Introduction to Computer Networks

TCP Connection Management (I)

https://pages.cs.wisc.edu/~mgliu/CS640/F22/

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Today

Last lecture

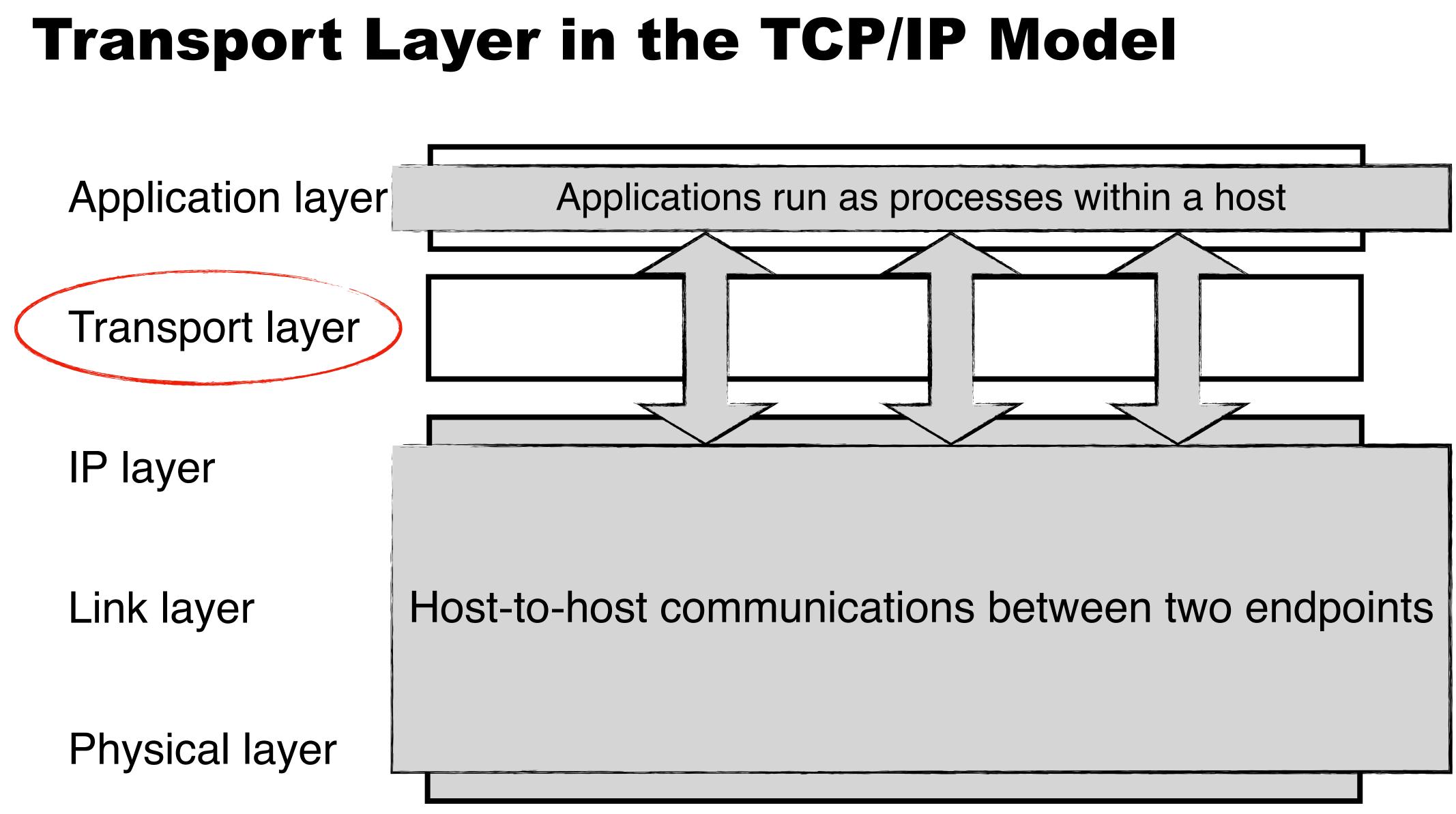
What functionalities does the transport layer provide?

Today

How to setup the TCP connection?

Announcements

• Lab4 is due 12/02/2022, 11:59 PM





Q: What functionalities does the transport layer provide?

A: Process-to-process communication channels

Q1: How to set up the process-to-process channel? Q2: How to multiplex concurrent channels over the physical link? Q3: How to control the transmission rate? Q4: How to achieve reliable delivery? Q5: How to share the in-network bandwidth resources?

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A: Process-to-process communication channels

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UDP Issues

#1: Arbitrary communication

Senders and receivers can talk to each other in any ways

#2: No reliability guarantee

- Packets can be lost/duplicated/reordered during transmission
- Checksum is not enough

#3: No resource management

- Each communication channel works as an exclusive network resource owner
- No adaptiveness support for the physical networks and applications



Q: What is the goal of TCP connection management?



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Client <--> Server



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On-demand communication

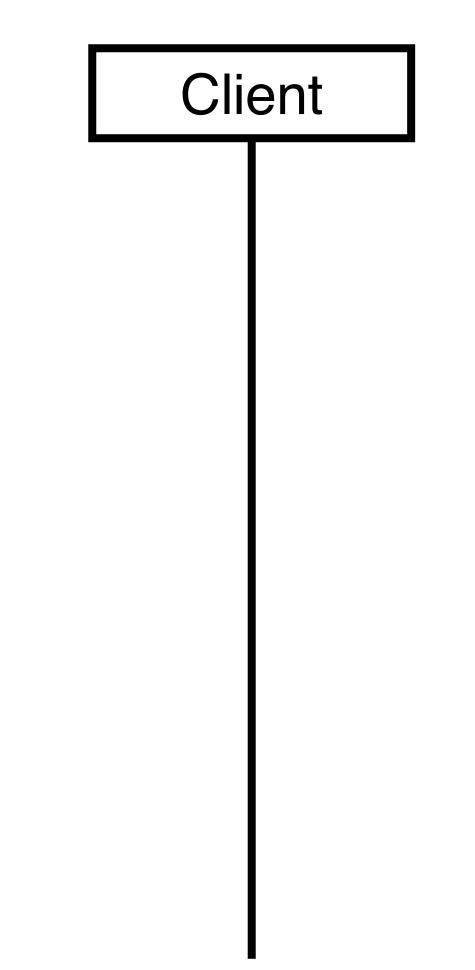
Client and server agree on the start of byte steams for two directions

