

An Analysis of Persistent Memory Use with WHISPER

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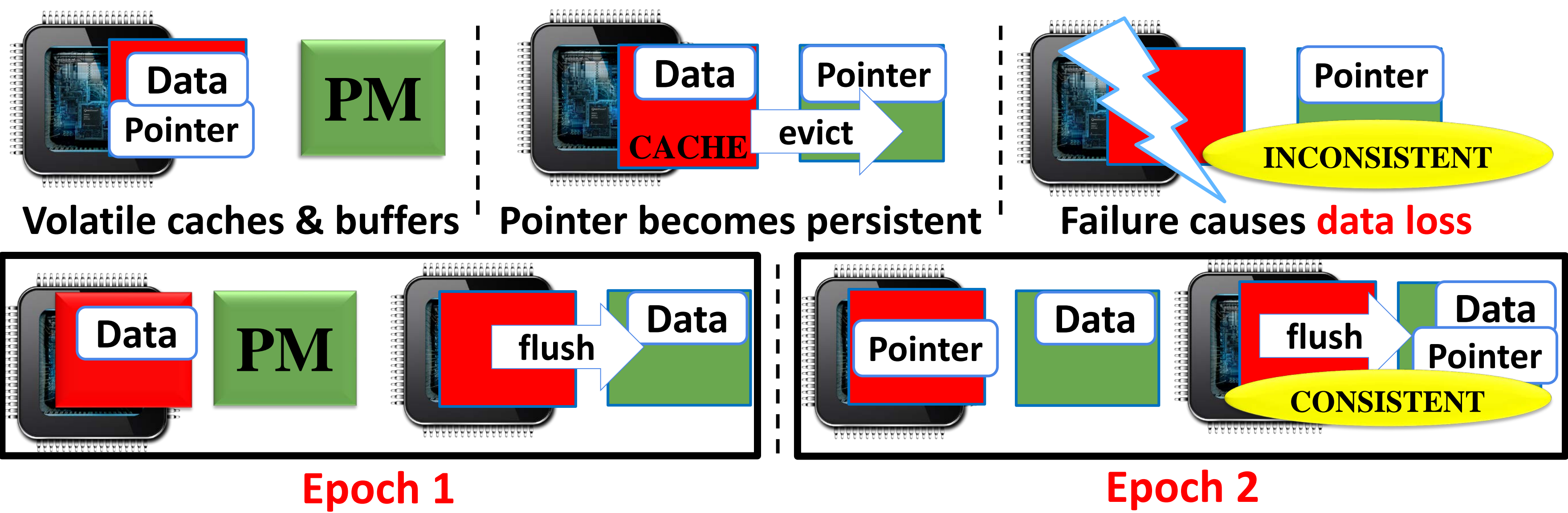
