Topology Planning for Long Distance Wireless Mesh Networks

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Under Guidance of Dr. Bhaskaran Raman

Outline

- Motivation & Background
- Problem statement, Uniqueness
- Problem formulation
 - Definitions, dependences
- Solution approach
- Evaluation
- Conclusions

Rural Connectivity

• Goal:

- Provide 100% connectivity to rural India, as 74% Indian population rural.
- What Technology to use ?

WiFi-based Rural Networks

Cost of land-line telephony: \$400 per line --> \$200 per line 400 million lines for India ==> \$80 billion

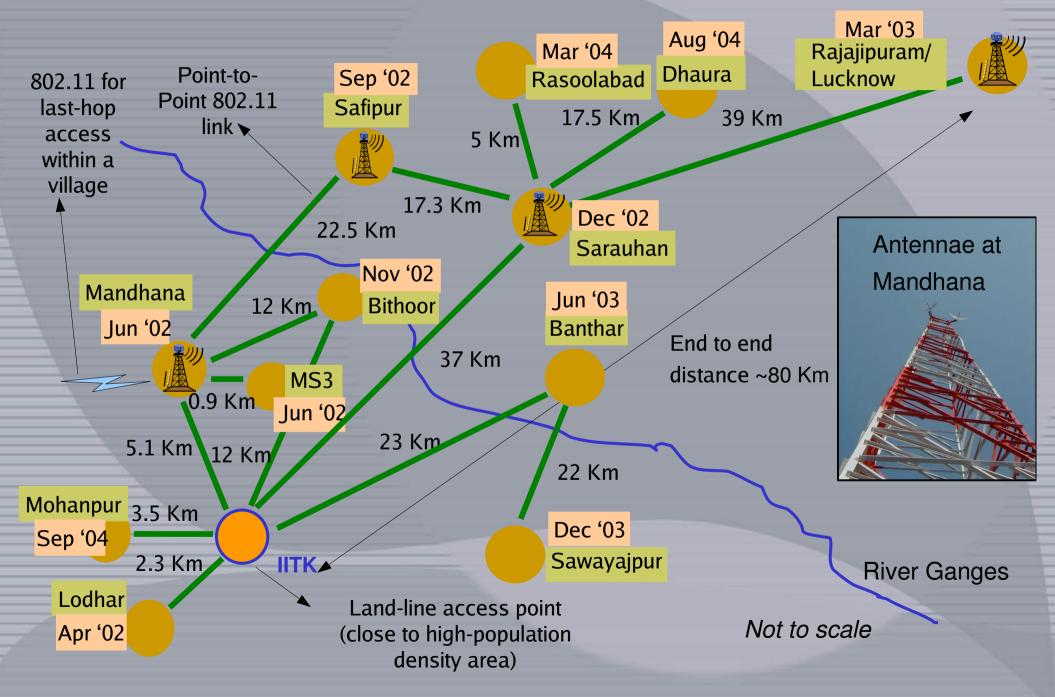
Cellular technology is *value-priced* (expensive for rural deployment)



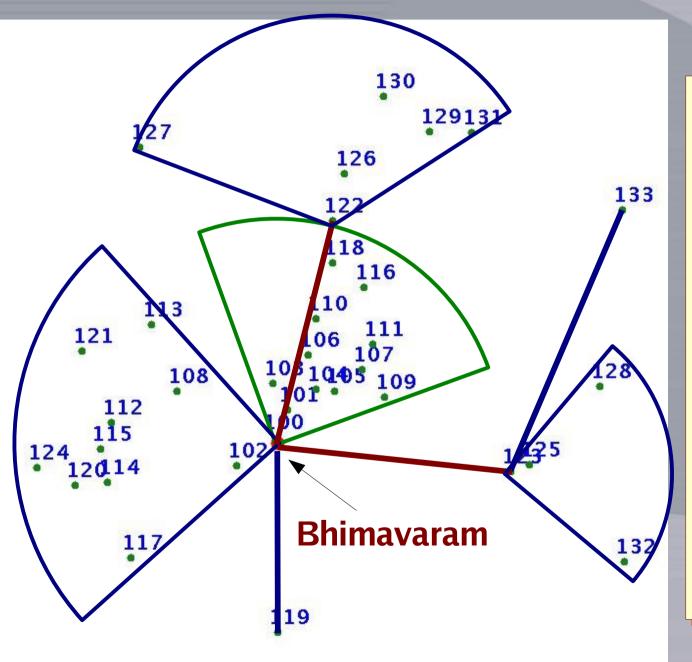
WiMAX (IEEE 802.16) yet to hit the market Unclear if it will be inexpensive enough for rural areas

In contrast: WiFi equipment *cost-priced* Rs 2-5K per WiFi radio Inexpensive enough for rural deployment

Digital Gangetic Plains: Testbed



The Ashwini Project



- Byrraju foundation,
 West Godavari,
 Andhra Pradesh
- To connect 34 villages
- Video-based health, education services

A WiFi Network in Djurslands, Denmark

www.DjurslandS.net



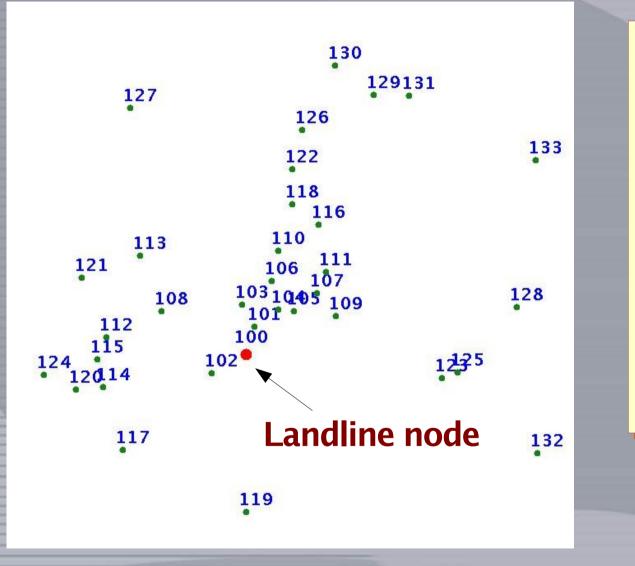
Problem Statement Motivation

• Problem:

- India has order of 6 lakh villages.
- No automated method exists to plan such networks.
- Current methods highly cost inefficient and ad-hoc.
- Last Mile Connectivity problem:
 - 85% of the villages are within 20 Km radius of a fiber Point-of-Presence (PoP).

(source:A. Jhunjhunwala et. al, Role of Wireless technologies in connecting rural India, Indian Journal of Radio & Space Physics'05)

Problem Statement



- Given: a set of village nodes, a single landline node
- Requirement: connect all villages to the landline in a network

Primary concern: cost and bandwidth guarantees.

