## CS354: Machine Organization and Programming

## Lecture 6

Wednesday the September $16^{\text {th }} 2015$

## Section 2

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## Class Announcements

1. How many of you attended the WACM tutorial and found it useful?
2. Assignment 1 released - due before 9AM on Sep 30 .

- You can find partners using Piazza too.
- Start Early! Much much harder than Assign 0!

3. Make sure you don't change your files to add very small changes like formatting, comments etc. after deadline. You get points deducted even if it is a small change.

## Lecture Overview

1. Doubly Linked Lists
2. Data Representation (Unsigned, 2's complement)
3. Signed <-> Unsigned Conversions
4. Integer Arithmetic (Addition)

## Example C Program on

 Singly Linked List
## Deleting a node in a Singly Linked List without copying



## Doubly Linked List

1. In order to delete a node in a singly linked list without copying values, a pointer to the previous node is also needed.
2. Doubly linked lists allow inserts and deletes in constant number of operations with only the node's address.
3. Doubly linked lists are easier to manipulate they allow fast and easy sequential access to the list in both directions.
struct node \{
int theint;
struct node *next;
struct node *previous;
\};


DOUBLY LINKED

For convenience, name this user-defined type:
typedef struct node \{
int theint;
struct node *next;
\} Node;

Now, declarations have less (keyboard) typing:

Node one, two, three;
Node *head;

## Some code, to show pointers and such. . .

one.theint = 1;
one.next = \&two;
one.next->next $=$ \&three;
three.next = NULL;
head
head = \&one;


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one. next $=$ \&two;
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one.theint $=1$;
one.next = \&two;
one.next->next $=$ \&three;
three.next = NULL; head
head = \&one;

int value = 1;
Node *ptr;
ptr = head;
while (ptr != NULL) \{ ptr->theint = value * 11; value++; ptr = ptr->next;
\}

## Some code, to show pointers and such. . .

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one.theint $=1$;
one.next = \&two;
one.next->next $=$ \&three;
three.next = NULL; head
head = \&one;



Why is this now incorrect?

With the correct code, what happens when this code is executed?

$$
\begin{aligned}
& \text { ptr }=\text { three.next; } \\
& \text { ptr }=\text { ptr->next; }
\end{aligned}
$$

## Some code, to show pointers and such. . .

```
one.theint = 1;
one.next = &two;
one.next->next = &three;
three.next = NULL;
head
head = &one;
```



With the correct code, what happens when this code is executed?

## Runtime error:

 NULL pointer dereference
## In Linux: <br> Segmentation <br> fault <br> (core dumped)

## Example C Program on

 Doubly Linked List
## Detailed Example C Program on Singly Linked List

## Bits, Nibbles, Bytes, Words

1. Bits represented using "high \& low voltages", "magnetic domain oriented clockwise or anticlockwise" etc.
2. 4 Bits $==$ Nibble ; 8 Bits $==$ Byte ; 16/32/64 Bits $==$ Word (depending on architecture);
3. Group of bits collected together with some interpretation is more useful than individual bits.


## Word size

1. It is the nomimal size of integers and pointer data
2. Determines the maximum size of virtual address space
3. $w$ bit word can address a virtual memory of size $\left(2^{w}\right)$ ranging from 0 to $2^{w}-1$.
4. Modern computers have 64 bit words. (Theoretically: $2^{64}=16$ Exabytes.)

## Byte encodings

- Byte $=8$ bits
- Binary $00000000_{2}$ to $11111111_{2}$
- Decimal: $0_{10}$ to $255_{10}$
- Hexadecimal $00_{16}$ to $\mathrm{FF}_{16}$
- Base 16 number representation
- Use characters ' 0 ' to ' 9 ' and ' A ' to ' F '
- Write FA1D37B16 in C as
- 0xFA1D37B
- 0xfa1d37b

| $\lambda^{e^{t}}$ |  |  |
| :--- | :---: | :---: |
| $e^{i n}$   <br> 0 0 0000 <br> 1 1 0001 <br> 2 2 0010 <br> 3 3 0011 <br> 4 4 0100 <br> 5 5 0101 <br> 6 6 0110 <br> 7 7 0111 <br> 8 8 1000 <br> 9 9 1001 <br> A 10 1010 <br> B 11 1011 <br> C 12 1100 <br> D 13 1101 <br> E 14 1110 <br> F 15 1111 |  |  |

## Byte Ordering/Endianness

1. Ordering of bytes within a word
2. Little endian - least significant byte comes first
3. Big endian - most significant byte comes first


Little-endian


## Representations

1. Unsigned encodings - positive integers
2. Two's complement - signed integers
3. Floating point - real numbers
4. Because of limited number of bits to encode a number, some operations can "overflow" when results are too large.

