CS 540 Introduction to Artificial Intelligence Search III: Advanced Search (aka Optimization)

University of Wisconsin-Madison
Spring 2023

## Outline

Homeworks:

- Homework 10 due Thursday
- Course evaluation due Friday

Class roadmap:

| Tuesday, May 2 | Advanced Search |
| :--- | :--- |
| Thursday, May 4 | Ethics and Review |
|  |  |
| Friday, May 12 <br> $5: 05-7: 05 \mathrm{pm}$ | Final Exam |

## Advanced Search Overview

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Problem Setting

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How is a search problem defined?

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## Genetic <br> Algorithms

Fitness
Population
Cross-over
Mutation

What is difference between two?

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- Advanced Search \& Hill-climbing
- More difficult problems, basics, local optima, variations


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- Hill Climbing
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- More difficult problems, basics, local optima, variations
- Hill Climbing
- Basic algorithm, local optima
- Genetic Algorithms
- Basics of evolution, fitness, natural selection


## Search vs. Optimization

Before: wanted a path from start state to goal state

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- States $s$ have values $f(s)$

- Want: Find $s$ with optimal value $f(s)$ (i.e, optimize over states)
- Challenging settings: too many states for previous search approaches, but maybe not a differentiable function for gradient descent.


## Examples: n Queens

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- State: configuration of the board
- $f(s)$ : \# of non-conflicting queens



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- $f(s): \#$ satisfied clauses

| $R(x, a, d) \wedge R(y, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(z, c, 0)$ |
| :--- |
| $R(0, a, d) \wedge R(0, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(0, c, 0)$ |
| $R(0, a, d) \wedge R(0, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(1, c, 0)$ |
| $R(0, a, d) \wedge R(1, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(0, c, 0)$ |
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| $R R(, a, d) \wedge R(1, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(0, c, 0)$ |
| $R(1, a, d) \wedge R(1, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(1, c, 0)$ |


| $R(-x, a, b) \wedge R(b, y, c) \wedge R(c, d,-z)$ |
| :--- |
| $R(1, a, b) \wedge R(b, 0, c) \wedge R(c, d, 1)$ |
| $R(1, a, b) \wedge R(b, 0, c) \wedge R(c, d, 0)$ |
| $R(1, a, b) \wedge R(b, 1, c) \wedge R(c, d, 1)$ |
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| $R(0, a, b) \wedge R(b, 0, c) \wedge R(c, d, 0)$ |
| $R(0, a, b) \wedge R(b, 1, c) \wedge R$ |
| $R(0, a, b) \wedge R(b, 1, c) \wedge R(c, d, 1)$ |

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- As we'll see, needs a careful choice



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$$
\stackrel{S}{f(s)=6}
$$



$$
f=6
$$

Neighborhood of $s$

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$$
A-B-C-D-E-F-G-H-A
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$$
\begin{aligned}
& A \vee \neg B \vee C \\
& \neg A \vee C \vee D \\
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(A=T, B=T, C=T, D=T, E=T) & \neg A \vee C \vee D \\
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- Q: terminate? When no neighbor has better value

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Note the local optima. How do we handle them?


## Escaping Local Optima

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- Probability of selecting a neighbor should be proportional to the value of that neighbor.


## Hill Climbing: Variations

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## Break \& Quiz

Q 1.1: Hill climbing and stochastic gradient descent are related by
(i) Both head towards optima
(ii) Both require computing a gradient
(iii) Both will find the global optimum for a convex problem (problem where all optima have the same value).

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- B. (i), (ii)
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Q 2.2: Which of the following would be better to solve with hill climbing rather than A* search?
i. Finding the smallest set of vertices in a graph that involve all edges
ii. Finding the fastest way to schedule jobs with varying runtimes on machines with varying processing power
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- A. (i) (No, (ii) better: huge number of states, don't care about path)
- B. (ii) (No, (i) complete graph might have too many edges for A*)
- C. (i) and (ii)
- D. (ii) and (iii) (No, (iii) is good for A*: few successors, want path)


## Genetic Algorithms

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Optimization approach based on nature

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- Crossover: exchange between parents' codes
- Mutation: rarer random process
- Happens at individual level



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- New terminology: state is 'individual'
- Value $f(s)$ is now the 'fitness'



## Genetic Algorithms Setup I

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(3 275241 1)


## Genetic Algorithms Setup II



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Goal of genetic algorithms: optimize using principles inspired by mechanism for evolution

- Analogous to natural selection, cross-over, and mutation



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## Genetic Algorithms Pseudocode