Problem Uniqueness

- Coverage only at village nodes (unlike cellular coverage)
- Line-of-Sight requirement
- Focus on cost optimality
 - Cost dominated by towers

)		Tower/mast height (m)	Cost (x1000 Rs.)
	_ine-of-Sight (LOS)	10	4
		15	6
		21	36
		24	41
		27	48
Tower	Mast	30	82

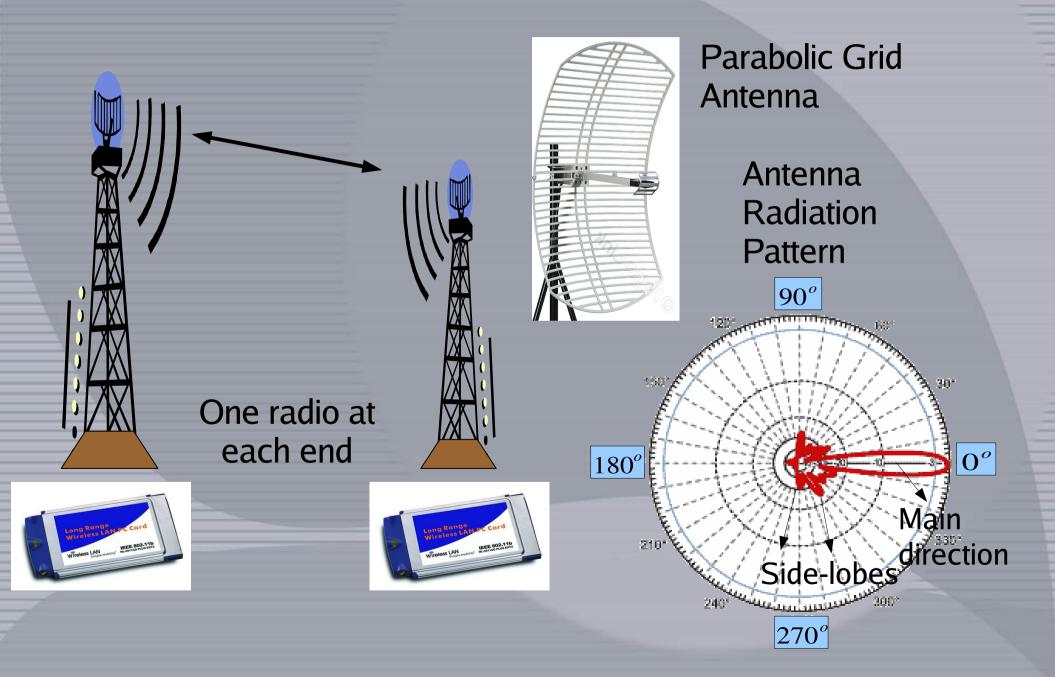
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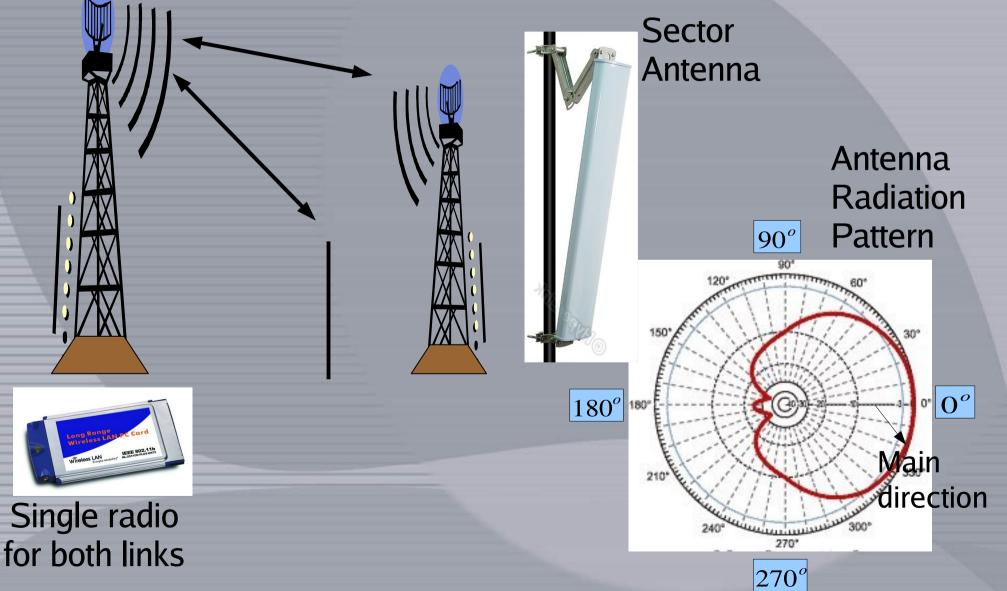
Problem Formulation

- Assumptions:
 - Antenna towers to be placed only at villages
 - Tree topology
- Application requirement:
 - Throughput per village: say, 384Kbps (for video)
- Definitions...
- Dependences...

Definition: Point-to-Point Link



Definition: Point-to-MultiPoint Link

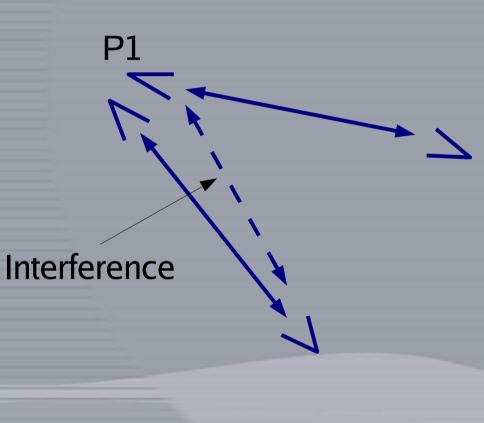


Notation: Top View

Point-to-Point (P2P) links

Point-to-MultiPoint (P2MP) link

Definitions: Transmit Power & Interference



Definitions: 2P and TDMA MAC

2P: The two links can operate simultaneously

TDMA: The two links operate in turns

Dependence: Throughput depends on the MAC

2P: The two links can operate simultaneously

TDMA: The two links operate in turns

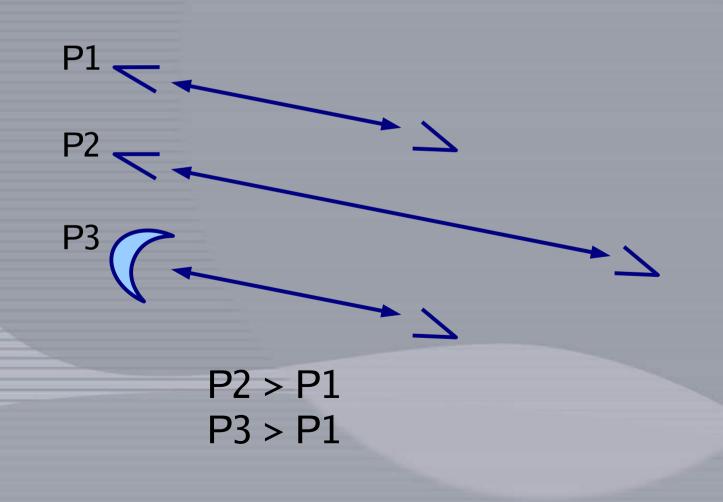
Dependence: Throughput depends on Link/Antenna Type

Point-to-Point (P2P) links

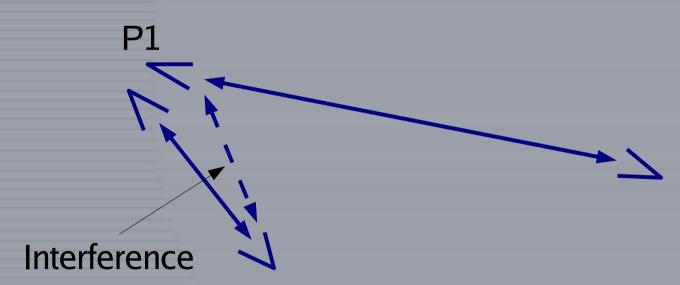
Point-to-MultiPoint (P2MP) link

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Dependence: Transmit Power (required) depends on Link (length) & Antenna Type



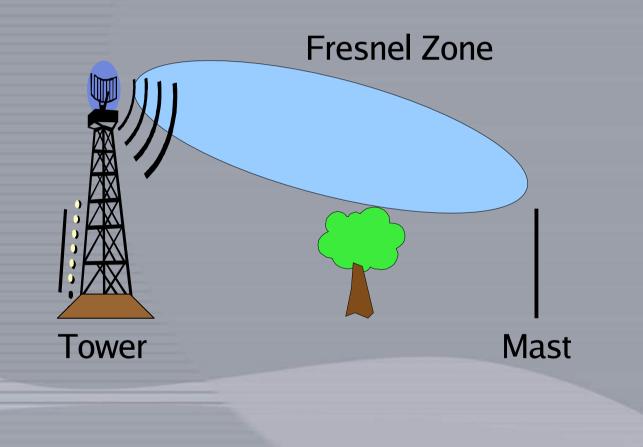
Dependence: MAC (feasibility) depends on the Transmit Powers