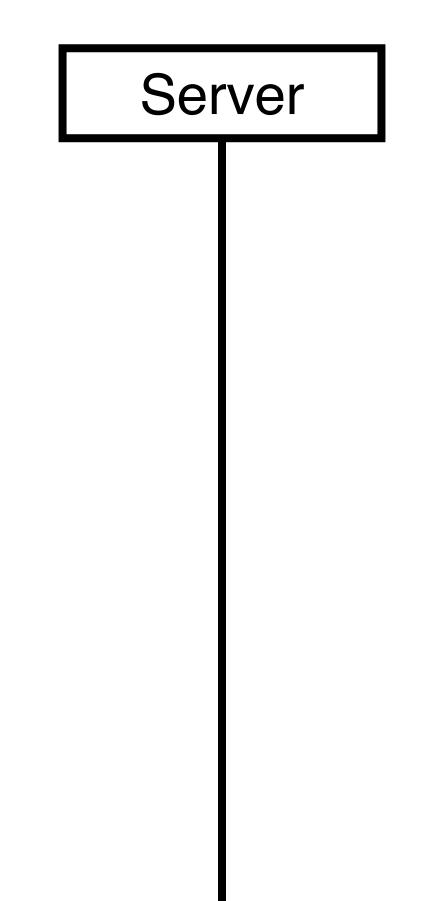
Client <--> Server



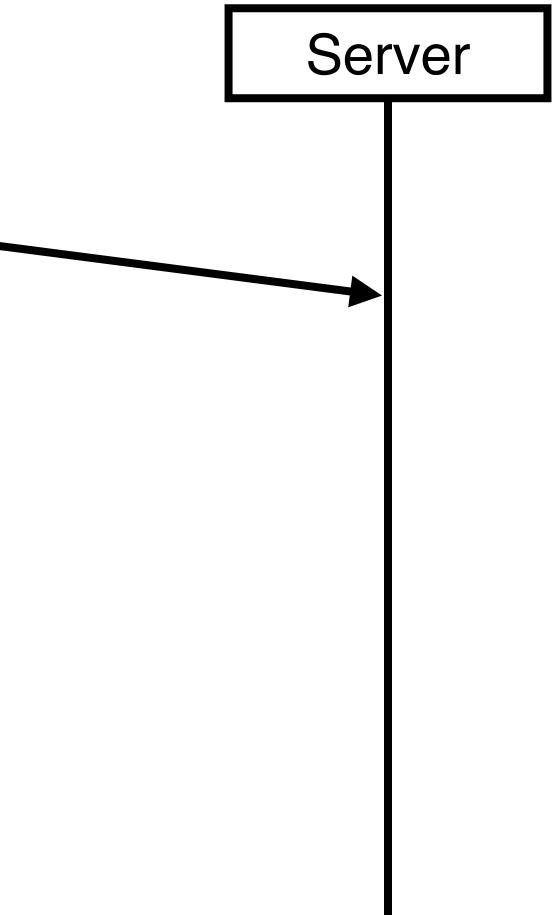
TCP Connection Establishment

Let's start with a naive approach





TCP Connection Establishment Let's start with a naive approach Client My (client) byte stream starts with a sequence number = X



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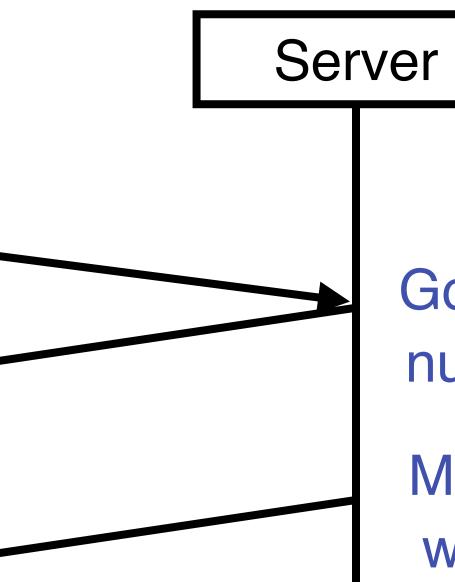




Got it, I acknowledge the sequence number of your next byte is = X + 1



TCP Connection Establishment Let's start with a naive approach Client My (client) byte stream starts with a sequence number = X



Got it, I acknowledge the sequence number of your next byte is = X + 1

My (server) byte stream starts with a sequence number = Y

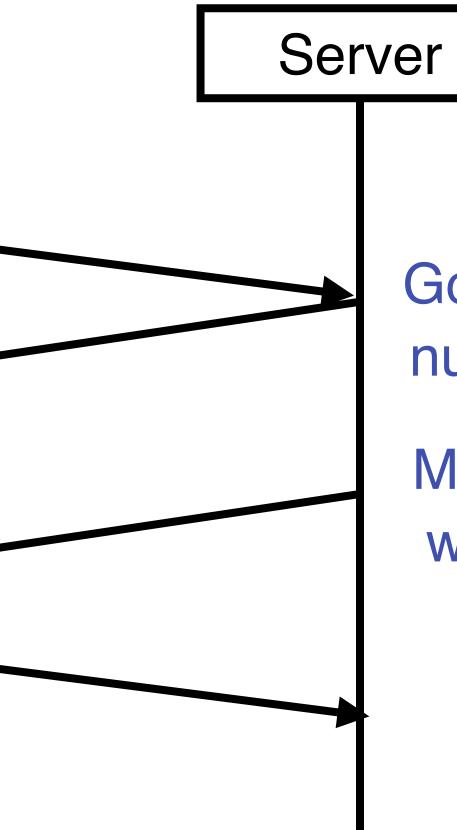


TCP Connection Establishment Let's start with a naive approach

Client

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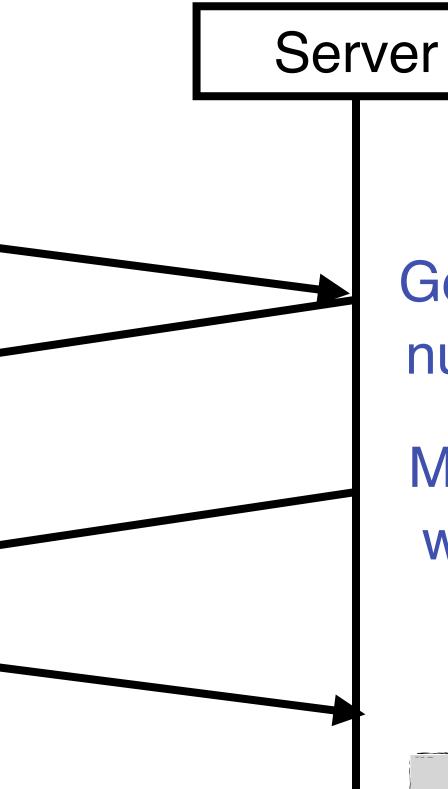


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Could we optimize a little bit?



