

CS 640 Introduction to Computer Networks

Lecture24

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Network security (continued)

- Key distribution
- Secure Shell
 - Overview
 - Authentication
 - Practical issues
- Firewalls
- Denial of Service Attacks
 - Definition
 - Examples

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Key Distribution – a first step

- How can we be sure a key belongs to the entity that purports to own it?
- Solution = certificates
 - special type of digitally signed document:
"I certify that the public key in this document belongs to the entity named in this document, signed X."
 - X is the name of the entity doing the certification
 - Only useful to the entity which knows the public key for X
 - Certificates themselves do not solve key distribution problem but they are a first step
- Certified Authority (CA)
 - administrative entity that issues certificates
 - useful only to someone that already holds the CA's public key
 - can trust more than one CA

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Key Distribution (cont)

- Chain of Trust
 - if X certifies that a certain public key belongs to Y , and Y certifies that another public key belongs to Z , then there exists a chain of certificates from X to Z
 - someone that wants to verify Z 's public key has to know X 's public key and follow the chain
 - X.509 is a standard for certificates
- Certificate Revocation List
 - Means for removing certificates
 - Periodically updated by CA

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Key Distribution (cont.)

- PGP (Pretty Good Privacy) provides email encryption and authentication
- Uses “web of trust” instead of “chain of trust”
 - You assign various levels of trust to public keys (e.g. if you got the key when you met face to face you trust it a lot)
 - People certify others' public keys
 - You trust a public key if it has enough “chains of trust”
 - The more disjoint paths in the trust graph the better
 - The shorter the paths the better
 - The more you trust the heads of the paths the better

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Secure Shell (SSH) Overview

- SSH is a **secure** remote virtual terminal application
 - Provides encrypted communication over an insecure network
 - Assumes eavesdroppers can hear *all* communications between hosts
 - Provides different methods of authentication
 - Encrypts data exchanged between hosts
 - Intended to replace insecure programs such as telnet, rsh, etc.
 - Includes capability to securely transfer file
 - SCP
 - Can forward X11 connections and TCP ports securely
- Very popular and widely used
 - Not invulnerable!

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SSH authentication

- Client authenticates server
 - The client caches the public keys of all servers it talks to
 - User can add new keys to the cache
 - Otherwise the user is warned when first connecting to a given server
- Server authenticates client
 - Through user's password
 - Public RSA key the user puts ahead of time on the server
 - Other, riskier methods
- At connection setup server and client agree on a session key used to encrypt communication
 - Many algorithms allowed (IDEA, Blowfish, Triple DES, etc.)

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SSH in Practice

- Host public/private key is generated when SSH is installed
 - Public key must be in `~/.ssh/known_hosts` on remote systems
- `ssh-keygen` command is used to generate users public/private keys
 - Public key copied to `~/.ssh/authorized_keys` on remote systems
 - Each private key in `~/.ssh/identity` requires a pass phrase when used
 - *Ssh-agent* eliminates need for repeated typing of pass phrase
- Password authentication is vulnerable to guessing attacks
 - Server logs all unsuccessful login attempts
- X11 and port forwarding enable encrypted pipe through the Internet
 - Can be used to securely access insecure application eg. SMTP
 - Can be used to circumvent firewalls

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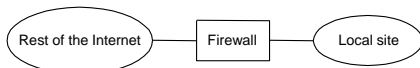
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Firewalls – overview

- Firewalls restrict communication between an organization's computers and the outside world
 - Keep the bad guys on the outside from exploiting vulnerabilities on the inside
 - Without restricting legitimate traffic
- NAT boxes implement a popular firewall policy
 - Allow internal clients to connect to outside servers
 - Do not allow inbound connections
- Two types of firewalls
 - Filter based (layer 4)
 - Proxy based (application layer)

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Firewalls



- Filter-Based Solution
 - Apply a set of rules to packets
 - Look at packet headers
 - Example of rules

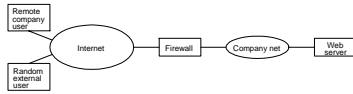
action	ourhost	port	theirhost	port	comment
block	*	*	BLASTER	*	We don't trust this system
allow	OUR_GW	25	*	*	Connects to our SMTP svr

- Default: forward or not forward?
- How dynamic?

