

The FBI created a fake malware-spreading newspaper article to trace a bomb threat

COMMENTS

By [Russell Brandom](#) on October 28, 2014 11:22 am [Email](#) [@russellbrandom](#)

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In 2007, the FBI was tra

Documents [uncovered by the Electronic Frontier Foundation](#) show that the FBI created a fake web page designed to look like a Seattle Times article, and used the page to spread tracking malware onto the suspect's computer. Creating dummy pages is a common way to spread malware — typically known as spoofing — but it's more common among criminals than law enforcement, and many are already interpreting the fake page as an attack on the press. "We are outraged that the FBI, with the apparent assistance of the U.S. Attorney's Office, misappropriated the name of *The Seattle Times*," a *Times* editor [told the paper](#). "Not only does that cross a line, it erases it." The Associated Press echoed the concern, saying, "this ploy violated AP's name and undermined AP's credibility."

Crypto: Passwords and RNGs

CS 642

Guest Lecturer: Adam Everspaugh

<http://pages.cs.wisc.edu/~ace>



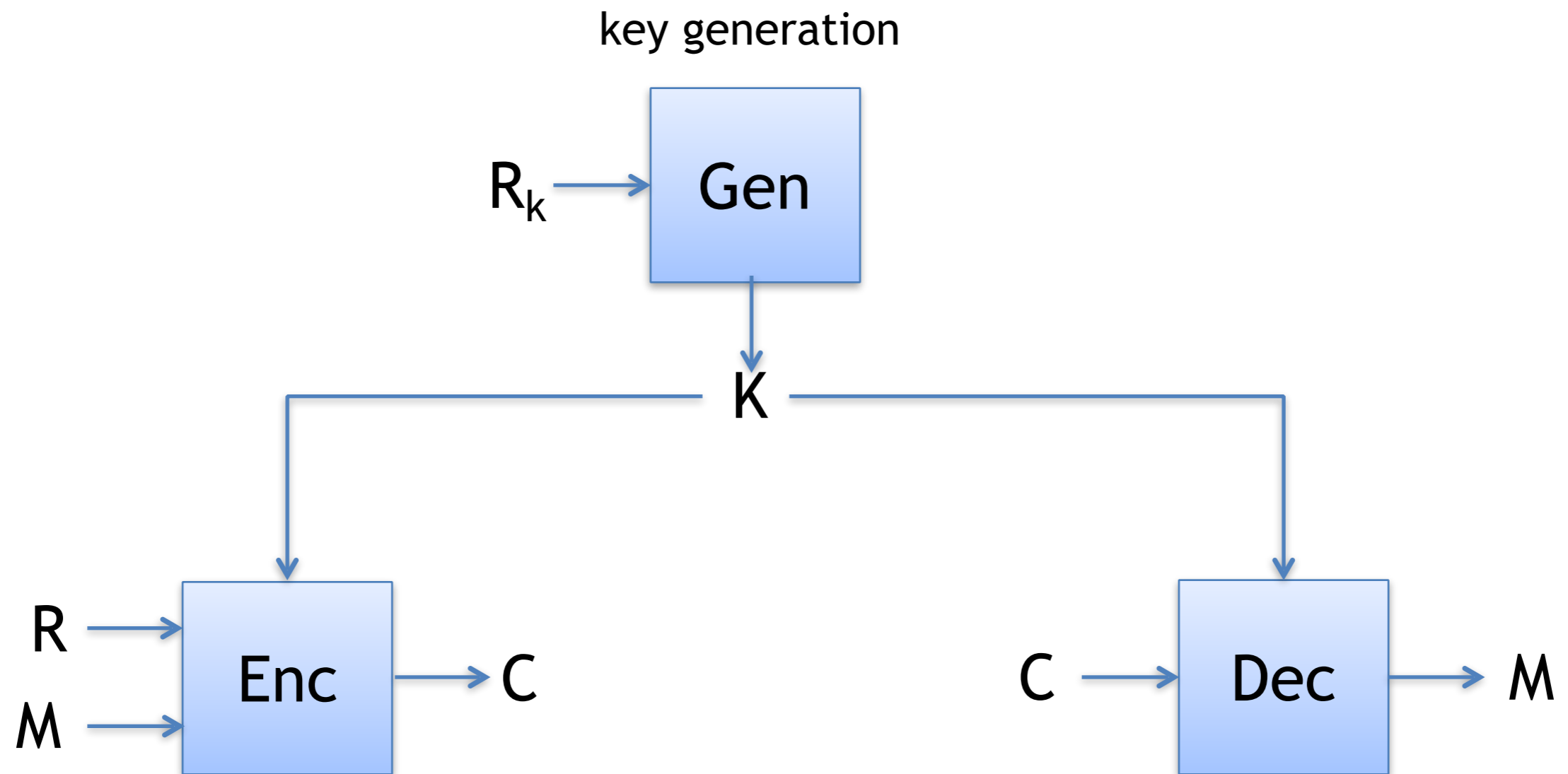
Topics

Password-based Crypto

Random Number Generators

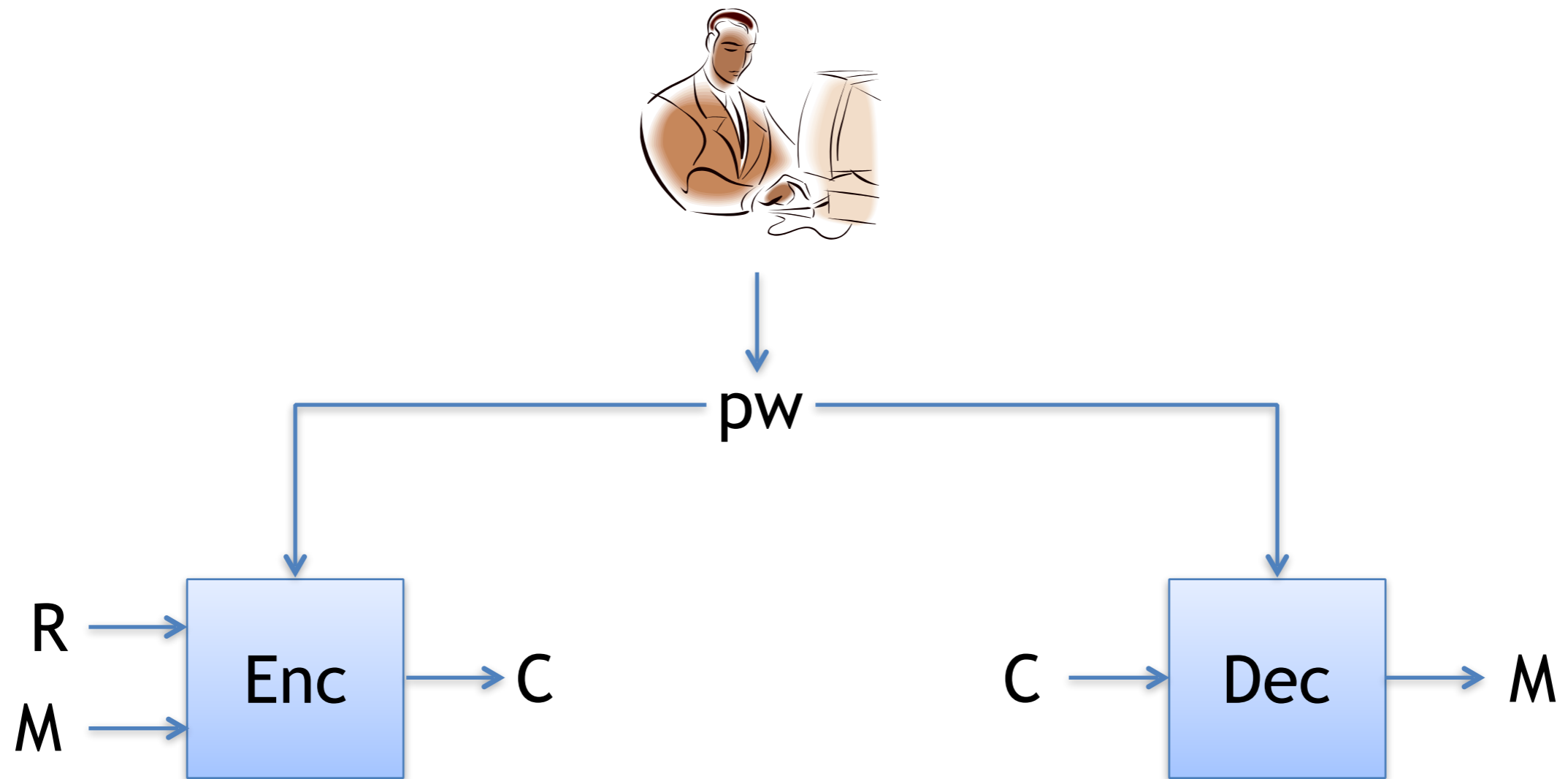


Symmetric Key Encryption



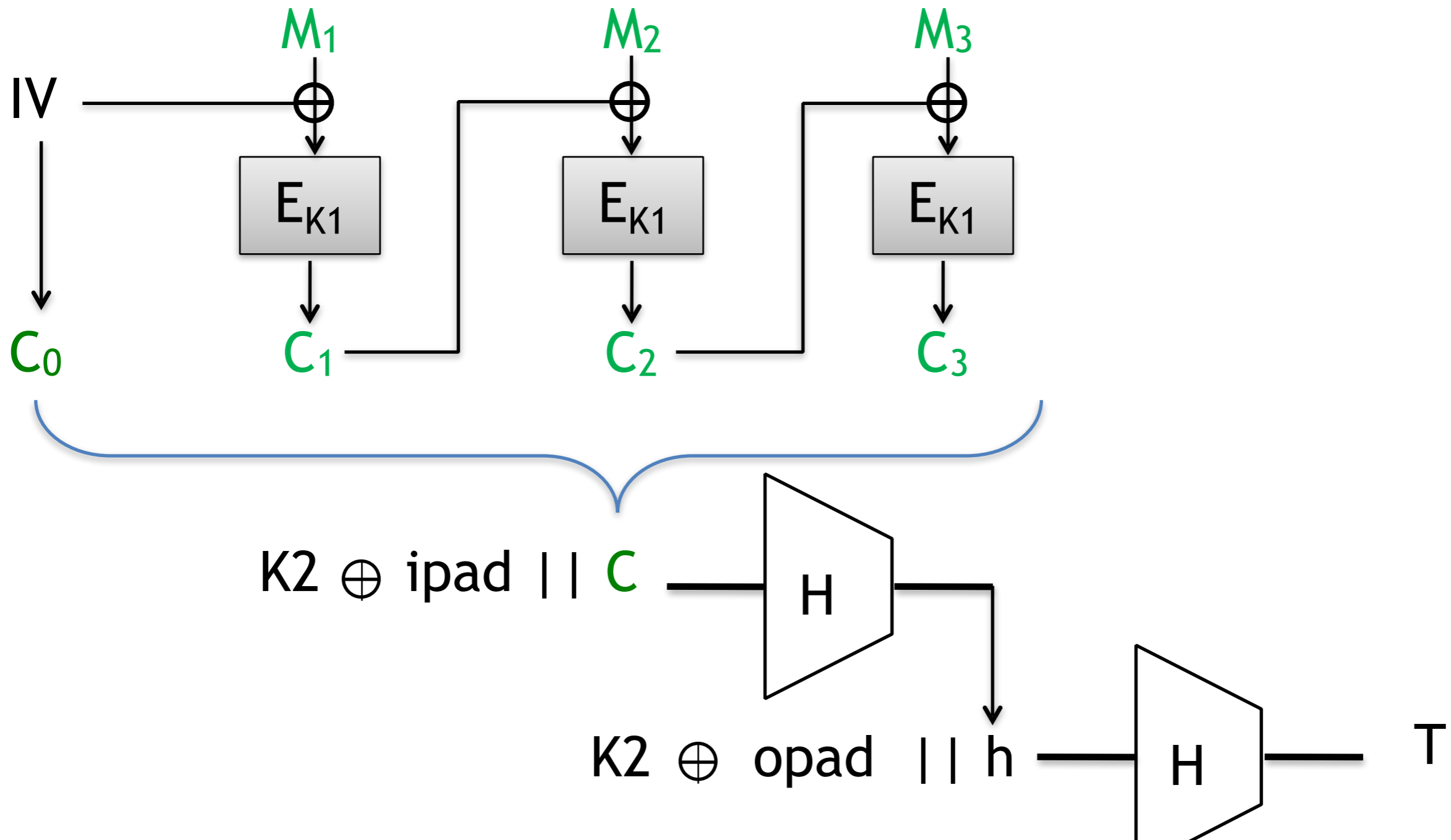
Correctness: $D_k(E_k(M,R)) = M$

Password-based Symmetric Encryption



Correctness: $D(\text{pw}, E(\text{pw}, M, R)) = M$

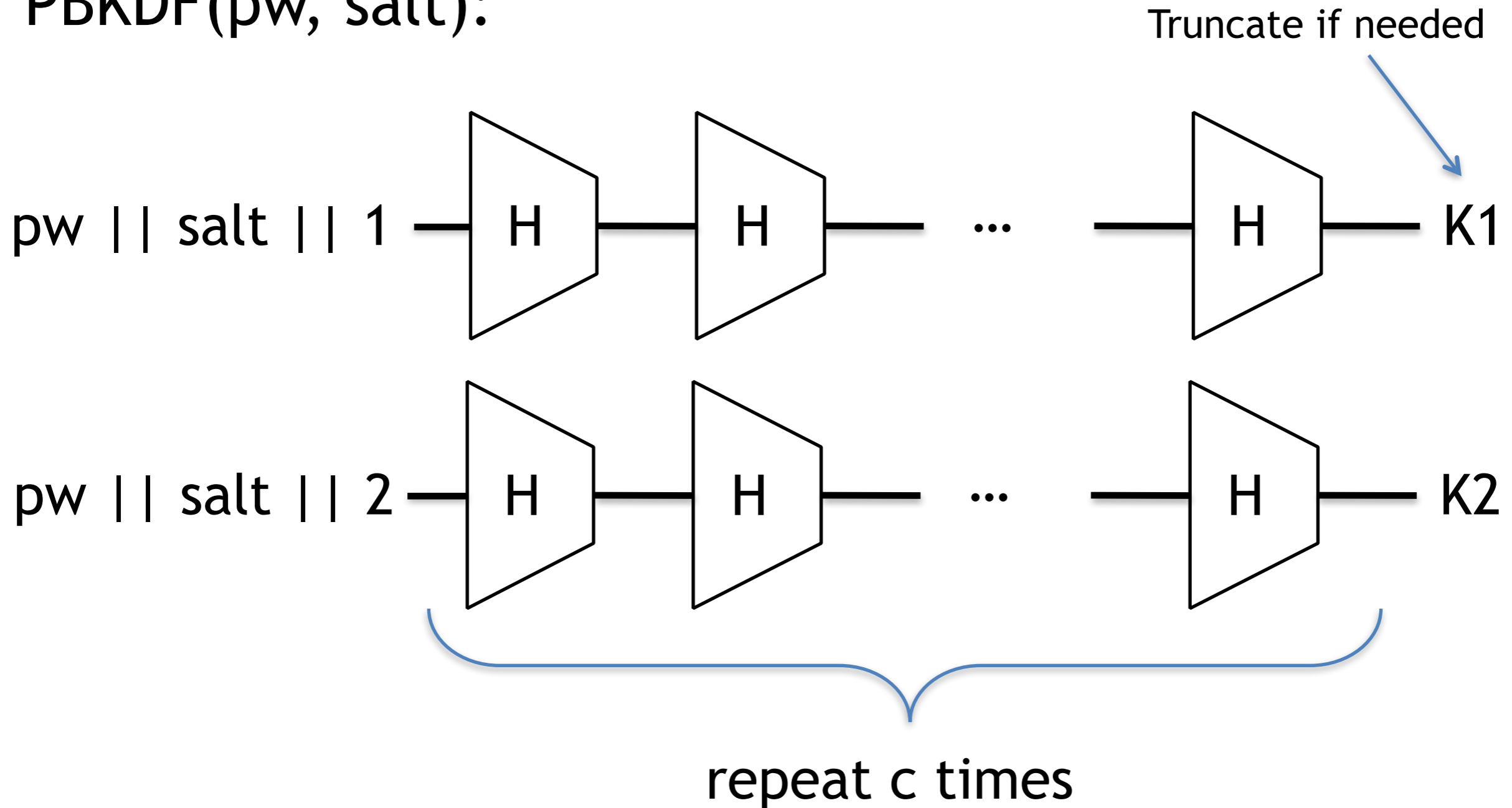
Encrypt-then-MAC with CBC and HMAC



How do we use this with a password?

