

# CS 354 - Machine Organization & Programming

## Tuesday December 10th, 2024

### Course Evals

<https://aefis.wisc.edu>

Course: CS354

Instructor: DEPPELER | MAHMOOD

**Final Exam - Thursday Dec 19th, 10:05 AM - 12:05 PM**

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**You must attend the exam room as assigned in the email you receive.**

Arrive early if possible with UW ID and #2 pencils. See additional exam info on course web site.

**All office hours, TA consulting, and Peer Mentoring end on Wed December 11th**

**Homework hw8:** DUE on or before Monday Dec 9 (late day Tuesday)

**Homework hw9:** DUE on or before Wednesday Dec 11 (NO LATE DAY)

**Project p6:** Due on last day of classes (NO LATE DAY or OOPS PERIOD). If you plan on getting help in labs, be sure to bring your own laptop in case there is no workstation available.

### Learning Objectives

- ◆ understand and describe how compiler resolves symbols across multiple source files
- ◆ understand and describe why relocation is necessary and how it occurs
- ◆ understand and describe what the Loader is and does

### This Week

Resolving Globals Symbol Relocation Executable Object File Loader What's next? take OS cs537 as soon as possible and Compilers cs536, too!	
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# Resolving Globals

## Confusing Globals

main.c

```
int m;  
int n = 11;  
short o;
```

```
extern int x;  
int y;  
static int z = 66;
```

//code continues...

fun1.c

```
int m = 22;  
int n;  
int o;
```

```
int x;  
static int y = 33;  
static int z = 77;
```

//code continues...

fun2.c

```
int m;  
extern int n;  
char o;
```

```
static int x = 33;  
static int y;  
int z;
```

//code continues...

\* *What happens if multiple definitions of an identifier exist?*

\* *Use `extern` to clearly indicate when*

\* *Use `static` to clearly indicate when*

**TEXTBOOK and OLD NOTES describe old rules for resolving global variables.**

**~~Strong and Weak Symbols~~ (no such thing any more, use `extern` when defined elsewhere)**

~~strong: function definitions and initialized global variables~~

~~weak: function declarations and uninitialized global variables~~

~~→ Which code statements above correspond to strong symbols?~~

## Rules for Resolving Globals

→ Which code statements above correspond to definitions?

Recall: `extern` is only a declaration

Note: `extern` vars must be defined in another file, otherwise undefined symbol linker error.

1. Multiple symbol defns in public gobal scope are not allowed

linker error -> mult defined symbol

Recall: `static` makes a global private, i.e., only visible within its source file)

2. Define one symbol in one file, and declare other with `extern`

Use `gcc -z muldefs` to ignore uninitialized symbols defined in multiple files.

