



# CS 540 Introduction to Artificial Intelligence

## **Search III: Advanced Search**

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# Announcements

- **Homeworks:**
  - HW8 due Tuesday
- **Class roadmap:**

Tuesday, Nov 16	Search II: Informed search		HW7 Due; HW8 Released
Thursday, Nov 18	Advanced Search and Review on Search		
<b>Everything below here is tentative and subject to change.</b>			
Tuesday, Nov 23	Games - Part I		HW8 Due; HW9 Released
Thursday, Nov 25	Happy Thanksgiving! (No class)		
Tuesday, Nov 30	Games - Part II		
Thursday, Dec 2	Reinforcement Learning I		HW9 Due; HW10 Released
Tuesday, Dec 7	Reinforcement Learning II		
Thursday, Dec 9	Review on Games and Reinforcement Learning		
Tuesday, Dec 14	Ethics and Trust in AI		HW10 Due

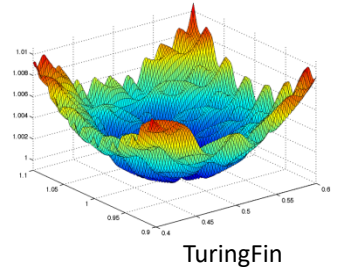
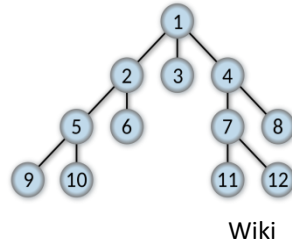
# Outline

- Advanced Search & Hill-climbing
  - More difficult problems, basics, local optima, variations
- Simulated Annealing
  - Basic algorithm, temperature, tradeoffs
- Genetic Algorithms
  - Basics of evolution, fitness, natural selection

# Search vs. Optimization

Before: wanted a **path** from start state to goal state

- Uninformed search, informed search



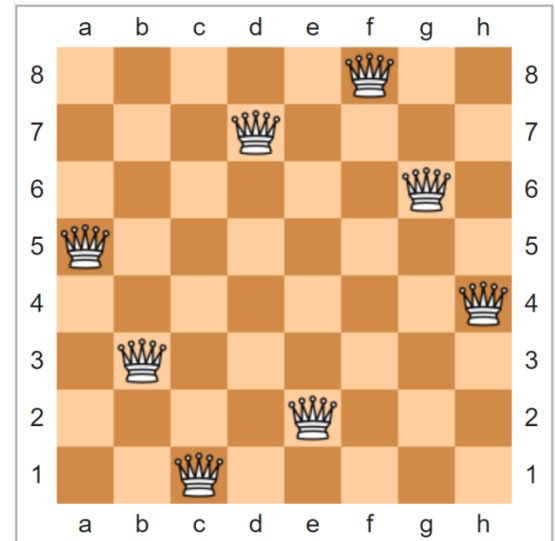
**New setting:** optimization

- States  $s$  have values  $f(s)$
- Want:  $s$  with optimal value  $f(s)$  (i.e, **optimize** over states)
- Challenging setting: **too many states** for previous search approaches, but maybe not a continuous function for SGD.

# Examples: $n$ Queens

A classic puzzle:

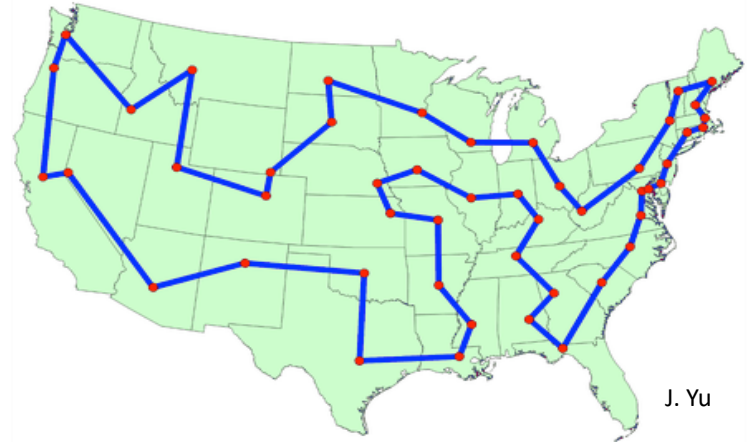
- Place 8 queens on a 8 x 8 chessboard so that no two have same row, column, or diagonal.
- Can generalize to  $n \times n$  chessboard.
- What are states  $s$ ? Values  $f(s)$ ?
  - State: configuration of the board
  - $f(s)$ : # of conflicting queens



# Examples: TSP

Famous graph theory problem.

