

CS/ECE 252: INTRODUCTION TO COMPUTER ENGINEERING

UNIVERSITY OF WISCONSIN—MADISON

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Examination 4
In Class (50 minutes)
Wednesday, May 3rd, 2017
Weight: 17.5%

NO BOOK(S), NOTE(S), CALCULATORS OR ELECTRONIC DEVICES OF ANY SORT.
The exam has eleven pages. You must turn in the pages 1-9. Circle your final answers. Plan your time carefully since some problems are longer than others. Use the blank sides of the exam for scratch work.

LAST NAME: _____
FIRST NAME: _____
ID#: _____

Problem	Maximum Points	Points Earned
1	4	
2	5	
3	5	
4	5	
5	5	
6	4	
7	4	
Total	32	

1. The following LC-3 assembly code contains assembly syntax errors. Identify and fix at least 4 such errors. (4 points)

```
.ORIG x3000
    AND R5, R5, ZERO
    LD R5, STRING
NEXT  ADD R5, R5, #32
      BRz NEXT
MAIN  LD R4, MAIN
      SUB R4, R4, #1
      ST R4, STRING
NEXT  HALT
ZERO  .FILL #0
      .BLKW 4
      .BLKW 3
STRING .STRINGZ "ABC"
      .END
```

```
.ORIG x3000
    AND R5, R5, ZERO
    LD R5, STRING
NEXT  ADD R5, R5, #32
      BRz NEXT

MAIN  LD R4, MAIN
      SUB R4, R4, #1
      ST R4, STRING
NEXT  HALT
ZERO  .FILL #0
      .BLKW 4

.BLKW 3
STRING .STRINGZ "ABC"
      .END
```

2. a) Fill in the symbol table for following LC-3 assembly code. You may not need to fill all rows. (5 points)

```

        .ORIG x3000
        AND R3, R3, #0
        LD R4, VAL1
LOOP    BRZ DONE
        JSR INC1
        JSR INC2
DONE    ST R3, ANS
        OUT
        HALT
INC1    ADD R3, R3, #1
        RET
INC2    ADD R4, R4, #-1
        RET ; Storage area for variables below:
ANS     .BLKW #4
VAL0    .STRINGZ "CS"
VAL1    .STRINGZ "252"
        .END

```

SYMBOL	Value (in hex)
LOOP	0x3002
DONE	0x3005
INC1	0x3008
INC2	0x300A
ANS	0x300C
VAL0	0x3010
VAL1	0x3013

b) Convert the instruction stored at memory location 0x3006 into binary.

0xF021 (TRAP x21)

3. An LC-3 assembly program is given below: (5 points)

```
.ORIG x3000
LD R0, DATA
ADD R0, R0, #10
PRINT1 OUT
    AND R0, R0, #0
    ADD R0, R7, R0
PRINT2 OUT
    HALT
DATA .FILL 0xFFFF
```

a. What is the output (in hex) after the OUT statement at the symbol PRINT1 finishes execution?

0x0009

b. What is the output (in hex) after the OUT statement at the symbol PRINT2 finishes execution?
Explain your answer.

0x3003

c. Complete the following incomplete code snippet that uses the memory-mapped LC-3 registers KBSR and KBDR to take input from the keyboard instead of the GETC instruction. Your code should store the value entered from keyboard in register R1.

(Assume KBSR is mapped to address xFE00 and KBDR is mapped to address xFE02.)

```
ECHO: _____, ____, KBSR
      _____
      _____, R1, KBDR
```

```
KBSR .FILL xFE00
KBDR .FILL xFE02
```

```
ECHO: LDI, R0, KBSR
      BRnp ECHO
      LDI, R1, KBDR
```

```
KBSR .FILL xFE00
KBDR .FILL xFE02
```

4. Short answer questions (5 points)

a. Briefly describe what happens in the **linking** phase of an assembly program.

Linking is resolving symbols between independent object files.

b. How are the Display Data Register (DDR) and Display Status Register (DSR) used when TRAP x21 (OUT) is called?

When the monitor is ready to display another character, DSR[15] is set to 1. DSR[7:0] is displayed and DSR[15] is set to 1. Any other data written to DDR is ignored while DSR[15] is zero.

c. What is a **service routine** in LC-3? Give an example.

A service routine is a function that performs a specific operation (optionally as part of the operating system). In LC-3, traps are service routines. Any trap will work here as an example.

d. **Briefly** describe the difference between synchronous and asynchronous I/O.

Synchronous I/O events occur at fixed, predictable rates. CPU reads every X seconds. Asynchronous I/O is unpredictable. Can use flag to achieve I/O. Example is keyboard input.

