# 802.11, Power Management, and Energy implications on smartphones

Shravan Rayanchu

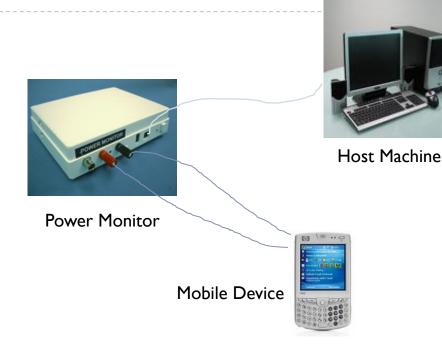
# Smartphones

- Data connection options:
  - 3G / WiFi
- Which one should we use?
  - ▶ 3G speeds: Advertized 2 14 Mbps (shared)
    - <u>"Downlink</u> Between 700 Kbps and 1.7 Mbps, <u>Uplink</u> Between 500 Kbps and 1.2 Mbps" (AT&T Media Newsroom)
  - WiFi speeds
    - IIa/g : up to 54 Mbps, IIn: up to 600 Mbps

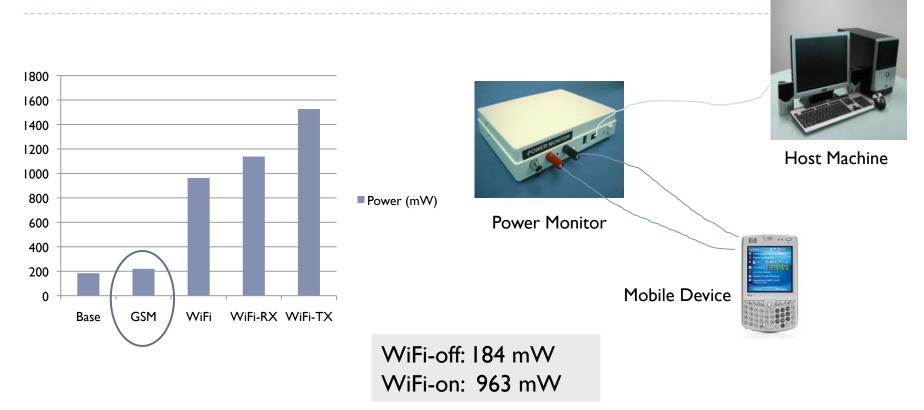


# How much energy does WiFi consume?

- Profiling Experiments
- CPU: 104 MHz
- TX/RX : I Mbps PHY
- HP iPAQ 6965



# How much energy does WiFi consume?



Turning on WiFi interface increases power consumption by 5 times!

Ex: HTC Tilt 8900 (base: 155-475 mW, w/WiFi: 1120 mW)

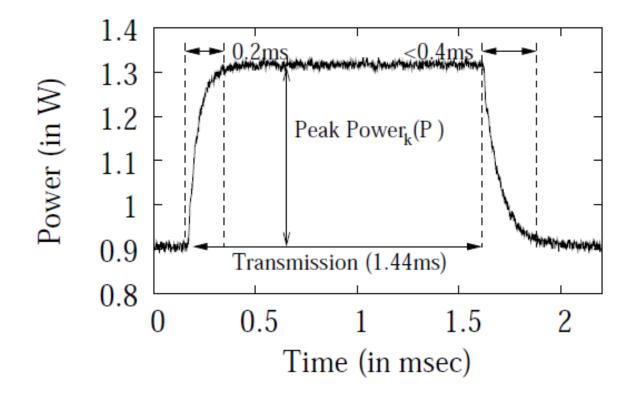
# Designing a Wireless Network for Energy Efficiency

#### Association

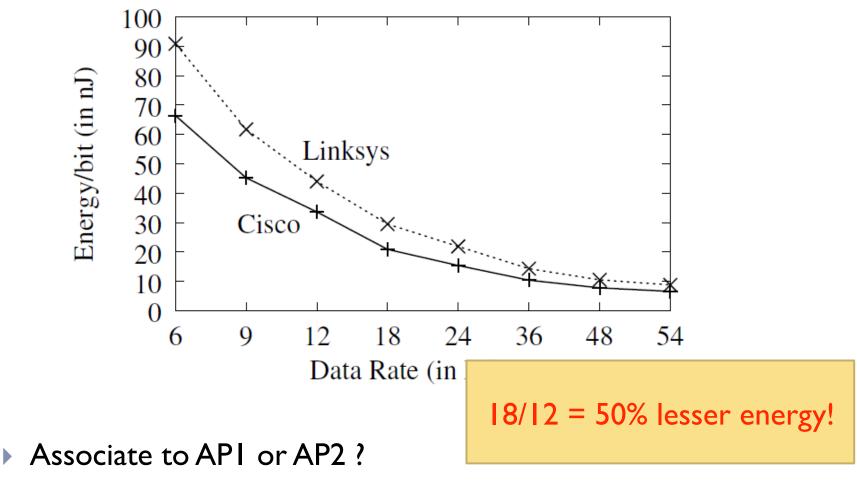
- Currently clients pick an AP with best signal strength
- More control to AP:
  - $\hfill\square$  Load on AP / Channel free time / No. of clients / data rate
- How do wireless parameters effect the energy consumption?
  - Channel Load
  - Data rate
  - Losses
  - Channel width

