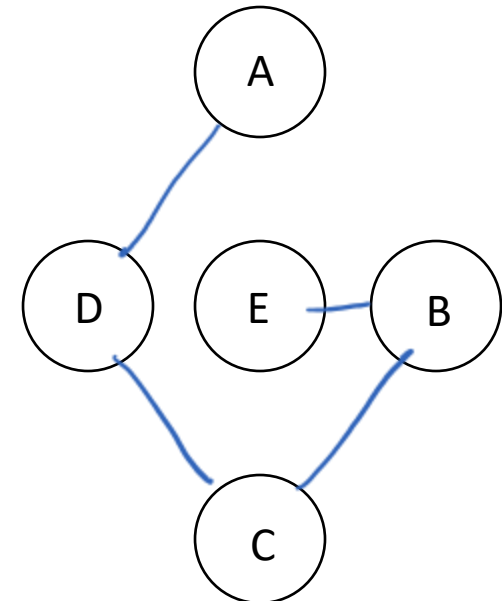
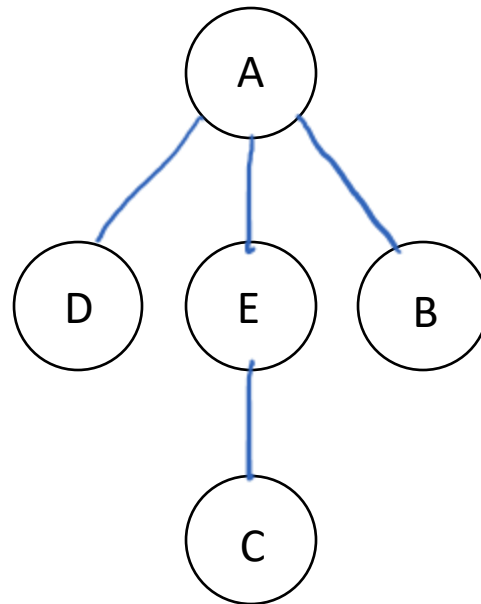
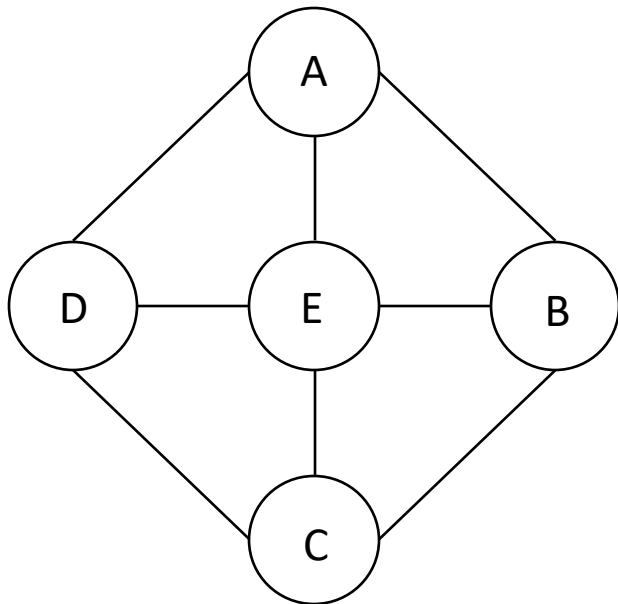


Spanning Trees

Definition

For an undirected graph G with V nodes: If T is a subgraph of G and T contains all nodes of G , is connected, and has exactly $(V-1)$ edges, then T is a spanning tree of G .



Depth First Spanning Trees

DFT(v):

mark v as visited

for each unvisited successor u of v :

add $v-u$ to spanning tree

DFT(u)

Breadth First Spanning Trees

BFT(v):

q = new Queue()

mark v as visited

q.enqueue(v)

while (!q.isEmpty()):

 c = q.dequeue()

 for each unvisited successor u of c:

 add edge c-u to spanning tree

 mark u as visited

 q.enqueue(u)

