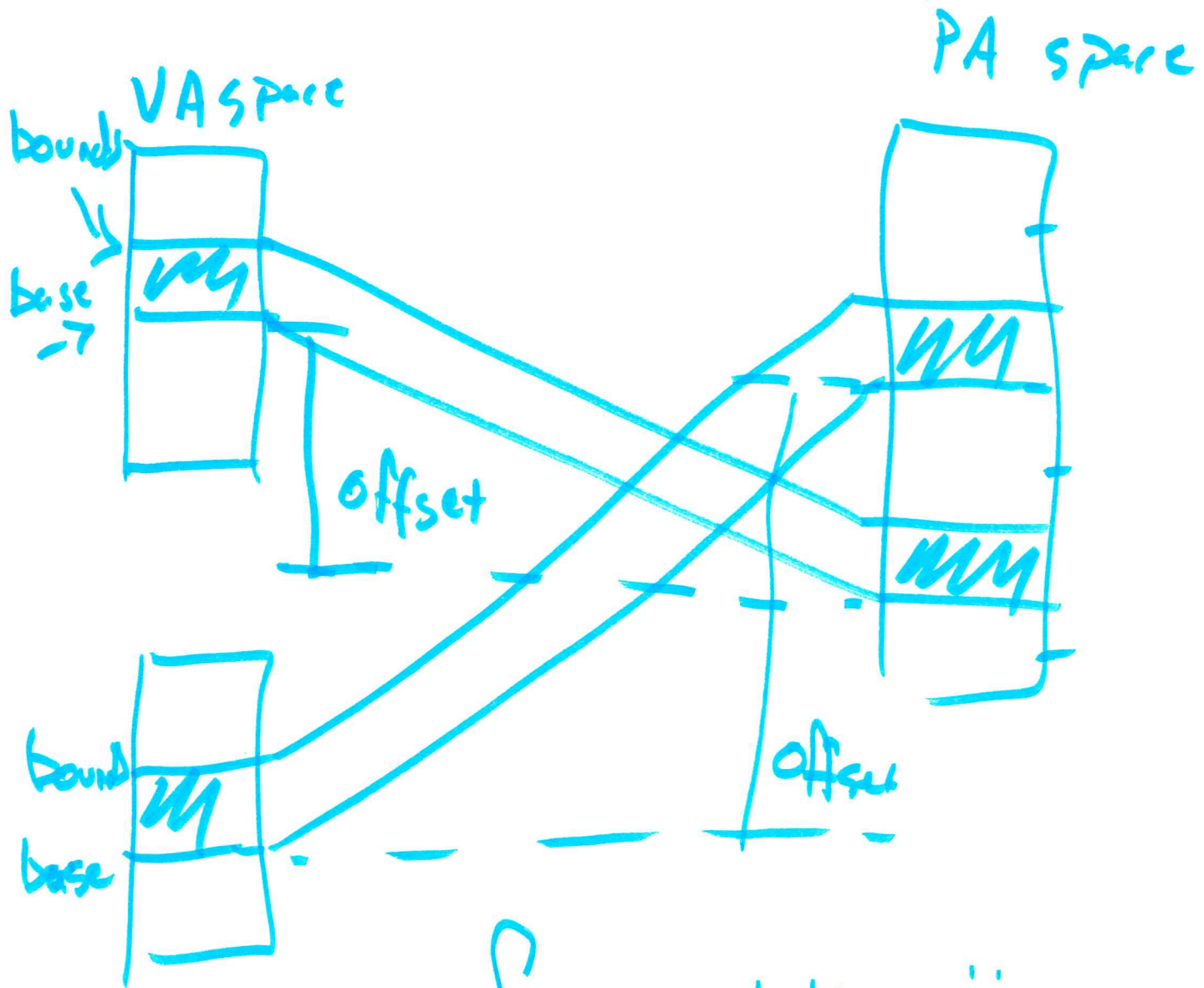


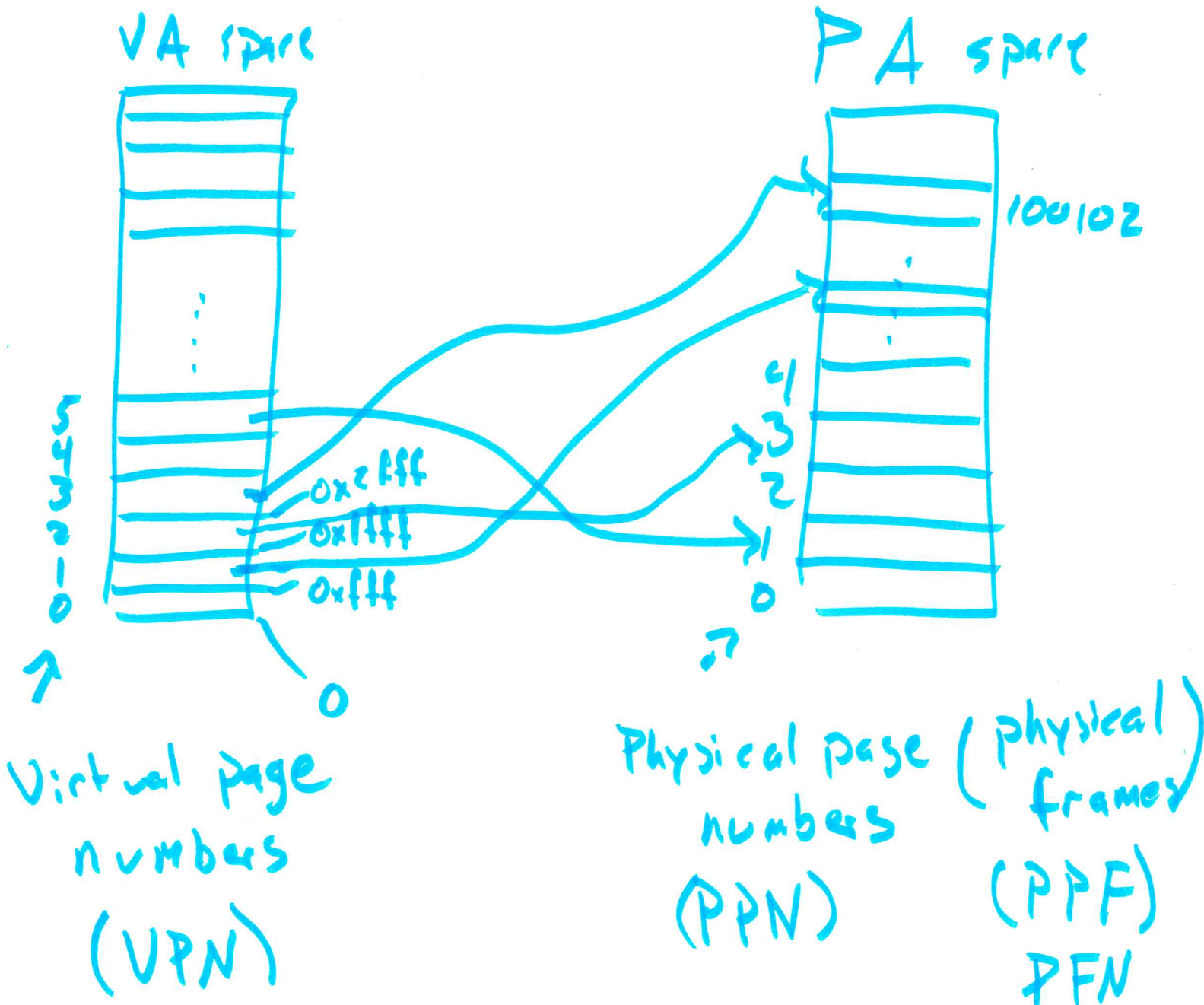
# Page-based Virtual Memory



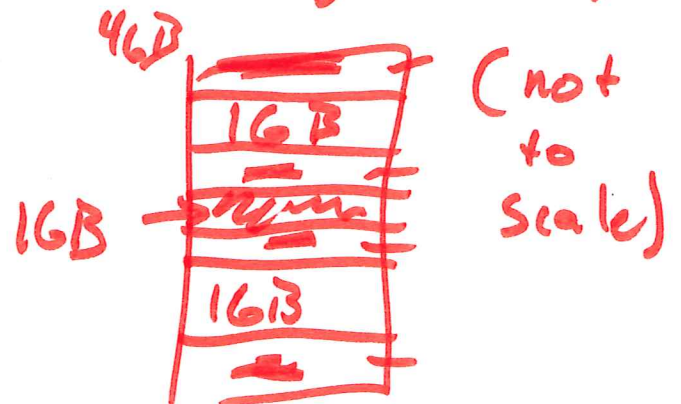
Fragmentation is

## Paging

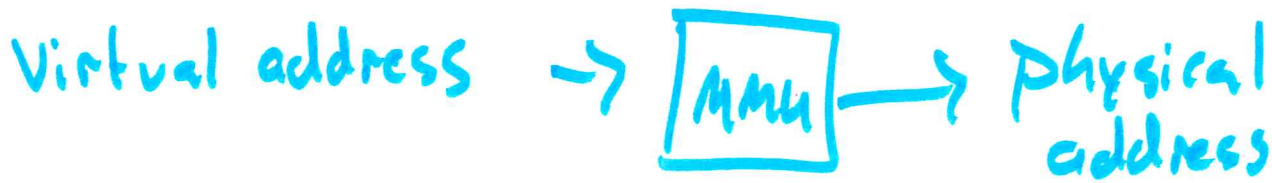
- break address up into same-sized chunks (4KB)



- Fragmentation problem (w/ segmentation)



# How to translate from VA to PA



## Page table

- MMU looks up translation in page table
- one page table per process

VPN → valid  
 VP → Present bit  
 PPN → indexed by the

VPN	VP	PPN
4	1	17
3	1	0
0	1	12

VPN → entries: Page table entry (PTE)

- includes:

- PPN mapping
- Valid entry?
- Permissions: read, write, execute

where on disk

Present bit → is this on disk or in memory?  
 Physical

Memory as large cache for disk

Virtual memory makes it possible

page faults → OS look up if page is valid  
(possibly on disk)

