

Abstract Data Types

ML also provides abstract data types in which the implementation of the type is *hidden* from users.

The general form is

```
abstype name = implementation  
with
```

```
    val and fun definitions
```

```
end;
```

Users may access the **name** of the abstract type and the **val** and **fun** definitions that follow the **with**, but the **implementation** may be used only with the body of the **abstype** definition.

Example

```
abstype 'a stack =  
  stk of 'a list  
with  
  val Null = stk([])  
  fun empty(stk([])) = true  
    | empty(stk(_::_)) = false  
  fun top(stk(h::_)) = h  
  fun pop(stk(_:t)) = stk(t)  
  fun push(v,stk(L)) =  
    stk(v::L)  
end  
type 'a stack  
val Null = - : 'a stack  
val empty = fn : 'a stack -> bool  
val top = fn : 'a stack -> 'a  
val pop =  
  fn : 'a stack -> 'a stack  
val push = fn :  
  'a * 'a stack -> 'a stack
```

Local value and function definitions, not to be exported to users of the type can be created using the **local** definition mechanism described earlier:

```
local
```

```
    val and fun definitions
```

```
in
```

```
    exported definitions
```

```
end;
```

```
abstype 'a stack =  
  stk of 'a list  
with  
  local  
    fun size(stk(L))=length(L);  
  in  
    val Null = stk([])  
    fun empty(s) =  
      (size(s) = 0)  
    fun top(stk(h::_)) = h  
    fun pop(stk(_:t)) = stk(t)  
    fun push(v,stk(L)) =  
      stk(v::L)  
  end  
end  
type 'a stack  
val Null = - : 'a stack  
val empty = fn : 'a stack -> bool  
val top = fn : 'a stack -> 'a  
val pop = fn :  
  'a stack -> 'a stack  
val push = fn :  
  'a * 'a stack -> 'a stack
```

Why are abstract data types useful?

Because they hide an implementation of a type from a user, allowing implementation changes without any impact on user programs.

Consider a simple implementation of queues:

```
abstype 'a queue =  
  q of 'a list  
with  
  val Null = q([])  
  fun front(q(h::_)) = h  
  fun rm(q(_::t)) = q(t)  
  fun enter(v,q(L)) =  
    q(rev(v::rev(L)))  
end  
type 'a queue  
val Null = - : 'a queue  
val front = fn : 'a queue -> 'a
```

```
val rm =  
  fn : 'a queue -> 'a queue  
val enter =  
  fn : 'a * 'a queue -> 'a queue
```

This implementation of queues is valid, but somewhat inefficient. In particular to enter a new value onto the rear end of a queue, we do the following:

```
fun enter(v,q(L)) =  
  q(rev(v::rev(L)))
```

We reverse the list that implements the queue, add the new value to the head of the reversed queue *then* reverse the list a second time.

A more efficient (but less obvious) implementation of a queue is to store it as two lists. One list represents the “front” of the queue. It is from this list that we extract the front value, and from which we remove elements. The other list represents the “back” of the queue (in reversed order). We add elements to the rear of the queue by adding elements to the front of the list. From time to time, when the front list becomes null, we “promote” the rear list into the front list (by reversing it). Now access to both the front and the back of the queue is fast and direct. The new implementation is:

```
abstype 'a queue =  
  q of 'a list * 'a list  
with  
  val Null = q([],[])  
  fun front(q(h::_,_)) = h  
  | front(q([],L)) =  
    front(q(rev(L),[]))  
  fun rm(q(_::t,L)) = q(t,L)  
  | rm(q([],L)) =  
    rm(q(rev(L),[]))  
  fun enter(v,q(L1,L2)) =  
    q(L1,v::L2)  
end  
type 'a queue  
val Null = - : 'a queue  
val front = fn :  
  'a queue -> 'a  
val rm = fn :  
  'a queue -> 'a queue  
val enter = fn :  
  'a * 'a queue -> 'a queue
```

From the user's point of view, the two implementations are *identical* (they export exactly the same set of values and functions). Hence the new implementation can replace the old implementation without any impact at all to the user (except, of course, performance!).

Exception Handling

Our definitions of stacks and queues are incomplete. Reconsider our definition of stack:

```
abstype 'a stack =  
  stk of 'a list  
with  
  val Null = stk([])  
  fun empty(stk([])) = true  
    | empty(stk(_::_)) = false  
  fun top(stk(h::_)) = h  
  fun pop(stk(_::_t)) = stk(t)  
  fun push(v, stk(L)) =  
    stk(v::_L)  
end
```

What happens if we evaluate

```
top(Null);
```

We get a “match failure” because our definition of **top** is incomplete!

In ML we can *raise* an exception if an illegal or unexpected operation occurs. Asking for the top of an empty stack ought to raise an exception since the requested value does not exist.

ML contains a number of predefined exceptions, including

Match Empty Div Overflow
(exception names by convention begin with a capital letter).

Predefined exception are raised by illegal values or operations. If they are not caught, the run-time system prints an error message.

```
fun f(1) = 2;  
val f = fn : int -> int  
f(2);  
uncaught exception nonexhaustive  
match failure  
hd [];  
uncaught exception Empty  
1000000*1000000;  
uncaught exception overflow  
(1 div 0);  
uncaught exception divide by zero  
1.0/0.0;  
val it = inf : real  
(inf is the IEEE floating-point  
standard “infinity” value)
```

User Defined Exceptions

New exceptions may be defined as
`exception name;`

or

`exception name of type;`

For example

`exception IsZero;`

`exception IsZero`

`exception NegValue of real;`

`exception NegValue of real`

Exceptions May be Raised

The `raise` statement raises (throws) an exception:

`raise exceptionName;`

or

`raise exceptionName(expr);`

For example

`fun divide(a,0) = raise IsZero
 | divide(a,b) = a div b;`

`val divide =
 fn : int * int -> int`

`divide(10,3);`

`val it = 3 : int`

`divide(10,0);`

`uncaught exception IsZero`

```
val sqrt = Real.Math.sqrt;  
val sqrt = fn : real -> real  
fun sqroot(x) =  
  if x < 0.0  
  then raise NegValue(x)  
  else sqrt(x);  
val sqroot = fn : real -> real  
sqroot(2.0);  
val it = 1.41421356237 : real  
sqroot(~2.0);  
uncaught exception NegValue
```

Exception Handlers

You may catch an exception by defining a *handler* for it:

```
(expr) handle exception1 => val1  
           || exception2 => val2  
           || ... ;
```

For example,

```
(sqroot ~100.0)  
  handle NegValue(v) =>  
    (sqrt (~v));  
val it = 10.0 : real
```

Stacks Revisited

We can add an exception, **EmptyStk**, to our earlier stack type to handle **top** or **pop** operations on an empty stack:

```
abstype 'a stack = stk of 'a list
with
  val Null = stk([])
  exception EmptyStk
  fun empty(stk([])) = true
    | empty(stk(_::_)) = false
  fun top(stk(h::_)) = h
    | top(stk([])) =
        raise EmptyStk
  fun pop(stk(_::_t)) = stk(t)
    | pop(stk([])) =
        raise EmptyStk
  fun push(v, stk(L)) =
    stk(v::L)
end
```

```
type 'a stack
val Null = - : 'a stack
exception EmptyStk
val empty = fn : 'a stack -> bool
val top = fn : 'a stack -> 'a
val pop = fn :
  'a stack -> 'a stack
val push = fn : 'a * 'a stack ->
  'a stack

pop(Null);
uncaught exception EmptyStk
top(Null) handle EmptyStk => 0;
val it = 0 : int
```