# Symbol Relocation

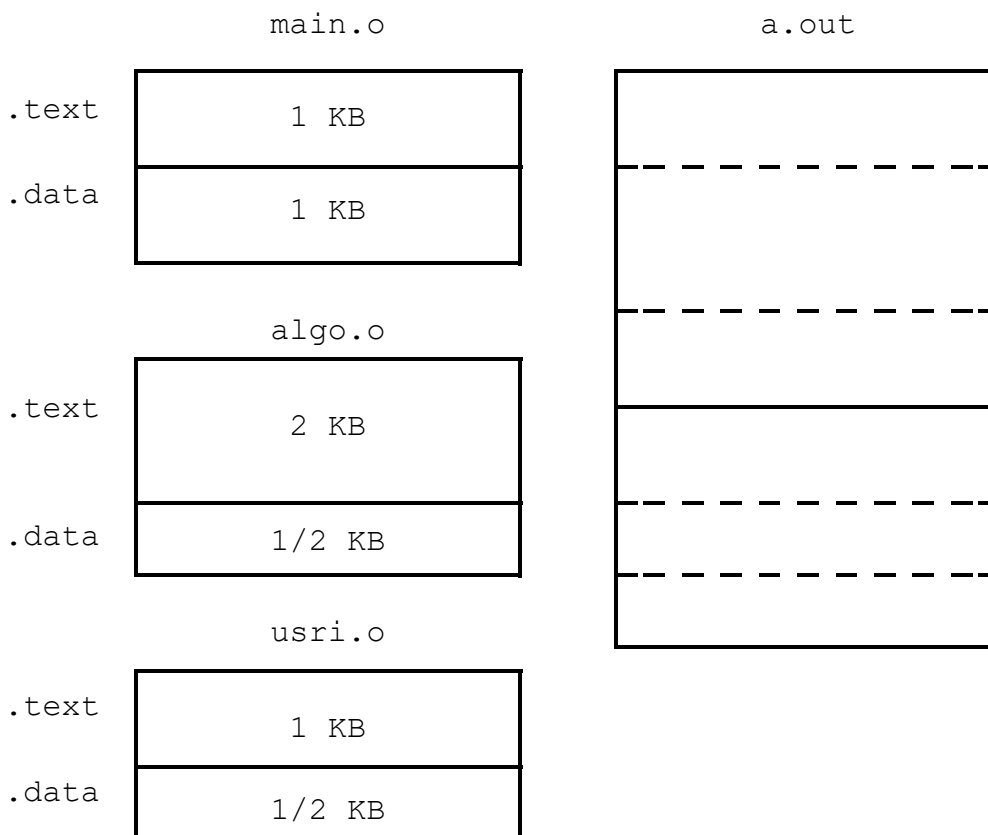
What? Symbol relocation

How?

1. Merge the same sections of ROFs into one aggregate for each section type.
2. Assign virtual addresses to each aggregate section and each symbol definition.
3. Update symbol references listed in ROF relocation sections (.rel.text, .rel.data).

Example

Consider the .text and .data sections of 3 object files below combined into an executable:



address = 1 start of section + 2 offset to subsection + 3 offset within subsection  
calc  
cnt

# Executable Object File (EOF)

**What?** An EOF, like an ROF, is

## Executable and Linkable Format

ELF Header

+ entry point = addr of 1st instr

+ Segment Header Table

+ info for each segment to be loaded into mem during execution  
offset in file

alignment

page size

size in file and size in mem

run-time permissions

ELF Header
Segment Header Table
.init
.text
.rodata
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.symtab
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.strtab
Section Header Table

→ Why aren't there relocation sections (.rel.text or .rel.data) in EOF?

since we've assumed static linking, all symbol relocations are done

➤ Why is the data segment's size in memory larger than its size in the EOF?

# Loader

## What? The *loader*

- ◆ is kernel code that
- ◆ can be invoked by

## Loading

1. “copies” code and data segments from EOF into memory
2. starts program executing by jumping to its entry point

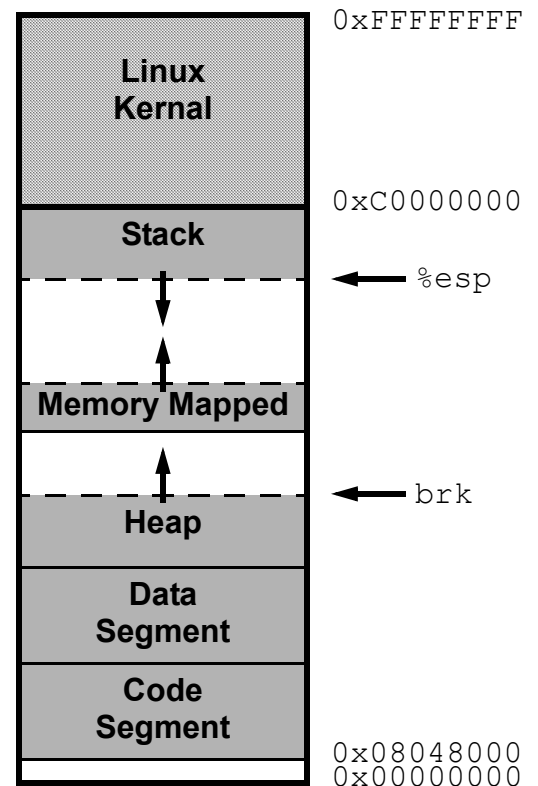
## Execution - the final story

1. shell creates a child process with `fork()`
2. child process invokes loader with `execve()`
3. loader creates the new runtime memory image
  - a. deletes curr segments code, data, heap, stack
  - b. creates new segments
  - c. heap and stack initialized to size 0
  - d. EOF's code and data segments are mapped in page table into page-sized chunks based on Segment Hdr Table BUT THESE ARE NOT COPIED INTO MEM except some header info

