

# CS 744: BIG DATA SYSTEMS

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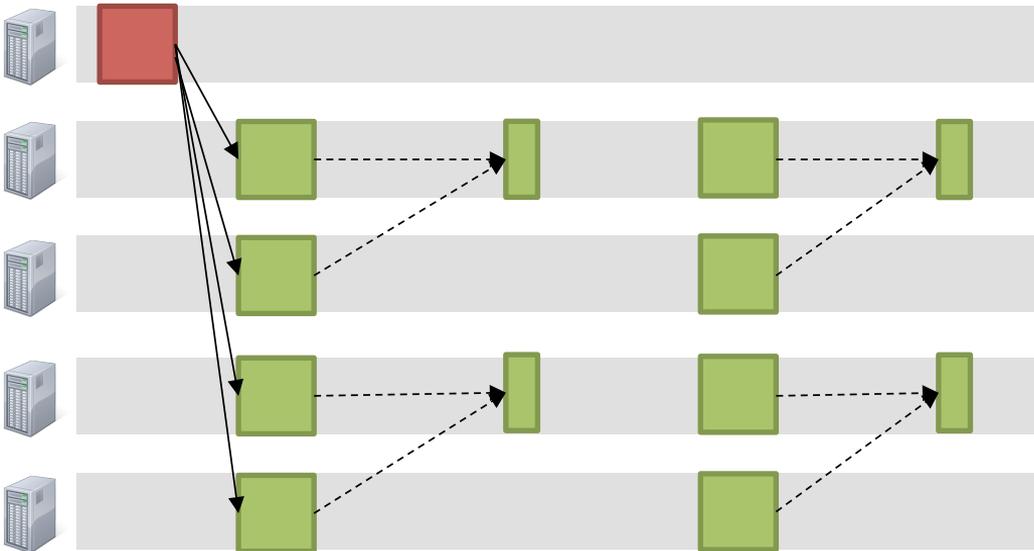
Fall 2018

# ADMINISTRIVIA

- Assignment 2, Midterm grades this week
- Course Projects: round 2 meetings next Friday
- Next Tuesday: Guest speaker for first part

**WHAT WE KNOW SO FAR**

# CONTINUOUS OPERATOR MODEL



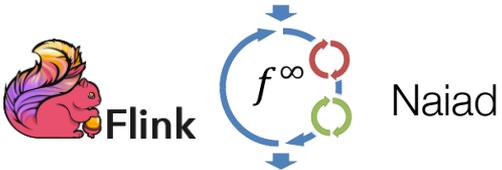
 Driver       Control Message  
 Task       Network Transfer