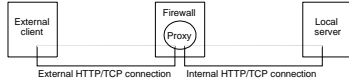
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Proxy-Based Firewalls

- Problem: complex policy
- Example: web server



- Solution: proxy



- Design: transparent vs. classical
- Limitations: attacks from within

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Denial of Service (DoS) Attacks

- A general form of attacking inter-networked systems
 - Based on overloading end systems (or network)
 - Result is severe reduction in performance or complete shutdown of target systems
- Focus of attack can be links, routers (CPU) or end hosts
- Flooding attacks pretty common nowadays
- Other most general form of attack is a break-in
 - Port scans
 - Buffer overflows
 - Password cracking...

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Overloading a System

- The goal of DoS is to drown legitimate traffic in a sea of garbage traffic
 - Clients experience delays due to congestion
 - Dropped packets lead to exponential backoff in timeouts
 - Routers can become overloaded
- Servers become overloaded by increased number of connect requests
 - TCP connection setup requires state on server
 - Server is required to respond to SYN from clients
 - Clients don't respond to server's response

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IP Spoofing

- Insert a different source IP address in TCP and IP headers
 - DoS attackers spoof for two reasons
 - They don't want to be discovered
 - Spoofing can add additional load
- If attacker spoofs a legitimate IP address
 - Reset can be triggered by either attacked host or actual IP host
 - Frees resources immediately on server
 - Carefully chosen sequence #s block new connections from host
- Attackers spoof with random IP addresses
 - Server response to client SYN will be lost
 - Server will not free resources for 75 seconds (typically)
 - SYN cookies on allow server kernel not to keep state

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Key Elements of DoS Attack

- Expansion in required work
 - Easy for me, harder for you
 - Expansion in IP spoofing
 - Me: generate SYNs as fast as possible (microseconds)
 - You: Timeout a SYN open every 75 seconds
- Best effort protocols
 - Drop tail queues
 - No source specificity
 - Clients can be starved or slowed to crawl

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DoS Attack Characteristics

- Expansion makes a only a few systems necessary
 - DDoS: attack from as many places as possible
 - Enables better utilization of network resources
 - Helps to prevent countermeasures
 - Helps to obscure attackers
- DoS software readily available
 - Most found in IRC chat rooms
- DoS attacks frequently preceded by break-ins to install DoS software onto “zombies”
 - Enables even more anonymity for attacker

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Things making DoS Attacks easy

- Lots of systems
- Large networks
- Naïve users with high speed Internet access
- Savvy bad guys
- Lots of free DoS software
- Poor operating and management policies
- Hugely complex software (on endhosts) with lots of well publicized holes
- Lack of means for stopping attacks

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Dealing with DoS Attacks

- Don't reserve state until receipt of client ACK
 - DOS attackers using spoofing don't send these
 - Otherwise they would have to keep state
 - Use of crypto to avoid saving state
 - Send one-use key with server response to SYN
 - Response ACK must return key
- Intrusion detection tools
 - Cut off an attack at a firewall if you recognize it
 - Bro, Snort
- IP traceback methods
- There are lots of companies in this space!

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Example of (D)DoS

- Code Red Worm
 - Released and identified on July 19, 2001
 - Infected over 250k systems in 9 hours
 - Takes advantage of hole in IIS on Win NT or Win 2k
 - And the fact that most people don't know IIS ON is default
 - Infected systems are completely compromised
 - Code Red installs itself in OS kernel
 - Small and efficient
 - V1 could be eliminated by reboot
 - Spends half its time trying to infect other systems, and half its time DoS'ing the White House and Pentagon

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