



#### Outline

- Classic Virtualization
- Software Virtualization
- Intel/AMD Hardware Virtualization
- Comparison and Results
- Discussion

## **Classic Virtualization**

- Popek and Goldberg's Criteria:
  - 1. Fidelity run any software
- 2. **Performance** run it fairly fast
- 3. Safety VMM manages all hardware
- Trap-and-Emulate only real solution until recently



![](_page_1_Figure_2.jpeg)

![](_page_1_Figure_3.jpeg)

![](_page_1_Figure_4.jpeg)

# Can x86 Trap and Emulate?

#### • No

- Even with 4 execution modes!
- Key problem: dual-purpose instructions don't trap
- Classic Example: popf instruction
  - Same instruction behaves differently depending on execution mode
  - User Mode: changes ALU flags
  - Kernel Mode: changes ALU and system flags
  - Does not generate a trap in user mode

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![](_page_2_Figure_16.jpeg)

# **VMWare's Binary Translation**On-the-fly Only need to translate OS code Makes SPEC run fast by default Most instruction sequences don't change Instructions that do change: Indirect control flow: call/ret, jmp PC-relative addressing Privileged instructions Adaptive Translation

– "Innocent until proven guilty"

# Performance Advantages of BT

- Translation sequences can be faster than native:
  - cli vs. vpu.flags.IF := 0
- Avoid privilege instruction traps
  - Example: rdtsc
    - Trap-and-emulate: 2030 cycles
    - Callout-and-emulate: 1254 cycles
    - BT emulation: 216 cycles (but TSC value is stale)

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#### AMD SVM and Intel VT

- Extensions to x86-32 and x86-64
  - Allows classic trap-and-emulate!
  - Hardware VM modes to reduce traps
  - Details:
    - VMCB virtual machine control block
    - VMX mode for running guest OSs
    - Vmrun instruction to enter VMX mode
    - Many instructions and events cause VMX exits
    - Control fields in VMCB can change VMX exit behavior

## Hardware VM Example: syscall

- 1. VMM fills in VMCB exception table for Guest OS
- Sets bit in VMCB not exit on syscall exception
- 2. VMM executes vmrun
- 3. Application invokes syscall
- CPU → CPL #0, *does not trap*, vectors to VMCB exception table

![](_page_4_Figure_1.jpeg)

![](_page_4_Figure_2.jpeg)

![](_page_4_Figure_3.jpeg)

![](_page_4_Figure_4.jpeg)

![](_page_5_Figure_1.jpeg)

# **Opportunities**

- Faster Microarchitecture implementations
- Intel Core Duo already much faster than P4
- Hardware VMM algorithms
- Software/Hardware Hybrid VMM
- Hardware MMU
  - Virtualize DMA

# • Is BT really faster for things that matter? - Process-based Apache on Linux?

- Who configures a system to constantly page?
- VMWare is done, why bother with Hardware VM support?
  - Simplicity of VMM w/ Hardware support
  - New applications
- Will next-gen hardware make binary translation unnecessary?