Long-lived operators

Mutable State

Distributed Checkpoints  
High overhead for Fault Recover

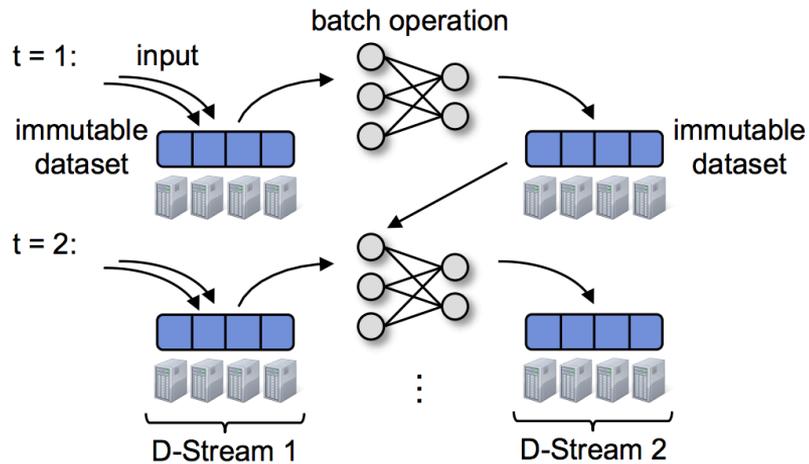
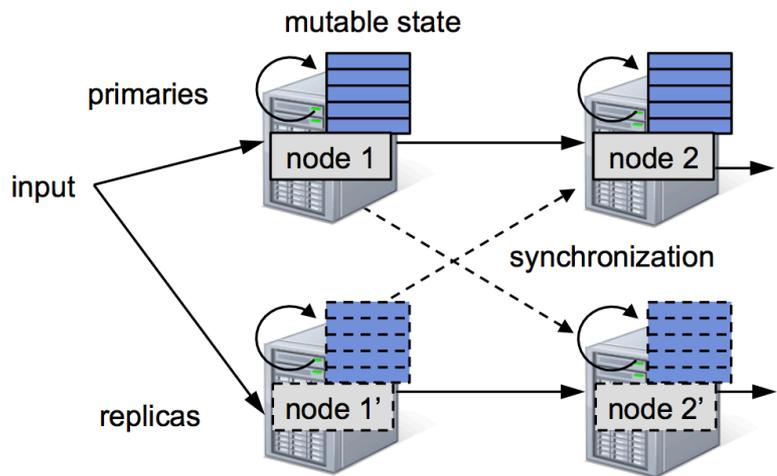
Stragglers ?



# GOALS

1. Scalability to hundreds of nodes
2. Minimal cost beyond base processing (no replication)
3. **Second-scale** latency
4. **Second-scale recovery** from faults and stragglers