Signal to Interference Ratio should be above threshold

$$P_R \quad I_R \ge SIR_{regd}$$

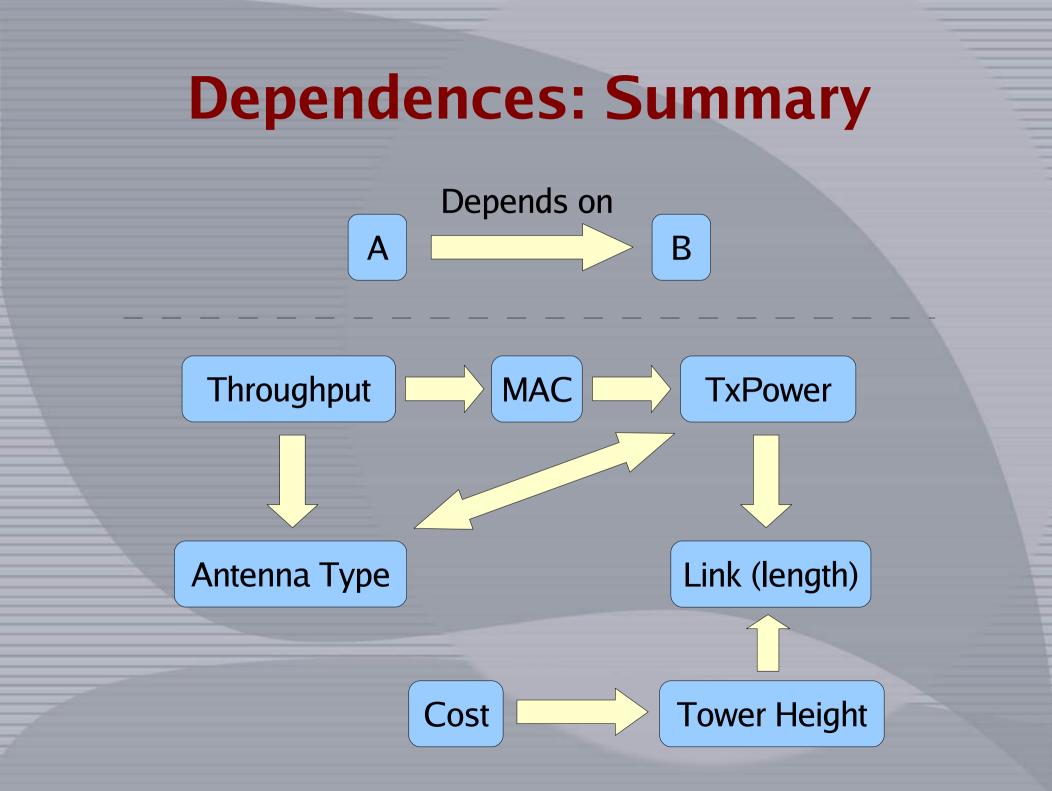
Dependence: Tower Height(s) (required) depends on Link (length)

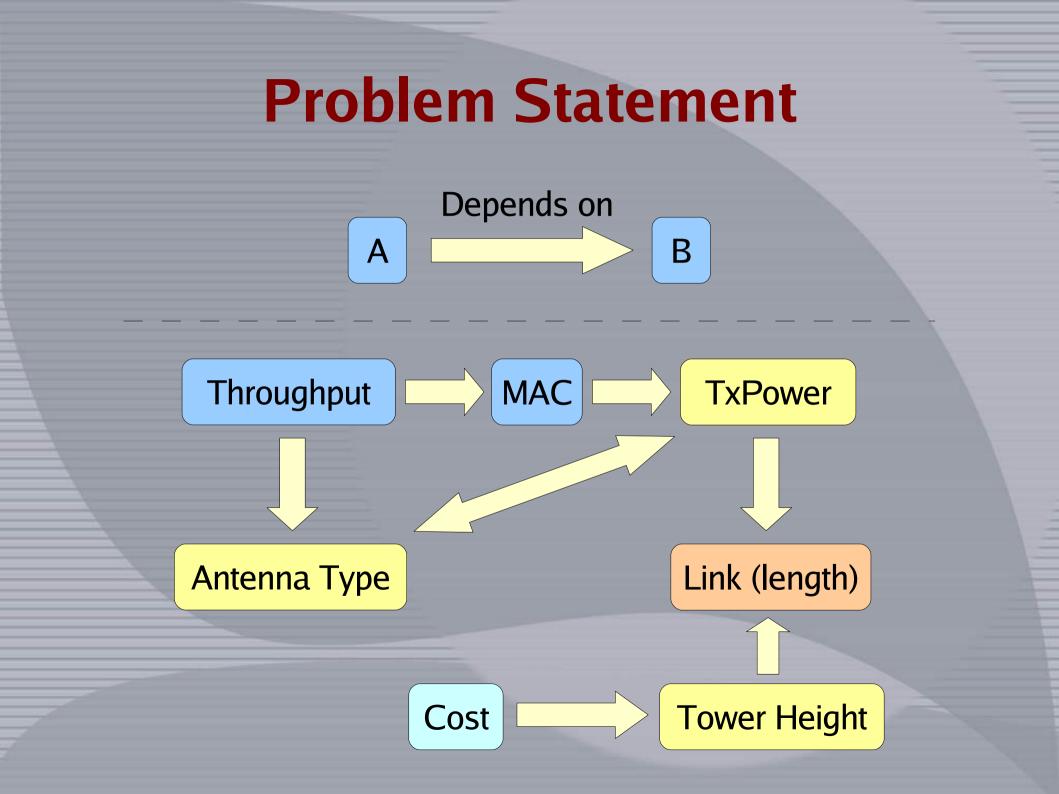


Dependence: Cost depends on Tower Height

- Cost takes quantum jumps due to change in underlying tower design
 - Increases superlinearly with height