5. The following LC-3 assembly code implements a **subroutine**. After taking input, it outputs a lowercase letter (a-z) if the input was uppercase (A-Z), and a "N" otherwise. (5 points)

```
SUBR GETC

    ADD R5, R0, 0
    NOT R5, R5
    ADD R5, R5, 1

    LD R1, LBOUND
    LD R2, UBOUND

    ADD R3, R5, R2
    BRn PNOT
    ADD R3, R5, R1
    BRp PNOT

    LD R3, DIFF

    ADD R0, R0, R3
    OUT
    BRnzp FINISH

PNOT LD R0, N
    OUT
FINISH RET

UBOUND .FILL x5A ; ASCII value for "Z"
LBOUND .FILL x41 ; ASCII value for "A"
N      .FILL x4E ; ASCII value for "N"
DIFF   .FILL x20
```

a. Is this code able to successfully return? Explain why or why not.

No. R7 is overwritten by GETC and OUT.

b. Add lines to the code above to make this subroutine **callee-saved**. You should not modify any existing lines. **Clearly indicate which lines you have added and where.**

Before GETC:
ST R0, SAVE0
ST R1, SAVE1
ST R2, SAVE2
ST R3, SAVE3

ST R5, SAVE5
ST R7, SAVE7

Before FINISH:

LD R0, SAVE0
LD R1, SAVE1
LD R2, SAVE2
LD R3, SAVE3
LD R5, SAVE5
LD R7, SAVE7

After DIFF

SAVE0 .FILL x0
SAVE1 .FILL x0
SAVE2 .FILL x0
SAVE3 .FILL x0
SAVE5 .FILL x0
SAVE7 .FILL x0

6. Examine the code below. You may assume that at the start of the program, all registers are set to 0. (4 points)

```
.ORIG x3000
GETC
AND R3, R3, 0
ADD R3, R3, R0

LD R1, ONE
NOT R1, R1
ADD R1, R1, 1

PRINT LD R0, C
OUT
LD R0, A
OUT
LD R0, T
OUT

ADD R0, R3, R1
BRZ FINISH

LD R0, S
OUT

FINISH HALT
C .FILL x63 ; ASCII 'c'
A .FILL x61 ; ASCII 'a'
T .FILL x74 ; ASCII 't'
S .FILL x73 ; ASCII 's'
ONE .FILL x31 ; ASCII '1'
```

a. If the input is the decimal number "1", what is output on the screen at the end of the program?

cat

b. OUT only prints out one character at a time. Instead of printing individually, we decide to replace the code stored at memory locations **x3006 to x300B** with the following, much shorter block of code:

```
PRINT LD R0, C
PUTS
```

Will this output the same result as a)? Why or why not? **You must explain your answer for credit.**

No. It'll output a bunch of garbage (whatever is stored at memory location x63 until it hits a null terminator). This is because PUTS uses the register value as an address, not as a value.

7. Multiple choice questions. Circle **one** answer for each question. (4 points)

(i) Which of the following can **not** be used multiple times in a single assembly program?

- a) .FILL
- b) .ORIG**
- c) .BLKW
- d) .STRINGZ

(ii) Assume that a LC-3 processor receives interrupts from 3 I/O devices (A, B and C) simultaneously. The priority levels for the interrupts are given below:

A) PL4 B) PL2 C) PL6

Assuming that no other interrupts come in, which of the above interrupts is serviced **last**?

- a) A
- b) B**
- c) C
- d) Any selected at random

(iii) Our program begins at memory location x4000. We want to load the value x4020 into R3. Which LC-3 instruction can we use to accomplish this in a **single** line?

- a) LEA**
- b) LD
- c) ST
- d) LDI

(iv) Which register is used to store **input data** after IN is called?

- a) R1
- b) R0**
- c) R7
- d) R4

LC-3 Instruction Set (Entered by Mark D. Hill on 03/14/2007; last update 03/15/2007)