# DISCRETIZED STREAMS



# DISCRETIZED STREAMS (DSTREAMS)

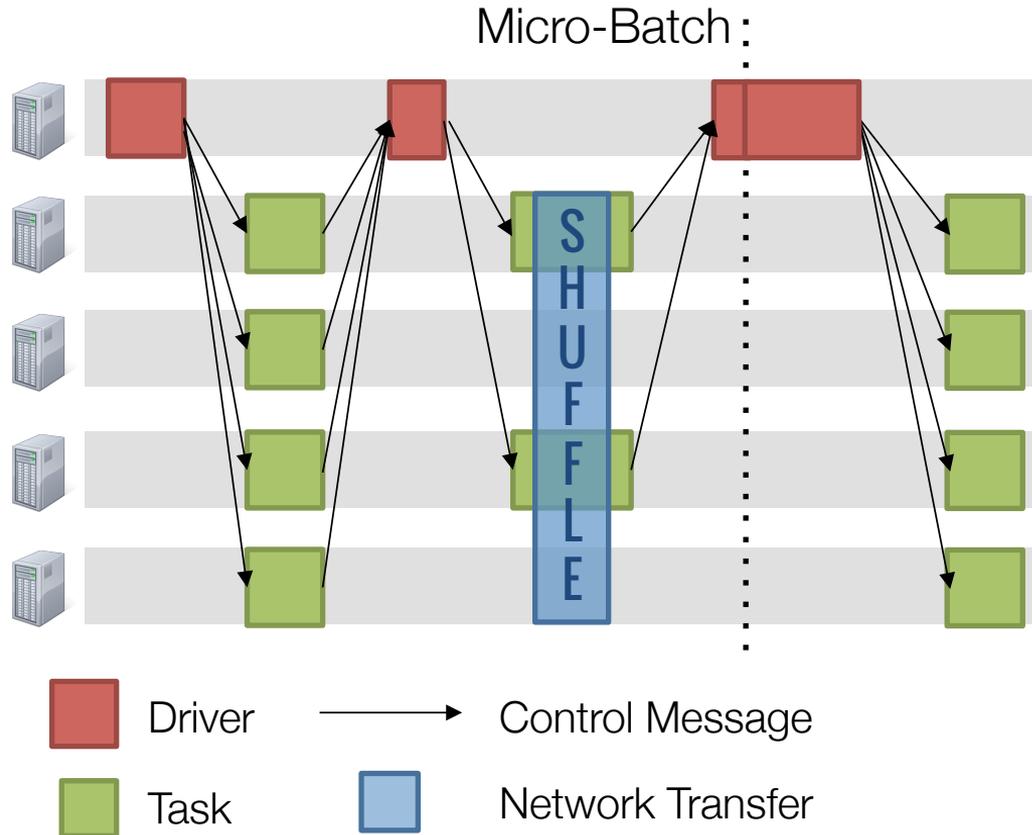
## Approach

- Use *short, stateless, deterministic tasks*
- Store **state** across tasks as in-memory RDDs
- Fine-grained tasks → Parallel recovery / speculation

## Model

- Chunk inputs into a number of **micro-batches**
- Processed via parallel operations (i.e., map, reduce, groupBy etc.)
- Save intermediate state as RDD / write output to external systems