- Get a graph  $G = (V, E)$ . **Goal:** a path that visits each node exactly once and returns to the initial node (a **tour**).
  - State: a particular tour (i.e., ordered list of nodes)
  - $f(s)$ : total weight of the tour (e.g., total miles traveled)



# Examples: Satisfiability

## Boolean satisfiability (e.g., 3-SAT)

- Recall our logic lecture. Conjunctive normal form

$$(A \vee \neg B \vee C) \wedge (\neg A \vee C \vee D) \wedge (B \vee D \vee \neg E) \wedge (\neg C \vee \neg D \vee \neg E) \wedge (\neg A \vee \neg C \vee E)$$

- Goal: find if satisfactory assignment exists.
- State: assignment to variables

–  $f(s)$ : # satisfied clauses

$$\begin{array}{l}
 \overline{R(x,a,d) \wedge R(y,b,d) \wedge R(a,b,e) \wedge R(c,d,f) \wedge R(z,c,0)} \\
 R(0,a,d) \wedge R(0,b,d) \wedge R(a,b,e) \wedge R(c,d,f) \wedge R(0,c,0) \\
 R(0,a,d) \wedge R(0,b,d) \wedge R(a,b,e) \wedge R(c,d,f) \wedge R(1,c,0) \\
 R(0,a,d) \wedge R(1,b,d) \wedge R(a,b,e) \wedge R(c,d,f) \wedge R(0,c,0) \\
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 \end{array}$$

$$\begin{array}{l}
 \overline{R(\neg x,a,b) \wedge R(b,y,c) \wedge R(c,d,\neg z)} \\
 R(1,a,b) \wedge R(b,0,c) \wedge R(c,d,1) \\
 R(1,a,b) \wedge R(b,0,c) \wedge R(c,d,0) \\
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 \end{array}$$

# Hill Climbing

One approach to such optimization problems.

- Basic idea: move to a neighbor with a better  $f(s)$
- **Q:** how do we define **neighbor**?
  - Not as obvious as our successors in search
  - Problem-specific
  - As we'll see, needs a careful choice

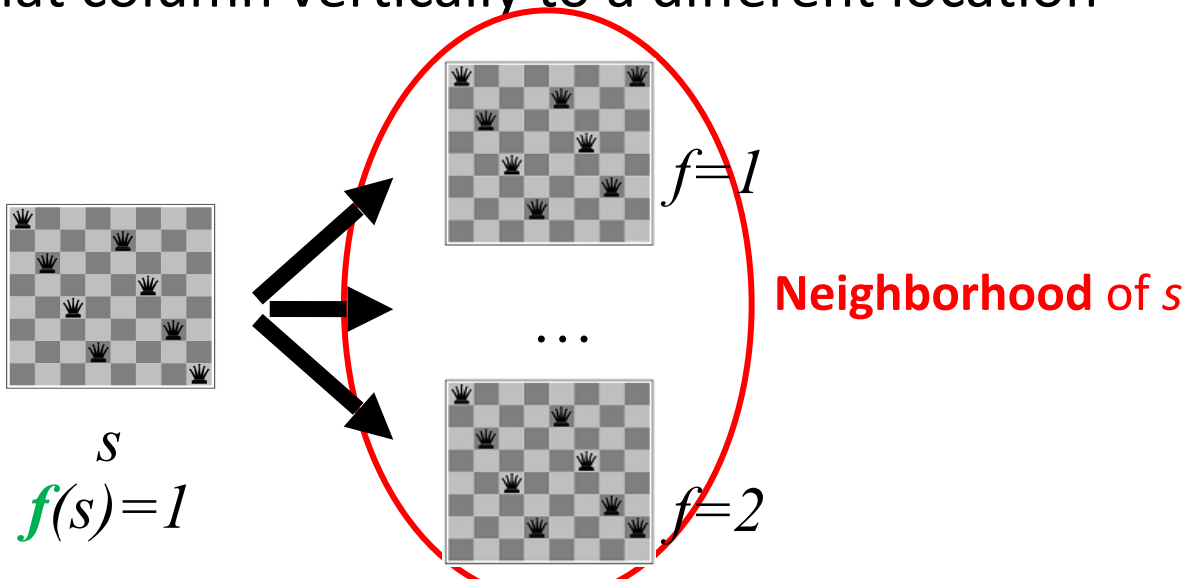




# Defining Neighbors: n Queens

In n Queens, a simple possibility:

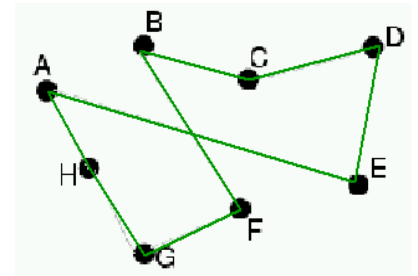
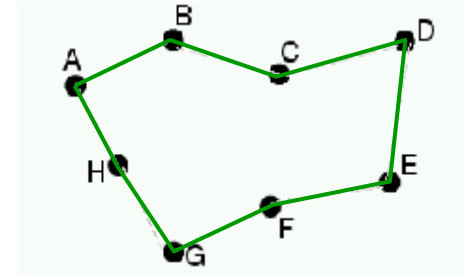
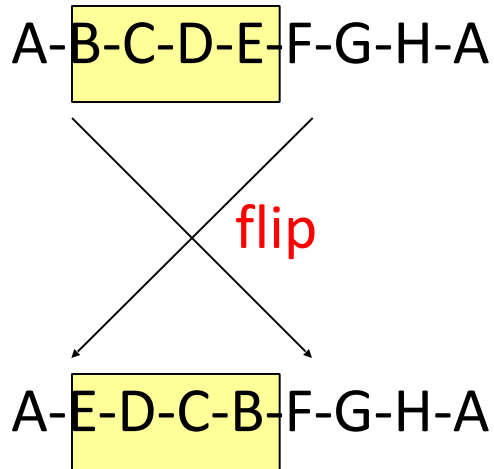
- Look at the **most-conflicting column** (ties? right-most one)
- Move queen in that column vertically to a different location