### Cost of a packet transmission

Data : 1064 byte packet, Cisco wireless card



### Impact of Data Rate



API uses a fixed rate of 12 Mbps, AP2 uses 18 Mbps

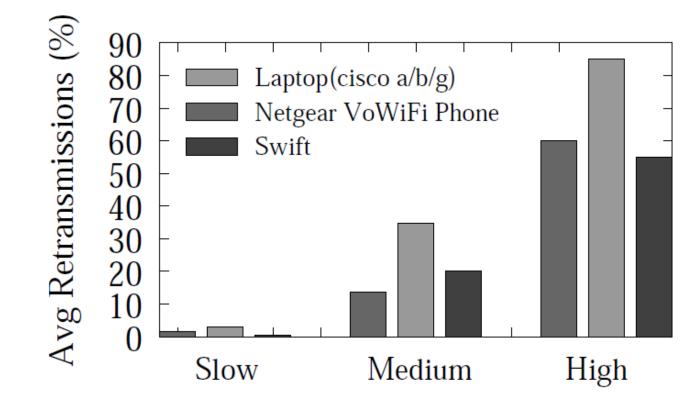
### Impact of Loss rates

- API : 12 Mbps, loss rate = 0%
- AP2 : 18 Mbps, loss rate = 40%
- > TX time for API  $\rightarrow$  p/I2 = 0.08333p
- TX time for AP2  $\rightarrow$  p/18 \* 1/0.6 = 0.092p

10% more energy at 18 Mbps!

# Impact of Mobility

• Imperfect rate adaptation during high mobility  $\rightarrow$  losses



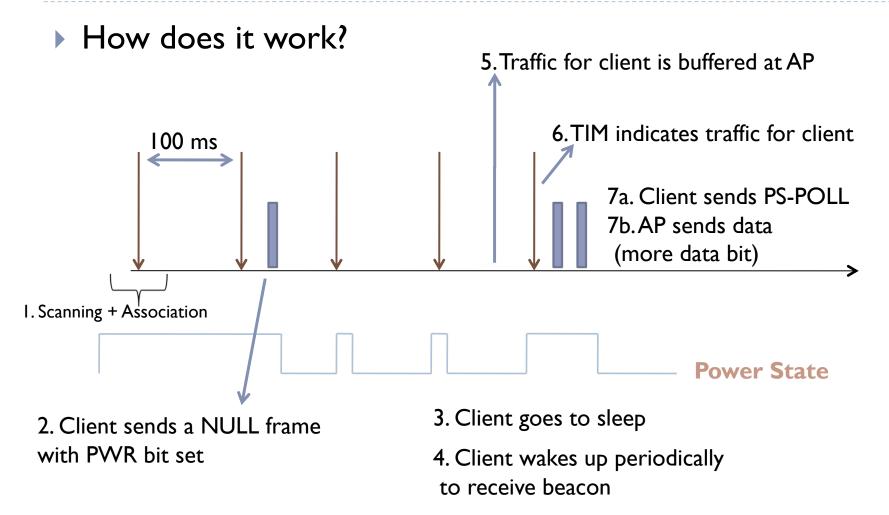
# WiFi Power Management

- What can we do to save power?
  - Switch off the radio?
  - Low power state ~ power consumed is as low as being off

Let us design a power management protocol

- Sleep: Inform the AP
  - Save the state (Association info)
- When should the client wake up?
  - Upload traffic ?
  - Download traffic?
    - $\square$  AP has to inform the client how? (Client wakes up periodically)
    - $\hfill\square$  Client has to inform the AP that it is awake  $\hfill$  .

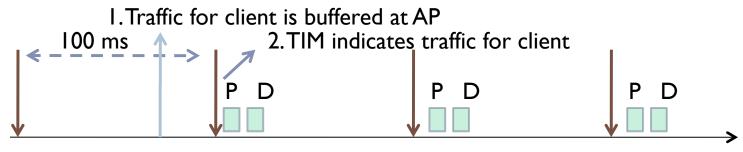
# IEEE 802.11 PSM



# IEEE 802.11 PSM

#### PSM Static

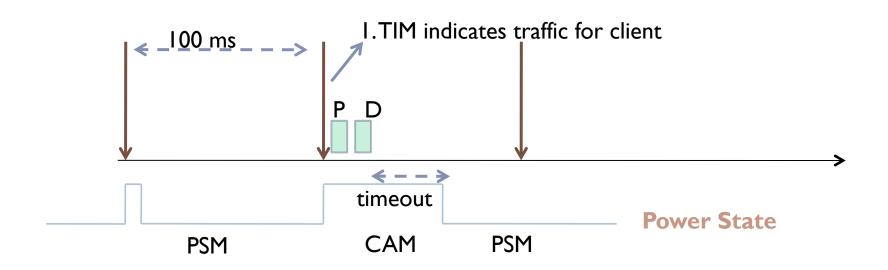
- What are the benefits of PSM?
- Can PSM static be harmful?
  - Increased latency! (Max latency tolerable ?)