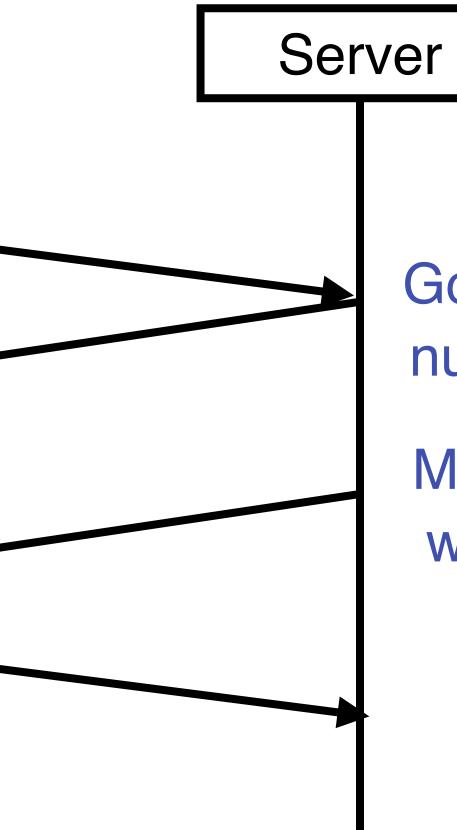


TCP Connection Establishment Let's start with a naive approach

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My (client) byte stream starts with a sequence number = X

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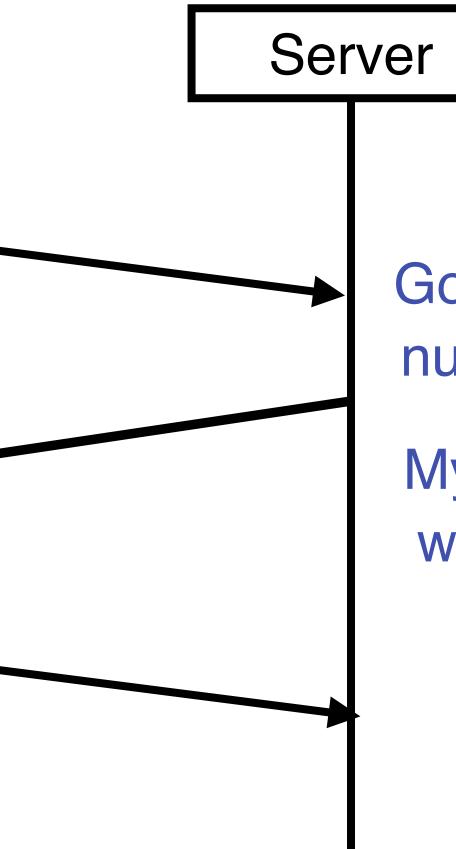


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My (server) byte stream starts with a sequence number = Y



TCP Connection Establishment Let's start with a naive approach Client My (client) byte stream starts with a sequence number = X Got it, I acknowledge the sequence number of your next byte is = Y + 1

