

Lecture 1

1-20-26

first version of definition:

Def An algorithm $A: \mathcal{X}^n \rightarrow \mathcal{Y}$ is

(ϵ, δ) -DP if $\forall x, x' \in \mathcal{X}^n$ such that

$x \neq x'$ differ in one entry and

events $E \subseteq \mathcal{Y}$,

$$\Pr[A(x) \in E] \leq e^\epsilon \Pr[A(x') \in E] + \delta.$$

Here: \mathcal{X} = space of data points

\mathcal{X}^n = datasets of n examples

\mathcal{Y} = output space for algorithm

$$\epsilon, \delta \geq 0$$

Another way to state definition.

Def (Adjacency) Two datasets $x = (x_1, \dots, x_n)$ and $x' = (x'_1, \dots, x'_n)$ are adjacent, written $x \sim x'$, if $\exists i^* \in [n]$ such that $\forall i \neq i^*, x_i = x'_i$.

Def ((ϵ, δ) -indistinguishability) Two distributions P, Q are (ϵ, δ) -indistinguishable, written $P \approx_{(\epsilon, \delta)} Q$, if \forall events E

$$P(E) \leq e^\epsilon Q(E) + \delta$$
$$Q(E) \leq e^\epsilon P(E) + \delta.$$

Def Algorithm A is (ϵ, δ) -DP if $\forall x \sim x'$, $A(x) \approx_{(\epsilon, \delta)} A(x')$.

With different notions of adjacency and/or different notions of "closeness of distributions", get different flavors of DP.

$\epsilon, \delta \Rightarrow$ approx

$\delta = 0 \Rightarrow$ pure

DP is worst-case over datasets.