One PS-POLL for every data packet!
 Energy = RX (data) + TX( PS-POLL)

#### PSM Adaptive:

- PSM + CAM
- CAM  $\rightarrow$  PSM: time out (how much?)
- ▶ PSM  $\rightarrow$  CAM: TIM bit, threshold packets etc.



Timeout

► Longer → more energy, but lesser latency for subsequent packets

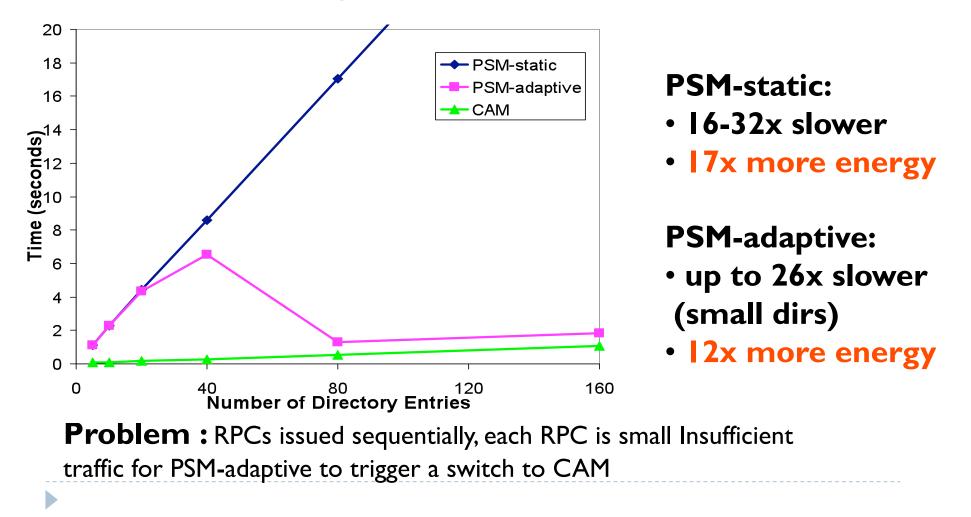
### Research paper

#### Self-Tuning Wireless Network Power Management

Manish Anand et. al, Mobicom'03

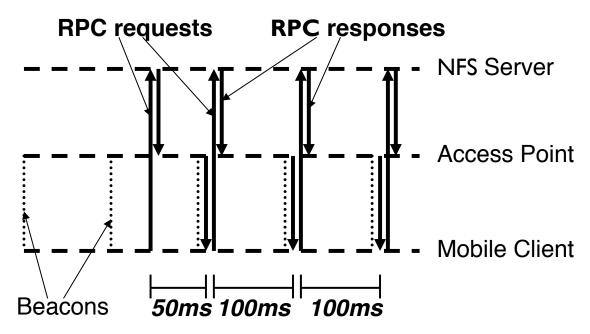
# Effect of Power Management on NFS

Time to list a directory on handheld with Cisco 350 card



# What's Going On?

NFS issues RPCs one at a time .....



- Get file listing
- Issues 2 RPCs (lookup, getattr for each file)
- Each RPC delayed 100ms cumulative delay is large

- Affects apps with **sequential request/response** pairs

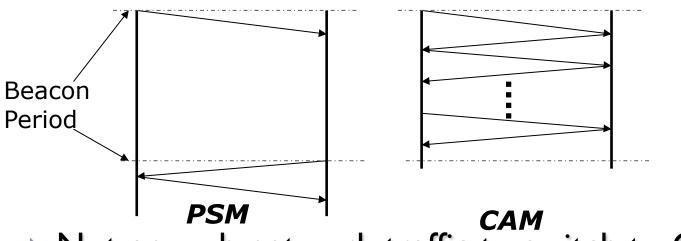
#### Applications:

- Low latency / Foreground
  - CAM is better
- Think time, Background
  - PSM useful
- Untuned power management during interactive episodes dominates the benefits realized during idle periods.

