## CS 640 Introduction to Computer Networks

Lecture 11

CS 640

# Routing – the big picture

- Internet divided into Autonomous Systems (ASes)
  - corresponds to an administrative domain
  - examples: University, company, backbone network
  - assign each AS a 16-bit number
- Two-level route propagation hierarchy
  - interior gateway protocol (RIP, OSPF)
  - exterior gateway protocol (Internet-wide standard)

CS 640

## Overview

- · Forwarding vs Routing
  - forwarding: to select an output port based on destination address and routing table
  - routing: process by which routing table is built
- · Network as a Graph

• Problem: Find best path between two nodes

- Factors
  - static: topologydynamic: load

CS 640

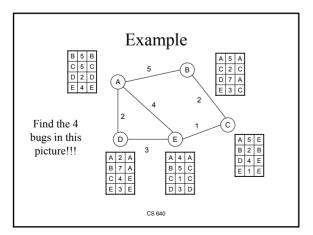
## Families of routing algorithms

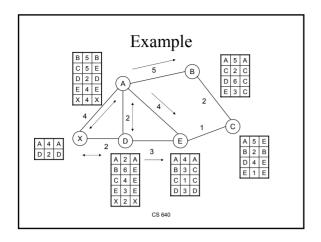
- · Distance vector
  - Tell your neighbors about everybody you know of
  - Lower memory
  - RIP: Route Information Protocol
    - · based on hop-count
- · Link state
  - Tell everybody about your neighbors
  - Most used today
  - OSPF: Open Shortest Path First

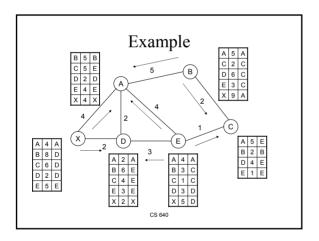
CS 640

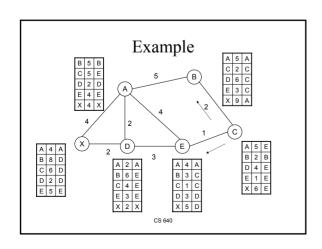
#### Distance Vector

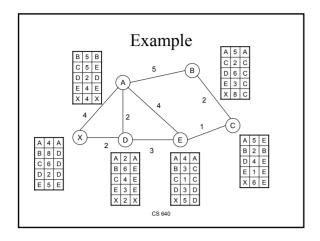
- Each node maintains a set of triples
  - (Destination, Cost, NextHop)
- Neighbors exchange updates
  - periodically (on the order of several seconds)
  - whenever table changes (called *triggered* update)
- Each update is a list of pairs: (Dest, Cost)
- Update local table if receive a "better" route
  - smaller cost
  - came from next-hop
- Refresh existing routes; delete if they time out  $_{\text{CS 640}}$