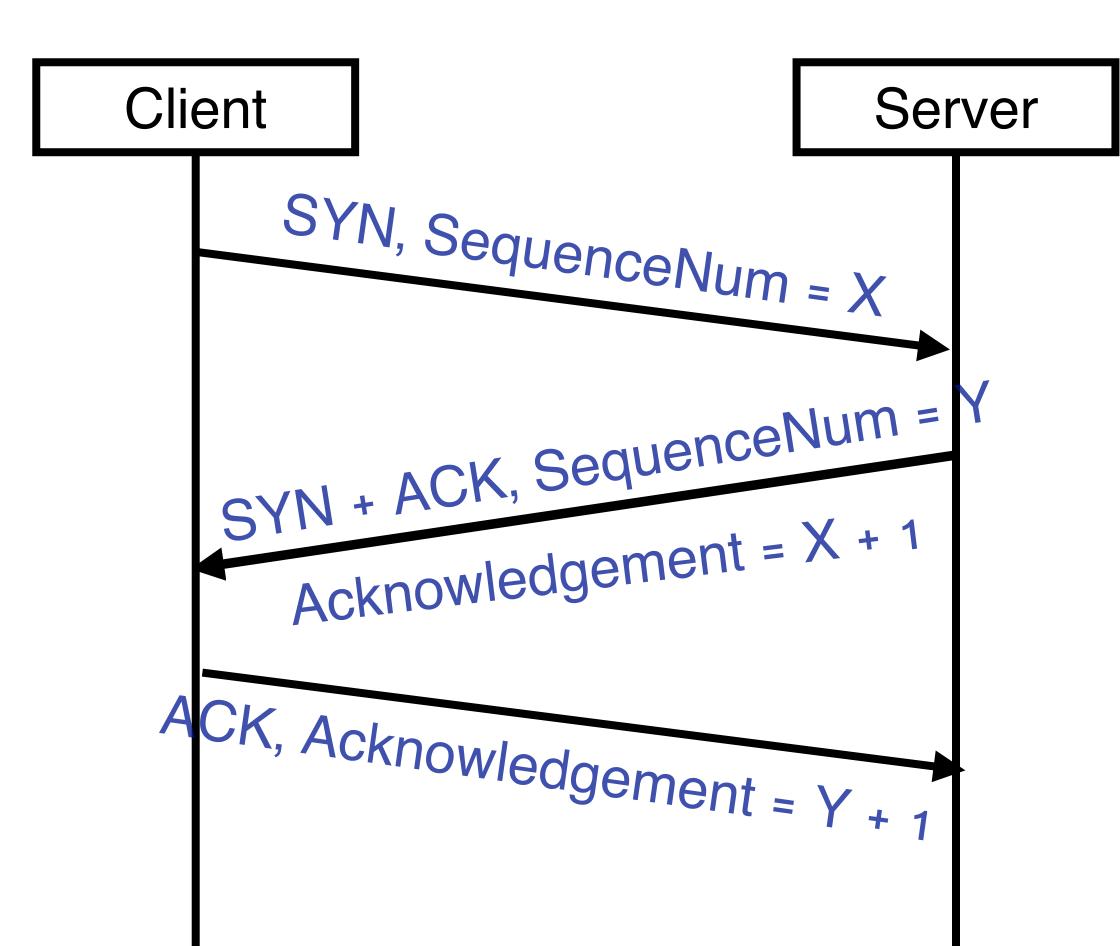


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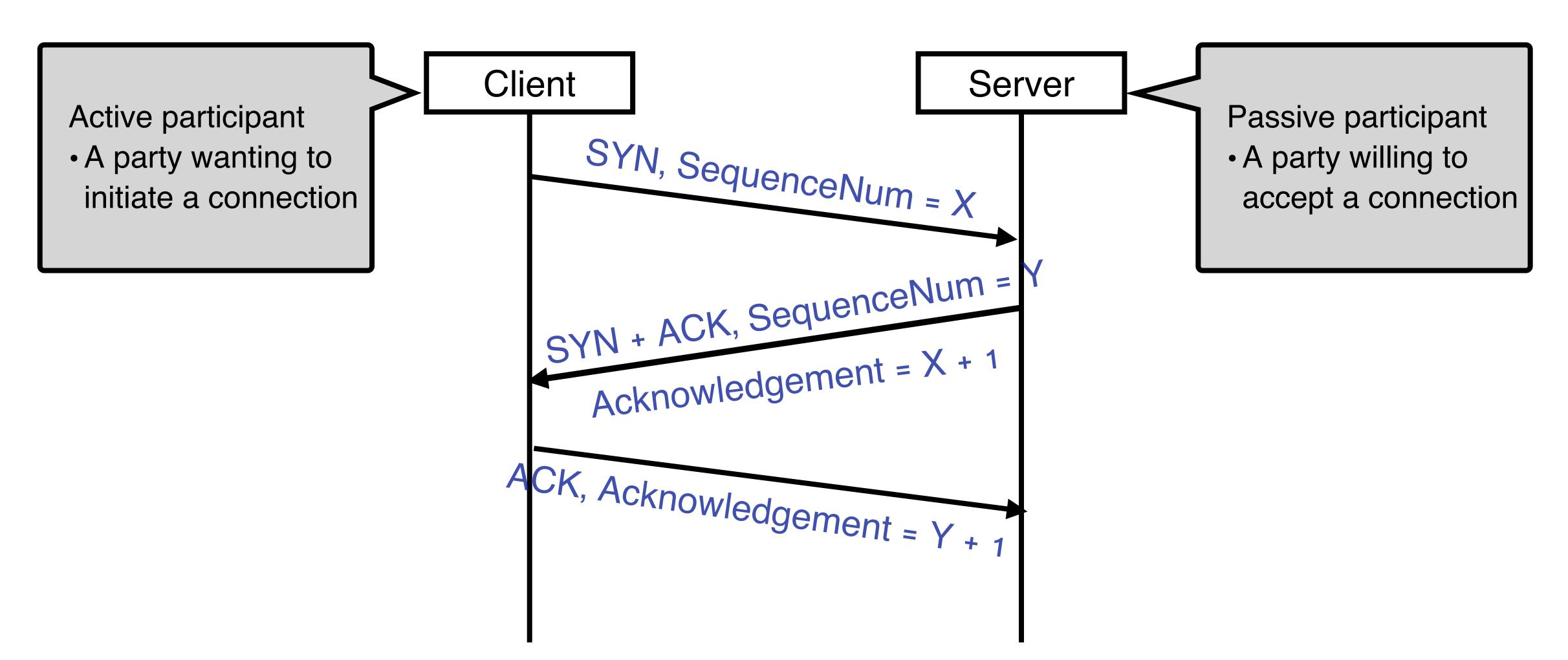