### STPM: Overview

- "One size fits all" approach does not work!
- Self tuning power management
  - Characteristics of network interface, mobile computer
    - Time & energy costs of switching power modes, Base power usage
  - Intent and access patterns of applications
    - Latency, transfer size etc
    - API to disclose hints about interface usage

# Know Application Intent Application: NFS File access



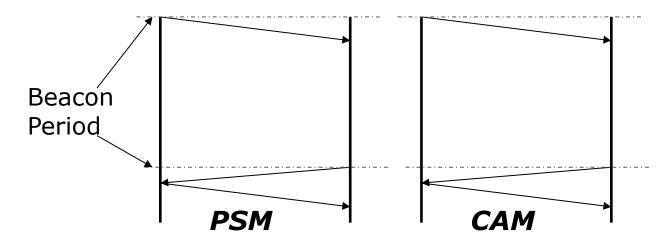
Not enough network traffic to switch to CAM

Data rate is dependent on the power management

Best Policy: Use CAM during activity period

# **Know Application Intent**

• Application: Stock Ticker that is receiving 10 packets per second



Data rate is not dependent on power management

#### STPM allows applications to disclose hints about

- When data transfer are occurring
- How much data will be transferred (optional)
- Max delay on incoming packets
- Best Policy: Use PSM

### Power Modes: to switch or not to switch?

|            | Cisco Aironet 350 |                  | Orinoco Silver   |                  |
|------------|-------------------|------------------|------------------|------------------|
|            | Time              | Energy           | Time             | Energy           |
| PSM to CAM | $0.40\mathrm{s}$  | $0.51~{ m J}$    | $0.23\mathrm{s}$ | $0.24\mathrm{J}$ |
| CAM to PSM | $0.41\mathrm{s}$  | $0.53~{ m J}$    | $0.26\mathrm{s}$ | $0.31\mathrm{J}$ |
| Disable    | $0.00~{ m s}$     | 0.00 J           | N/A              | N/A              |
| Enable     | $0.39\mathrm{s}$  | $0.51\mathrm{J}$ | N/A              | N/A              |

- Transition incurs cost : time, energy
- Measuring transition time
  - Initiation transition, immediately perform a ping
- Short RPC: transition cost > latency reduction achieved in CAM
- Break even size: For transfer greater than this, performance benefit of CAM > transition cost

### Proactive Vs Reactive

Transition cost of changing power mode: 200 -600 ms.

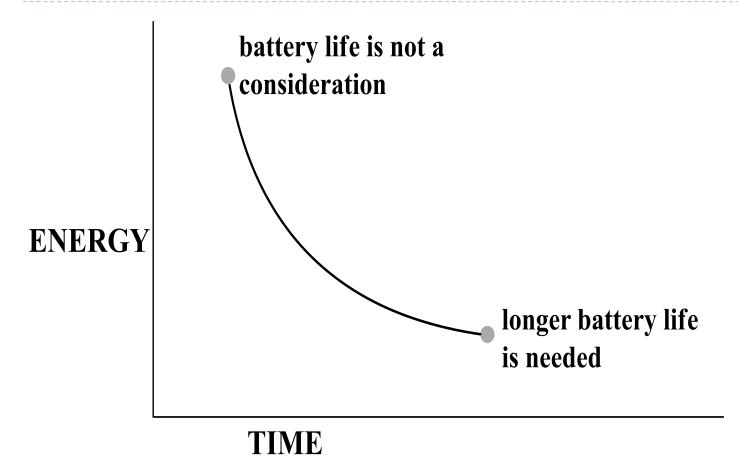
#### Large transfers: use a reactive strategy

- If transfer large enough, should switch to CAM
- Break even point depends on card characteristics (supported data rates, transition cost)
- STPM calculates this dynamically
- Many applications (like NFS) only make short transfers: be proactive
  - Benefit of being in CAM small for each transfer
  - But if many transfers, can amortize transition cost
  - Need foreknowledge: Build empirical distribution of n/w transfers
  - Switches to CAM when it predicts many transfers likely in future

# Respect the Critical Path

- Many applications are latency sensitive
  - NFS file accesses
  - Interactive applications
  - Performance and Energy critical
- Other applications are less sensitive to latency
  - Prefetching, asynchronous write back (Coda DFS)
  - Multimedia applications (with client buffering)
  - Only energy conservation critical
- Applications disclose the nature of transfer: foreground or background

