

Today: Beating Global Sensitivity

- Local notions of sensitivity
- Propose-Test-Release
- Gaussian Mean Estimation

Local Notions of Sensitivity

Def Global sensitivity of $f: \mathcal{X}^n \rightarrow \mathbb{R}$

$$\text{is } \Delta_f = \max_{x \sim x'} |f(x) - f(x')|$$

Privacy guarantees must be worst-case

In many cases, global sensitivity is large because of pathological or atypical instances.

Example Median of data in $[0, R]$,

$$\Delta_{\text{med}} = R, \text{ only happens when half/half } 0 \text{ \& } R$$

Example Mean of $x_1, \dots, x_n \in \mathbb{R}$.

$$\Delta_{\text{mean}} = \infty \quad \text{b/c outliers}$$

But if $x_1, \dots, x_n, x_{n+1} \stackrel{\text{iid}}{\sim} \mathcal{N}(\mu, \sigma^2)$ then

$$\left(\frac{1}{n} \sum_{i=1}^n x_i \right) - \left(\frac{1}{n} \sum_{i=2}^{n+1} x_i \right) = \frac{1}{n} (x_1 - x_{n+1}) \\ \sim \mathcal{N}\left(0, \frac{2\sigma^2}{n}\right)$$

Want to take advantage of when data is not worst case.

Def The local sensitivity of f at x is

$$LS_f(x) \triangleq \max_{x'} |f(x) - f(x')| \\ \text{s.t. } x' \sim x$$

Q: Does $f(x) + \text{Lap}\left(\frac{LS_f(x)}{\epsilon}\right)$ give DP?

A: No! (median, 2 points on $\frac{\mathbb{R}}{2}$)

Def The down sensitivity of f at x is

$$DS_f(x) \triangleq \max_{x'} |f(x) - f(x')|$$

s.t. $x' \subseteq x$
 $|x'| = n-1$

Ex Mean of Gaussian data has small down-sensitivity, $LS = \infty$

Instance optimality, b/c for any pair $x \sim x'$ definitely must hide difference.

I know four techniques for beating global sensitivity

- ① Inverse sensitivity mechanism
- ② Privately estimate $LS_f(x)$
- ③ Smooth sensitivity
- ④ Propose - test - Release

Skipped today

If someone asks

$$SS_f^\epsilon(x) = \max_{x'} L_{S_f}(x') \cdot e^{-\epsilon d(x, x')}$$

Can do: $f(x) + \frac{SS_f^\epsilon(x)}{\epsilon} \cdot Z$ for ϵ -DP

$$f(x) + \frac{SS_f^\epsilon(x)}{\epsilon} \cdot Z e^{\sigma Y}$$

Cauchy

σ scale gives (ϵ, δ) -DP

Std Laplace

Std normal

Propose - Test - Release

Simplest form: "guess" that local sensitivity is small on x & all its neighbors.

We'll see more general views, as a wrapper around a "goal" algorithm.

Not obviously implementable on a computer, but many efficient examples.

Algorithm Propose-Test-Release

Input: $\epsilon, \delta \in (0, 1)$, $x \in \mathcal{X}^n$, $A: \mathcal{X}^n \rightarrow \mathcal{Y}$

Output: $\tilde{y} \in \mathcal{Y}$ or "FAIL"

- 1) $\text{SAFE} = \{x' \in \mathcal{X}^n : \forall x'' \sim x', A(x') \approx_{(\epsilon, \delta)} A(x'')\}$
- 2) $\text{UNSAFE} = \mathcal{X}^n \setminus \text{SAFE}$
- 3) $k \leftarrow \min_{x' \in \text{UNSAFE}} d(x, x')$
- 4) If $\tilde{k} = k + \text{Lap}(\frac{1}{\epsilon}) \leq \frac{\log(1/\delta)}{\epsilon} + 1$
- 5) Return "FAIL"
- 6) Else return $\tilde{y} = A(x)$

(Privacy proof)

Gaussian mean in d dimensions

Assume bound on \mathbb{R}

Iteratively refine it

For $x \sim \mathcal{N}(\mu, \mathbb{I})$, w.p. $1 - \beta$

$$\|x - \mu\|_2 \leq \sqrt{d} + \sqrt{\log \frac{2}{\beta}}$$