

# CS 354 - Machine Organization & Programming

## Tuesday Mar 19 and Thursday Mar 21, 2024

Midterm Exam - Thurs April 4, 7:30 - 9:30 pm

- ◆ UW ID and #2 required
- ◆ closed book, no notes, no electronic devices (e.g., calculators, phones, watches)  
see "Midterm Exam 2" on course site Assignments for topics
- ◆ Exam room information will be sent via email by Friday

A09 GF18 and p4B Worksheet\_completed.pdf

Homework hw4: DUE on or before Monday, 3/18

Homework hw5: will be DUE on or before Monday, April 8

Project p4A: DUE on or before Thursday, Mar 21

Project p4B: DUE on or before Friday, Apr 5

### Learning Objectives

- ◆ explain low-level details of program execution
- ◆ identify and describe assembly language data formats
- ◆ identify IA-32 registers, by name, size, and common usage
- ◆ identify size and type of operand by name and syntax
- ◆ interpret basic assembly language instructions: mov, push, pop, leal, arithmetic
- ◆ interpret basic assembly language control instructions: cmp, test, set, jmp, br
- ◆ interpret and trace sequence of assembly code instructions
- ◆ interpret and explain memory addressing modes by name and syntax
- ◆ able to encode target for control instructions

### This Week

Memory Mountain (from L08) C, Assembly, & Machine Code - L16-10 Low-level View of Data Registers Operand Specifiers & Practice L18-7 Instructions - MOV, PUSH, POP	Instruction - LEAL Instructions - Arithmetic and Shift Instructions - CMP and TEST, Condition Codes Instructions - SET & Jumps Encoding Targets & Converting Loops
<p><b>Next Week:</b> Stack Frames and Exam 2 B&amp;O 3.7 Intro - 3.7.5, 3.8 Array Allocation and Access 3.9 Heterogeneous Data Structures</p>	

# C, Assembly, & Machine Code

C Function	Assembly (AT&T)	Machine (hex)
int accum = 0; int sum(int x, int y) { int t = x + y; accum += t; return t; }	sum: pushl %ebp movl %esp, %ebp movl 12(%ebp), %eax addl 8(%ebp), %eax addl %eax, accum popl %ebp ret	55 89 e5 8b 45 0C 03 45 08 01 05 ?? ?? ?? ?? 5D C3

## C

- ◆
- ◆
- ◆

→ What aspects of the machine does C hide from us?

## Assembly (ASM)

- ◆
- ◆

→ What ISA (Instruction Set Architecture) are we studying?

→ What does assembly remove from C source?

→ Why Learn Assembly?

- 1.
- 2.
- 3.

## Machine Code (MC) is

- ◆
- ◆

→ How many bytes long is an IA-32 instructions?

## Low-Level View of Data

### C's View

- ◆
- ◆

### Machine's View

\* *Memory contains bits that do not*

→ How does a machine know what it's getting from memory?

- 1.
- 2.

### Assembly Data Formats

C	IA-32	Assembly Suffix	Size in bytes
char	byte		
short	word		
int	double word		
long int	double word		
char*	double word		
float	single precision		
double	double prec		
long double	extended prec		

\* *In IA-32 a word*

# Registers

What? Registers

## General Registers

bit 31

16 15

8 7

0

%eax	%ax	%ah	%al
%ecx	%cx	%ch	%cl
%edx	%dx	%dh	%dl
%ebx	%bx	%bh	%bl
%esi	%si		
%edi	%di		
%esp	%sp		
%ebp	%bp		