Tower/mast height (m)	Cost (x1000 Rs.)
10	4
15	6
21	36
24	41
27	48
30	82
45	220



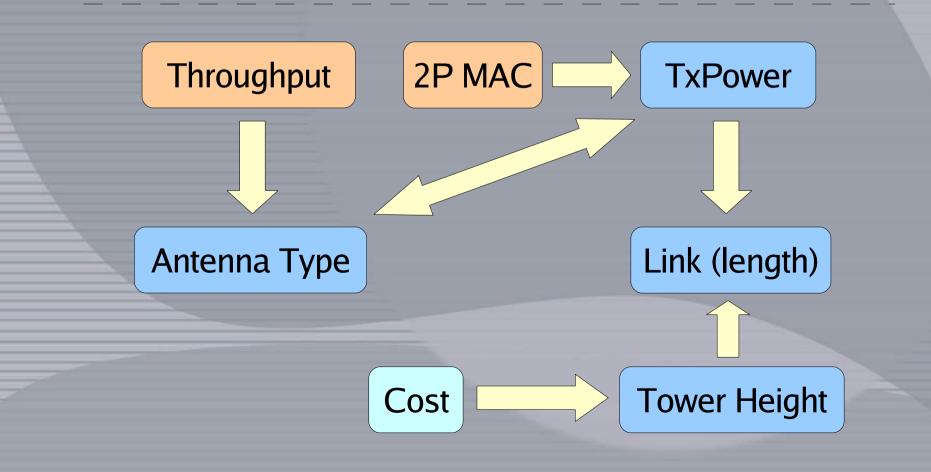


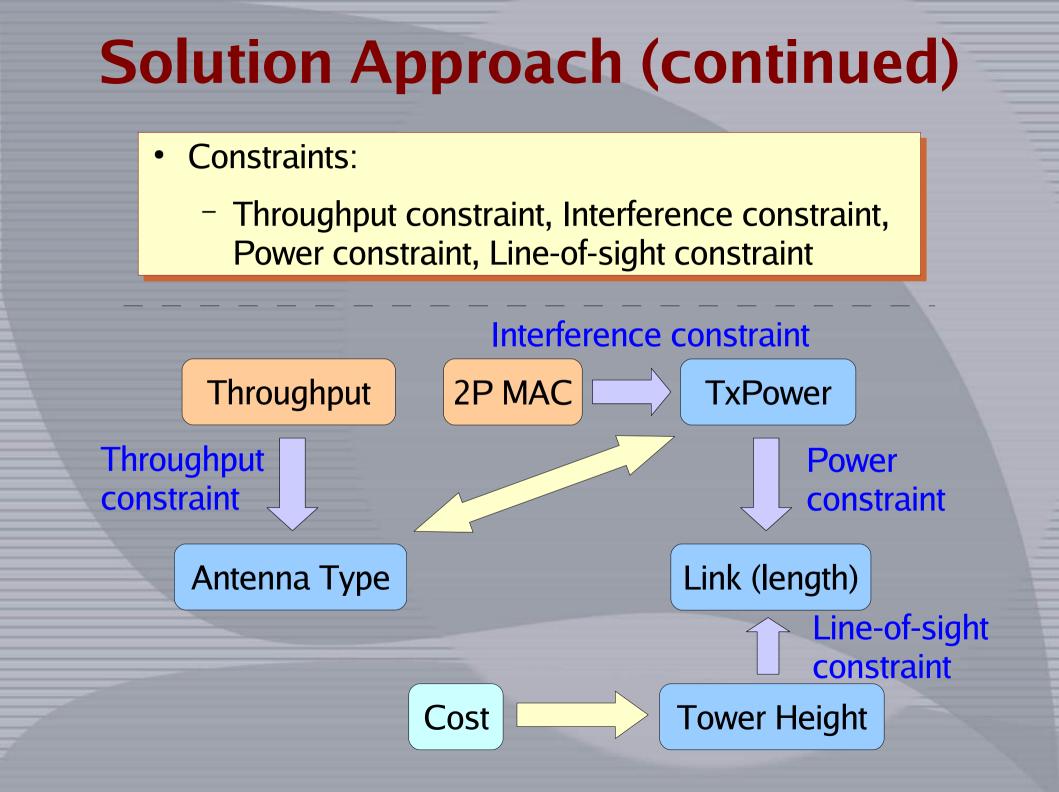
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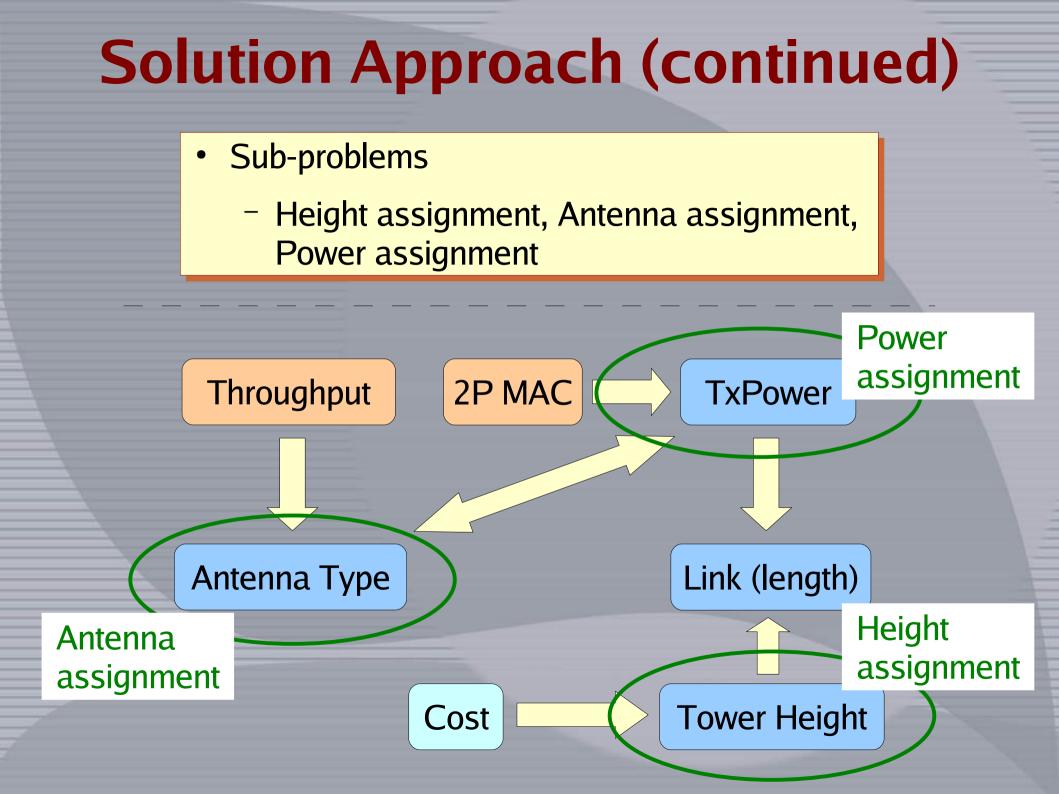
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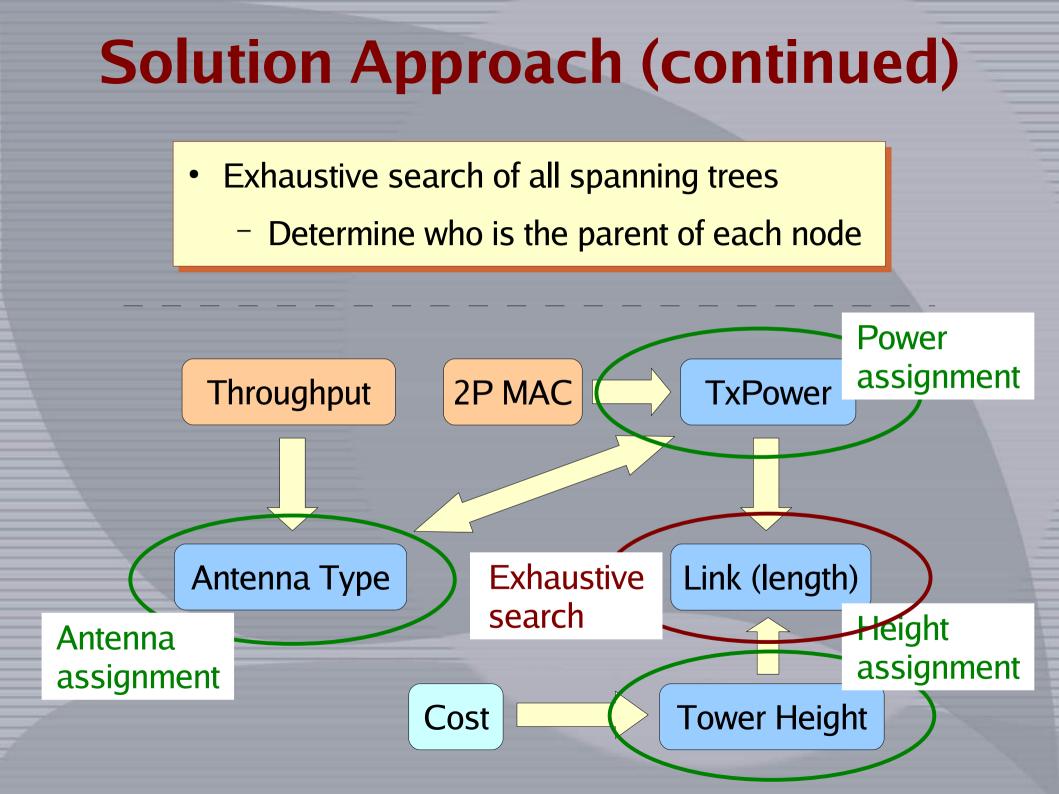
Solution Approach (Overview)