→ OS will allocate a ~~physical~~  
Physical page → possibly  
moving a page to disk

→ OS will copy requested page  
into memory

→ OS updates the page table(s)

## Swapping

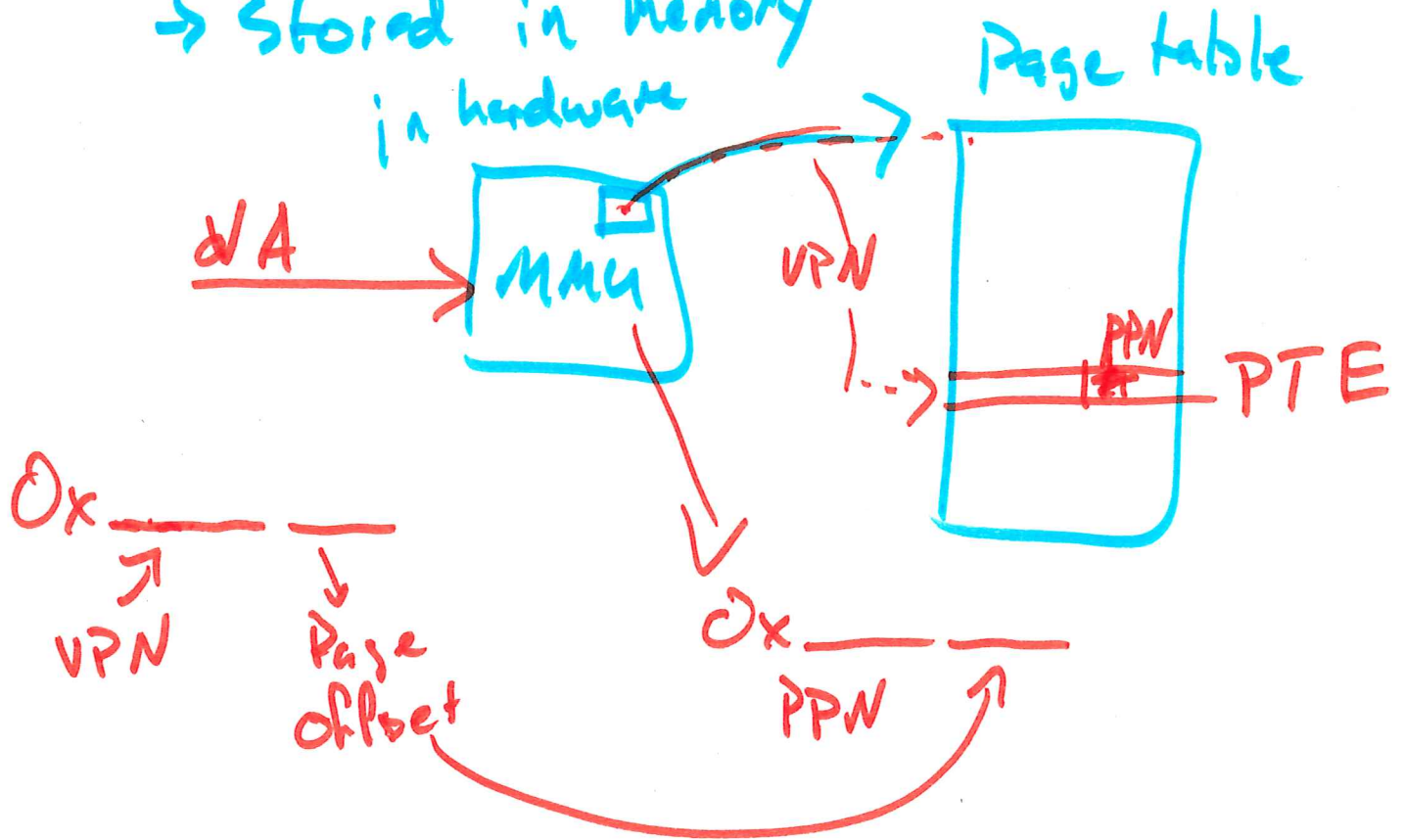
Page faults → caused by "page misses"

~~the~~ MMU looks up VPN in the  
page table → in memory (valid &  
present)  
↳ "page hit"

→ not in memory page miss  
page fault

Where is the page table?

→ Stored in memory  
in hardware



MMU has a pointer to the page table → physical address (CR3)

Each process has unique page table and unique PT pointer

Multi-level page table

Page table

4KB pages  
2<sup>12</sup>

VPN

	V	P	PPN	
F	1	1	5	} stack
E	1	1	7	
D	1	1	4	
C	1	1	8	
B				
A	1	1	A	
9	1	1	3	→ heap
8	0	0	0	←
7				
6	1	1	1	} global data
5	1	1	2	
4	1	1	6	
3				
2				
1	1	1	0	
0				→ code

0x8A14 → 0x6A14  
 ↳ offset  
 ↳ VPN

~~Invalid~~  
 If the MMU finds an invalid entry → Page fault  
 or not present