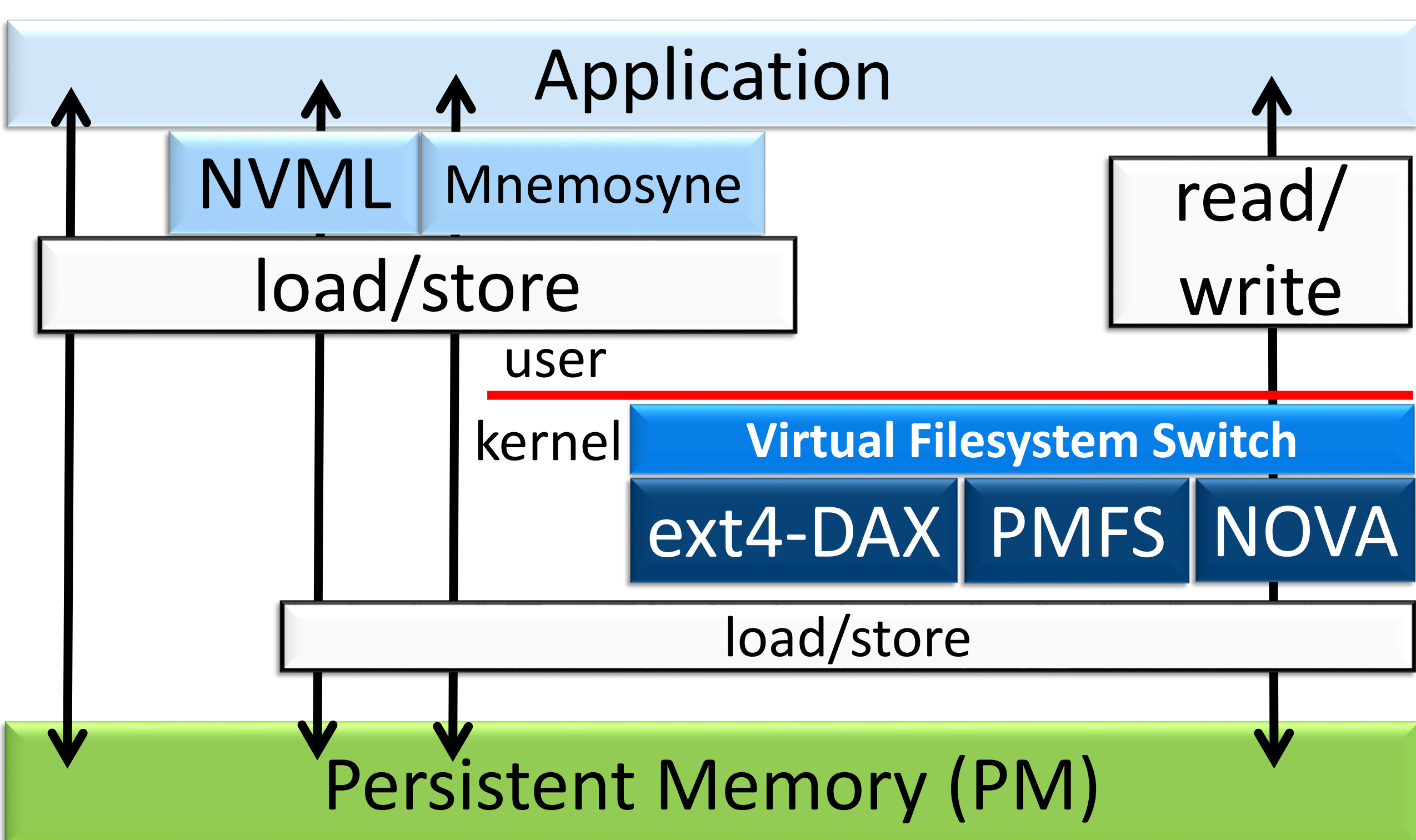
MOTIVATION: A STANDARD BENCHMARK SUITE FOR PERSISTENT MEMORY (PM)