Password-based Key Derivation (PBKDF)

PBKDF(pw, salt):



PBKDF + Symmetric Encryption yields PW-Based Encryption

Enc(pw,M,R):

salt || R' = R

K = PBKDF(pw,salt)

C = Enc'(K,M,R')

Return (salt,C)

Here Enc'/Dec' is a typical symmetric encryption scheme (CBC+HMAC)

Dec(pw,C):

salt || C' = C

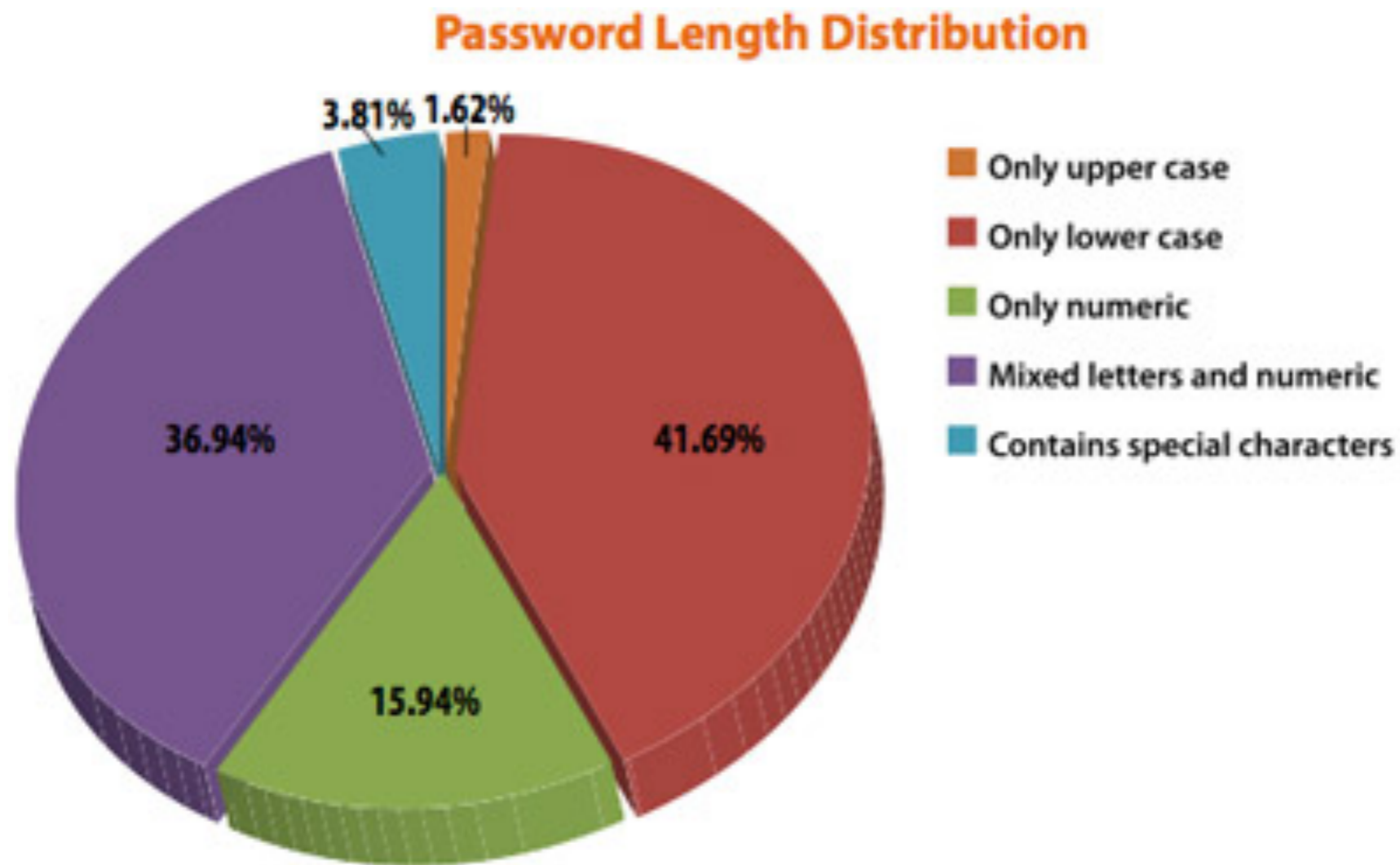
K = PBKDF(pw,salt)

M = Dec'(K,C')

Return M

Attacks?