PC': incremented PC. setcc(): set condition codes N, Z, and P. mem[A]:memory contents at address A.
 SEXT(immediate): sign-extend immediate to 16 bits. ZEXT(immediate): zero-extend immediate to 16 bits.

```

15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
+-----+-----+-----+-----+-----+-----+
| 0 0 0 1 | DR | SR1 | 0 | 0 0 | SR2 |
+-----+-----+-----+-----+-----+
ADD DR, SR1, SR2 ; Addition
DR ← SR1 + SR2 also setcc()

+-----+-----+-----+-----+-----+-----+
| 0 0 0 1 | DR | SR1 | 1 | imm5 |
+-----+-----+-----+-----+-----+
ADD DR, SR1, imm5 ; Addition with Immediate
DR ← SR1 + SEXT(imm5) also setcc()

+-----+-----+-----+-----+-----+-----+
| 0 1 0 1 | DR | SR1 | 0 | 0 0 | SR2 |
+-----+-----+-----+-----+-----+
AND DR, SR1, SR2 ; Bit-wise AND
DR ← SR1 AND SR2 also setcc()

+-----+-----+-----+-----+-----+-----+
| 0 1 0 1 | DR | SR1 | 1 | imm5 |
+-----+-----+-----+-----+-----+
AND DR, SR1, imm5 ; Bit-wise AND with Immediate
DR ← SR1 AND SEXT(imm5) also setcc()

+-----+-----+-----+-----+-----+-----+
| 0 0 0 0 | n | z | p | PCOffset9 | GO | ((n and N) OR (z AND Z) OR (p AND P))
+-----+-----+-----+-----+-----+-----+
BRx, label (where x={n,z,p,zp,np,nz,nzp}); Branch
if(GO is true) then PC←PC' + SEXT(PCOffset9)

+-----+-----+-----+-----+-----+-----+
| 1 1 0 0 | 0 0 0 | BaseR | 0 0 0 0 0 0 |
+-----+-----+-----+-----+-----+
JMP BaseR ; Jump
PC ← BaseR

+-----+-----+-----+-----+-----+-----+
| 0 1 0 0 | 1 | PCOffset11 |
+-----+-----+-----+-----+-----+
JSR label ; Jump to Subroutine
R7 ← PC', PC ← PC' + SEXT(PCOffset11)

+-----+-----+-----+-----+-----+-----+
| 0 1 0 0 | 0 | 0 0 | BaseR | 0 0 0 0 0 0 |
+-----+-----+-----+-----+-----+
JSRR BaseR ; Jump to Subroutine in Register
temp ← PC', PC ← BaseR, R7 ← temp

+-----+-----+-----+-----+-----+-----+
| 0 0 1 0 | DR | PCOffset9 |
+-----+-----+-----+-----+-----+
LD DR, label ; Load PC-Relative
DR ← mem[PC' + SEXT(PCOffset9)] also setcc()

+-----+-----+-----+-----+-----+-----+
| 1 0 1 0 | DR | PCOffset9 |
+-----+-----+-----+-----+-----+
LDI DR, label ; Load Indirect
DR←mem[mem[PC'+SEXT(PCOffset9)]] also setcc()

+-----+-----+-----+-----+-----+-----+
| 0 1 1 0 | DR | BaseR | offset6 |
+-----+-----+-----+-----+-----+
LDR DR, BaseR, offset6 ; Load Base+Offset
DR ← mem[BaseR + SEXT(offset6)] also setcc()

+-----+-----+-----+-----+-----+-----+
| 1 1 1 0 | DR | PCOffset9 |
+-----+-----+-----+-----+-----+
LEA, DR, label ; Load Effective Address
DR ← PC' + SEXT(PCOffset9) also setcc()

+-----+-----+-----+-----+-----+-----+
| 1 0 0 1 | DR | SR | 1 | 1 1 1 1 1 |
+-----+-----+-----+-----+-----+
NOT DR, SR ; Bit-wise Complement
DR ← NOT(SR) also setcc()

+-----+-----+-----+-----+-----+-----+
| 1 1 0 0 | 0 0 0 | 1 1 1 | 0 0 0 0 0 0 |
+-----+-----+-----+-----+-----+
RET ; Return from Subroutine
PC ← R7

+-----+-----+-----+-----+-----+-----+
| 1 0 0 0 | 0 0 0 0 0 0 0 0 0 0 0 0 0 0 |
+-----+-----+-----+-----+-----+
RTI ; Return from Interrupt
See textbook (2nd Ed. page 537).

+-----+-----+-----+-----+-----+-----+
| 0 0 1 1 | SR | PCOffset9 |
+-----+-----+-----+-----+-----+
ST SR, label ; Store PC-Relative
mem[PC' + SEXT(PCOffset9)] ← SR

+-----+-----+-----+-----+-----+-----+
| 1 0 1 1 | SR | PCOffset9 |
+-----+-----+-----+-----+-----+
STI, SR, label ; Store Indirect
mem[mem[PC' + SEXT(PCOffset9)]] ← SR

+-----+-----+-----+-----+-----+-----+
| 0 1 1 1 | SR | BaseR | offset6 |
+-----+-----+-----+-----+-----+
STR SR, BaseR, offset6 ; Store Base+Offset
mem[BaseR + SEXT(offset6)] ← SR

+-----+-----+-----+-----+-----+-----+
| 1 1 1 1 | 0 0 0 0 | trapvect8 |
+-----+-----+-----+-----+-----+
TRAP ; System Call
R7 ← PC', PC ← mem[ZEXT(trapvect8)]

+-----+-----+-----+-----+-----+-----+
| 1 1 0 1 |
+-----+-----+-----+-----+-----+
; Unused Opcode
Initiate illegal opcode exception
15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

```

Assembler Directives

<i>Opcode</i>	<i>Operand</i>
.ORIG	address
.END	
.BLKW	n
.FILL	n
.STRINGZ	n-character string

Trap Codes

<i>Code</i>	<i>Equivalent</i>
HALT	TRAP x25
IN	TRAP x23
OUT	TRAP x21
GETC	TRAP x20
PUTS	TRAP x22