ<sup>≠</sup> evaluation (we can add disequalities) with relations that are unary? For example, such a query would be  $q() = R(x), S(y), T(z), x \neq y, y \neq z$ .
2. [15%] For the following conjunctive queries, determine if they are acyclic or not. If the query is acyclic, construct the join tree of the query as well.
  - (a)  $q_1() : -R(x, t_1, y), S(y, t_2, z), T(z, t_3, x), U(x, y, z)$
  - (b)  $q_2() : -R(x, y, z), S(y, z, w), T(x, y, w), U(x, z, w)$
  - (c)  $q_3() : -R(x, y, z), S(x, y), T(y, z), U(z, x)$
3. [10%] For the following CQs, construct a GHD with the smallest possible ghw.
  - (a) The  $k$ -th spider-web query  $W_k$ :
$$W_k(x) : -R(x, z_1), R(x, z_2), \dots, R(x, z_k), R(z_1, z_2), R(z_2, z_3), \dots, R(z_k, z_1)$$
  - (b)  $V_k() : -R(x_1, x_2), R(x_2, x_3), \dots, R(x_k, x_1), R(x_1, z_2), R(z_2, z_3), \dots, R(z_k, x_1)$ , for any  $k \geq 1$ . This query can be drawn as two cycles of length  $k$  connected at variable  $x_1$ .
4. [10%] Among all CQs with  $n$  atoms in the body, what is the maximum possible ghw as a function of  $n$ ? Describe a class of queries that achieves this ghw.

## DELIVERABLES

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Submit a single PDF document using Canvas (Homework 1). It is strongly suggested to use  $\text{\LaTeX}$  to write your solution.