# COMPUTATION MODEL: MICRO-BATCHES

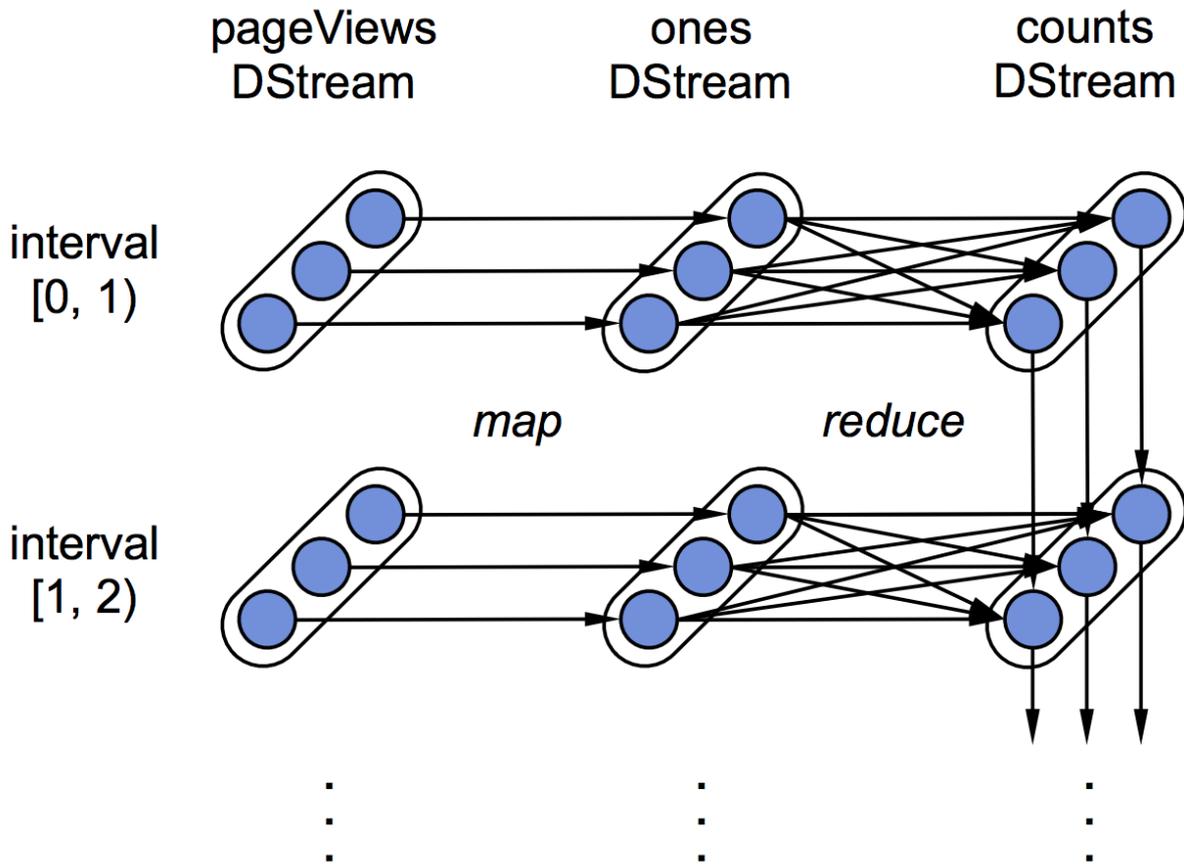


# EXAMPLE

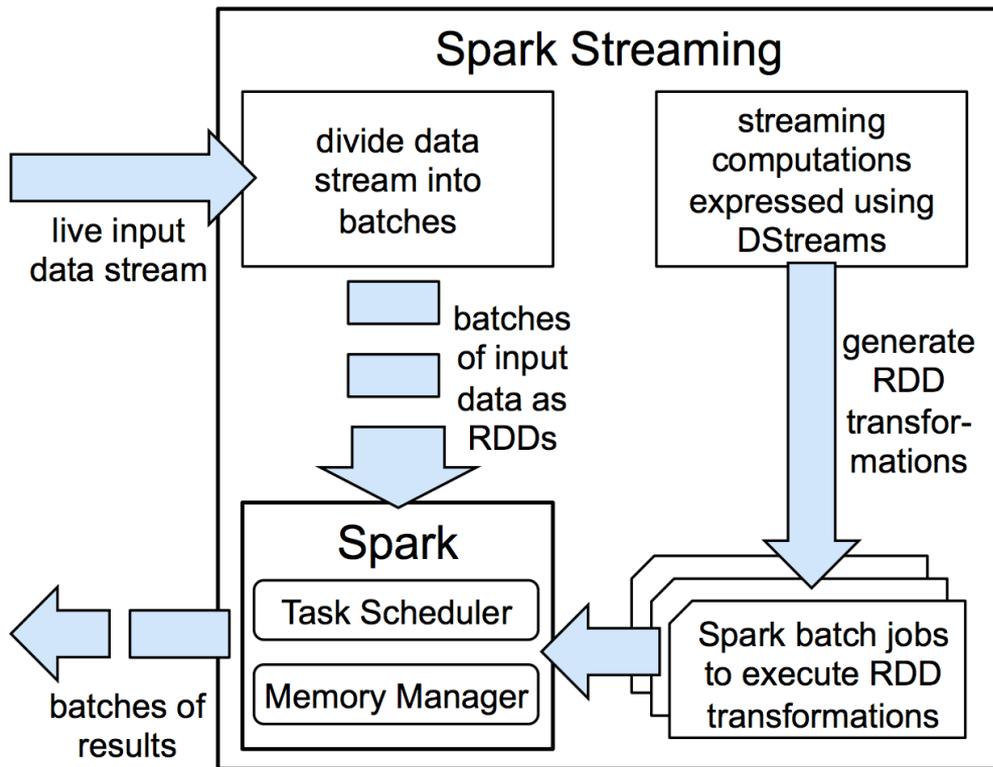
```
pageViews =  
  readStream(http://...,  
            "1s")
```

```
ones = pageViews.map(  
  event =>(event.url, 1))
```

```
counts =  
  ones.runningReduce(  
    (a, b) => a + b)
```