# Defining Neighbors: TSP

For TSP, can do something similar:

- Define neighbors by small changes
- Example: 2-change: A-E and B-F



# Defining Neighbors: SAT

For Boolean satisfiability,

- Define neighbors by flipping one assignment of one variable

Starting state: TFTTT

(A=**F**, B=F, C=T, D=T, E=T)

(A=T, B=**T**, C=T, D=T, E=T)

(A=T, B=F, C=**F**, D=T, E=T)

(A=T, B=F, C=T, D=**F**, E=T)

(A=T, B=F, C=T, D=T, E=**F**)

$A \vee \neg B \vee C$

$\neg A \vee C \vee D$

$B \vee D \vee \neg E$

$\neg C \vee \neg D \vee \neg E$

$\neg A \vee \neg C \vee E$

# Hill Climbing Neighbors

## Q: What's a **neighbor**?

- **Vague definition.** For a given problem structure, neighbors are states that can be produced by a small change
- **Tradeoff!**
  - Too small? Will get stuck.
  - Too big? Not very efficient
- **Q:** how to pick a neighbor? Greedy
- **Q:** terminate? When no neighbor has bigger value



# Hill Climbing Algorithm

## Pseudocode:

1. Pick initial state  $s$
2. Pick  $t$  in **neighbors**( $s$ ) with the largest  $f(t)$
3. if  $f(t) \leq f(s)$  THEN stop, return  $s$
4.  $s \leftarrow t$ . goto 2.

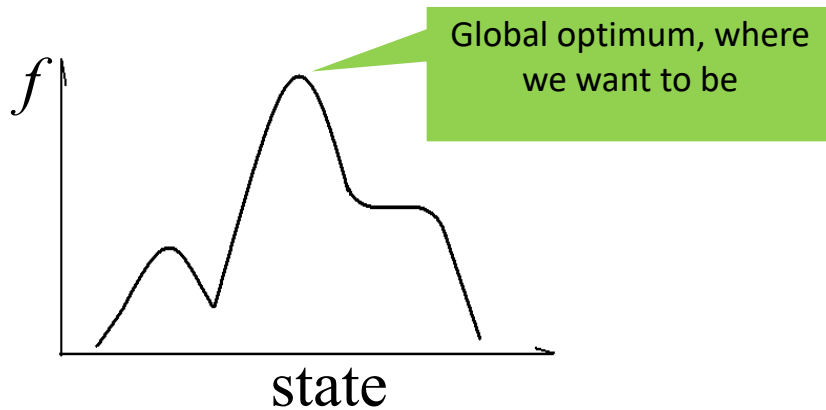
Maximization problem;  
flip for minimization.

What could happen? **Local optima!**

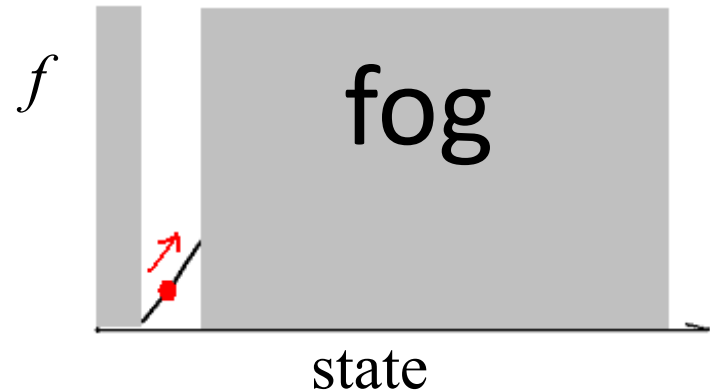


# Hill Climbing: Local Optima

Q: Why is it called hill climbing?