Password Distribution



From an Imperva study of released RockMe.com password database (2010)

Dictionary Attack

- Given a (message, ciphertext) pair:
- Enumerate a dictionary D of possible passwords, in order of likelihood
- Test each candidate password

DictionaryAttack(D,M,C):

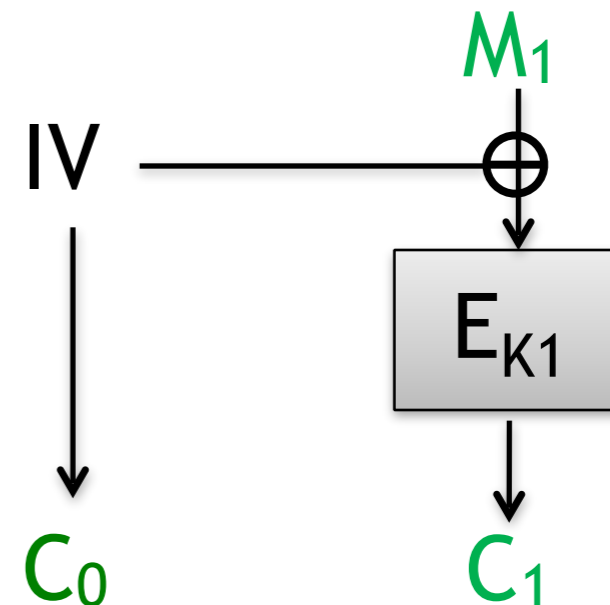
$R \parallel C' = C$

for pw^* in D :

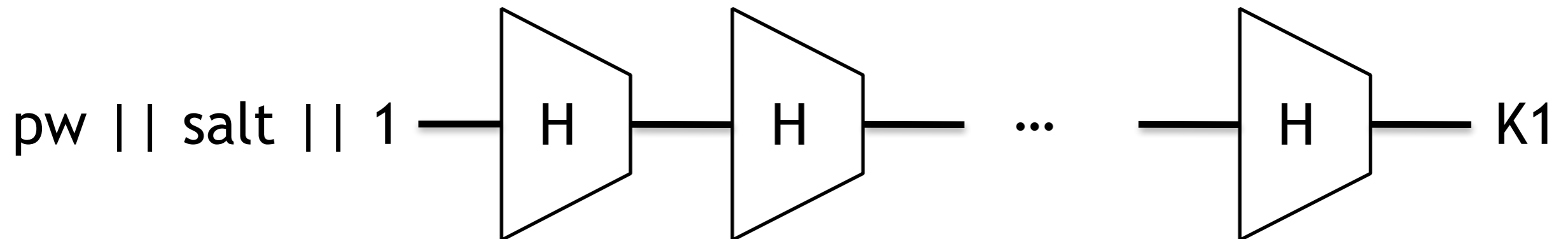
$C^* = \text{Enc}(pw^*, M, R)$

if $C^* == C'$:

return pw^*



PBKDF Slows Down Dictionary Attacks



Iterating c times should slow down attacks by factor of c

Salts:

- Different derived keys, even if same password

- Slows down attacks against multiple users

- Prevents precomputation attacks, if salts chosen randomly

How Fast Are Dictionary Attacks?

- openssl speed sha1
- Assume: 4 cores @ 2.2M hashes per second

	Size of Dictionary	Computation time c=1	Computation time c=4096
6 digit PIN	10	0.11 seconds	7.8 minutes
6 alphanumerics (lowercase)	36	4.1 minutes	11.7 days
8 alphanumerics (mixed case)	62	287 days	3,222 years