- Fix 2P MAC: more efficient
- Throughput, 2P MAC feasibility are constraints
- Cost minimization is an objective

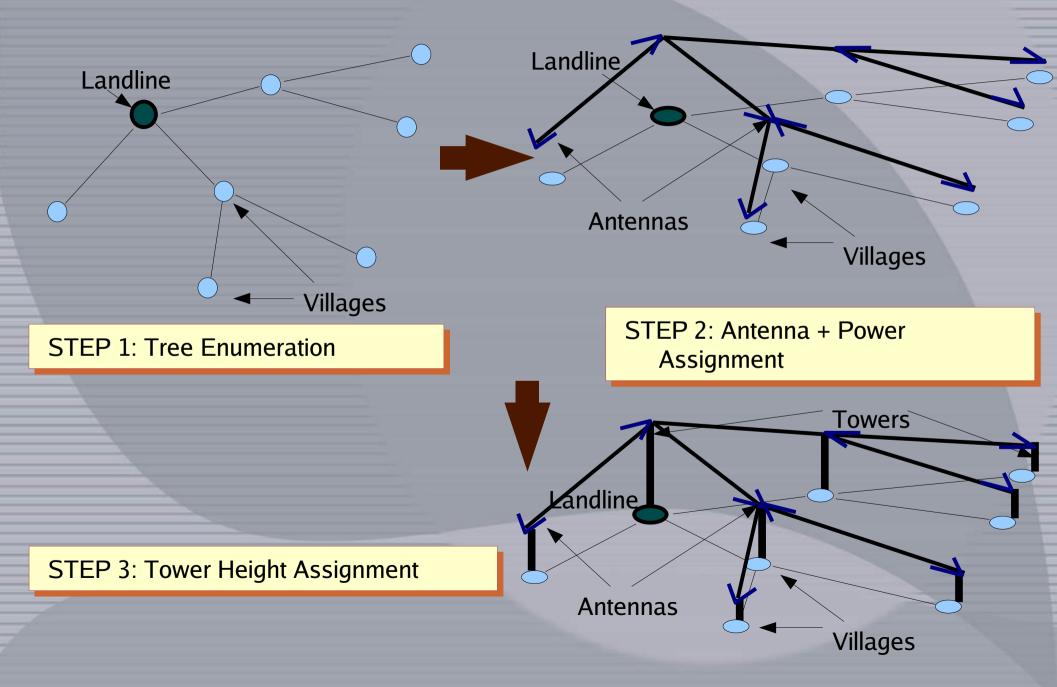








Solution Methodology

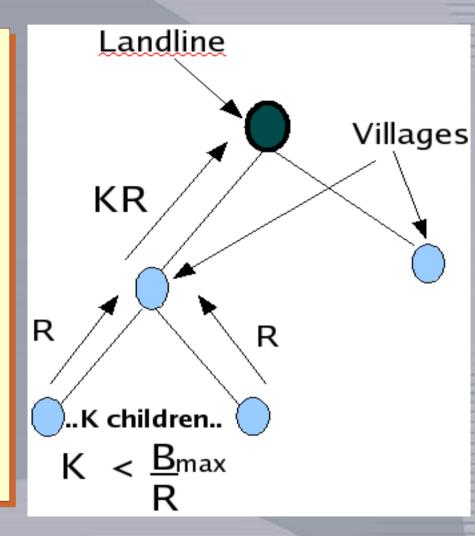


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 - Tree Enumeration
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Tree Enumeration

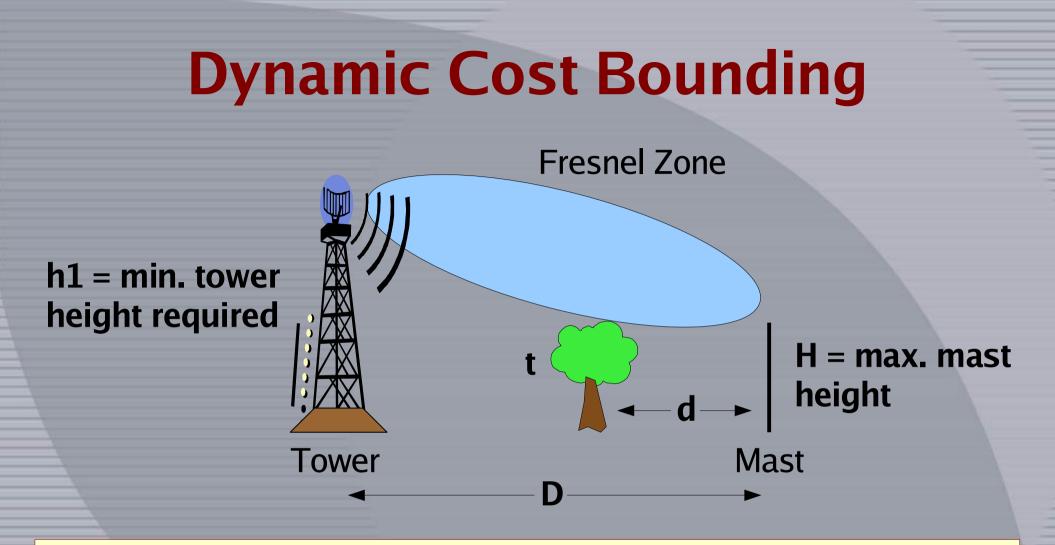
- Exhaustive Enumeration of all spanning trees
 - Throughput check: MAC dependant
 - Other domain based relaxations



Optimizations on Exhaustive Search

Domain-knowledge based optimizations

- Eliminate "long" links to begin with
- Tree depth restriction
- Dynamic cost bounding

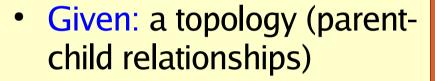


- Observation-1: height of level-2 tower determined by children set
- Observation-2: given a link-length, can lower-bound tower height
- Implication: can lower-bound the cost of a sub-tree
 - Can pre-compute lower-bounds for efficiency

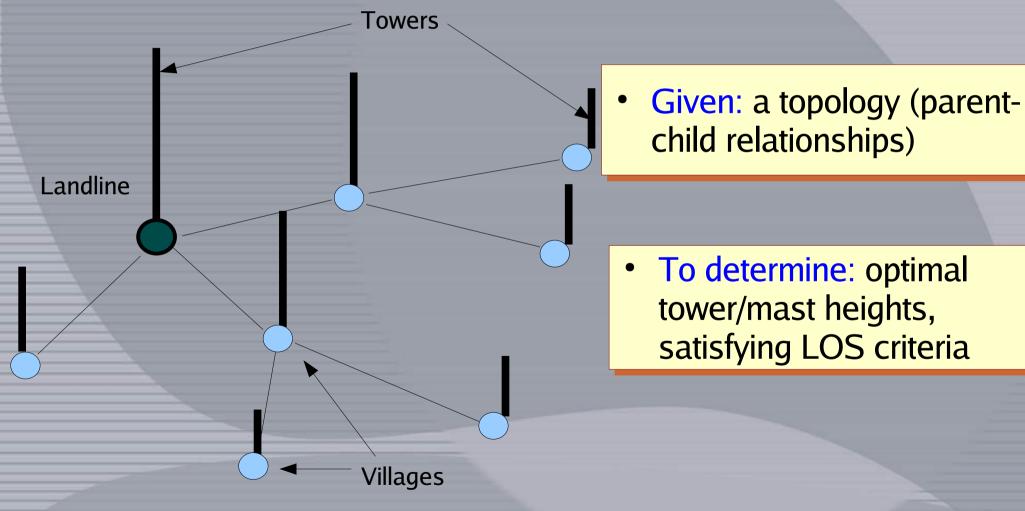
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Height Assignment: Problem Statement