# Embrace Performance/Energy Tradeoff

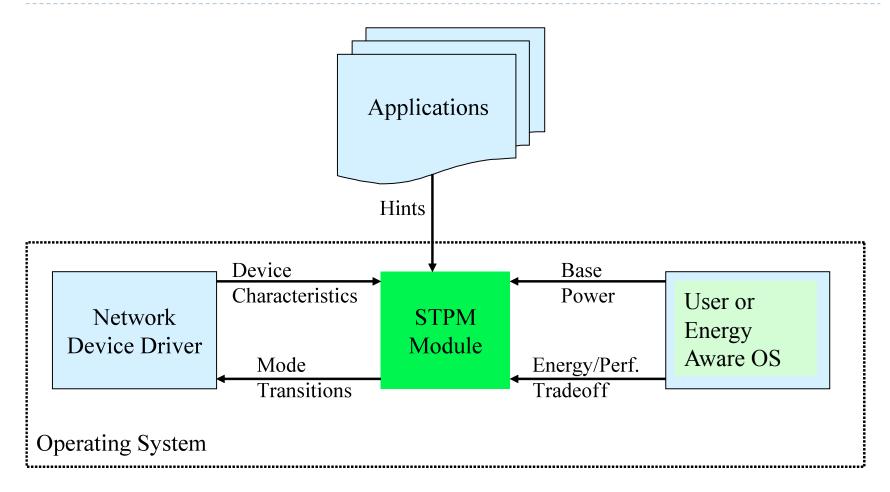


- Trade off: Energy Vs Performance
- STPM lets user specify relative priorities using a tunable knob

# Adapt to the operating environment

- Must consider base power of the mobile computer
- Consider mode that reduces network power from 2W to IW
   Delays interactive application by 10%
- On handheld with base power of 2 Watts:
  - Reduces power 25% (from 4W to 3W)
  - Energy reduced 17.5% (still pretty good)
- On laptop with base power of 15 Watts:
  - Reduces power by only 5.9%
  - Increases energy usage by 3.5%
  - Battery lasts longer, user gets less work done

### STPM Architecture



# STPM: API

### TransferHintBegin:

- Foreground / Background
- Expected transfer size (default: small)

# ListenHintBegin:

Max. delay application is willing to tolerate

# HintEnd

Removes hints when application terminates

# SetKnob,

• 0 to 100 (min, max performance)

# SetBasePower

### Network Power Costs

| Card            | Cisco            | Orinoco          |
|-----------------|------------------|------------------|
| Mode            | Aironet 350      | Silver           |
| $P_{disabled}$  | $0.24\mathrm{W}$ | N/A              |
| $P_{PSM\_idle}$ | $0.39\mathrm{W}$ | $0.19\mathrm{W}$ |
| $P_{PSM\_recv}$ | $1.42\mathrm{W}$ | $2.22\mathrm{W}$ |
| $P_{PSM\_send}$ | $2.48\mathrm{W}$ | $2.70\mathrm{W}$ |
| $P_{CAM\_idle}$ | $1.41\mathrm{W}$ | $1.21\mathrm{W}$ |
| $P_{CAM\_recv}$ | $2.61\mathrm{W}$ | $2.25\mathrm{W}$ |
| $P_{CAM\_send}$ | $3.69\mathrm{W}$ | $2.67\mathrm{W}$ |

Using Digital Multimeter, measure the power drawn by mobile computer

- w/o network card (gives **base** power)
- with card (**Tx**, **Rx**, **Idle**, **disabled**)

# Transition to CAM

STPM switches from PSM to CAM when:

- Case I:Application specifies max delay < beacon period
- Case 2: Disclosed transfer size > break-even size
  Expected time to perform transfer:

$$T = L + B_s / DR_s + B_r / DR_r$$

Transition to CAM

• Expected energy to perform transfer  $E = L \times (P_i + P_b) + B_s / DR_s \times (P_s + P_b) + B_r / DR_r \times (P_r + P_b)$ 

Switch from PSM to CAM if

1.  $T_{CAM} + T_{transition} < T_{PSM}$ 

2.  $E_{CAM}$  +  $E_{transition}$  <  $E_{PSM}$ 

Conflicting Cases: Use knob to resolve

 $C = (T) \times knob + (E) \times (100 - knob)$ 

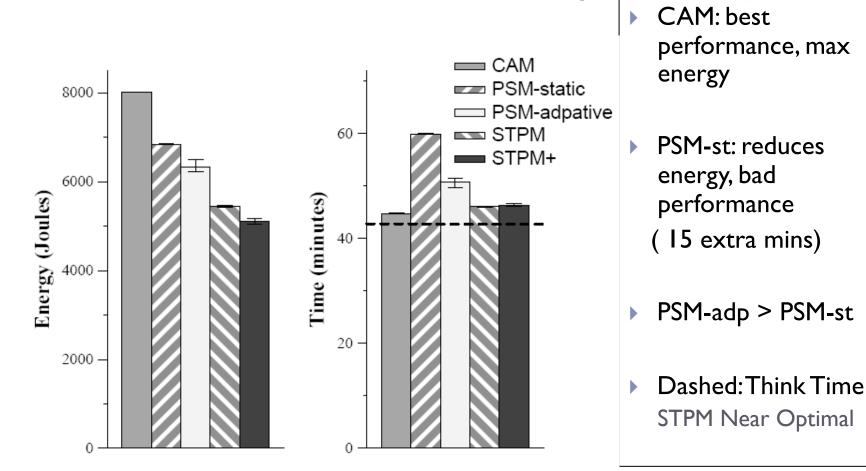
### Transition to CAM

Case 3: Many forthcoming transfers are likely

- **Run**: Transfers closely correlated in time
- **Run Length**: Number of Transfers
- To predict forthcoming transfers STPM generates an empirical distribution of run lengths