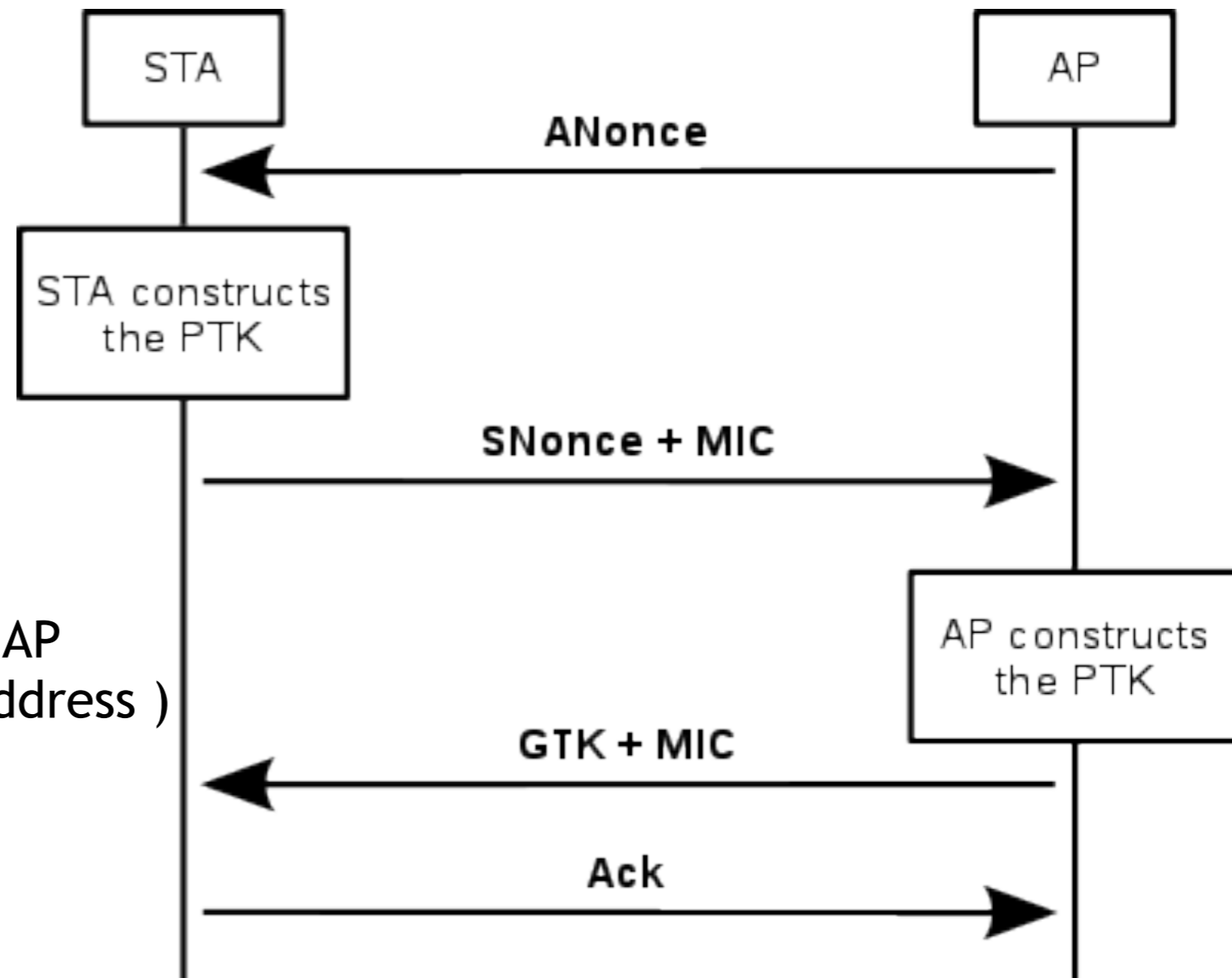
802.11 WPA Authentication



$PMK = PBKDF(pw, ssid || ssidlength)$
with $c = 4096$

$PTK = H(PMK || ANonce || SNonce || AP$
MAC address || STA MAC address)

$MIC = HMAC-MD5(PTK, M2)$



Observe just one handshake by another party, and attacker can mount *offline* dictionary attack against the password

Attacking WPA Passwords

Wifi AP



$PMK = PBKDF(pw, ssid || ssidlength)$
with $c = 4096$

$PTK = H(PMK || ANonce || SNonce || AP$
MAC address || STA MAC address)

$MIC = HMAC-MD5(PTK, M2)$

DictionaryAttack(D, MIC, ANonce, SNonce, SSID, M2):

for pw^* in D:

$PMK^* = PBKDF(pw^*, ssid || ssidlength)$

$PTK^* = H(PMK^* || ANonce || \dots)$

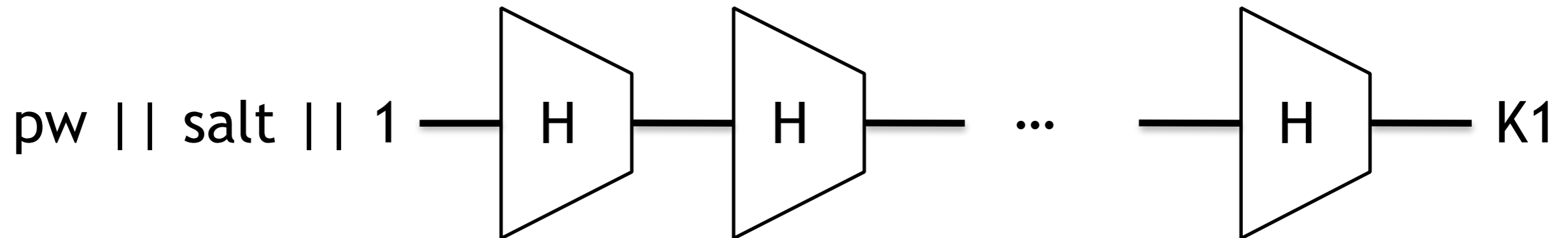
$MIC^* = HMAC-MD5(PTK^*, M2)$

If $MIC^* == MIC$:

 return pw^*

return None

Recap: Password-based Crypto



- Allows use of passwords in existing crypto schemes
- **Gain:**
 - Increases attackers computations
 - Prevents precomputation
- **Cost:**
 - Increased computation
- **Limitation:**
 - Strength of key still limited to strength of password
 - Don't make it easy for attacker to mount offline dictionary attacks



INTERMISSION

Uses for Secure Random Numbers

Cryptography

- Keys
- Nonces, initial values (IVs), salts

System Security

- TCP Initial Sequence Numbers (ISNs)
- ASLR
- Stack Canaries



Where can we get secure random numbers?



OSX/Linux

- `cat /dev/urandom`
- `xxd -l 1024 -p /dev/urandom`
- `openssl rand 256 -hex`

Intel HW RNG

- **OSX:** `sysctl -a | grep RDRAND`
- **Linux:** `cat /proc/cpuinfo | grep rdrand`



