

1)

Claim: # of edges in case of n nodes, $f(n) = n - 1$

Proof: Base case: $n = 1$, # of edges $f(1) = 1 - 1 = 0$

Inductive step: Assume # of edges for n nodes $f(n) = n - 1$

For $n + 1$ nodes # of nodes, $f(n + 1) = f(n) + 1 = n - 1 + 1 = n$

2)

Max. height of BST -- $n-1$

Min. height of BST -- $\text{floor}(\log_2 n)$

Max. # of leaves – $\text{floor}((n+1)/2)$

Min. # of leaves – 1

3)

Yes. Insertion to BST always creates a leaf node.

4)

$$2 * (4! + 2 * 3! + 2 * 2!) = 80$$

or

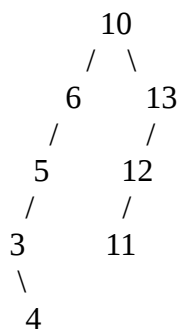
$$[C(6,3) * C(1,1) * C(2,1) * C(1,1)] * [C(3,3) * C(1,1) * C(2,1) * C(1,1)] = 80$$

or

$$1 * C(6,3) * 2 * 2 = 80$$

1; since 4 must be the first insertion. 6 choose 3; to pick which 3 of the following 6 insertions will be the set {1, 2, 3}. 2; since the insertion ordering for subtree {1, 2, 3} can be 213 or 231. 2; since the insertion ordering for subtree {5, 6, 7} can be 657 or 675.

5)



6)

```
public void inorderPrint ( BinaryTreeNode <K > node ) {
// Create a stack to hold nodes
    Stack < BinaryTreeNode <K > > stack = new Stack < BinaryTreeNode <K > >();
    while (! stack . isEmpty () || node != null ) {
        if ( node != null ) {
            stack . push ( node );
            node = node . getLeft () ;
        } else {
            node = stack . pop () ;
            System . out . print ( node . getKey () + " " ); // space as delimiter
            node = node . getRight () ;
        }
    }
    System. out . println () ; // newline after all keys are printed
}
```

EC:

```
public void prune ( BinaryTreeNode <K > node , int depth ) {  
    // Handle base case  
    if ( depth == 0 ) {  
        node . setLeft ( null ) ;  
        node . setRight ( null ) ;  
        return ;  
    }  
    // Recurse into subtrees , if they exist  
    if ( node . getLeft () != null ) {  
        prune ( node . getLeft () , depth -1 ) ;  
    }  
    if ( node . getRight () != null ) {  
        prune ( node . getRight () , depth -1 ) ;  
    }  
}
```