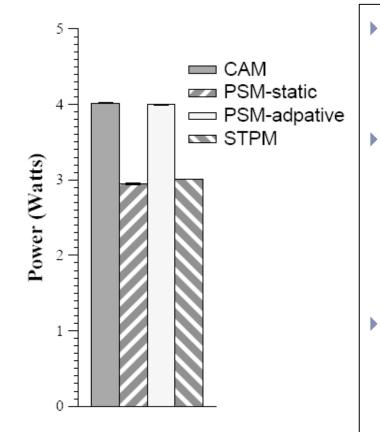
# Results: CODA on handheld

Workload: 45 minute interactive software development activity



STPM: 21% less energy, 80% less time than 802.11b power mgmt.

# Results: Xmms



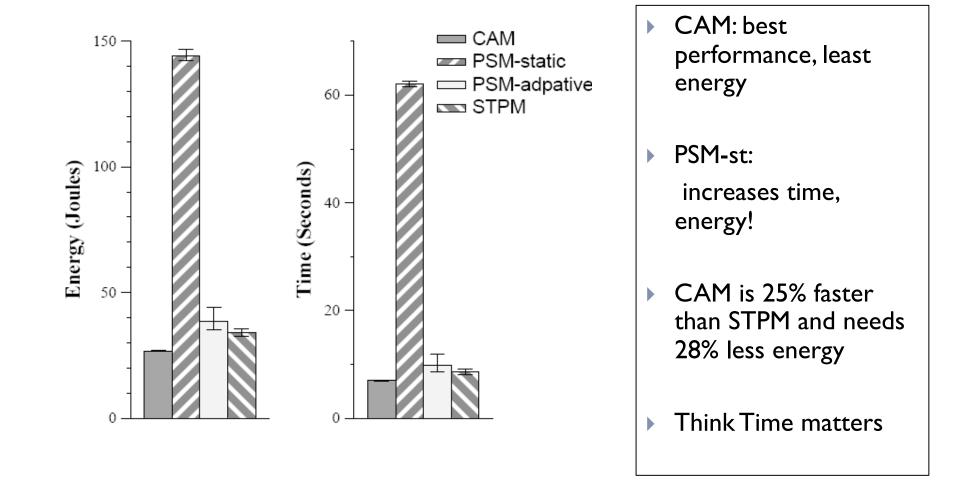
 Client buffering determines if power management can be used

#### STPM

- 25% less power than CAM
- > 2% more energy than PSM-st
- No audio packets dropped
- PSM-adaptive performs badly
  - Assumes latency is critical

#### Workload: 128Kb/s streaming MP3 audio from an Internet server

### Results: remote X



# Research Paper

- NAPman: Network Assisted Power Management for WiFi Devices
  - Eric Rozner et. al, MobiSys 2010
  - Revisiting PSM

Effect of "background" traffic on PSM clients

# Effect of background traffic



#### **Experiment Setup**

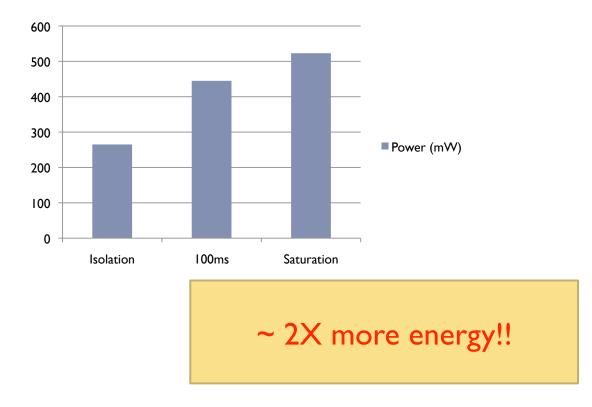
Linksys WRT54GL AP 2 HP iPAQ clients

Auto-rate algorithm at AP Packet size: 1440 bytes PSM-adaptive: 30ms timeout, RX thresh: 10

- I. PSM client (near)
  - II Mbps PHY
  - Periodic traffic: one packet every 100ms
- 2. CAM client (far)
  - I Mbps PHY
  - Traffic :
    - I. Periodic (one packet/100ms)
    - 2. Saturated