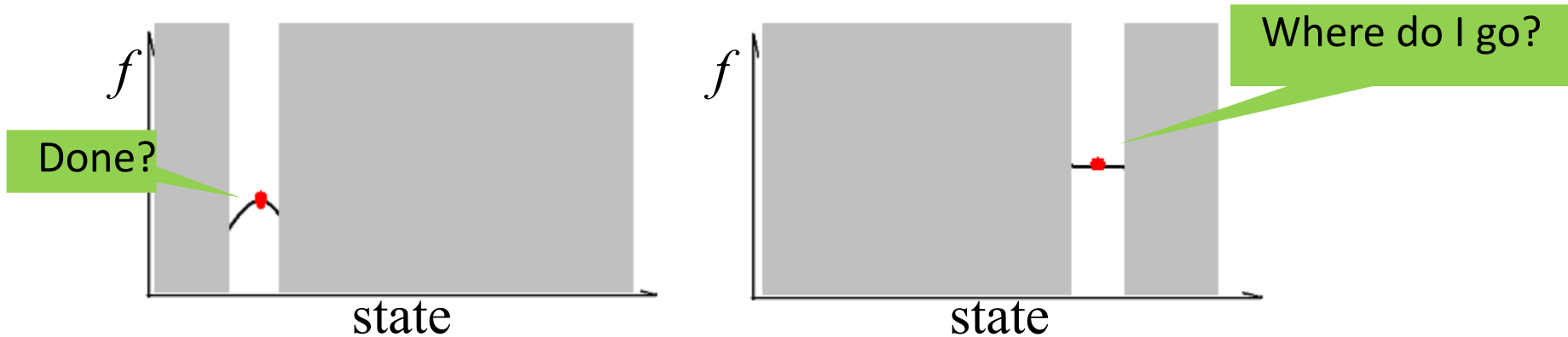
Left: What's actually going on.



Right: What we get to see.

# Hill Climbing: Local Optima

Note the **local optima**. How do we handle them?

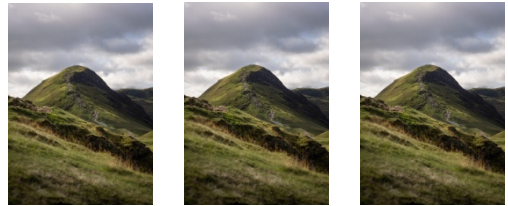


# Escaping Local Optima

## Simple idea 1: random restarts

- Stuck: pick a random new starting point, re-run.
- Do  $k$  times, return best of the  $k$ .

Might get same optima each time.



## Simple idea 2: reduce greed

- “Stochastic” hill climbing: randomly select between neighbors
- Probability proportional to the value of neighbors



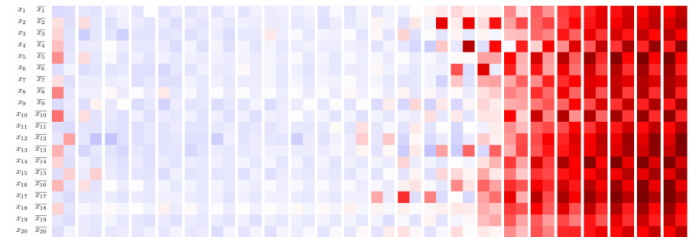
# Hill Climbing: Variations

**Q:** neighborhood too large?

- Generate random neighbors, **one at a time**. Take the better one.

**Q:** relax requirement to always go up?

- Often useful for harder problems
- 3SAT algorithm: Walk-SAT



# Break & Quiz

**Q 1.1:** Hill climbing and SGD are related by

- (i) Both head towards optima
- (ii) Both require computing a gradient
- (iii) Both will find the global optimum for a convex problem

- A. (i)
- B. (i), (ii)
- C. (i), (iii)
- D. All of the above

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# Break & Quiz

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- (i) Both head towards optima
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- A. (i) (No: (iii) also true since convexity->local optima are global)
- B. (i), (ii) (No: (ii) is false. Hill-climbing looks at neighbors only.)
- C. (i), (iii)
- D. All of the above (No: (ii) false, as above.)

# Simulated Annealing

A more sophisticated optimization approach.

- **Idea:** move quickly at first, then slow down
- Pseudocode:

Pick initial state  $s$

For  $k = 0$  through  $k_{\max}$ :

$T \leftarrow \text{temperature}( (k+1)/k_{\max} )$

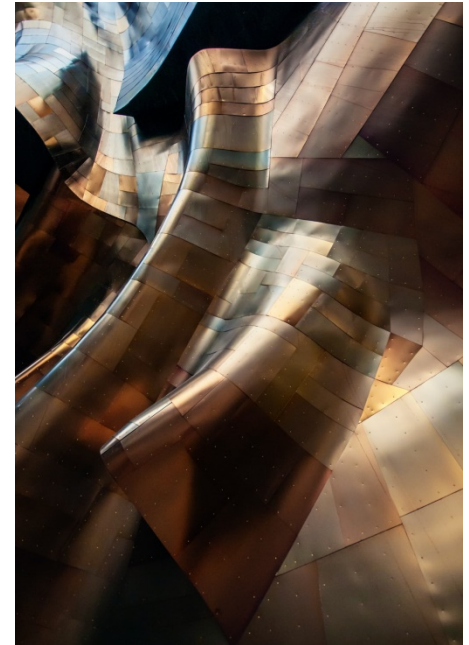
Pick a random neighbour,  $t \leftarrow \text{neighbor}(s)$

If  $f(s) \leq f(t)$ , then  $s \leftarrow t$

Else, with prob.  $P(f(s), f(t), T)$  then  $s \leftarrow t$

**Output:** the final state  $s$

The interesting bit



# Simulated Annealing: Picking Probability

How do we pick probability  $P$ ? Note 3 parameters.