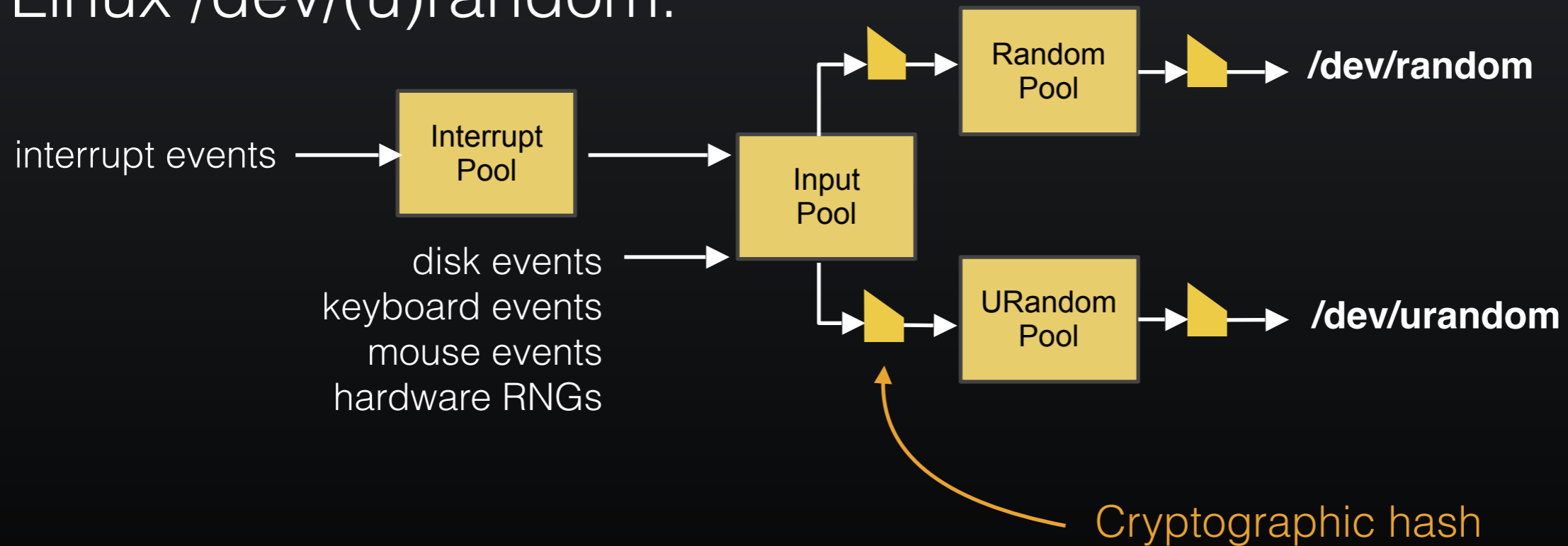
Operating System Random Number Generators



Linux RNG



Linux /dev/(u)random:



RNG Failures



RNG Failures

Predictable Output

Repeated Output

Outputs from a small range (not-statistically uniform)

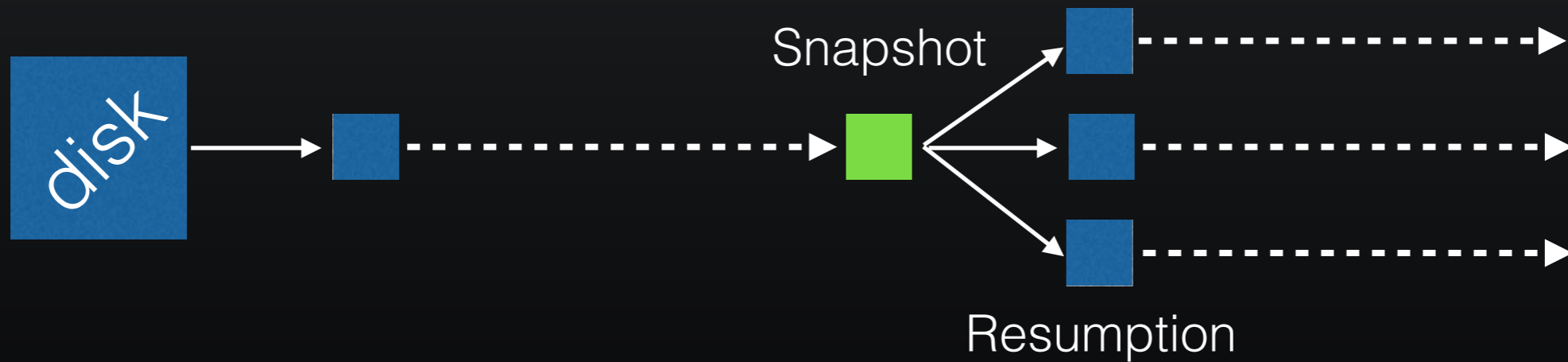
Broken Windows RNG: [DGP 2007]

Broken Linux RNG: [GPR 2008], [LRSV 2012], [DPRVW 2013], [EZJSR 2014]

Factorable RSA Keys: [HDWH 2012]

Taiwan National IDs: [BCCHLS 2013]

Virtual Machine Snapshots



Security Problems with VM Resets

VM Reset Vulnerabilities [Ristenpart, Yilek 2010]

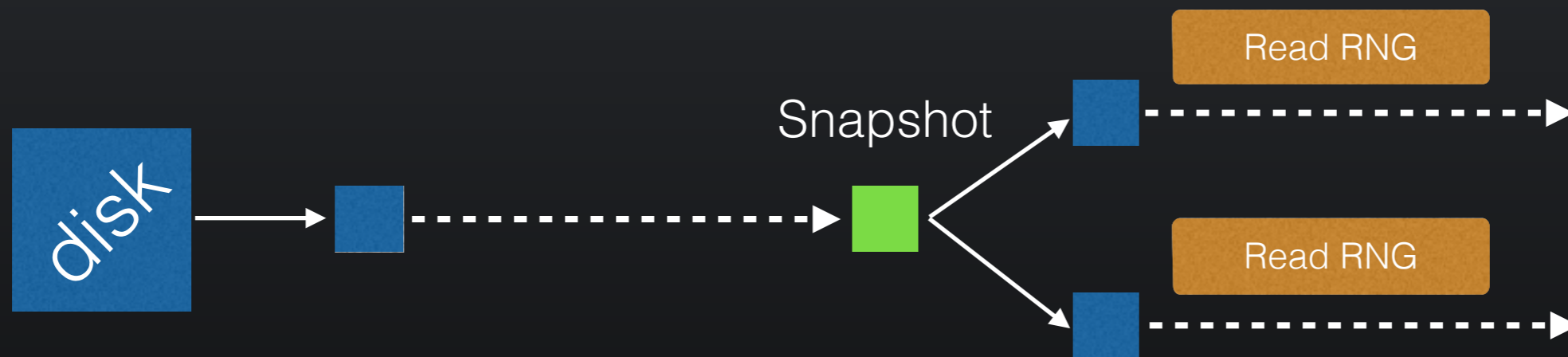


**Firefox and Apache reused random values for TLS
Attacker can read previous TLS sessions, recover private
keys from Apache**