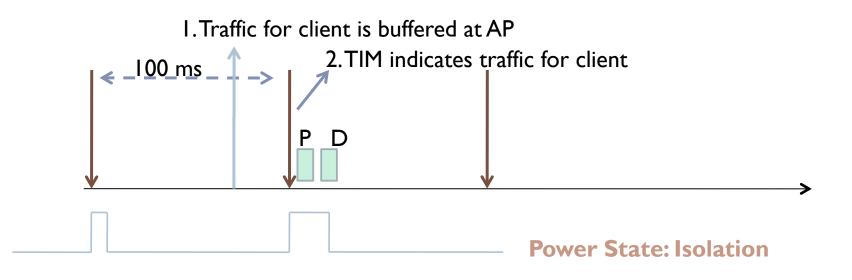
# Effect of background traffic

Total packets: 1000 (~ 100 sec)

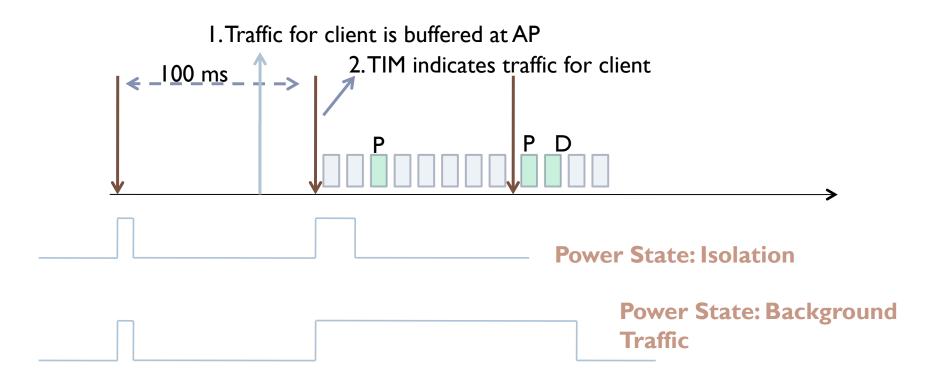


### NAPman

#### Basic PSM problem and solution

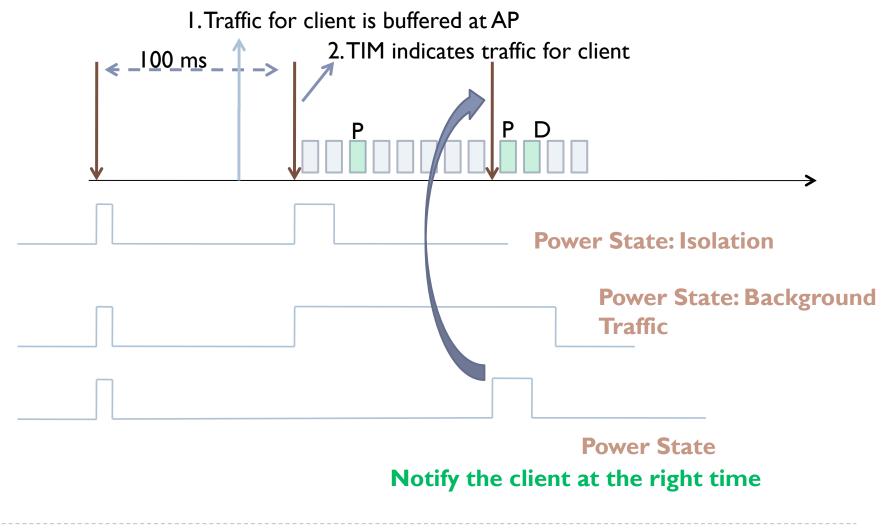


#### NAPman

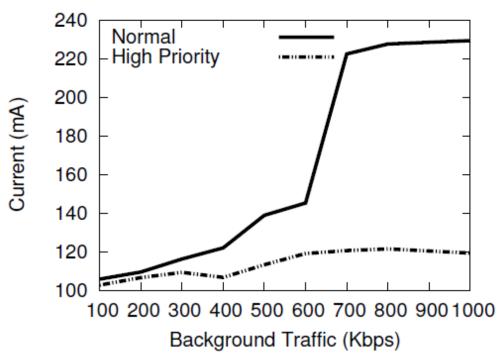




### NAPman

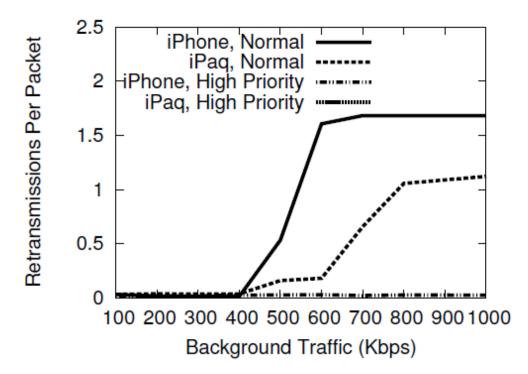


# Priority queuing



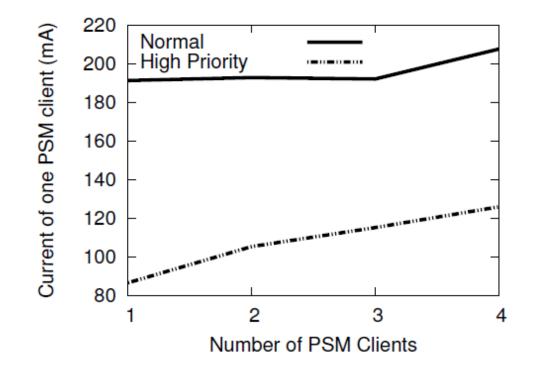
- AP sets the TIM bit
- On receiving PS-POLL, the data packet is scheduled using a higher priority queue
- > PSM client immediately receives the packet and then goes to sleep

