Optimality conditions and complementarity, variational inequalities, operators and graphs

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Funded by DOE-MACS Grant with Argonne National Laboratory

Auckland, April 26, 2018

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## The PIES Model (Hogan) - Optimal Power Flow (OPF)

$$\begin{array}{ll} \min_{x} c(x) & \mbox{cost} \\ \mbox{s.t.} \ Ax \geq q & \mbox{balance} \\ Bx = b, x \geq 0 & \mbox{technical constr} \end{array}$$

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•  $q = d(\pi)$ : issue is that  $\pi$  is the multiplier on the "balance" constraint

- Such multipliers (LMP's) are critical to operation of market
- Can try to solve the problem iteratively (shooting method):

 $\pi^{new} \in \mathsf{multiplier}(OPF(d(\pi)))$ 

Alternative: Form KKT of QP, exposing  $\pi$  to modeler

$$L(x,\mu,\lambda) = c(x) + \mu^{T}(d(\pi) - Ax) + \lambda^{T}(b - Bx)$$

$$0 \le -\nabla_{\mu}L = Ax - d(\pi) \qquad \bot \quad \mu \ge 0$$
  
$$0 = -\nabla_{\lambda}L = Bx - b \qquad \bot \quad \lambda$$
  
$$0 \le \quad \nabla_{x}L = \nabla c(x) - A^{T}\mu - B^{T}\lambda \quad \bot \quad x \ge 0$$

- empinfo: dualvar  $\pi$  balance
- Fixed point: replaces  $\mu \equiv \pi$

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- empinfo: dualvar  $\pi$  balance
- Fixed point: replaces  $\mu \equiv \pi$
- LCP/MCP is then solvable using PATH

$$z = \begin{bmatrix} \pi \\ \lambda \\ x \end{bmatrix}, \quad F(z) = \begin{bmatrix} & & A \\ & & B \\ -A^T & -B^T \end{bmatrix} z + \begin{bmatrix} -d(\pi) \\ -b \\ \nabla c(x) \end{bmatrix}$$

• Existence, uniqueness, stability from variational analysis

## Convex subdifferentials



- Assume f is convex, then  $f(z) \ge f(x) + \nabla f(x)^T (z - x)$ (linearization is below the function)
- Incorporate constraints by allowing f to take on +∞ if constraint is violated f : ℝ<sup>n</sup> → (-∞, +∞]
- $\partial f(x) = \{g: f(z) \ge f(x) + g^T(z-x), \forall z\},\$ the subdifferential of f at x

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- $\partial f(x) = \{g: f(z) \ge f(x) + g^T(z-x), \forall z\},\$ the subdifferential of f at x
- If f is differentiable and convex, then  $\partial f(x) = \{\nabla f(x)\}$
- e.g.  $f(z) = \frac{1}{2}z^TQz + p^Tz$ , then  $\partial f(x) = \{Qx + p\}$
- $x^*$  solves min f(x) if and only if  $0 \in \partial f(x^*)$

## Indicator functions and normal cones

$$\psi_{\mathcal{C}}(z) = egin{cases} 0 & ext{if } z \in \mathcal{C} \ \infty & ext{else} \end{cases}$$

 $\psi_{\mathcal{C}}$  is a convex function when  $\mathcal{C}$  is a convex set



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If 
$$x \in C$$
, then  
 $g \in \partial \psi_C(x)$   
 $\iff \psi_C(z) \ge \psi_C(x) + g^T(z - x), \ \forall z$   
 $\iff 0 \ge g^T(z - x), \ \forall z \in C$ 

Normal cone to C at x.

$$N_{\mathcal{C}}(x) := \partial \psi_{\mathcal{C}}(x) = \begin{cases} \{g : g^{T}(z - x) \leq 0, \forall z \in \mathcal{C} \} & \text{if } x \in \mathcal{C} \\ \emptyset & \text{if } x \notin \mathcal{C} \end{cases}$$

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#### Some calculus

•  $f_i : \mathbb{R}^n \mapsto (-\infty, \infty], i = 1, \dots, m$ , proper, convex functions  $F = f_1 + \cdots + f_m$ m assume  $\bigcap \operatorname{rint}(\operatorname{dom}(f_i)) \neq \emptyset$  then (as sets) i=1 $\partial F(x) = \partial f_1(x) + \cdots + \partial f_m(x), \ \forall x$ •  $C = \bigcap C_i$ , then  $\psi_C = \psi_{C_i} + \cdots + \psi_{C_m}$ , so  $N_C = N_{C_i} + \cdots + N_{C_m}$  $x^*$  solves  $\min_{x \in \mathcal{C}} f(x) \iff x^*$  solves  $\min_{x} (f + \psi_{\mathcal{C}})(x)$  $\iff 0 \in \partial (f + \psi_{\mathcal{C}})(x^*) \iff 0 \in \nabla f(x^*) + N_{\mathcal{C}}(x^*)$ 