- Decrease with time
- Decrease with gap  $|f(s) - f(t)|$

Pick initial state  $s$

For  $k = 0$  through  $k_{\max}$ :

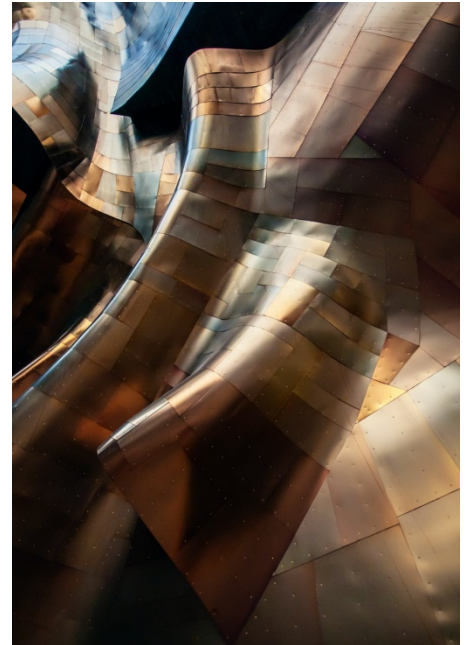
$T \leftarrow \text{temperature}( (k+1)/k_{\max} )$

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Else, with prob.  $P(f(s), f(t), T)$  then  $s \leftarrow t$

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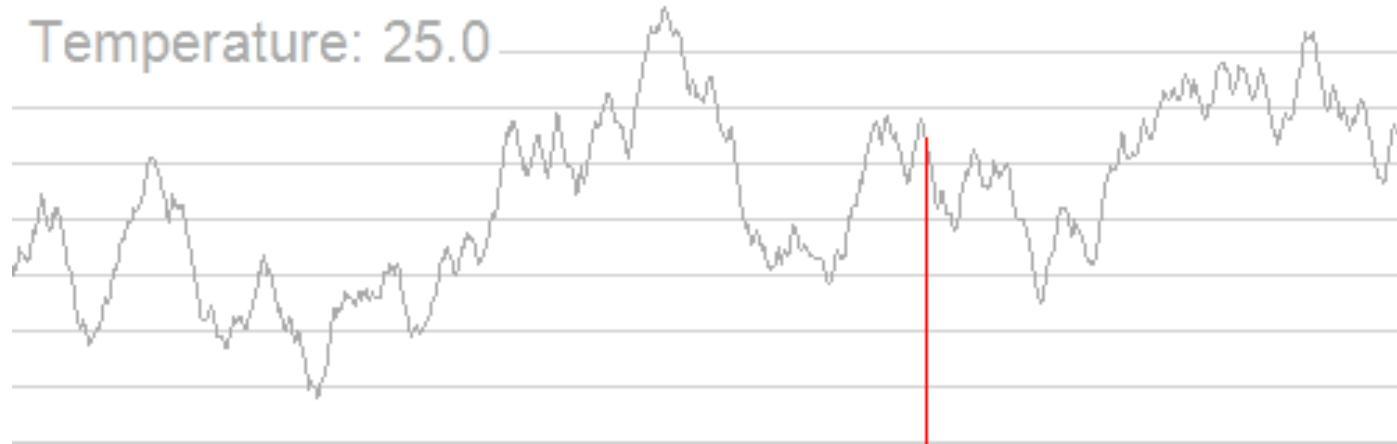
# Simulated Annealing: Picking Probability

How do we pick probability P? Note 3 parameters.

- Decrease with time
- Decrease with gap  $|f(s) - f(t)|$ :  $\exp\left(-\frac{|f(s) - f(t)|}{Temp}\right)$
- Temperature cools over time.
  - So: high temperature, accept any  $t$
  - But, low temperature, behaves like hill-climbing
  - Still,  $|f(s) - f(t)|$  plays a role: if big, replacement probability low.

# Simulated Annealing: Visualization

What does it look like in practice?