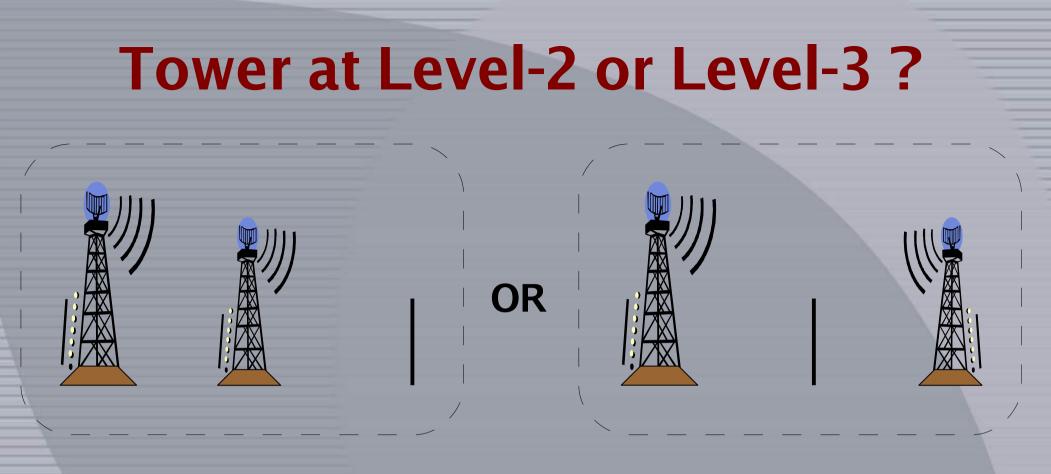


Height Assignment: Problem Statement



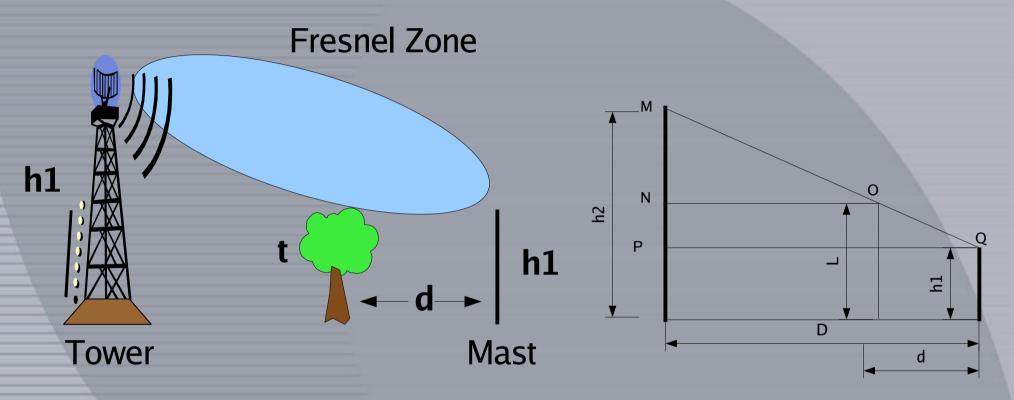
Height Assignment: Simplifications

- 2-hop topology only
 - One hop ~ 10km ==> 20km radius ==> 40km dia
 - Accommodates significant # practical scenarios
- Tower at central location: say 50m
 - Typically in a town with reasonably tall buildings
- Assumption: No link between two masts (tree obstructions)
- Assumption: tower cost is linear in height
 - But we distinguish between towers and masts
 - Cost is a piece-wise linear function of height



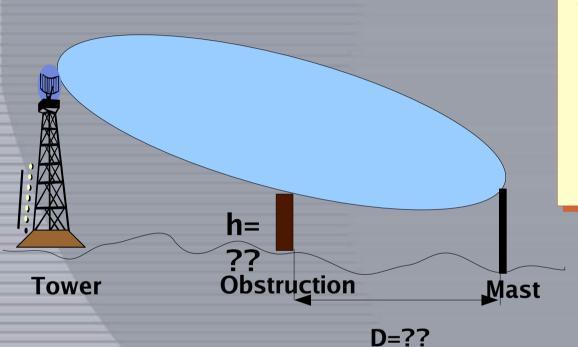
- Observation-1: tower heights can be interchanged in a link, retaining the same cost
 - Note: does not hold if terrain uncertainties are considered
- Observation-2: # level-3 nodes (leaves) > # level-2 nodes
- Implication: towers at level-2 and masts at level-3

The LP Formulation



- Linear equations for obstruction clearance
- Linear cost optimization function

Finding Obstruction Height

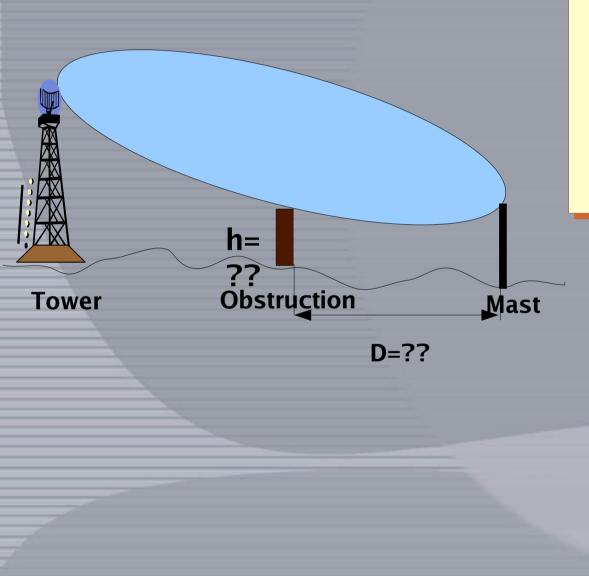


• How to estimate the

Maximum
 Obstruction Height

- And it's location ?

Finding Obstruction Height



- How to estimate the
 - Maximum
 Obstruction Height
 - And it's location ?

- Use freely available Satellite data (ftp://e0srp01u.ecs.nasa.gov)
- Interpolate to estimate.

Outline

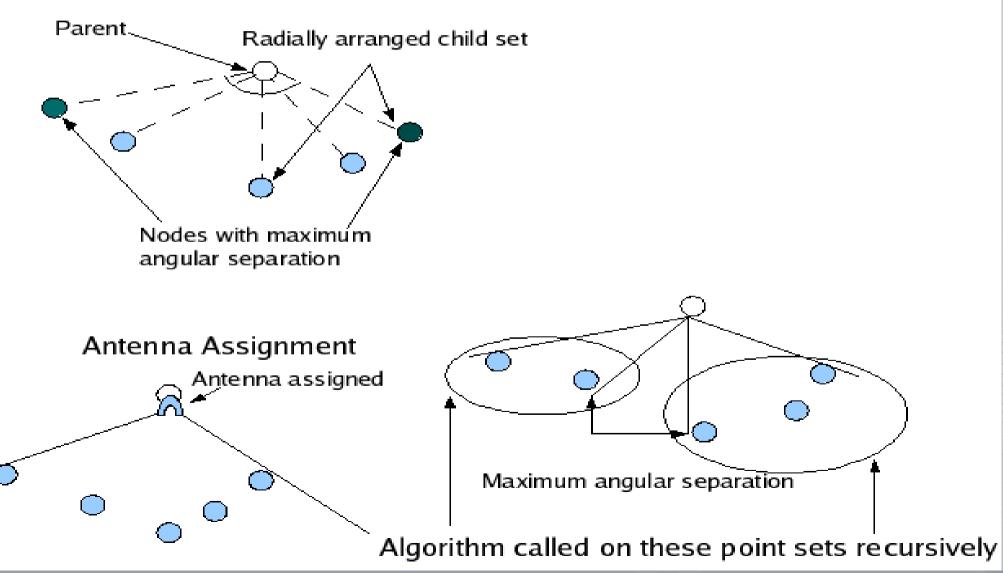
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Antenna Assignment

- Problem Statement:
 - Given a node and its children
 - What antenna types?
 - How many?, and
 - In which directions to use?
 - So that interference is minimised.
 - Similar to the minimum set-cover problem
 - Solved locally: For a node and its child set
 - Chicken egg problem
 - Power assignment takes care that all the links are working anyway
 - Child always has a high gain directional antenna