Linux RNG after VM Reset



Not-So-Random Numbers in Virtualized Linux
[Everspaugh, et al, 2014]



Experiment:

- Boot VM in Xen or VMware
- Capture snapshot
- Resume from snapshot, read from `/dev/urandom`

Repeat: 8 distinct snapshots
20 resumptions/snapshot

/dev/urandom outputs after resumption

Linux RNG is **not** reset secure:
7/8 snapshots produce mostly identical outputs

1E6DD331

8CC97112

2A2FA7DB

DBBF058C

26C334E7

F17D2D20

CC10232E

...

Reset 1

1E6DD331

8CC97112

2A2FA7DB

DBBF058C

26C334E7

F17D2D20

CC10232E

...

Reset 2

1E6DD331

8CC97112

2A2FA7DB

DBBF058C

26C334E7

45C78AE0

E678DBB2

...

Reset 3

Reset insecurity and applications

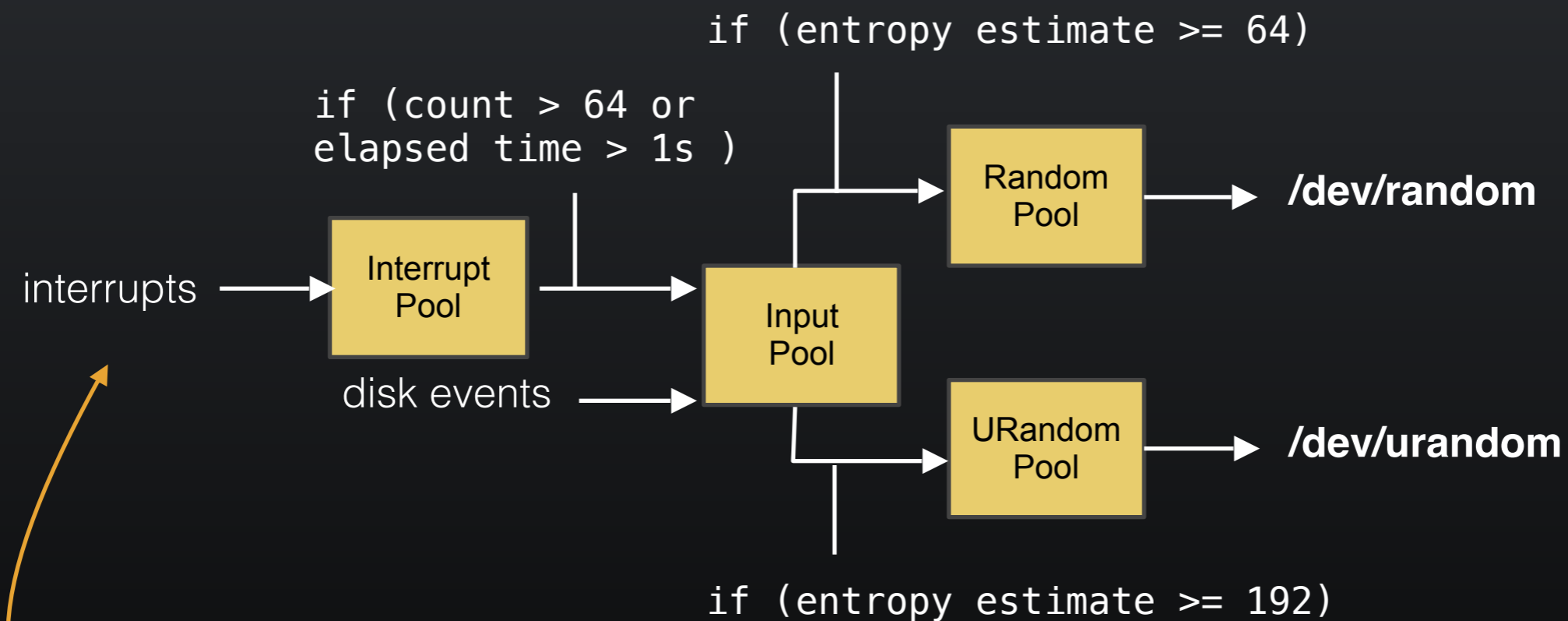
Generate RSA key on resumption:

```
openssl genrsa
```

30 snapshots; 2 resets/snapshot (ASLR Off)

- 27 trials produced **identical** private keys
- 3 trials produced unique private keys

Why does this happen?



Buffering and thresholds prevent new inputs from impacting outputs

Linux /dev/(u)random

What about other platforms?

FreeBSD

/dev/random produces **identical** output stream
Up to 100 seconds after resumption



Microsoft Windows 7

Produces **repeated** outputs indefinitely

rand_s (stdlib)

CryptGenRandom (Win32)

RngCryptoServices (.NET)

RNG Recap

- **RNGs are critical for security**
 - Keys, nonces, etc
- **Building good RNGs is hard**
- **OS provides a strong RNG**
 - e.g.: /dev/urandom
- **Intel CPUs provide an RNG**
 - RDRAND instructions



RNG

/dev/urandom