# Simulated Annealing: Picking Parameters

- Have to balance the various parts., e.g., cooling schedule.
  - Too fast: becomes hill climbing, stuck in local optima
  - Too slow: takes too long.
- Combines with variations (e.g., with random restarts)
  - Probably should try hill-climbing first though.
- Inspired by cooling of metals
  - We'll see one more alg. inspired by nature



# Break & Quiz

**Q 2.1:** Which of the following is likely to give the best cooling schedule for simulated annealing?

- A.  $\text{Temp}_{t+1} = \text{Temp}_t * 1.25$
- B.  $\text{Temp}_{t+1} = \text{Temp}_t$
- C.  $\text{Temp}_{t+1} = \text{Temp}_t * 0.8$
- D.  $\text{Temp}_{t+1} = \text{Temp}_t * 0.0001$

# Break & Quiz

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- C.  $\text{Temp}_{t+1} = \text{Temp}_t * 0.8$**
- D.  $\text{Temp}_{t+1} = \text{Temp}_t * 0.0001$

# Break & Quiz

**Q 2.1:** Which of the following is likely to give the best cooling schedule for simulated annealing?

- A.  $\text{Temp}_{t+1} = \text{Temp}_t * 1.25$  (No, temperate is increasing)
- B.  $\text{Temp}_{t+1} = \text{Temp}_t$  (No, temperature is constant)
- C.  $\text{Temp}_{t+1} = \text{Temp}_t * 0.8$
- D.  $\text{Temp}_{t+1} = \text{Temp}_t * 0.0001$  (Cools too fast---basically hill climbing)

# Break & Quiz

**Q 2.2:** Which of the following would be better to solve with simulated annealing than A\* search?

- i. Finding the smallest set of vertices in a graph that involve all edges
- ii. Finding the fastest way to schedule jobs with varying runtimes on machines with varying processing power
- iii. Finding the fastest way through a maze

- A. (i)
- B. (ii)
- C. (i) and (ii)
- D. (ii) and (iii)

# Break & Quiz

**Q 2.2:** Which of the following would be better to solve with simulated annealing than A\* search?

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- **C. (i) and (ii)**
- D. (ii) and (iii)

# Break & Quiz

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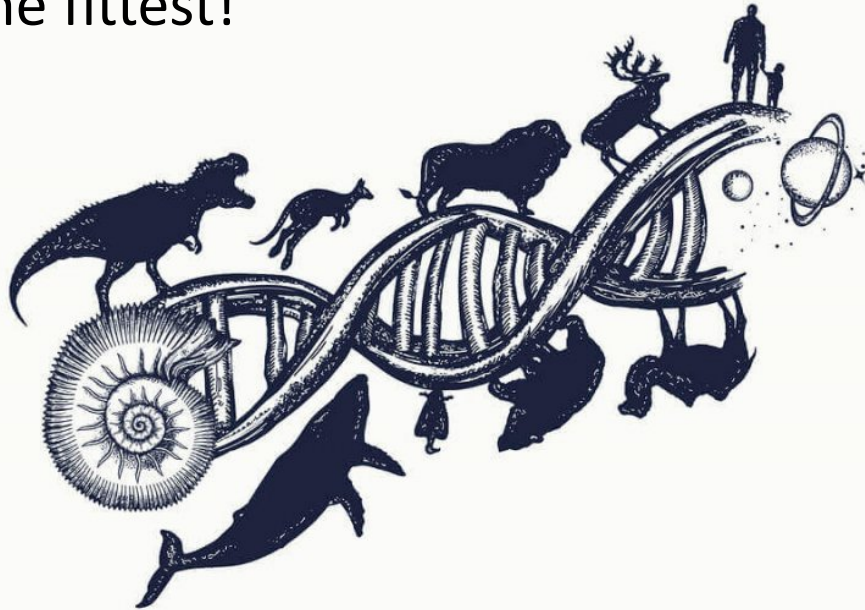
- i. Finding the smallest set of vertices in a complete graph (i.e., all nodes connected)
- ii. Finding the fastest way to schedule jobs with varying runtimes on machines with varying processing power
- iii. Finding the fastest way through a maze

- A. (i) (No, (ii) better: huge number of states, don't care about path)
- B. (ii) (No, (i) complete graph might have too many edges for A\*)
- C. (i) and (ii)
- D. (ii) and (iii) (No, (iii) is good for A\*: few successors, want path)

# Genetic Algorithms

Another optimization approach based on nature

- Survival of the fittest!

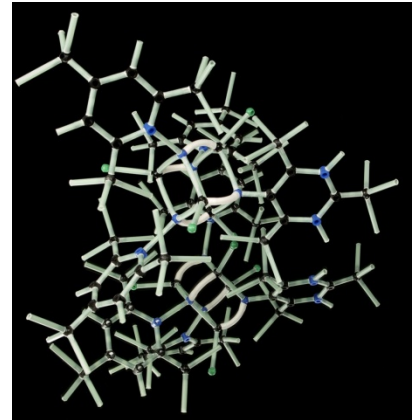




# Evolution Review

Encode genetic information in DNA (four bases)

- A/C/T/G: nucleobases acting as symbols ACCCATGT
- Two types of changes
  - Crossover: exchange between parents' codes
  - Mutation: rarer random process
    - Happens at individual level