## Arithmetic Operations

> Arithmetic Operations
> addition
> subtraction
> multiplication
> division
> Each of these operations on the integer representations:
> unsigned
> two's complement

## Addition

## One bit of binary

 addition
a


Addition Truth Table

| Carry In | a | b | Carry <br> Out | Sum Bit |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 |

## Unsigned Representation

$B 2 \mathrm{U}_{\mathrm{w}}\left(x_{\text {vec }}\right)=\operatorname{Sum}_{\mathrm{i}=0->\mathrm{w}-1} \mathrm{x}_{\mathrm{i}} \cdot 2^{\mathrm{i}}$

B2 $_{4}([0101])=0.2^{3}+1.2^{2}+0.2^{1}+1.2^{0}=5$
$B 2 \mathrm{U}_{\mathrm{w}}$ is a bijection:

- associates a unique value to each bit vector of length w
- each integer between 0 and $2^{\mathrm{w}}-1$ has a unique binary representation as a bit vector of length w


## Unsigned Addition

Of two unsigned w bit values X \& Y $X+Y$ equals:

- $\mathrm{X}+\mathrm{Y}$, if $(\mathrm{X}+\mathrm{Y})<2^{\mathrm{w}}$
$-\mathrm{X}+\mathrm{Y}-2^{\mathrm{w}}$, if $2^{\mathrm{w}}<=(\mathrm{X}+\mathrm{Y})<2^{\mathrm{w}+1}$


## Addition

> Unsigned and 2's complement use the same addition algorithm
> Due to the fixed precision, throw away the carry out from the msb

$$
\begin{array}{r}
00010111 \\
+\quad 10010010
\end{array}
$$

## Addition

> Unsigned and 2's complement use the same addition algorithm
> Due to the fixed precision, throw away the carry out from the msb

$$
\begin{array}{r}
00010111 \\
+\quad 10010010 \\
\hline 10101001
\end{array}
$$

## Two's complement Representation

B2T $\mathrm{w}_{\mathrm{w}}\left(x_{\text {vec }}\right)=-\mathrm{x}_{\mathrm{w}-1} 2^{\mathrm{w}-1}+\operatorname{Sum}_{\mathrm{i}=0->\mathrm{w}-2} \mathrm{X}_{\mathrm{i}} \mathrm{2}^{\mathrm{i}}$

B2 $\mathrm{T}_{4}([1011])=-1.2^{3}+0.2^{2}+1.2^{1}+1.2^{0}=-5$
$B 2 \mathrm{~T}_{\mathrm{w}}$ is a bijection:

- associates a unique value to each bit vector of length w
- each integer between $-2^{\mathrm{w}-1}$ and $2^{\mathrm{w}-1}-1$ has a unique binary representation as a bit vector of length w


## Range of Values for Unsigned and 2's Complement (16 bits)

|  | Decimal | Hex | Binary |  |
| :--- | ---: | ---: | ---: | ---: |
| UMax | 65535 | FF FF | 11111111 | 11111111 |
| TMax | 32767 | $7 F ~ F F$ | 01111111 | 11111111 |
| TMin | -32768 | $80 \quad 00$ | 10000000 | 00000000 |
| -1 | -1 | FF FF | 11111111 | 11111111 |
| 0 | 0 | $00 \quad 00$ | 00000000 | 00000000 |

\#include <limits.h> declares constants, e.g., ULONG_MAX, LONG_MAX, LONG_MIN (Values platform specific)

4-bit Unsigned and 2's complement Integers

| $X$ | $\mathrm{~B} 2 \mathrm{U}(X)$ | $\mathrm{B} 2 \mathrm{~T}(X)$ |
| :---: | :---: | :---: |
| 0000 | 0 | 0 |
| 0001 | 1 | 1 |
| 0010 | 2 | 2 |
| 0011 | 3 | 3 |
| 0100 | 4 | 4 |
| 0101 | 5 | 5 |
| 0110 | 6 | 6 |
| 0111 | 7 | 7 |
| 1000 | 8 | -8 |
| 1001 | 9 | -7 |
| 1010 | 10 | -6 |
| 1011 | 11 | -5 |
| 1100 | 12 | -4 |
| 1101 | 13 | -3 |
| 1110 | 14 | -2 |
| 1111 | 15 | -1 |