## Just one variant:

1. Let $s_{1}, \ldots, s_{N}$ be the current population
2. Let $p_{i}=f\left(s_{i}\right) / \Sigma_{j} f\left(s_{j}\right)$ be the reproduction probability
3. for $k=1 ; k<N ; k+=2$

- parent1 = randomly pick according to $p$
- parent2 = randomly pick another
- randomly select a crossover point, swap strings of parents 1,2 to generate children $t[k], t[k+1]$

4. for $k=1 ; k<=N ; k++$

- Randomly mutate each position in $t[k]$ with a small probability (mutation rate)

5. The new generation replaces the old: $\{s\} \leftarrow\{t\}$. Repeat

## Reproduction: Proportional Selection

Reproduction probability: $p_{i}=f\left(s_{i}\right) / \Sigma_{j} f\left(s_{j}\right)$

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| B | 20 | $40 \%$ |
| C | 11 | $22 \%$ |
| D | 8 | $16 \%$ |
| E | 6 | $12 \%$ |

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- Example: $\Sigma_{j} f\left(s_{j}\right)=5+20+11+8+6=50$
- $p_{1}=5 / 50=10 \%$

| Individual | Fitness | Prob. |
| :--- | :--- | :--- |
| A | 5 | $10 \%$ |
| B | 20 | $40 \%$ |
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## Example: Scheduling Courses

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## Example: Scheduling Courses

Let's run through an example:

- 5 courses: A,B,C,D,E
- 3 time slots: Mon/Wed, Tue/Thu, Fri/Sat
- Students wish to enroll in three courses
- Goal: maximize student enrollment

| Courses | Students |
| :---: | :---: |
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## Example: Scheduling Courses

Let's run through an example:

- State: course assignment to time slot

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
| B C D | 4 |
| B D E | 10 |
| C D E | 5 |

## Example: Scheduling Courses

Let's run through an example:

- State: course assignment to time slot

| $M$ | $M$ | $F$ | $T$ | $M$ |
| :---: | :---: | :---: | :---: | :---: |
| $A$ | $B$ | $C$ | $D$ | $E$ |


| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
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| :---: | :---: | :---: | :---: | :---: |
| $A$ | $B$ | $C$ | $D$ | $E$ |
|  | $=M M F T M$ |  |  |  |


| Courses | Students |
| :---: | :---: |
| A B C | 2 |
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| $A$ | $B$ | $C$ | $D$ | $E$ |
| $=M M F T M$ |  |  |  |  |

- Here:

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
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## Example: Scheduling Courses

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- State: course assignment to time slot

| M | M | F | T | M |
| :---: | :---: | :---: | :---: | :---: |
| A | B | C | D | E |

- Here:
- Courses A, B, E scheduled Mon/Wed

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
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## Example: Scheduling Courses

Let's run through an example:

- State: course assignment to time slot

| $M$ | $M$ | $F$ | $T$ | $M$ |
| :---: | :---: | :---: | :---: | :---: |
| $A$ | $B$ | $C$ | $D$ | $E$ |
| $=M M F T M$ |  |  |  |  |

- Here:
- Courses A, B, E scheduled Mon/Wed
- Course D scheduled Tue/Thu

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
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| A D E | 3 |
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## Example: Scheduling Courses

Let's run through an example:

- State: course assignment to time slot

| $M$ | $M$ | $F$ | $T$ | $M$ |
| :---: | :---: | :---: | :---: | :---: |
| $A$ | $B$ | $C$ | $D$ | $E$ |

- Here:
- Courses A, B, E scheduled Mon/Wed
- Course D scheduled Tue/Thu

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
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| B D E | 10 |
| C D E | 5 |

- Course C scheduled Fri/Sat


## Example: Scheduling Courses

## Example: Scheduling Courses

| Courses | Students | Can enroll? |
| :---: | :---: | :---: |
| A B C | 2 | No |
| A B D | 7 | No |
| A D E | 3 | No |
| B C D | 4 | Yes |
| B D E | 10 | No |
| C D E | 5 | Yes |

## Example: Scheduling Courses

Value of a state? Say MMFTM

| Courses | Students | Can enroll? |
| :---: | :---: | :---: |
| A B C | 2 | No |
| A B D | 7 | No |
| A D E | 3 | No |
| B C D | 4 | Yes |
| B D E | 10 | No |
| C D E | 5 | Yes |

- Here 4+5=9 students can enroll in desired courses


## Example: Scheduling Courses

First step:

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
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## Example: Scheduling Courses

First step:

- Randomly initialize and evaluate states

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
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## Example: Scheduling Courses

First step:

- Randomly initialize and evaluate states

```
MMFTM = 9
TTFMM = 4
FMTTF = 19
MTTTF = 3
```

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
| B C D | 4 |
| B D E | 10 |
| C D E | 5 |

## Example: Scheduling Courses

First step:

- Randomly initialize and evaluate states

```
MMFTM = 9
TTFMM = 4
FMTTF = 19
MTTTF = 3
```

- Calculate reproduction probabilities

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
| B C D | 4 |
| B D E | 10 |
| C D E | 5 |

## Example: Scheduling Courses

## First step:

- Randomly initialize and evaluate states

| MMFTM $=9$ | MMFTM $=26 \%$ |
| :--- | :--- |
| TTFMM $=4$ | TTFMM $=11 \%$ |
| FMTTF $=19$ | FMTTF $=54 \%$ |
| MTTTF $=3$ | MTTTF $=9 \%$ |

- Calculate reproduction probabilities

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
| B C D | 4 |
| B D E | 10 |
| C D E | 5 |

## Example: Scheduling Courses

Next steps:

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Next steps:

- Select parents using reproduction probabilities


## Example: Scheduling Courses

## Next steps:

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## Example: Scheduling Courses

## Next steps:

- Select parents using reproduction probabilities
- Perform crossover



## Example: Scheduling Courses

## Next steps:

- Select parents using reproduction probabilities
- Perform crossover


| FMTTF | EMF'M |
| :---: | :---: |
| MMFTM | MMTTF |
| MTTTF | MMTTF |
| FMTTF | FTTTF |

## Example: Scheduling Courses

## Next steps:

- Select parents using reproduction probabilities
- Perform crossover
- Randomly mutate new children


| FMTTF | FMFTM |
| :---: | :---: |
| MMFTM | MMTTF |
| MTTTF | MMTTF |
| FMTTF | FTTTF |

## Example: Scheduling Courses

## Next steps:

- Select parents using reproduction probabilities
- Perform crossover
- Randomly mutate new children
MMFTM $=26 \%$
TTFMM $=11 \%$
MTTTF $=54 \%$

| FMTTF | FMFTM | FMFTM | FMFTT |
| :---: | :---: | :---: | :---: |
| MMFTM | MMTTF | MMTTF | MMTTF |
| MTTTF | MMTTF | MMTTF | MMTFF |
| FMTTF | FTTTF | FTTTF | FTTTF |

## Example: Scheduling Courses

Continue:

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
| B C D | 4 |
| B D E | 10 |
| C D E | 5 |

## Example: Scheduling Courses

## Continue:

- Now, get our function values for updated population

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
| B C D | 4 |
| B D E | 10 |
| C D E | 5 |

## Example: Scheduling Courses

## Continue:

- Now, get our function values for updated population

$$
\begin{aligned}
& \mathrm{FMFTT}=11 \\
& \mathrm{MMTTF}=13 \\
& \mathrm{MMTFF}=4 \\
& \mathrm{FTTPF}=0
\end{aligned}
$$

| Courses | Students |
| :---: | :---: |
| A B C | 2 |
| A B D | 7 |
| A D E | 3 |
| B C D | 4 |
| B D E | 10 |
| C D E | 5 |

## Example: Scheduling Courses

## Continue:

- Now, get our function values for updated population
- Calculate reproduction probabilities

$$
\begin{aligned}
& \text { FMFTT }=11 \\
& \text { MMTTF }=13 \\
& \text { MMTFF }=4 \\
& \text { FTTTF }=0
\end{aligned}
$$

| Courses | Students |
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## Example: Scheduling Courses

## Continue:

- Now, get our function values for updated population
- Calculate reproduction probabilities

$$
\begin{array}{ll}
\text { FMFTT }=11 & \text { FMFTT }=39 \% \\
\text { MMTTF }=13 & \text { MMTTF }=46 \% \\
\text { MMTFF }=4 & \text { MMTFF }=14 \% \\
\text { FTTTF }=0 & \text { FTTTF }=0 \%
\end{array}
$$

| Courses | Students |
| :---: | :---: |
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## Variations \& Concerns

Many possibilities:

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## Some challenges

- Formulating a good state encoding
- Lack of diversity: converge too soon



## Variations \& Concerns

## Many possibilities:

- Parents survive to next generation
- Use ranking instead of exact value of $f(s)$ for reproduction probabilities (reduce influence of extreme $f$ values)


## Some challenges

- Formulating a good state encoding
- Lack of diversity: converge too soon
- Must pick a lot of parameters



## Summary

- Challenging optimization problems
- First, try hill climbing. Simplest solution
- Genetic algorithms
- Biology-inspired optimization routine