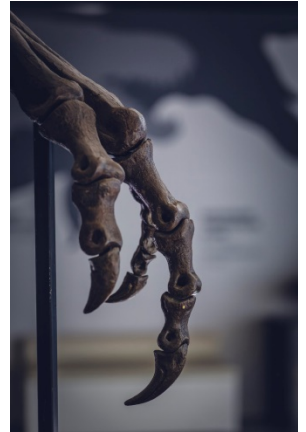
# Natural Selection

## Competition for resources

- Organisms better fit → better probability of reproducing
- Repeated process: fit become larger proportion of population

## Goal: use these principles for optimization

- New terminology: state  $s$  ‘**individual**’
- Value  $f(s)$  is now the ‘**fitness**’

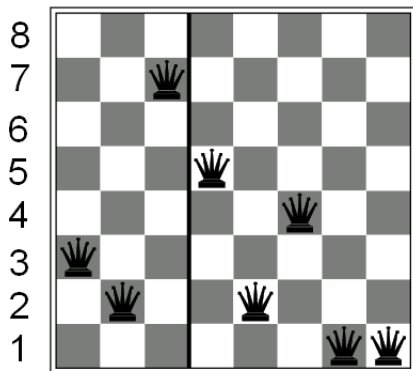


# Genetic Algorithms Setup I

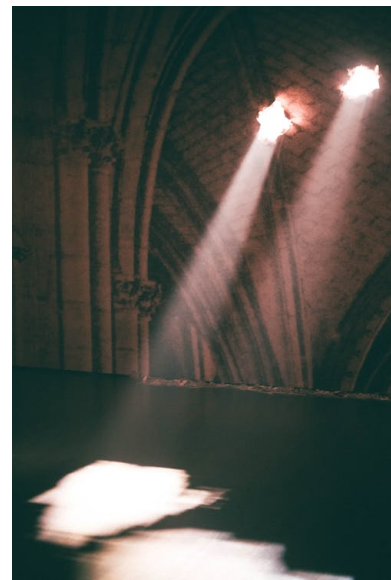
Keep around a fixed number of states/individuals

- A bit like beam search
- Call this the **population**

For our n Queens game example, an individual:



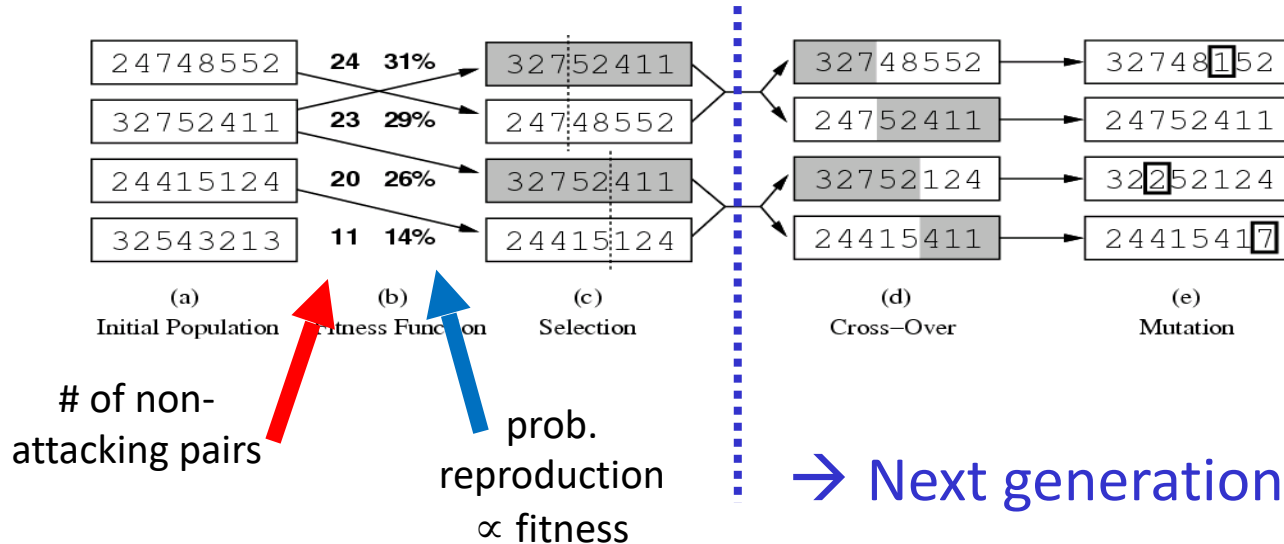
(3 2 7 5 2 4 1 1)



# Genetic Algorithms Setup II

Goal of genetic algorithms: optimize using principles inspired by mechanism for evolution

- E.g., analogous to **natural selection, cross-over, and mutation**



# Genetic Algorithms Pseudocode

Just one variant:

1. Let  $s_1, \dots, s_N$  be the current population
2. Let  $p_i = f(s_i) / \sum_j f(s_j)$  be the reproduction probability
3. for  $k = 1; k < N; k += 2$ 
  - parent1 = randomly pick according to  $p$
  - parent2 = randomly pick another
  - randomly select a crossover point, swap strings of parents 1, 2 to generate children  $t[k], t[k+1]$
4. for  $k = 1; k \leq N; k++$ 
  - Randomly mutate each position in  $t[k]$  with a small probability (mutation rate)
5. The new generation replaces the old:  $\{s\} \leftarrow \{t\}$ . Repeat