## Priority queuing



- Increased retransmissions  $\rightarrow$  reduced capacity!
- Priority queuing can resolve the issue
- Problems (i) Fairness (ii) Multiple PSM clients

# Multiple PSM clients create a problem ..



 Multiple PSM clients can nullify the effect of priority queuing

# Questions?

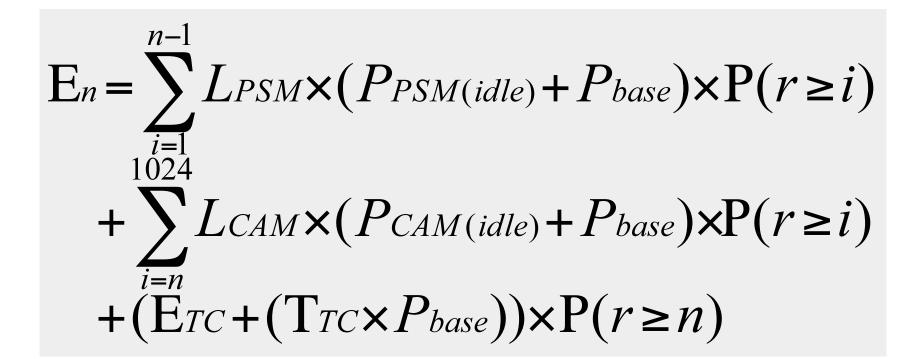
# Transition to CAM

Expected time to perform nth transfer

$$T_{n} = \sum_{i=1}^{n-1} L_{PSM} \times P(r \ge i) + \sum_{i=n}^{1024} L_{CAM} \times P(r \ge i) + T_{TC} \times P(r \ge n)$$

 $P(r \ge n)$  is the probability that *length of run*  $\ge n$ 

#### Transition to CAM



 $C_n = (T_n / T_{mean}) \times knob + (E_n / E_{mean}) \times (100 - knob)$ 

# Intuition: Using the Run-Length History



Switch when expected # of transfers remaining in run is high

# Evaluation

#### Client: iPAQ handheld with Cisco 350 wireless card

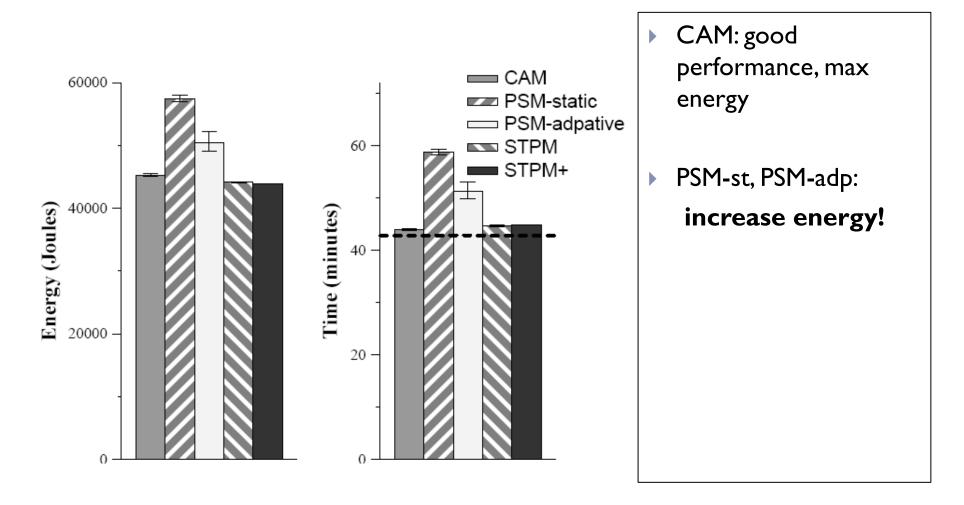
## Evaluate STPM vs. CAM, PSM-static, and PSM-adaptive:

- NFS distributed file system
- Coda distributed file system (file access, asynch writes)
- XMMS streaming audio (streaming data, buffering)
- Remote X (thin client display)

## Run DFSTrace tool generate access stats for STPM

- Use Mummert's file system trace (SOSP '95)
- File system operations (e.g. create, open, close)
- Captures interactive software development

# Results: CODA on laptop



STPM: 2.6% less energy, 15% more time than CAM