## Special cases and examples

- Normal cone is a cone
- $x \in int(\mathcal{C})$ , then  $N_{\mathcal{C}}(x) = \{0\}$

• 
$$\mathcal{C} = \mathbb{R}^n$$
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## Special cases and examples

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• 
$$C = \{z : a_i^T z \le b_i, i = 1, \dots, m\}$$
  
polyhedral

• 
$$\mathcal{N}_{\mathcal{C}}(x) = \left\{ \sum_{i=1}^{m} \lambda_i \mathbf{a}_i : 0 \le b_i - \mathbf{a}_i^T x \perp \lambda_i \ge 0 \right\}$$

•  $\perp$  makes product of items around it 0, i.e.

$$(b_i - a_i^T x)\lambda_i = 0, \ i = 1, \dots, m$$

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## Combining: KKT conditions

• Example: convex optimization first-order optimality condition:

$$x^* \text{ solves } \min_{x \in \mathcal{C}} f(x) \iff 0 \in \nabla f(x^*) + N_{\mathcal{C}}(x^*)$$
$$\iff 0 = \nabla f(x^*) + y, \ y \in N_{\mathcal{C}}(x^*)$$
$$\iff 0 = \nabla f(x^*) + y, \ y = A^T \lambda,$$
$$0 \le b - Ax^* \perp \lambda \ge 0$$
$$\iff 0 = \nabla f(x^*) + A^T \lambda,$$
$$0 \le b - Ax^* \perp \lambda \ge 0$$

• More generally, if  $\mathcal{C} = \{z : g(z) \leq 0\}$ , g convex, (with CQ)

$$\begin{aligned} x^* \text{ solves } \min_{x \in \mathcal{C}} f(x) & \Longleftrightarrow 0 \in \nabla f(x^*) + \mathcal{N}_{\mathcal{C}}(x^*) \\ & \Longleftrightarrow 0 = \nabla f(x^*) + \nabla g(x^*)\lambda, \\ & 0 \leq -g(x^*) \perp \lambda \geq 0 \end{aligned}$$

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Variational Inequality (replace  $\nabla f(z)$  with F(z))

•  $F: \mathbb{R}^n \to \mathbb{R}^n$ 

• Ideally:  $C \subseteq \mathbb{R}^n$  – constraint set; Often:  $C \subseteq \mathbb{R}^n$  – simple bounds

$$VI(F,C): 0 \in F(z) + N_{\mathcal{C}}(z)$$

- VI generalizes many problem classes
- Nonlinear Equations: F(z) = 0 set  $\mathcal{C} \equiv \mathbb{R}^n$
- Convex optimization:  $F(z) = \nabla f(z)$
- For NCP:  $0 \leq F(z) \perp z \geq 0$ , set  $\mathcal{C} \equiv \mathbb{R}^n_+$
- For MCP (rectangular VI), set  $C \equiv [I, u]^n$ .
- For LP, set  $F(z) \equiv \nabla f(z) = p$  and  $C = \{z : Az = a, Hz \leq h\}$ .

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VI:  $0 \in F(z) + \mathcal{N}_{\mathcal{C}}(z)$ 



#### Many applications where F is not the derivative of some f

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Optimality and Complementarity

Auckland, April 2018 10 / 17

# Other applications of complementarity

#### Complementarity can model fixed points and disjunctions

- Economics: Walrasian equilibrium (supply equals demand), taxes and tariffs, computable general equilibria, option pricing (electricity market), airline overbooking
- Transportation: Wardropian equilibrium (shortest paths), selfish routing, dynamic traffic assignment
- Applied mathematics: Free boundary problems
- Engineering: Optimal control (ELQP)
- Mechanics: Structure design, contact problems (with friction)
- Geology: Earthquake propagation

Good solvers exist for large-scale instances of Complementarity Problems

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## Complementarity Problems via Graphs



$$-y \in \mathcal{T}(z) \iff (z,-y) \in \mathcal{T} \iff 0 \leq y \perp z \geq 0$$

By approximating (smoothing) graph can generate interior point algorithms for example  $yz = \epsilon, y, z > 0$ 

 $0 \in F(z) + \mathcal{N}_{\mathbb{R}^n_+}(z) \iff (z, -F(z)) \in \mathcal{T}^n \iff 0 \leq F(z) \perp z \geq 0$ 

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Operators and Graphs  $(\mathcal{C} = [-1, 1], \mathcal{T} = \mathcal{N}_{\mathcal{C}})$ 

$$z_i = -1, -y_i \leq 0 ext{ or } z_i \in (-1, 1), -y_i = 0 ext{ or } z_i = 1, -y_i \geq 0$$