Program Counter      %eip

## Condition Code Registers

# Operand Specifiers

**What?** Operand specifiers are

- ◆ S
- ◆ D

**Why?**

**How?**

1. )	specifies an operand value that's		
	<b>specifier</b>	<b>operand value</b>	
	\$/imm	/imm	
2. )	specifies an operand value that's		
	<b>specifier</b>	<b>operand value</b>	
	%E <sub>a</sub>	R[%E <sub>a</sub> ]	
3. )	specifies an operand value that's		
	<b>specifier</b>	<b>operand value</b>	<b>effective address</b>
	imm	M[EffAddr]	imm
	(%E <sub>a</sub> )	M[EffAddr]	R[%E <sub>a</sub> ]
	imm(%E <sub>b</sub> )	M[EffAddr]	imm+R[%E <sub>b</sub> ]
	(%E <sub>b</sub> ,%E <sub>i</sub> )	M[EffAddr]	R[%E <sub>b</sub> ]+R[%E <sub>i</sub> ]
	imm(%E <sub>b</sub> ,%E <sub>i</sub> )	M[EffAddr]	imm+R[%E <sub>b</sub> ]+R[%E <sub>i</sub> ]
	imm(%E <sub>b</sub> ,%E <sub>i</sub> ,s)	M[EffAddr]	imm+R[%E <sub>b</sub> ]+R[%E <sub>i</sub> ]*s
	(%E <sub>b</sub> ,%E <sub>i</sub> ,s)	M[EffAddr]	R[%E <sub>b</sub> ]+R[%E <sub>i</sub> ]*s
	imm(,%E <sub>i</sub> ,s)	M[EffAddr]	imm+R[%E <sub>i</sub> ]*s
	(,%E <sub>i</sub> ,s)	M[EffAddr]	R[%E <sub>i</sub> ]*s

## Operands Practice

Given:

Memory Addr	Value	Register	Value
0x100	0x	%eax	0x
0x104	0x	%ecx	0x
0x108	0x	%edx	0x
0x10C	0x		
0x110	0x		

→ What is the value being accessed? Also identify the type of operand, and for memory types name the addressing mode and determine the effective address.

- | Operand              | Value | Type:Mode | Effective Address |
|----------------------|-------|-----------|-------------------|
| 1. (%eax)            |       |           |                   |
| 2. 0xF8 (, %ecx, 8)  |       |           |                   |
| 3. %edx              |       |           |                   |
| 4. \$0x108           |       |           |                   |
| 5. -4 (%eax)         |       |           |                   |
| 6. 4 (%eax, %edx, 2) |       |           |                   |
| 7. (%eax, %edx, 2)   |       |           |                   |
| 8. 0x108             |       |           |                   |
| 9. 259 (%ecx, %edx)  |       |           |                   |

## Instructions - MOV, PUSH, POP

**What?** These are instructions to

**Why?**

**How?**

instruction class	operation	description
MOV S, D		

MOVS S, D

MOVZ S, D

pushl S

popl D

### Practice with Data Formats

→ What data format suffix should replace the \_ given the registers used?

1. mov\_ %eax, %esp
2. push\_ \$0xFF
3. mov\_ (%eax), %dx
4. mov\_ (%esp, %edx, 4), %dh
5. mov\_ 0x800AFFE7, %bl
6. mov\_ %dx, (%eax)
7. pop\_ %edi

\* Focus on register type operands

## Operand/Instruction Caveats

### Missing Combination?

→ Identify each source and destination operand type combinations.

1. `movl $0xABCD, %ecx`
2. `movb $11, (%ebp)`
3. `movb %ah, %dl`
4. `movl %eax, -12(%esp)`
5. `movb (%ebx, %ecx, 2), %al`

→ What combination is missing?

### Instruction Oops!

→ What is wrong with each instruction below?

1. `movl %bl, (%ebp)`
2. `movl %ebx, $0xA1FF`
3. `movw %dx,%eax`
4. `movb $0x11, (%ax)`
5. `movw (%eax), (%ebx,%esi)`
6. `movb %sh, %bl`

## Instruction - LEAL

### Load Effective Address

```
leal S,D      D <-- &S
```

### LEAL vs. MOV

```
struct Point {  
    int x;  
    int y;  
} points[3];  
  
int y = points[i].y;           mov 4(%ebx,%ecx,8),%eax  
  
points[_].y;  
  
int *py = &points[i].y;       leal 4(%ebx,%ecx,8),%eax
```

### LEAL Simple Math

```
leal -3(%ebx), %eax          subl $3, %ebx  
                             movl %ebx, %eax
```

→ Suppose register %eax holds x and %ecx holds y.  
What value in terms of x and y is stored in %ebx for each instruction below?

1. leal (%eax,%ecx,8),%ebx
2. leal 12(%eax,%eax,4),%ebx
3. leal 11(%ecx),%ebx
4. leal 9(%eax,%ecx,4),%ebx

## Instructions - Arithmetic and Shift

### Unary Operations

INC D	D <-- D + 1
DEC D	D <-- D - 1
NEG D	D <-- -D
NOT D	D <-- ~D

### Binary Operations

ADD S,D	D <-- D + S
SUB S,D	D <-- D - S
IMUL S,D	D <-- D * S
XOR S,D	D <-- D ^ S
OR S,D	D <-- D   S
AND S,D	D <-- D & S

Given:

<b>0x100</b> 0xFF	<b>%eax</b> 0x100
<b>0x104</b> 0xAB	<b>%ecx</b> 0x1
<b>0x108</b> 0x10	<b>%edx</b> 0x2

→ What is the destination and result for each? (do each independently)

1. incl 4(%eax)
2. addl %ecx, (%eax)
3. addl \$32, (%eax,%edx,4)
4. subl %edx, 0x104

### Shift Operations

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#### logical shift

SHL k,D	D <-- D << K
SHR k,D	D <-- D >> K

#### arithmetic shift

SAL k,D	D <-- D << K
SAR k,D	D <-- D >> K

## Instructions - CMP and TEST, Condition Codes

### What?

- ◆
- ◆
- ◆

### Why?

### How?

CMP S2, S1                            CC <-- S1 - S2

TEST S2, S1                            CC <-- S1 & S2

- What is done by `testl %eax, %eax`

### Condition Codes (CC)

ZF: zero flag

CF: carry flag

SF: sign flag

OF: overflow flag

## Instructions - SET

### What?

set a byte register to 1 if a condition is true, 0 if false  
specific condition is determined from CCs

### How?

sete D	setz	D <-- ZF	== <u>equal</u>
setne D	setnz	D <-- ~ZF	!= <u>not equal</u>
sets D		D <-- SF	< 0 <u>signed</u> (negative)
setns D		D <-- ~SF	>= 0 <u>not signed</u> (nonnegative)

**Unsigned Comparisons:** t = a - b   if a - b < 0 => CF = 1   if a - b > 0 => ZF = 0

setb D	setnae	D <-- CF	< <u>below</u>
setbe D	setna	D <-- CF   ZF	<= <u>below or equal</u>
seta D	setnbe	D <-- ~CF & ~ZF	> <u>above</u>
setae D	setnb	D <-- ~CF	>= <u>above or equal</u>

**Signed (2's Complement) Comparisons**

setl D	setnge	D <-- SF ^ OF	< <u>less</u> (note l ISN'T size suffix)
setle D	setng	D <-- (SF ^ OF)   ZF	<= <u>less or equal</u>
setg D	setnle	D <-- ~(SF ^ OF) & ~ZF	> <u>greater</u>
setge D	setnl	D <-- ~(SF ^ OF)	>= <u>greater or equal</u>

Demorgan's Law:  $\sim(a \& b) \Rightarrow \sim a \mid \sim b$     $\sim(a \mid b) \Rightarrow \sim a \& \sim b$  note  $\sim$  bitwise not, ! logical not

**Example: a < b** (assume int a is in %eax, int b is in %ebx)

1. cmpl %ebx, %eax
2. setl %cl
3. movzbl %cl, %ecx

## Instructions - Jumps

**What?**

target:

**Why?**

**How? Unconditional Jump**

indirect jump:

```
jmp *Operand
```

direct jump:

```
jmp Label
```

**How? Conditional Jumps**

◆

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both:	je Label	jne Label	js Label	jns Label
unsigned:	jb Label	jbe Label	ja Label	jae Label
signed:	jl Label	jle Label	jg Label	jge Label

# Encoding Targets

## What?

## Absolute Encoding

## Problems?

- ◆ code is not
- ◆ code cannot be

## Solution?

IA-32:

→ What is the distance (in hex) encoded in the `jne` instruction?

Assembly Code	Address	Machine Code
<code>cmpl %eax, %ecx</code>		
<code>jne .L1</code>	<code>0x_B8</code>	<code>75 ??</code>
<code>movl \$11, %eax</code>	<code>0x_BA</code>	
<code>movl \$22, %edx</code>	<code>0x_BC</code>	
<code>.L1:</code>	<code>0x_BE</code>	

→ If the `jb` instruction is 2 bytes in size and is at `0x08011357` and the target is at `0x8011340` then what is the distance (hex) encoded in the `jb` instruction?

## Converting Loops

- Identify which C loop statement (for, while, do-while) corresponds to each goto code fragment below.

```
loop1:                                t = loop_condition
    loop_body
    t = loop_condition
    if (t) goto loop1:
                                         if (!t) goto done:
loop2:                                loop_body
    t = loop_condition
    if (t) goto loop2
done:                                     
```

  

```
loop_init
t = loop_condition
if (!t) goto done:
loop3:                                loop_body
    loop_update
    t = loop_condition
    if (t) goto loop3
done:                                     
```

*Most compilers (gcc included)*