- Address **lack of understanding** of PM usage
- Minimize use of ad hoc micro-benchmarks
- Study **diverse, real interfaces** for consistency and durability of data and metadata in PM
- Aid design of better PM runtimes
- Establish a **standard for evaluating** PM runtimes

WHISPER = Wisconsin-HP Labs Suite for Persistence
research.cs.wisc.edu/multifacet/whisper/

Application	System	Brief Description	(*Adapted to PM)
Echo*	Native	Multi-version KV store	(U of Washington)
N-store*	Native	Fast, in-memory relational DB	(Carnegie-Mellon)
Redis	NVML	In-memory KV store	(Intel)
C-tree, Hashmap	NVML	Multi-threaded microbenchmarks	(Intel)
Vacation*	Mnemosyne	OLTP travel reservation system	(U of Wisconsin)
Memcached*	Mnemosyne	Distributed in-memory KV store	(U of Wisconsin)
NFS	PMFS	Linux server/client for remote PM access	(Intel)
Exim	PMFS	Popular mail server using mailbox format	
MySQL	PMFS	Widely used RDBMS for OLTP	

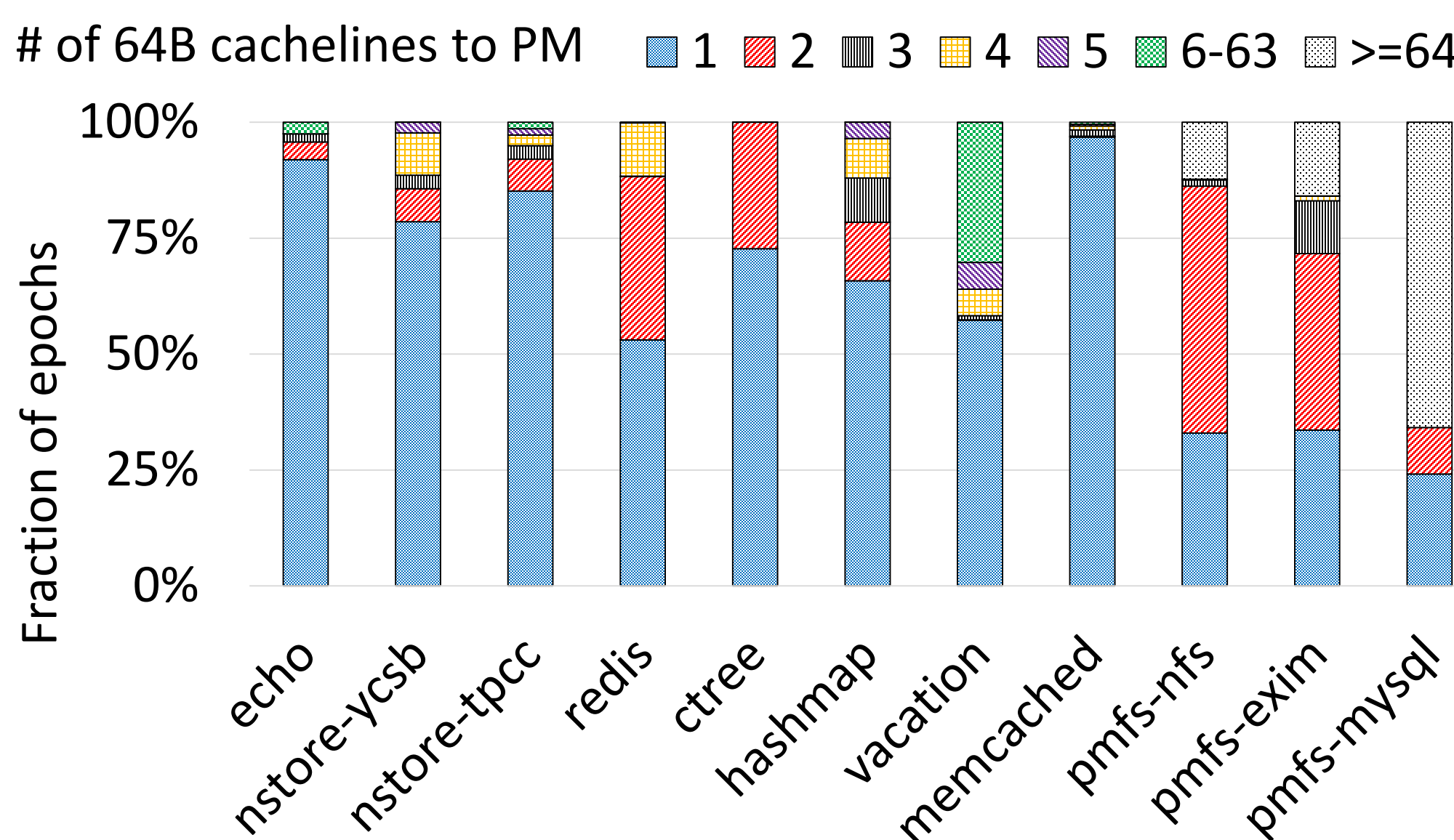
BACKGROUND : PM WORKLOADS, EPOCHS, CONSISTENCY & DURABILITY



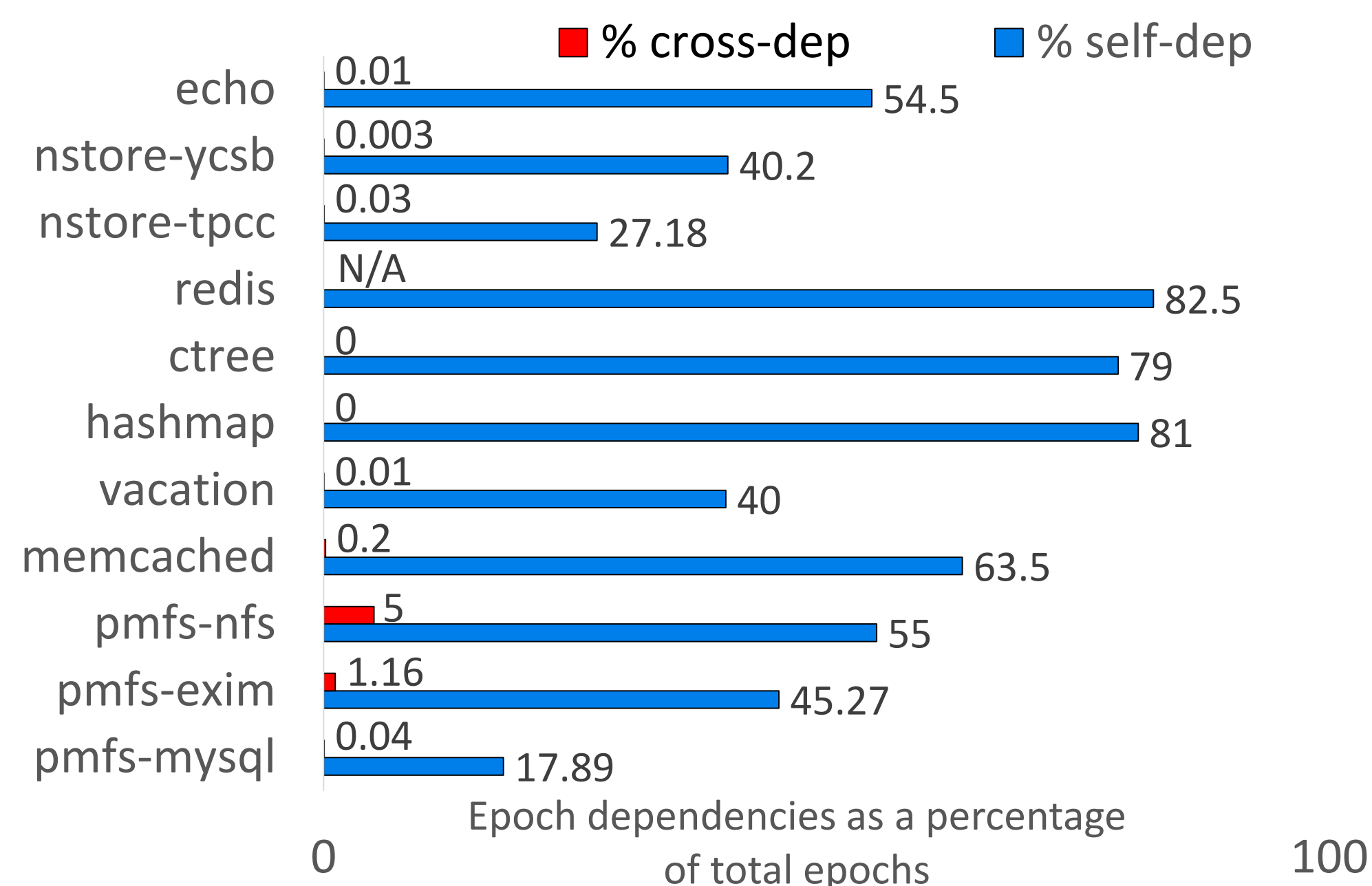
- Controlling order of writes to PM preserves consistency at all stages.
- Epoch** = Set of writes to PM guaranteed to be durable before any subsequent writes become durable.

SELECTED ANALYSIS RESULTS

- What % of accesses in WHISPER applications are to PM?
- How many epochs are there in a transaction to PM?
(**∴ Durability involves writes to PM & costly if enforced on each epoch**)
- How often do epochs from same or different threads write to the same cacheline to PM ?
(**∴ Epochs writing to same cache line depend on each other and stall execution**)



Takeaway: Enforce durability only after last epoch in a transaction as enforcing it on each epoch can slow down execution and impede transaction performance.