Three-Way Handshake



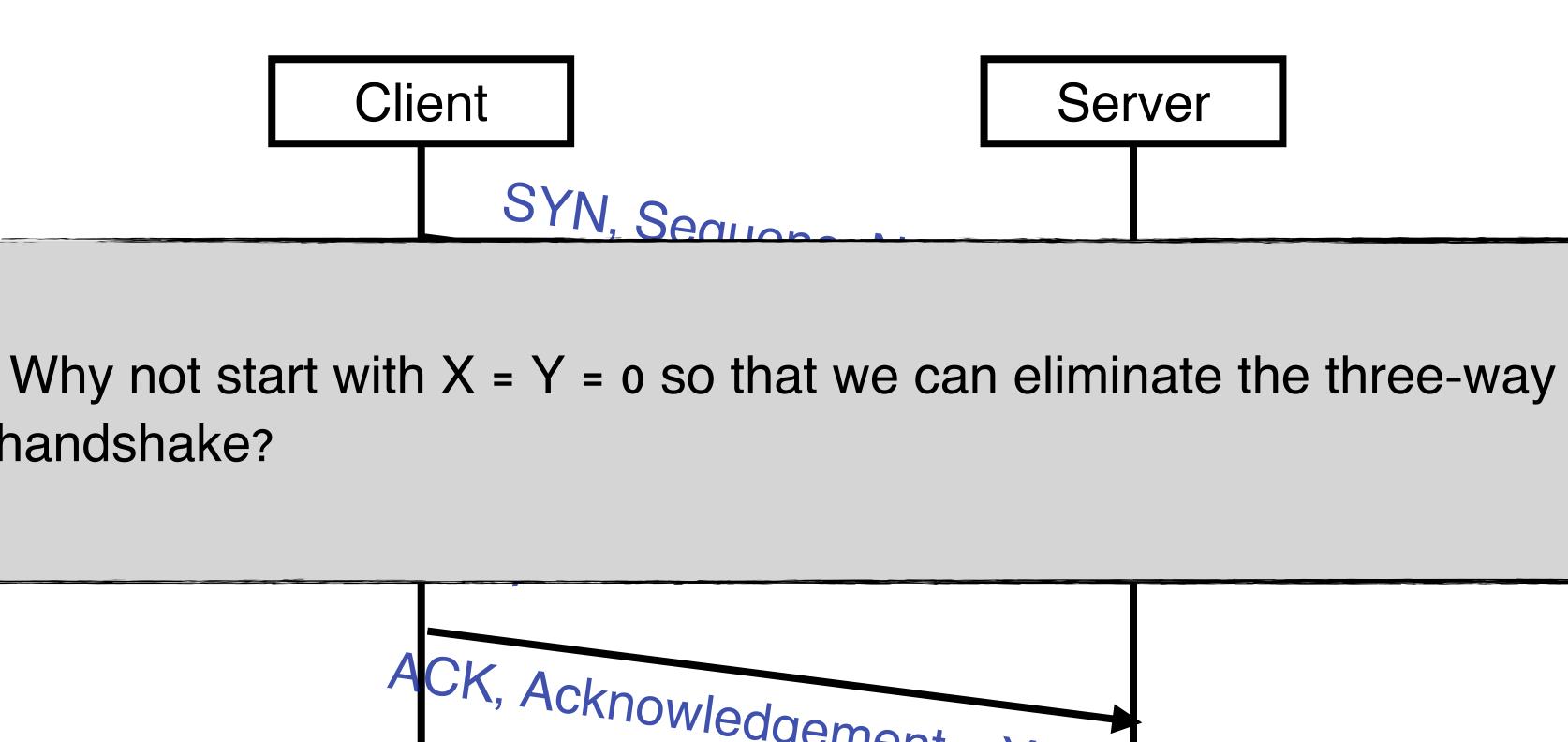


Three-Way Handshake

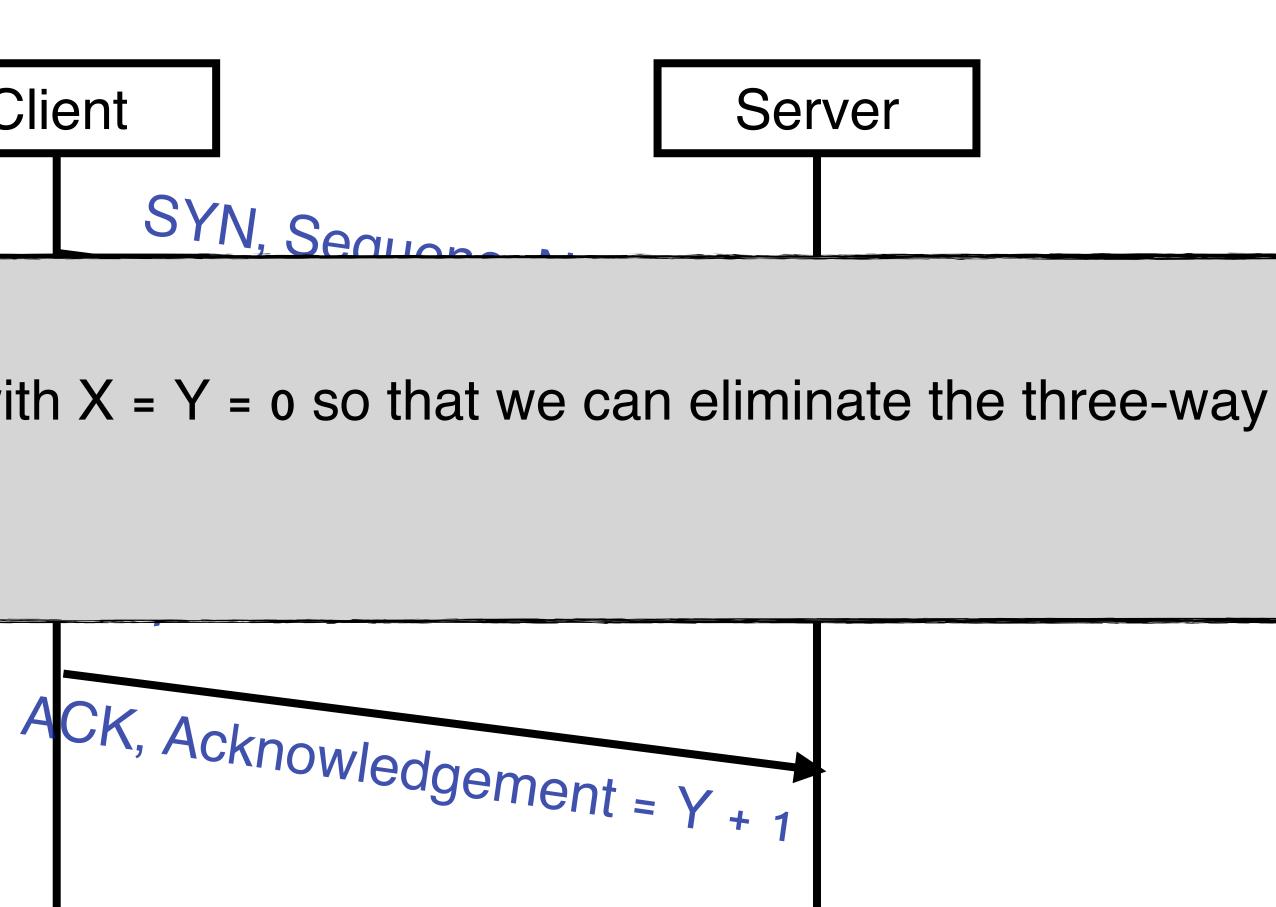




Three-Way Handshake



handshake?





