Finite-state machines

CS 536
Last time

A compiler is a recognizer of language $S$ (Source) a translator from $S$ to $T$ (Target) a program in language $H$ (Host)

For example, gcc: $S$ is C, $T$ is x86, $H$ is C
Why do we need a compiler?

- Processors can execute only binaries (machine-code/assembly programs)
- Writing assembly programs will make you want to reconsider your life choices
- Write programs in a nice(ish) high-level language like Java; compile to binaries
Last time
The scanner

Translates sequence of chars into sequence of tokens

Each time scanner is called it should:

find longest sequence of chars corresponding to a token

return that token
Special linkage between scanner and parser in most compilers

Conceptual organization
Scanner generator

Generates a scanner!!!

Needs one regular expression for each token

Needs regular expressions for things to ignore

Comments, whitespace, etc.

To understand how it works, we need FSMs

Finite state machines
FSMs: Finite State Machines

Aka finite automata

**Input:** string (seq of chars)

**Output:** accept / reject

i.e., input is legal in language
FSMs

Represent regular languages
Good enough for tokens in PLs
Example 1

single line comments with //

1 -> 2

2 -> 3

3 -> 4
Example 2

What language does this accept?

Can you find an equivalent, but smaller, FSM?
How an FSM works

curr_state = start_state

let in_ch = current input char

repeat

if there is edge out of curr_state with label in_ch into next_state

cur_state = next_state

in_ch = next char of input

o/w stuck // error condition

until stuck or input string is consumed

string is accepted iff entire string is consumed and final_states.contains(cur_state)
FSMs, formally

\[(Q, \Sigma, \delta, q, F)\]

- finite set of states
- the alphabet (characters)
- transition function \(\delta : Q \times \Sigma \rightarrow Q\)
- start state \(q \in Q\)
- final states \(F \subseteq Q\)
FSMs, formally

\[(Q, \Sigma, \delta, q, F)\]

FSM accepts string

\[x_1x_2x_3 \ldots x_n\]

\[\iff\]

\[\delta(\ldots \delta(\delta(q, x_1), x_2), x_3) \ldots , x_n) \in F\]

The language of FSM \(M\) is the set of all words it accepts, denoted \(L(M)\)
FSM example, formally

$$(Q, \Sigma, \delta, q, F)$$

\[Q = \{s_0, s_1\}\]
\[\Sigma = \{a, b, c\}\]
\[q = s_0\]
\[F = \{s_0\}\]
\[\delta = s_0, a \rightarrow s_1\]
\[s_1, b \rightarrow s_0\] anything else, machine is stuck
Coding an FSM

curr_state = start_state

done = false

while (!done)

ch = nextChar()

next = transition[curr_state][ch]

if (next == error || ch == EOF)

done = true

else

    curr_state = next

return final_states.contains(curr_state) &&

    next!=error
FSM types: DFA & NFA

Deterministic
no state has > 1 outgoing edge with same label

Nondeterministic
states may have multiple outgoing edges with same label
edges may be labelled with special symbol $\varepsilon$ (empty string)
$\varepsilon$-transitions can happen without reading input
NFA example

Equivalent DFA
Why NFA?

Much more compact

What does this accept?
An equivalent DFA needs $2^5$ states
Extra example

Hex literals
must start with 0x or 0X
followed by at least one hex digit (0-9,a-f,A-F)
can optionally have long specifier (l,L) at the end
A C/C++ identifier is a sequence of one or more letters, digits, or underscores. It cannot start with a digit.
Extra Example - Part 1

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What if you wanted to add the restriction that it can't end with an underscore?
Extra Example - Part 2

What if you wanted to add the restriction that it can't end with an underscore?
Recap

The scanner reads stream of characters and finds tokens

Tokens are defined using regular expressions, which are finite-state machines

Finite-state machines can be non-deterministic

Next time: understand connection between deterministic and non-deterministic FSMs