## 4. loader

### `_start:`

```
call __libc_init_first
call _init
call atexit
call main
call _exit
```



## CS354 Project Reminders

**p1 Building and running an executable from C source code**

**p2A reading a file with 2D array data and checking the contents meet requirements**

**p2B implementing algorithm to fill 2D array and writing 2D array to a file**

**p3A implementing alloc for a dynamically allocated memory space (heap)**

**p3B implementing free with immediate coalescing (heap)**

**p4A analyzing cache performance for large 2D array access in various sequences (strides)**

**p4B implementing a cache simulator for any sequence of memory access and cache config.**

**p5 disassembling Linux executable and tracing ASM to find input to open a safe.**

**p6 handle SIGINT SIGUSR1 SIGFPE signals, send signals via command line and syscalls**

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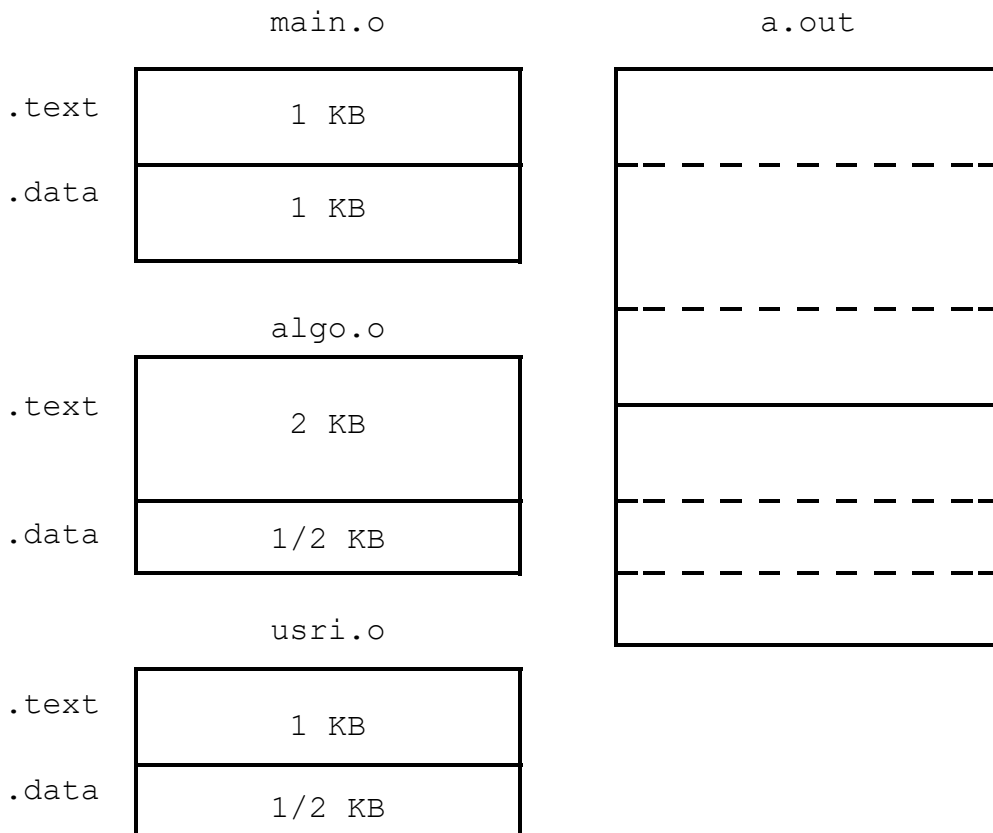
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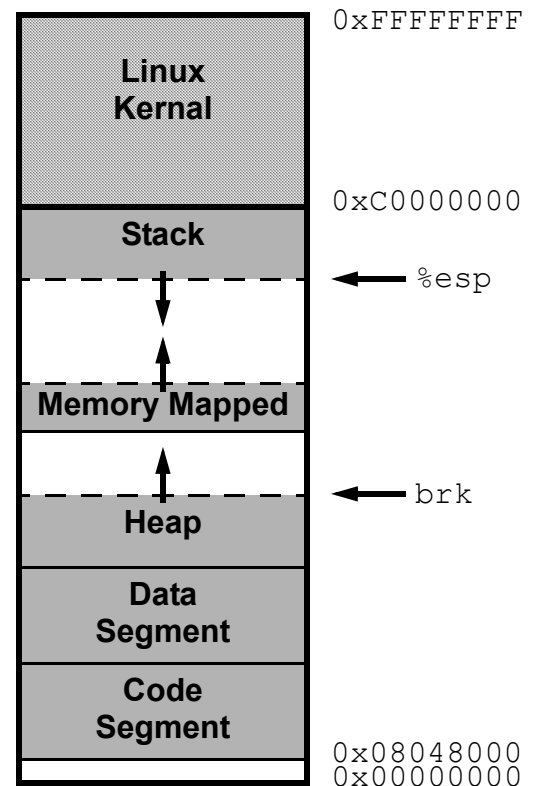
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