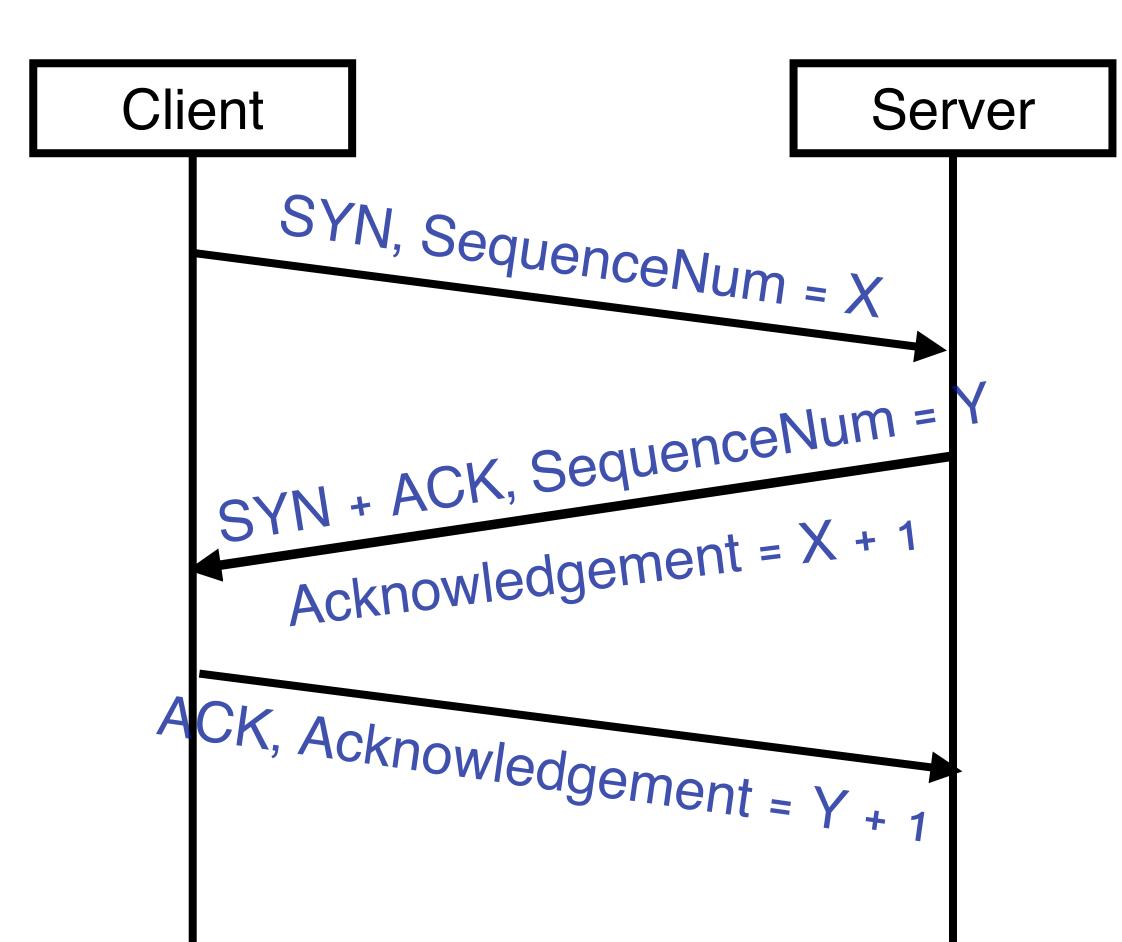
The Incarnation Issue

A connection (defined by a particular host and port pair) to be reused again

Solution: initial sequence number is randomly generated



How to implement this?



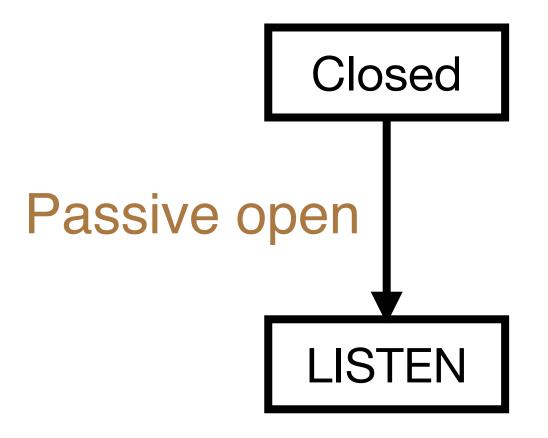


State Machine (event/action)