Heuristic Antenna Assignment Algorithm

INPUT:



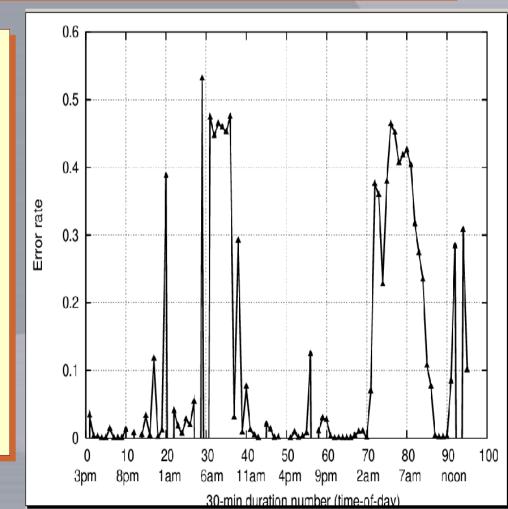
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Power Assignment: Motivation

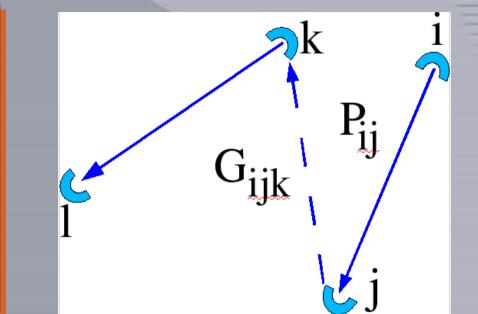
- Given: topology, antennas at nodes
- To determine: transmit powers for each radio

- Motivation: Direct causeeffect relationship between interference & packet error rate
 - Error rate as high as 50 %.
 - RTS/CTS not a remedy.



Power Assignment

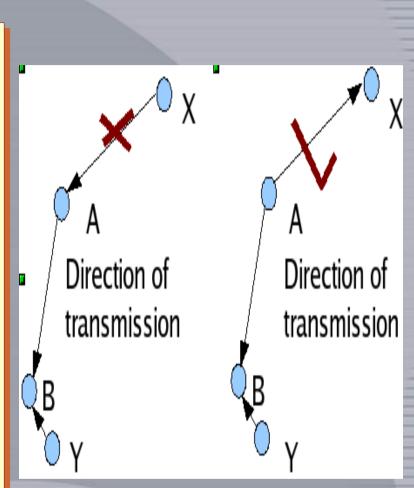
- P_{ij} = power transmitted by antenna at i for j, towards j.
- G_{ijk} = gain of antenna at i for j, towards k.
- $PL_{ii} = pathloss from i to j.$



- Maximum broadcast power criteria: $P_{ij} * G_{ijj} \leq P_{max}$
- Minimum received power criteria: $P_{ij} * G_{ij} * G_{jii} / PL_{ij} \ge P_{min}$
- SIR criteria: $P_{ij} * G_{ijj} * G_{jii} / PL_{ij} \ge SIR_{reqd} * \Sigma_{(k,l) \in R} P_{kl} * G_{jik} * G_{klj} / PL_{kj}$
 - R is the set of interfering links

Interfering links (for 2P)

- Denote by A⇒B, A transmitting towards B
- For 2P:
 - X⇒A, does not interfere
 - A⇒X and antenna different from A⇒B, interferes
 - Y⇒B, interferes
 - B⇒Y, interfers
 - If none of above satisfied, C⇒D interferes



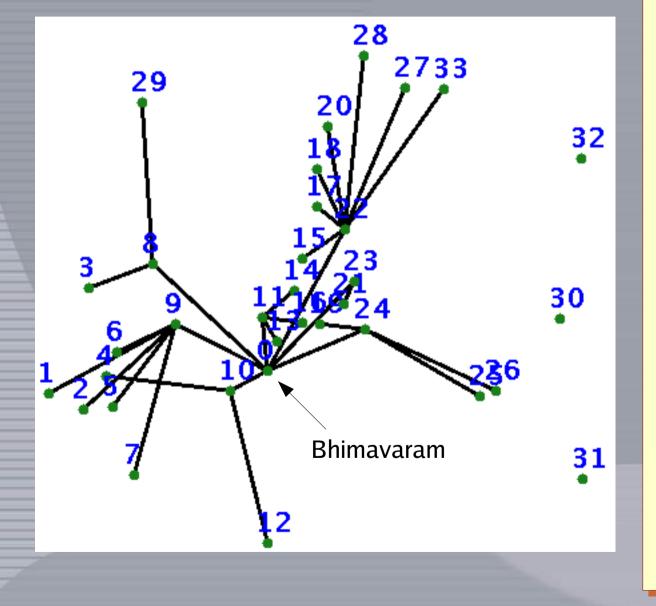
Putting it All Together

- Exhaustive enumeration of all spanning trees
 - Connected sub-trees at each stage
 - BFS-based enumeration
 - Eliminate "long" edges before starting enumeration
- For each sub-tree during enumeration:
 - Depth restriction check
 - Throughput check
 - Dynamic Cost Bounding
- For each spanning tree formed:
 - Height, antenna, power assignment

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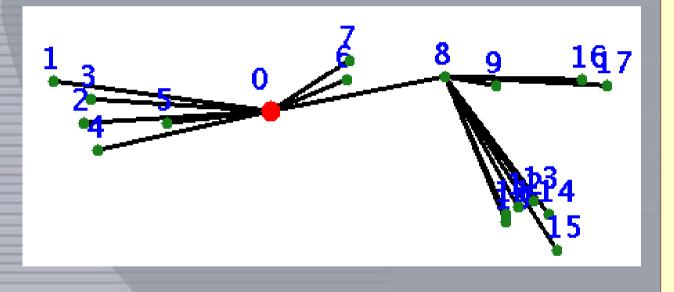
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Evaluation: Bhimavaram topology



- Ashwini Project: Byrraju
 foundation, West
 Godavari, Andhra Pradesh
- To connect 34 villages (result only for 31 nodes)
- Uses ONE wireless channel compared to THREE by current deployment.
- Careful topology planning led to 21% cost savings.
- Careful height assignment led to 15 times cost benefit over current deployment (undergoing tests).