# Reproduction: Proportional Selection

Reproduction probability:  $p_i = f(s_i) / \sum_j f(s_j)$

- **Example:**  $\sum_j f(s_j) = 5+20+11+8+6=50$
- $p_1=5/50=10\%$

Individual	Fitness	Prob.
A	5	10%
B	20	40%
C	11	22%
D	8	16%
E	6	12%



# Example: Scheduling Courses

Let's run through an example:

- **5 courses: A,B,C,D,E**
- *3 time slots: Mon/Wed, Tue/Thu, Fri/Sat*
- Students wish to enroll in three courses
- Goal: maximize student enrollment

Courses	Students
A B C	2
A B D	7
A D E	3
B C D	4
B D E	10
C D E	5

# Example: Scheduling Courses

Let's run through an example:

- State: course assignment to time slot

M	M	F	T	M
A	B	C	D	E

= MMF'TM

- Here:
  - Courses A, B, E scheduled Mon/Wed
  - Course D scheduled Tue/Thu
  - Course C scheduled Fri/Sat

Courses	Students
A B C	2
A B D	7
A D E	3
B C D	4
B D E	10
C D E	5



# Example: Scheduling Courses

Value of a state? Say MMFTM

Courses	Students	Can enroll?
A B C	2	No
A B D	7	No
A D E	3	No
B C D	4	Yes
B D E	10	No
C D E	5	Yes

- Here  $4+5=9$  students can enroll in desired courses

# Example: Scheduling Courses

First step:

- Randomly initialize and evaluate states

MMFTM = 9

MMFTM = 26%

TTFMM = 4

TTFMM = 11%

FMTTF = 19

FMTTF = 54%

MTTTF = 3

MTTTF = 9%

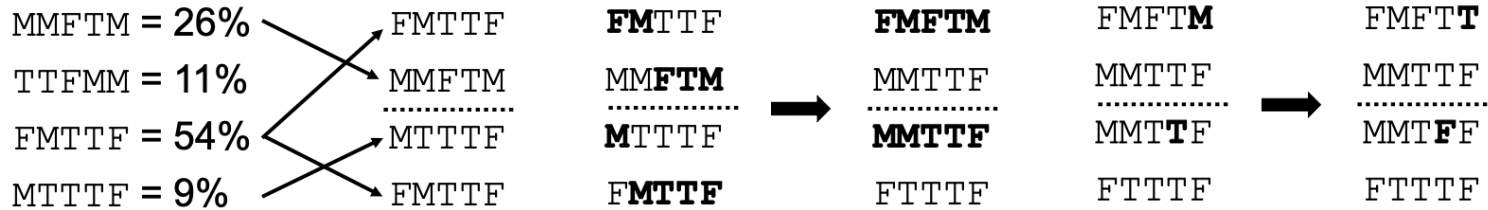
- Calculate reproduction probabilities

Courses	Students
A B C	2
A B D	7
A D E	3
B C D	4
B D E	10
C D E	5

# Example: Scheduling Courses

Next steps:

- Select parents using reproduction probabilities
- Perform crossover
- Randomly mutate new children



# Example: Scheduling Courses

Continue:

- Now, get our function values for updated population
- Calculate reproduction probabilities

$$FMFTT = 11 \quad FMFTT = 39\%$$

$$MMTTF = 13 \quad MMTTF = 46\%$$

$$MMTFF = 4 \quad MMTFF = 14\%$$

$$FTTTF = 0 \quad FTTTF = 0\%$$

Courses	Students
A B C	2
A B D	7
A D E	3
B C D	4
B D E	10
C D E	5

# Variations & Concerns

Many **possibilities**:

- Parents survive to next generation
- Ranking instead of exact value of  $f(s)$  for reproduction probabilities

Some **challenges**

- State encoding
- Lack of diversity: converge too soon
- Must pick a lot of parameters



# Summary

- Challenging optimization problems
  - First, try hill climbing. Simplest solution
- Simulated annealing
  - More sophisticated approach; helps with local optima
- Genetic algorithms
  - Biology-inspired optimization routine

# Search Summary

- Uninformed Search
  - Find path from initial to goal state.
  - Know cost to current state and successors
  - BFS, UCS, DFS, IDS
- Informed Search
  - Use heuristic to estimate how close states are to goal; focus on more promising states.
  - A\* Search is uses an admissible heuristic to guarantee completeness and optimality.
- Optimization / Advanced Search
  - States have values and many neighbors.
  - Just interested in final state, not path to goal.
  - Hill-climbing, simulated annealing, and genetic algorithms.



**Acknowledgements:** Adapted from materials by Jerry Zhu, Tony Gitter, Fred Sala (University of Wisconsin), Andrew Moore