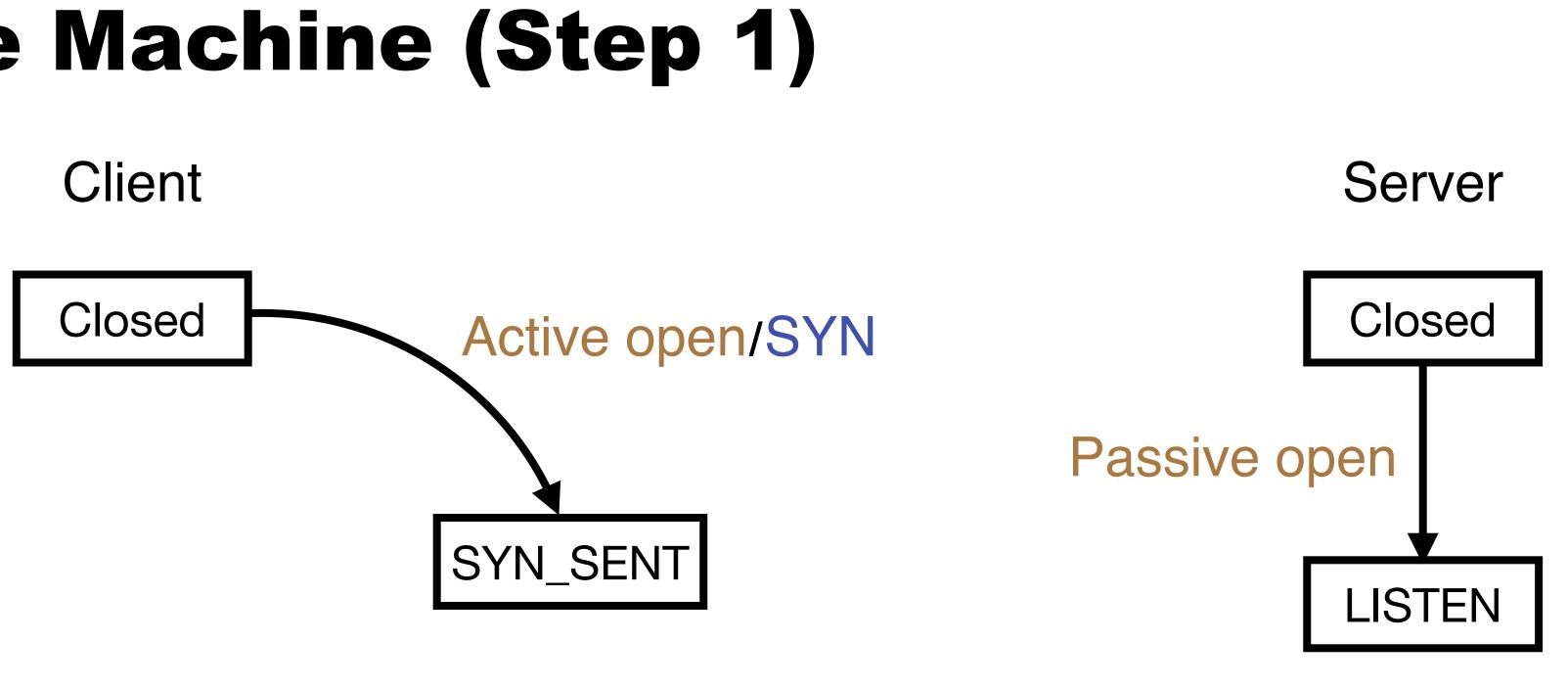
Client

Closed

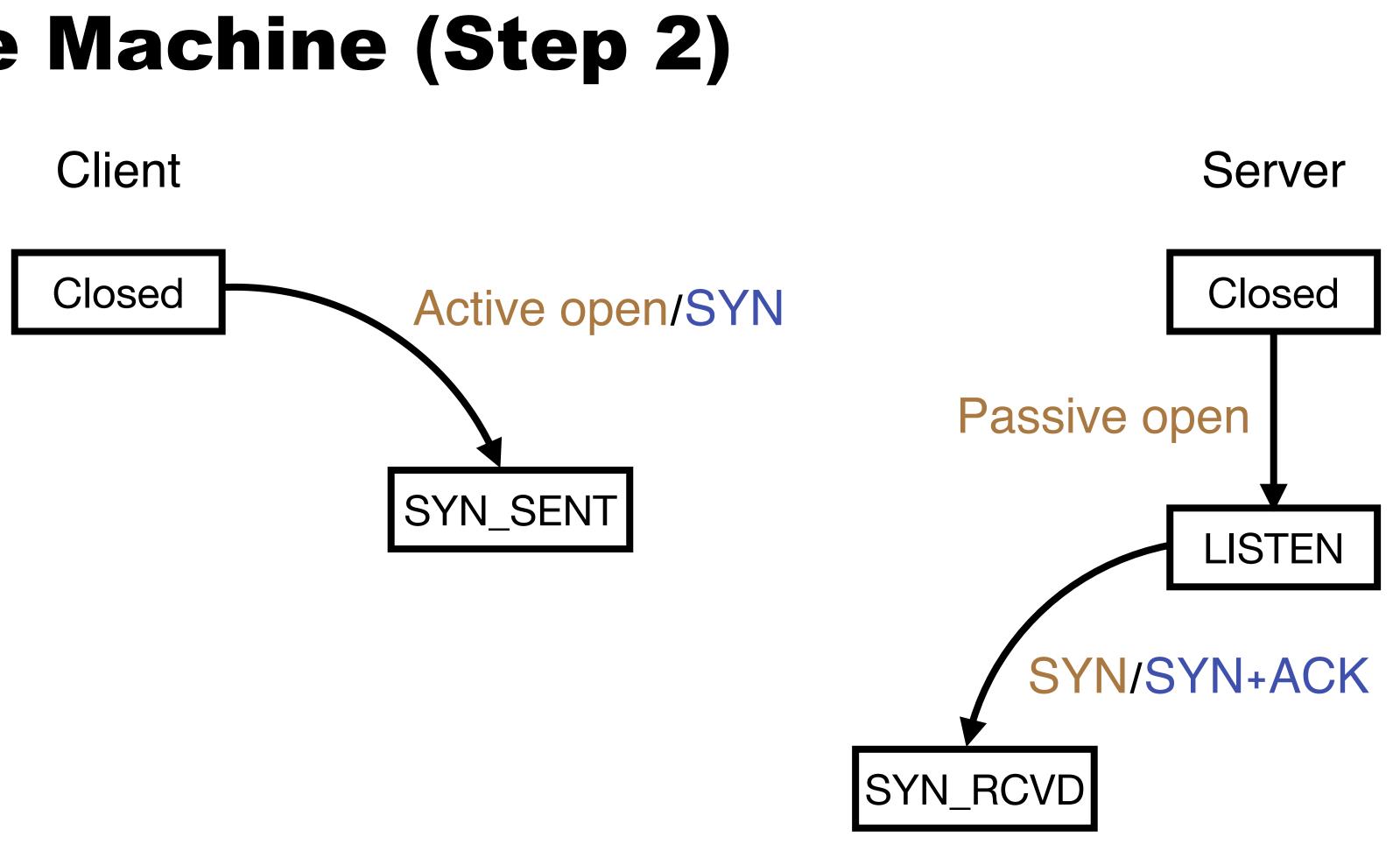