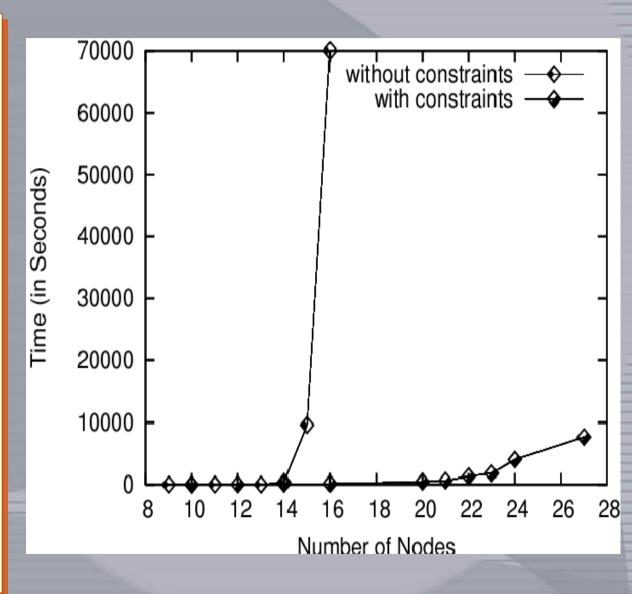
Evaluation: Amalapuram topology

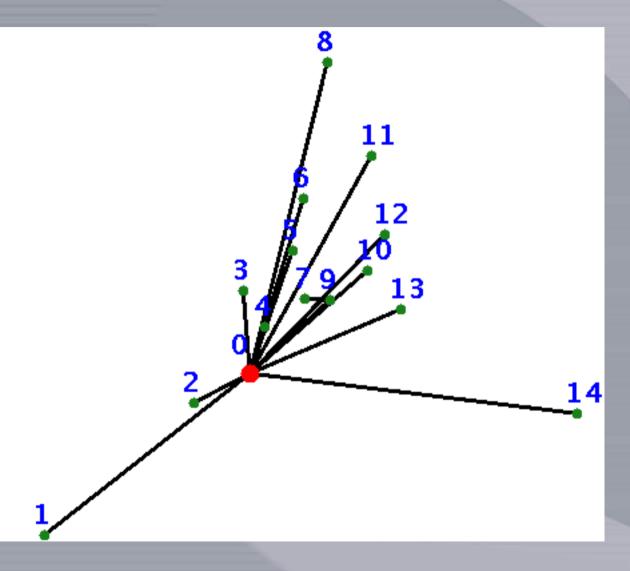


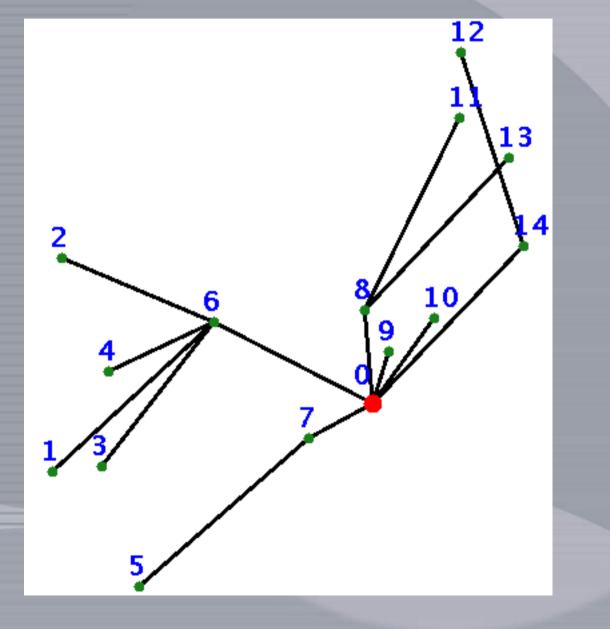
- Ashwini Project: Byrraju foundation, East Godavari, Andhra Pradesh
- To connect 18 villages
- Uses ONE wireless channel.

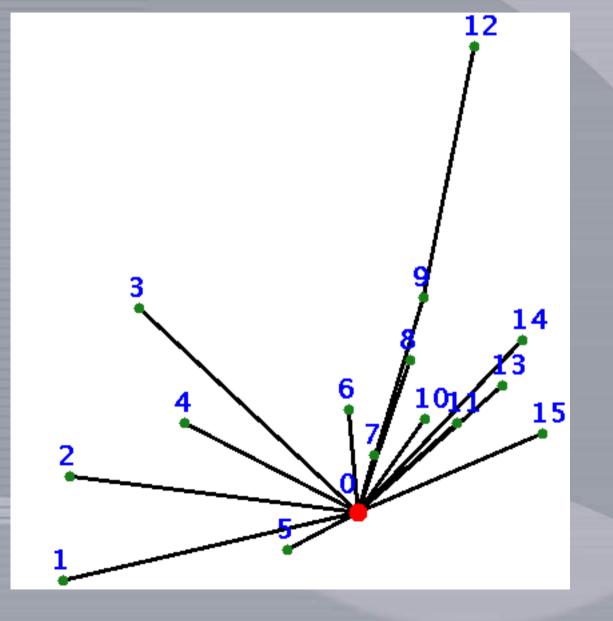
Evaluation

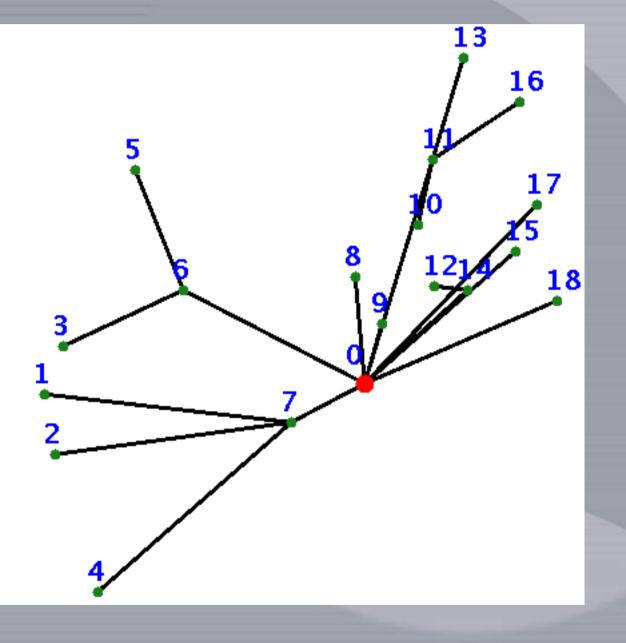
- Runtime Plot
- Observations
 - Antennas of max half power beamwidth 30 degree.
 - The two ends of a link are assigned same power values
 - Linearity of tower cost holds

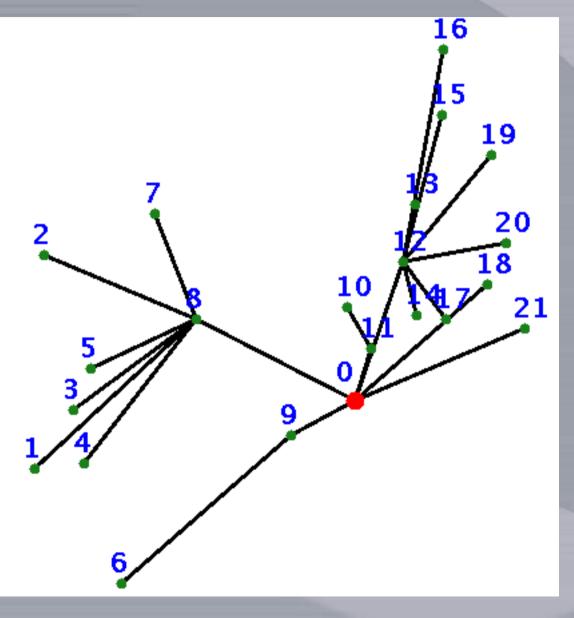


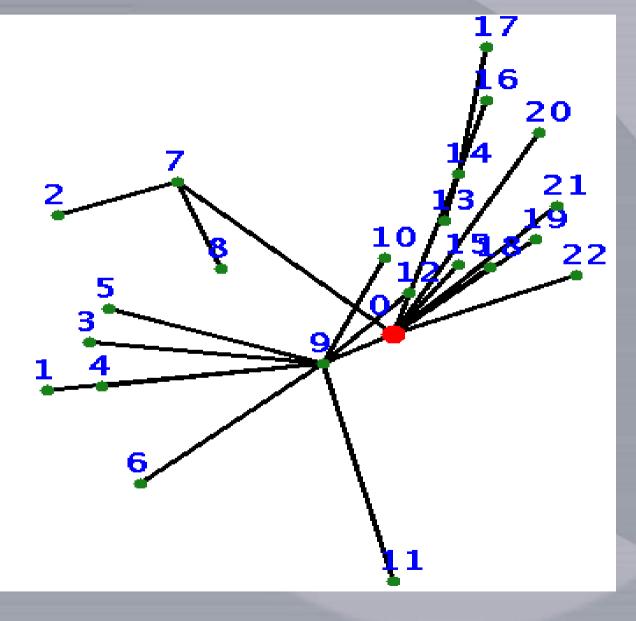












Conclusions

- Topology construction an important problem
- Unique problem thus far
- Challenging to formulate
- Our contributions:
 - Problem formulation
 - Overall approach
 - First-cut solution
- Lots of scope for further in-depth work

