# File System Checkers

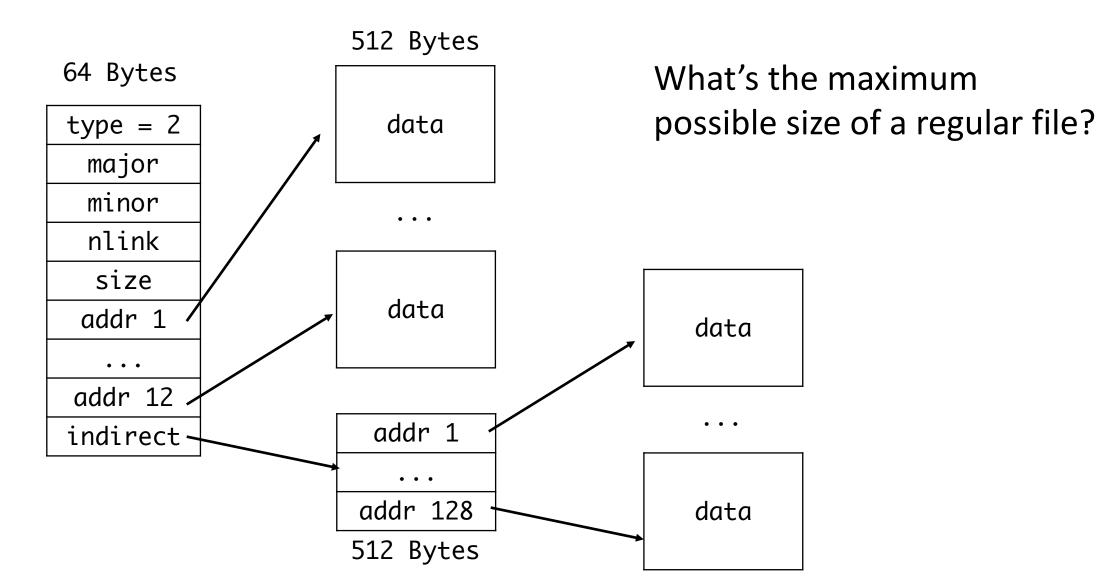
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### File System Layout in xv6

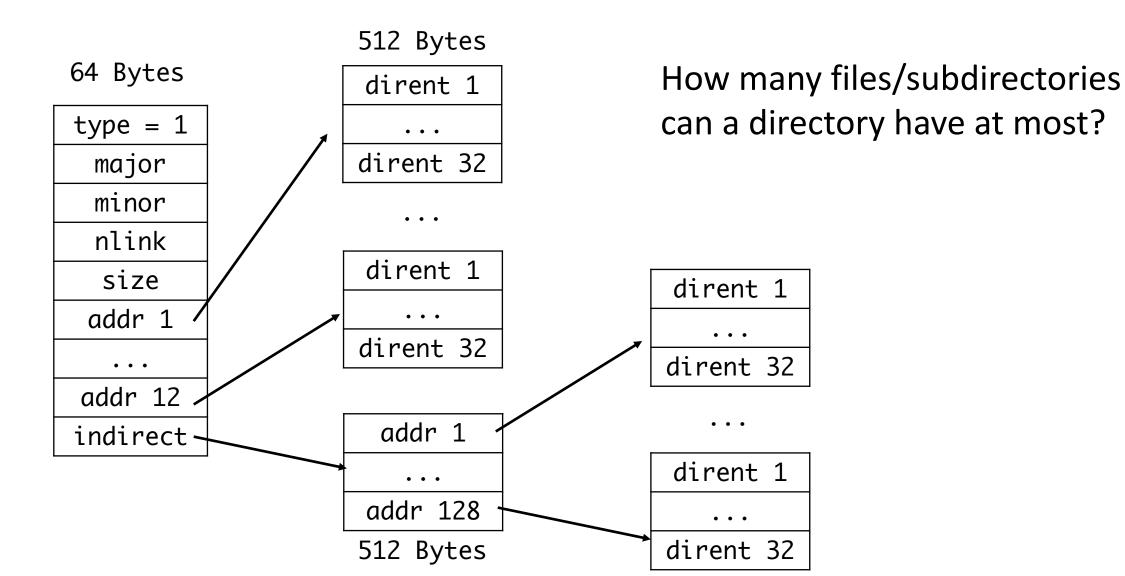
Unused | Superblock | Inodes ... | Bitmap | Data ...

```
50 struct dirent {
51  ushort inum;
52  char name[DIRSIZ];
53 };
```

# Inode of a regular file in xv6



# Inode of a directory in xv6



#### Bitmap in xv6

- Each bit in the bitmap is associated with a block, NOT an inode.
- Although the very first block is unused, it is always marked as 1 in the first bit of the bitmap, and so are all the blocks where the inodes and bitmap itself reside.
- Bitmap is grouped in byte.
- Intel x86 processors use little-endian.
- Example:
  - ff c2 => 1111 1111 1100 0010 7 6 5 4 3 2 1 0 | 15 14 13 12 11 10 9 8

#### Demos

- How to build your own file system image and reflect in xv6
  - Closer look at mkfs.c
- How xv6 files change update the image
  - Closer look at fs.img with xxd
- How to read image? mmap()!

#### How ln works

- Reference counts (number of links) for regular files match the number of times file is referred to in directories (i.e., hard links work correctly).
- No extra links allowed for directories (each directory only appears in one other directory).