# ARCHITECTURE



# DSTREAM API

## Output operations

save output to external database / filesystem

## Transformations

Stateless: map, reduce, groupBy, join

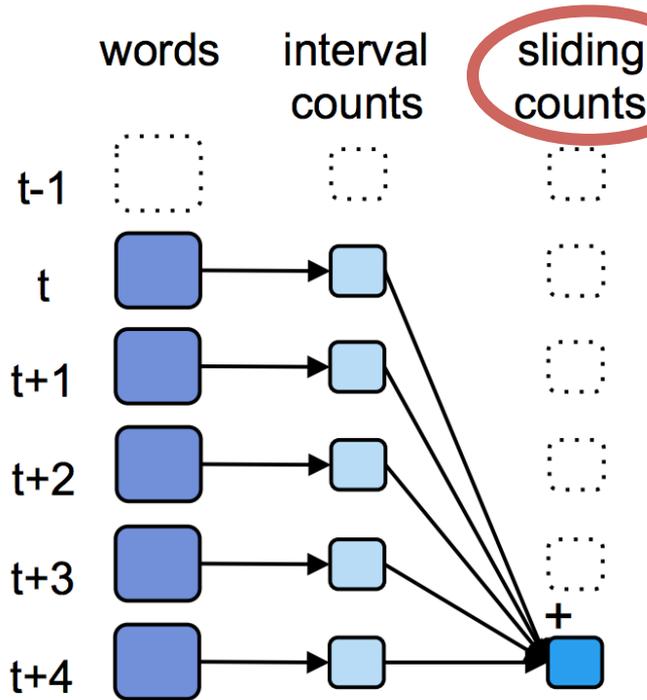
Stateful:

window("5s") → RDDs with data in [0,5), [1,6), [2,7)

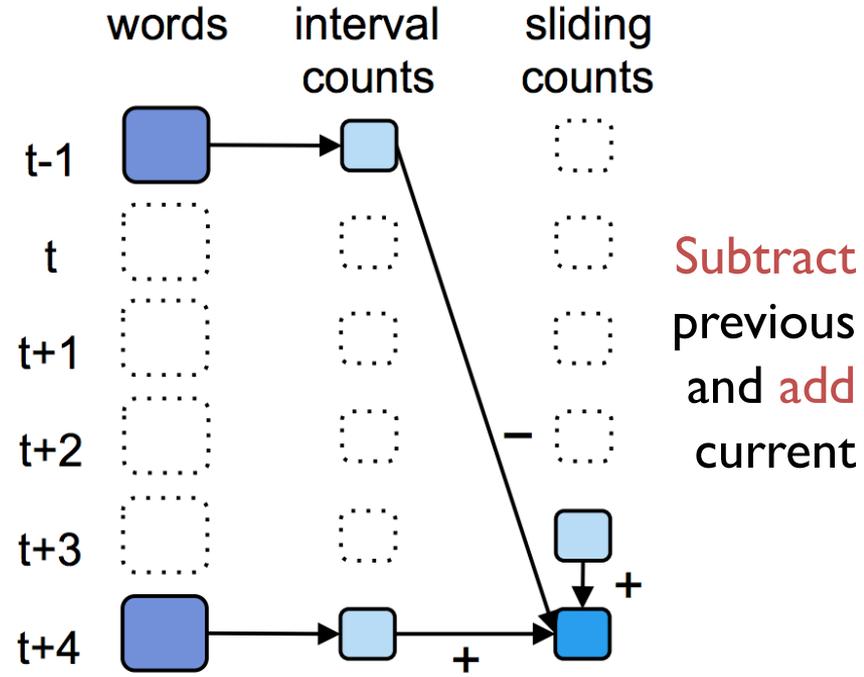
reduceByWindow("5s", (a, b) => a + b) → **incremental** aggregation