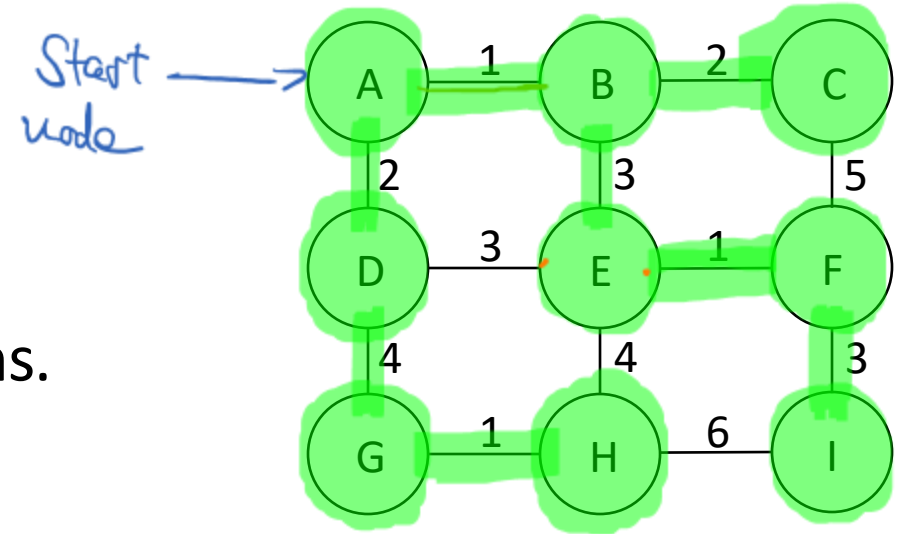
Minimum Spanning Trees

In a weighted graph:

Minimum spanning trees are the spanning trees with the lowest sum of edge weights.

Prim's Algorithm

For weighted, connected, and undirected graphs.



Start node

```

prim(v):
    pq = new PriorityQueue()
    mark v as visited
    pq.insert(v.outgoingEdges)
    while (!pq.isEmpty()):
        c = pq.removeMin()
        if (c.endNode is unvisited):
            mark c.endNode as visited
            add c to tree
            pq.insert(c.endNode.outgoingEdges to unvisited nodes)
    
```

$E \cdot \log E$

E iterations

$\log E$

$\log E$

p.q.:

~~AB:1, AD:2, BC:2, BE:3, DE:3,~~
~~DG:4, CF:5, EF:1, EH:4, FI:3,~~
~~IH:6, GH:1~~

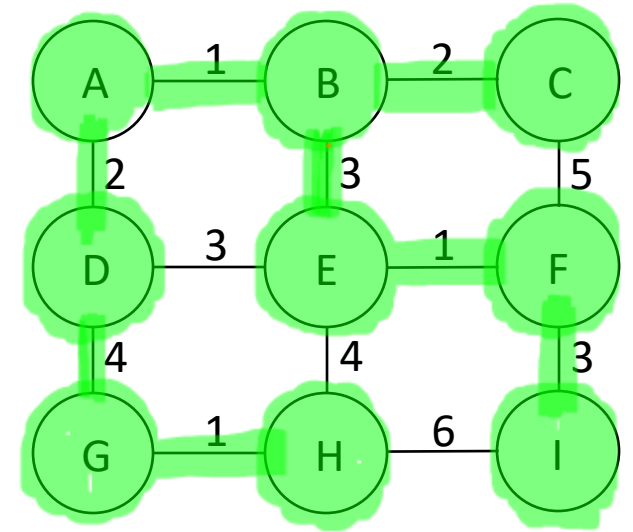
$\Sigma = 17$

edge weights

$$\Rightarrow E \cdot \log E + E \cdot 2 \cdot \log E = 3 \cdot E \cdot \log E$$

Kruskal's Algorithm

For weighted, connected, and undirected graphs.



```
kruskal(edgeList):
```

sort edgeList

init nodeSets with singleton sets

```
while (!edgeList.isEmpty()):
```

```
c = edgeList.removeFirst()
```

```
if (c.startNode and c.endNode in different nodeSets):
```

add c to tree

```
nodeSets.join(c.startNode, c.endNode)
```

edgeList

~~AB:1, EF:1, GH:1, BC:2, AD:2,~~
~~BE:3, DE:3, FI:3, DG:4, EH:4,~~
~~CF:5, HI:6~~

$$\sum_{\text{abgewandt}} = 17$$

$$\text{---} \hookrightarrow \log V \Rightarrow E \cdot \log E + V + E \cdot 2 \cdot \log V$$

Complexities

V: # nodes in graph

E: # edges in graph

Prim: $O(\cancel{3} \cdot E \cdot \log E)$

$$O(E \cdot \log V^2)$$

$$O(E \cdot \cancel{2} \cdot \log V)$$

Kruskal: $O(E \cdot \log E + \cancel{V} + 2 \cdot E \cdot \log V)$

$$O(E \cdot \log V^2 + 2 \cdot E \cdot \log V)$$

$$O(\cancel{4} \cdot E \cdot \log V)$$