

Mark Moir

Principal Investigator

Scalable Synchronization Research Group

Joint work with: Peter Damron, Yossi Lev, Victor Luchangco, Dan Nussbaum, Nir Shavit

Transactional Memory



© Sun Microsystems 2005



Outline

- Progress in software TM
- Developments in hardware TM
- Hybrid Hardware/Software TM
 - Key idea in abstract
 - Applied to DSTM
 - Applied to word-based STM
- Concluding remarks



Software Transactional Memory

- Lots of progress in:
 - Space overhead
 - Data transparency
 - Parallelism
 - Performance
- Work beginning on programming language/model/interface
- Still requires multiple expensive atomic instructions per transaction



Hardware Transactional Memory...

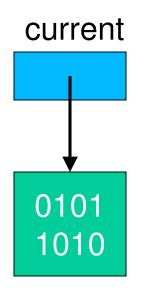
- (...including related things like SLE).
- Renewed interest.
- Two flavours of design:
 - Best effort: simple, but no guarantees
 - Unbounded: can commit any transaction, but more complicated



Hybrid Transactional Memory

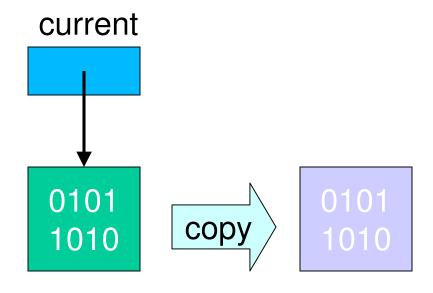
- Make HTM and STM txns "play together nicely"; try hardware first, resort to software when it fails
- Programmer writes txn code once; hybrid TM does the rest
- Hardware can be "best effort", without impacting programmers
- Transactional code can be developed and tested without hardware support





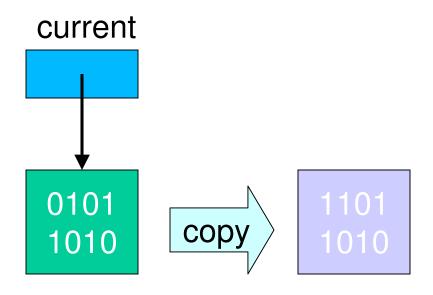
- Well-known software synchronization technique
- Copies of data and indication of which is "current".





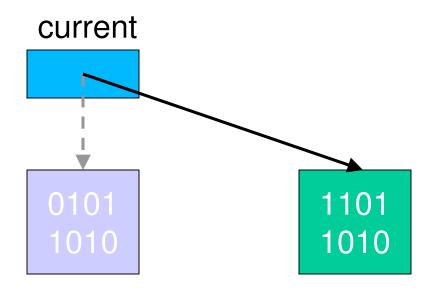
- Applying an operation:
 - make private copy





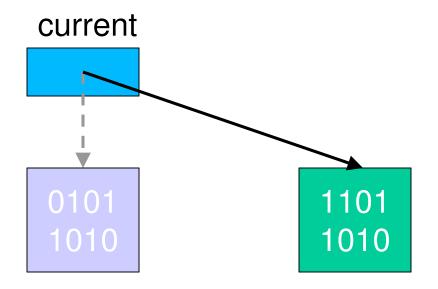
- Applying an operation:
 - make private copy
 - modify private copy (sequential code)





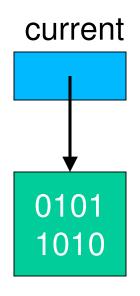
- Applying an operation:
 - make private copy
 - modify private copy (sequential code)
 - atomically make copy become current





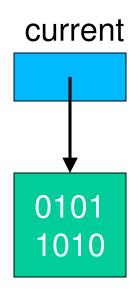
- Disadvantages:
 - no parallelism
 - copying overhead





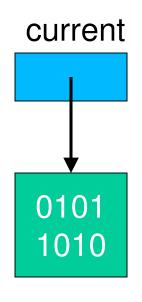
- Simplified (and wrong!)
 hybrid optimization
 approach:
 - use hardware TM to modify copy in-place, and
 - check current pointer doesn't change





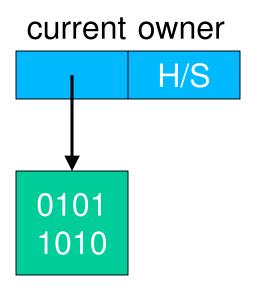
- Advantages:
 - Nonconflicting operations can succeed in parallel
 - no copying





- Advantages:
 - Nonconflicting operations can succeed in parallel
 - no copying
- Disadvantage:
 - → not correct ②
 - software operations might copy inconsistent data



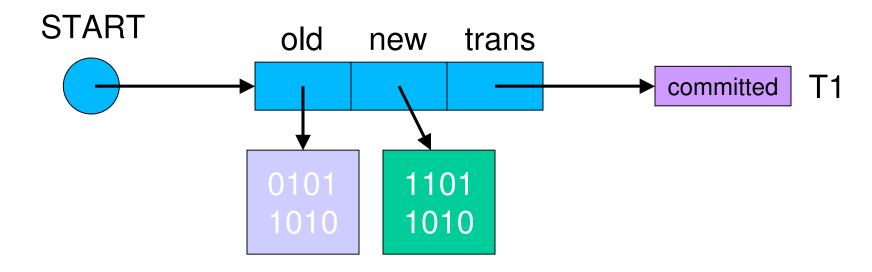


Solution:

- indicate "ownership" by software operation(s)
- software ops acquire ownership before copying
- hardware ops abort if object owned by software

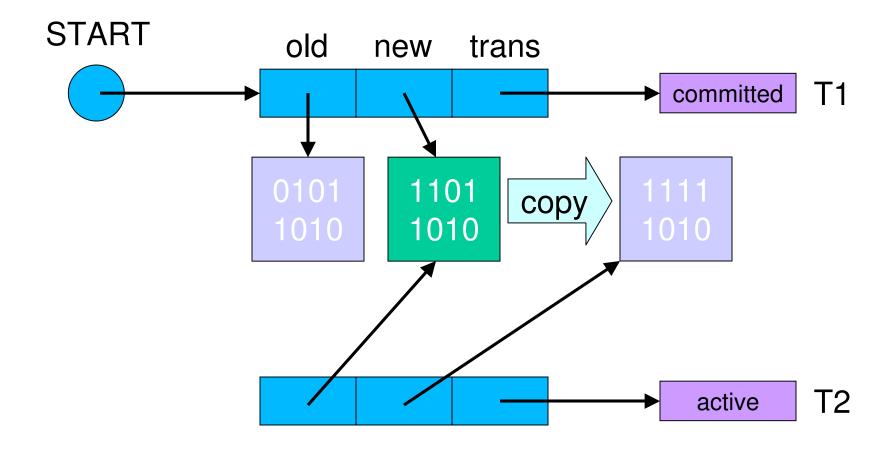


Example: DSTM



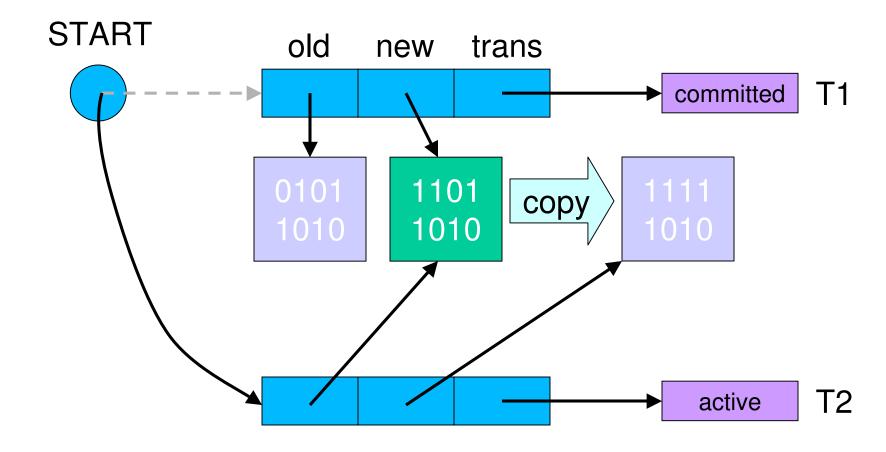


Example: DSTM



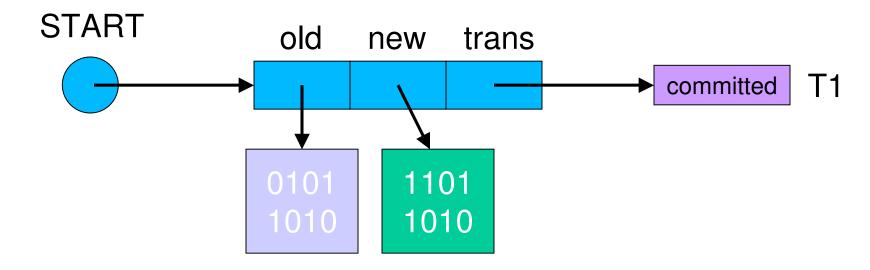


Example: DSTM





Hybrid DSTM



- Use hardware TM to check object not owned by active software txn, and if not modify current object in-place
- Avoids copying and expensive synchronization, allows intra-object parallelism



An Aside

- Level of indirection because of singleword CAS, multi-word locator
- H/W TM that guarantees (eventual) success of single-cache-line txn would remove indirection simply.
- Other simple guarantees useful too: e.g., txns that read one cache line and write one cache line



Word-based Hybrid TM

- TM useful at lower levels too, e.g. JVM, GC, etc.
- No object infrastructure, want to modify data in-place



Word-based Hybrid TM Prototype

- Similar in structure to Harris & Fraser STM (OOPSLA 2003)
- Ownership table; data transparency
- Compiler emits code for hardware txn and software txn (library calls), and retry/contention mgmt code
- Hardware transactions augmented to check ownership table for conflicting software txns (multiple levels of granularity possible)



Word-based Hybrid TM Prototype

- Shoe-horned into existing languages to allow experimentation with existing code, e.g. GC, JVM, etc.
- Programming model primitive, but workable
- Remaining challenges include:
 - better language integration
 - validation of hybrid approach



Concluding Remarks

- Hybrid Transactional Memory
 - Use best-effort hardware TM to boost performance of self-contained STM
 - Eases constraints on hardware designers
 - Best effort ok, but simple guarantees desirable
 - Eases path to adoption of transactional programming