# ASSOCIATIVE, INVERTIBLE

Add  
previous 5  
each time



(a) Associative only



Subtract  
previous  
and add  
current

(b) Associative & invertible

# OTHER ASPECTS

**Tracking State:** streams of (Key, Event) → (Key, State)

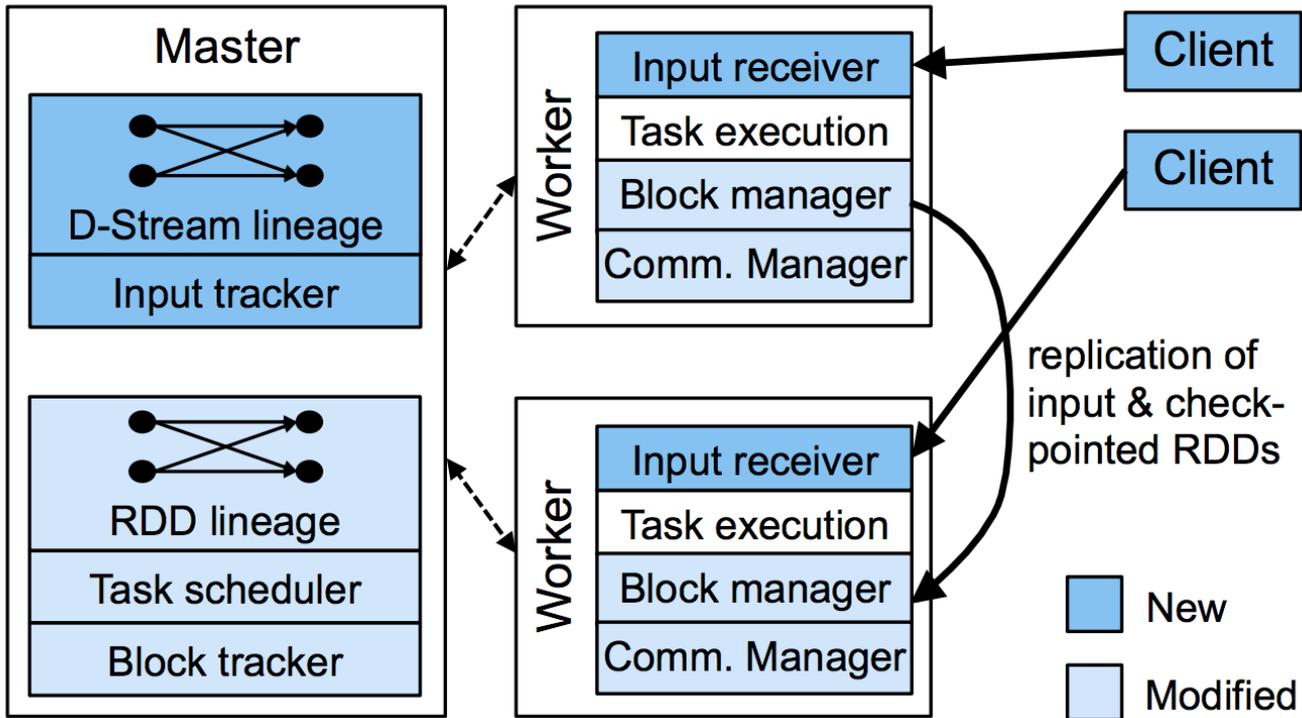
- Initialize: Create a State from the first event
- Update: Return new State given, old state and event
- Timeout for dropping old states.

```
events.track(  
  (key, ev) => 1,  
  (key, st, ev) =>  
    ev == Exit ?  
      null : 1,  
  "30s")
```

Unifying batch and stream

- Join DStream with static RDD
- Attach console and query existing RDDs
- Shared codebase, functions etc.