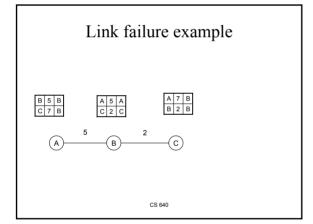


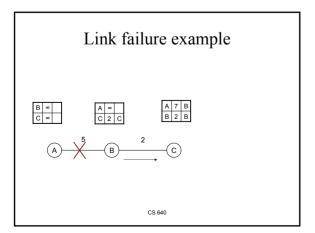


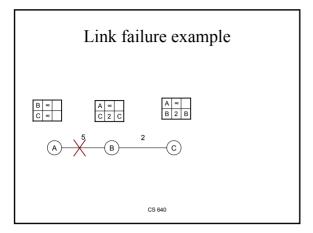


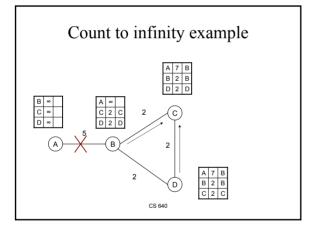


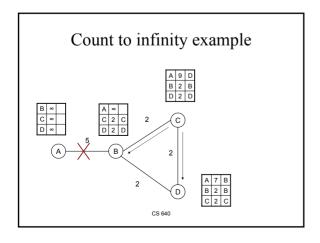


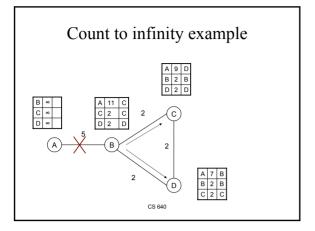


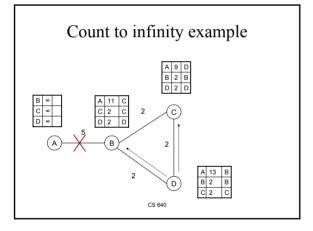












# Loop-Breaking Heuristics

- Set infinity to 16
- Split horizon
  - Don't advertise route to neighbor you heard it from
- Split horizon with poison reverse
  - Advertise it with ∞ cost

CS 640

#### Link State

- Strategy
  - send to all nodes (not just neighbors) information about directly connected links (not entire routing table)
- Link State Packet (LSP)
  - id of the node that created the LSP
  - cost of link to each immediate neighbor
  - sequence number (SEONO)
  - time-to-live (TTL) for this packet

CS 640

### Link State (cont)

- Reliable flooding
  - store most recent LSP from each node
  - forward **new** LSPs to all neighbors (except the one that sent it)
  - generate new LSP periodically
    - increment SEQNO
  - start SEQNO at 0 when reboot
  - decrement TTL of each stored LSP
    - discard when TTL=0

#### **Route Calculation**

- Dijkstra's shortest path algorithm
  - s denotes node performing calculation
    l (i, j) denotes non-negative cost (weight) for edge (i, j)

  - C(n) denotes cost of the path from s to node n
  - N denotes set of all nodes in the graph
    M denotes the set of nodes incorporated so far

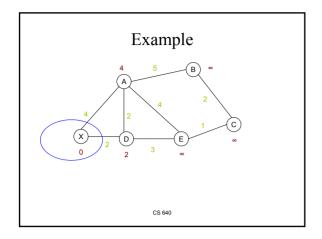
 $M = \{s\}$ for each n in  $N - \{s\}$ 

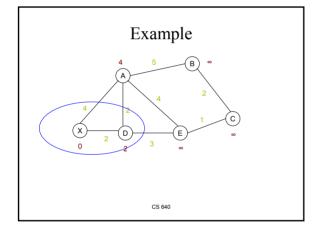
for each n in N - {s} C(n) = 1(s, n)while (N := M)  $M = M + \{w\} \text{ such that } C(w) \text{ is the min for all } w \text{ in } N - M$ for each n in (N - M)  $C(n) = \min(C(n), C(w) + 1(w, n))$ 

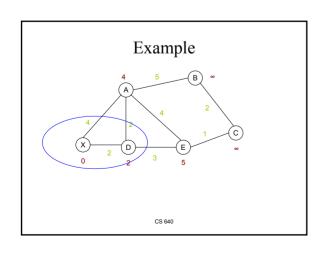
- Invariant of Dijkstra's algorithm
  - We have shortest path for nodes from M to s
- For nodes outside M we have shortest path that goes to s only using nodes in M as next hop

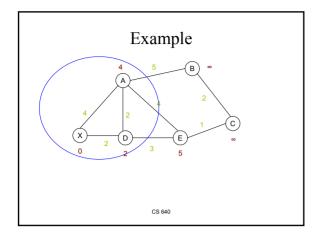
CS 640

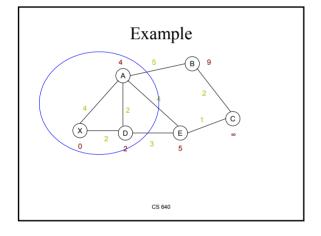
•	
•	
-	
-	

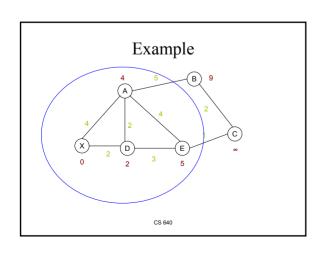












-	
	-