 $(\mathcal{I} + \mathcal{T})^{-1}(y)$  is the projection of y onto [-1, 1]:  $P_{\mathcal{C}}(y)$ 

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#### Generalized Equations

• Suppose  ${\mathcal T}$  is a maximal monotone operator

$$0 \in F(z) + \mathcal{T}(z)$$
 (GE)

- Define  $P_{\mathcal{T}} = (\mathcal{I} + \mathcal{T})^{-1}$
- If  $\mathcal{T}$  is polyhedral (graph of  $\mathcal{T}$  is a finite union of convex polyhedral sets) then  $P_{\mathcal{T}}$  is piecewise affine (continous, single-valued, non-expansive)

$$egin{aligned} 0 \in F(z) + \mathcal{T}(z) & \iff & z \in F(z) + \mathcal{I}(z) + \mathcal{T}(z) \ & \iff & z - F(z) \in (\mathcal{I} + \mathcal{T})(z) \iff P_{\mathcal{T}}(z - F(z)) = z \end{aligned}$$

Use in fixed point iterations (cf projected gradient methods)

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## Splitting Methods

• Suppose  $\mathcal{T}$  is a maximal monotone operator

$$0 \in F(z) + \mathcal{T}(z)$$
 (GE)

- Can devise Newton methods (e.g. SQP) that treat F via calculus and  ${\cal T}$  via convex analysis
- Alternatively, can split F(z) = A(z) + B(z) (and possibly T also) so we solve solve (GE) by solving a sequence of problems involving just

$$\mathcal{T}_1(z) = A(z)$$
 and  $\mathcal{T}_2(z) = B(z) + \mathcal{T}(z)$ 

where each of these is "simpler"

• Forward-Backward splitting (or ADMM):

$$z^{k+1} = (I + c_k T_2)^{-1} (I - c_k T_1) (z^k),$$

## Normal Map

• Suppose  ${\mathcal T}$  is a maximal monotone operator

$$0 \in F(z) + \mathcal{T}(z)$$
 (GE)

• Define  $P_{\mathcal{T}} = (I + \mathcal{T})^{-1}$  (continuous, single-valued, non-expansive)

$$0 \in F(z) + \mathcal{T}(z) \iff z \in F(z) + \mathcal{I}(z) + \mathcal{T}(z)$$
  
$$\iff z - F(z) = x \text{ and } x \in (\mathcal{I} + \mathcal{T})(z)$$
  
$$\iff z - F(z) = x \text{ and } P_{\mathcal{T}}(x) = z$$
  
$$\iff P_{\mathcal{T}}(x) - F(P_{\mathcal{T}}(x)) = x$$
  
$$\iff 0 = F(P_{\mathcal{T}}(x)) + x - P_{\mathcal{T}}(x)$$

This is the so-called Normal Map Equation

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## Conclusions

- Convexity separates easy optimization problems from hard ones
- Modern convex analysis extends linear programming to richer but still tractable settings
- Modern optimization within applications requires multiple model formats, computational tools and sophisticated solvers
- Variational inequalities and set valued analysis important tools for big data problems
- Modeling, optimization, statistics and computation embedded within the application domain is critical
- Many new settings available for deployment; need for more theoretic and algorithmic enhancements

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## Bimatrix Games: Golden Balls

- VI can be used to formulate many standard problem instances corresponding to special choices of M and C.
- Nash game: two players have I and J pure strategies.
- *p* and *q* (strategy probabilities) belong to unit simplex  $\triangle_I$  and  $\triangle_J$  respectively.
- Payoff matrices  $A \in R^{J \times I}$  and  $B \in R^{I \times J}$ , where  $A_{j,i}$  is the profit received by the first player if strategy *i* is selected by the first player and *j* by the second, etc.
- The expected profit for the first and the second players are  $q^T A p$  and  $p^T B q$  respectively.
- A Nash equilibrium is reached by the pair of strategies  $(p^*, q^*)$  if and only if

$$p^* \in \arg \min_{p \in riangle_I} \langle Aq^*, p 
angle$$
 and  $q^* \in \arg \min_{q \in riangle_J} \langle B^T p^*, q 
angle$ 

## Formulation using complementarity

The optimality conditions for the above problems are:

$$-Aq^* \in N_{ riangle_I}(p^*)$$
 and  $-B^T p^* \in N_{ riangle_J}(q^*)$ 

Therefore the corresponding VI is affine and can be written as:

$$0 \in \begin{bmatrix} 0 & A \\ B^{T} & 0 \end{bmatrix} \begin{bmatrix} p \\ q \end{bmatrix} + N_{\triangle_{I} \times \triangle_{J}} \left( \begin{bmatrix} p \\ q \end{bmatrix} \right).$$
(1)