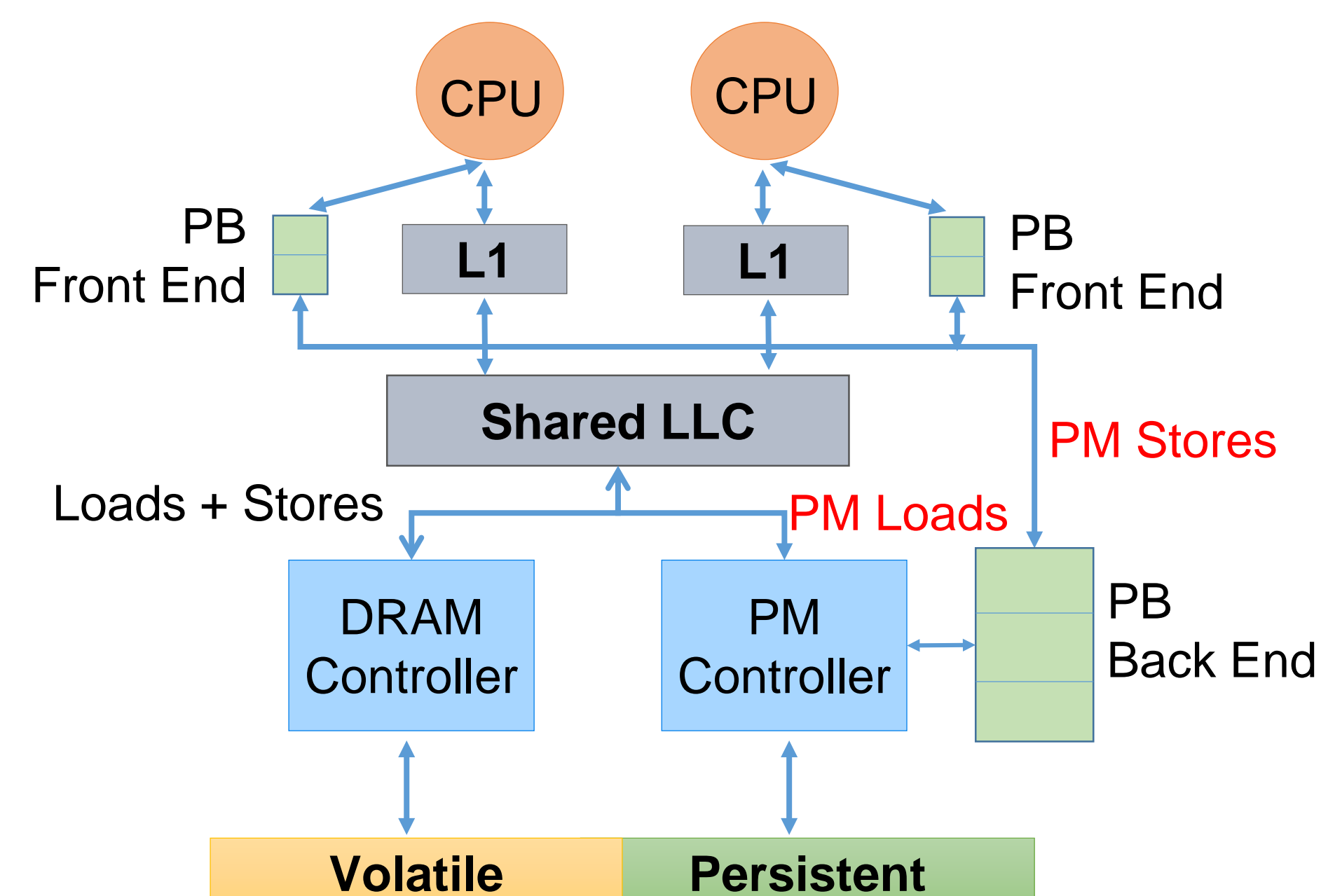
Takeaway: Design hardware with multi-versioned buffers and optimize data structures in software to minimize dependencies among epochs.

ANALYSIS SUMMARY

HOPS DESIGN

PERSIST BUFFERS IN HOPS

- 4% accesses are to PM, 96% are to DRAM**
- 5-50 epochs per transaction**
- Self-dependencies **common**
- Cross-dependencies are **rare**
- Volatile memory hierarchy (almost) unchanged by PBs
- Order Epochs without flushing (OFENCE and DFENCE)
- Allow multiple copies of same cacheline in PBS via timestamps
- Correct, conservative method using coherence, timestamps



24% faster than Intel extensions for PM