# SYSTEM IMPLEMENTATION



# OPTIMIZATIONS

## Network Communication

- Rewrote Spark's data plane to use asynchronous I/O

## Timestep Pipelining

- No barrier across timesteps unless needed

- Tasks from the next timestep scheduled before current finishes

## Checkpointing

- Async I/O, as RDDs are immutable

- Forget lineage after checkpoint

# FAULT TOLERANCE: PARALLEL RECOVERY

## Worker failure

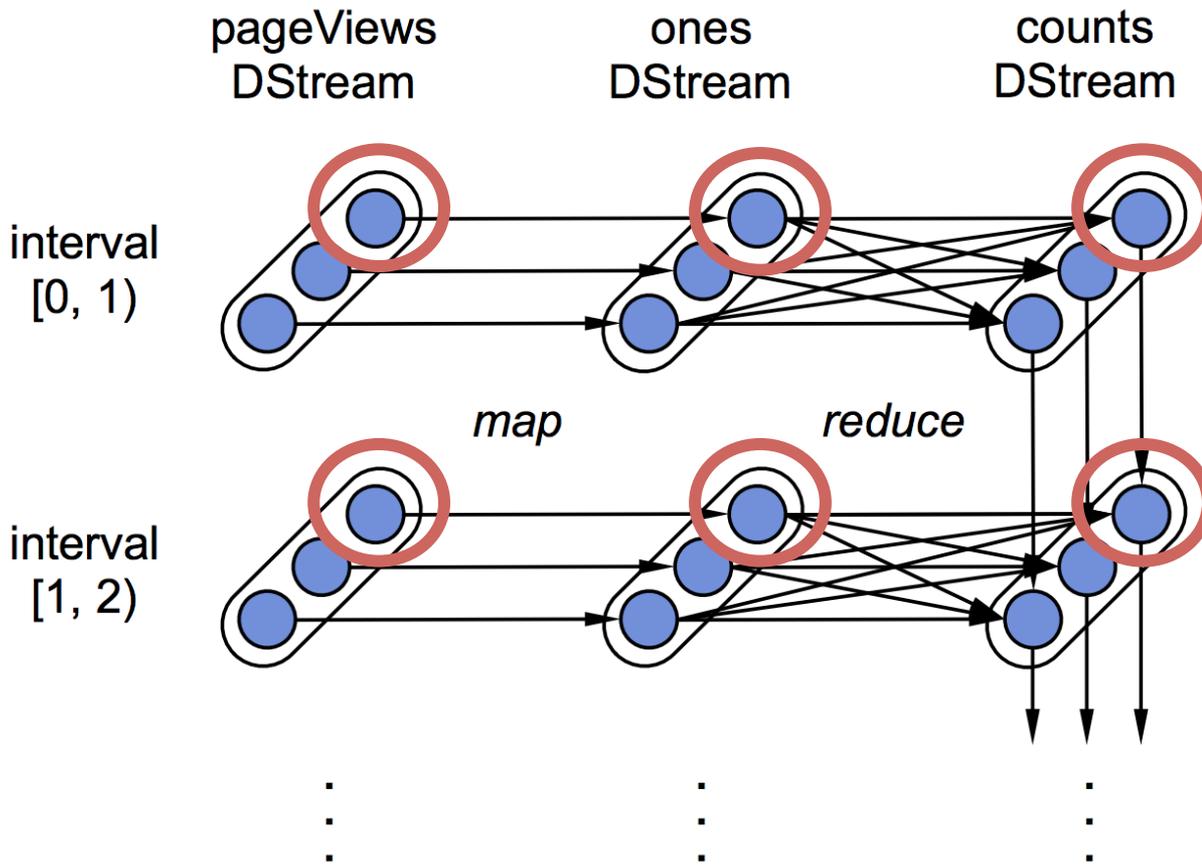
- Need to recompute state RDDs stored on worker
- Re-execute tasks running on the worker

## Strategy

- Run all independent recovery tasks in **parallel**
- Parallelism from partitions *in timestep* and *across timesteps*

# EXAMPLE

```
pageViews =  
  readStream(http://...,  
            "1s")  
  
ones = pageViews.map(  
  event =>(event.url, 1))  
  
counts =  
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```



# FAULT TOLERANCE

## Straggler Mitigation

Use speculative execution

Task runs more than 1.4x longer than median task → straggler

## Master Recovery

- At each timestep, write out **graph of DStreams** and Scala function objects
- **Workers connect** to a new master and report their RDD partitions
- Note: No problem if a given RDD is computed twice (determinism).

# DISCUSSION/SHORTCOMINGS

## Expressiveness

- Current API requires users to “think” in micro-batches

## Setting batch interval

- **Manual tuning**. Higher batch → better throughput but worse latency

## Memory usage

- LRU **cache** stores state RDDs in memory

# SUMMARY

Micro-batches: New approach to stream processing

Higher latency for fault tolerance, straggler mitigation

Unifying batch, streaming analytics