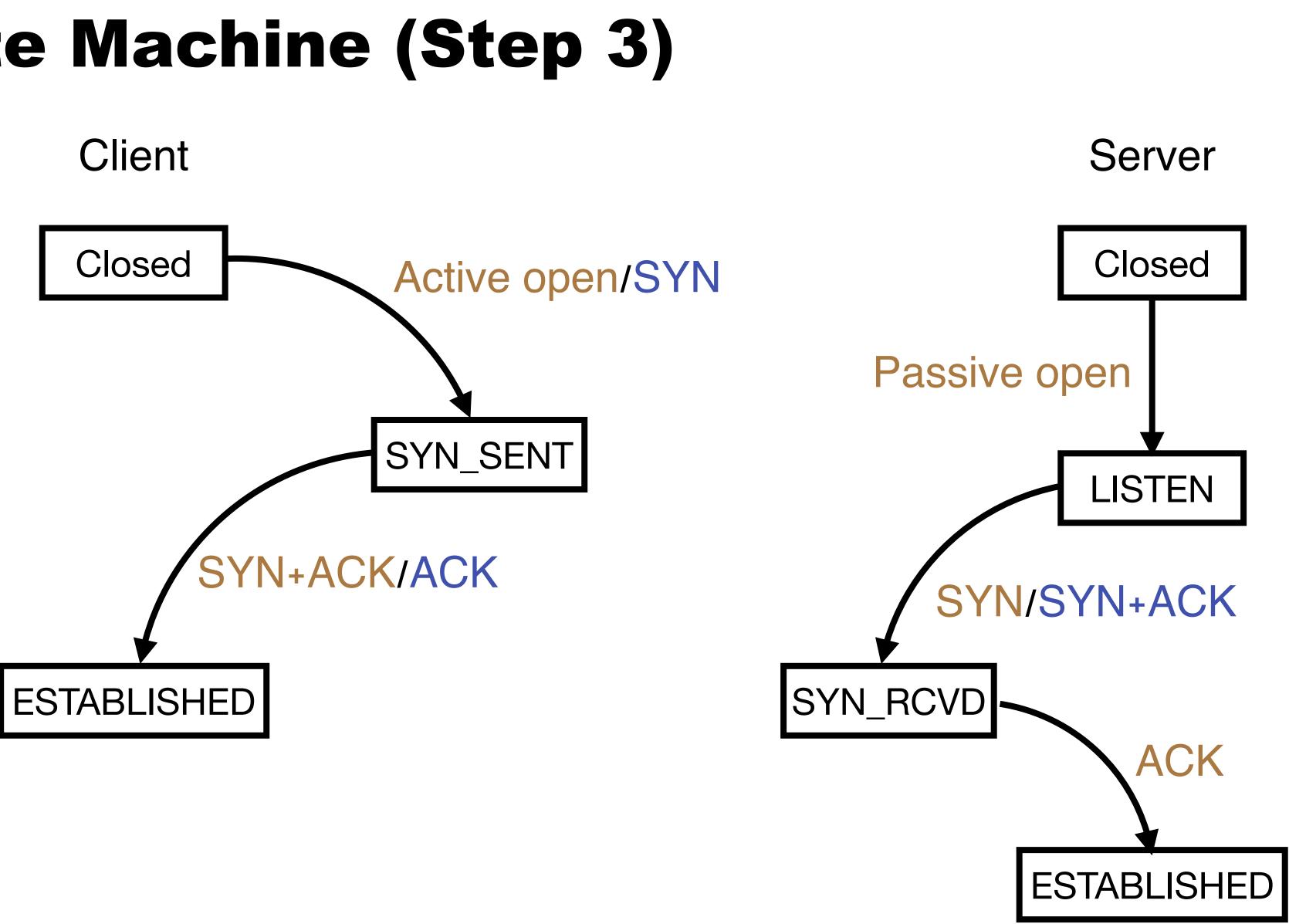
State Machine (Step 1)



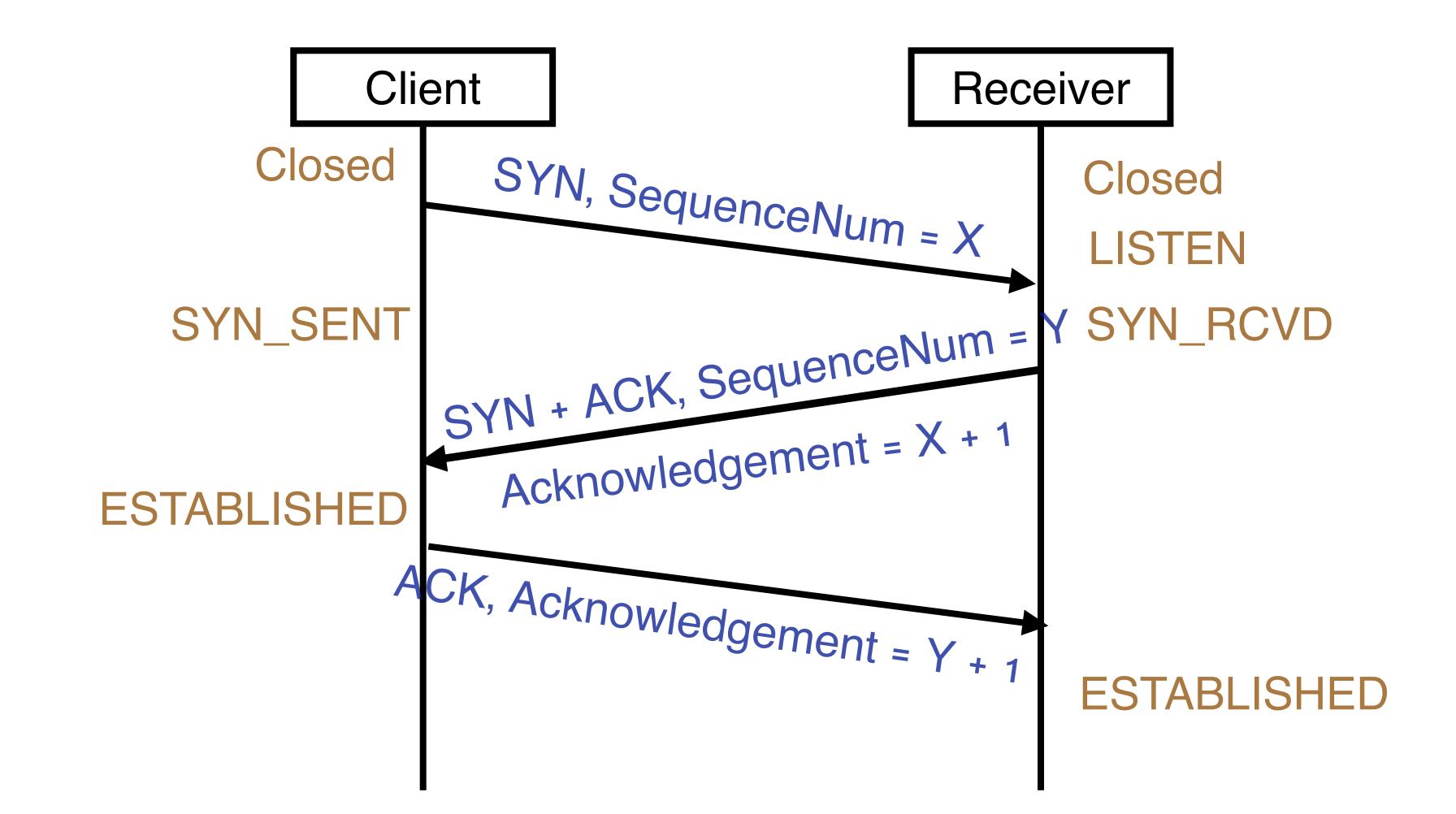
State Machine (Step 2)



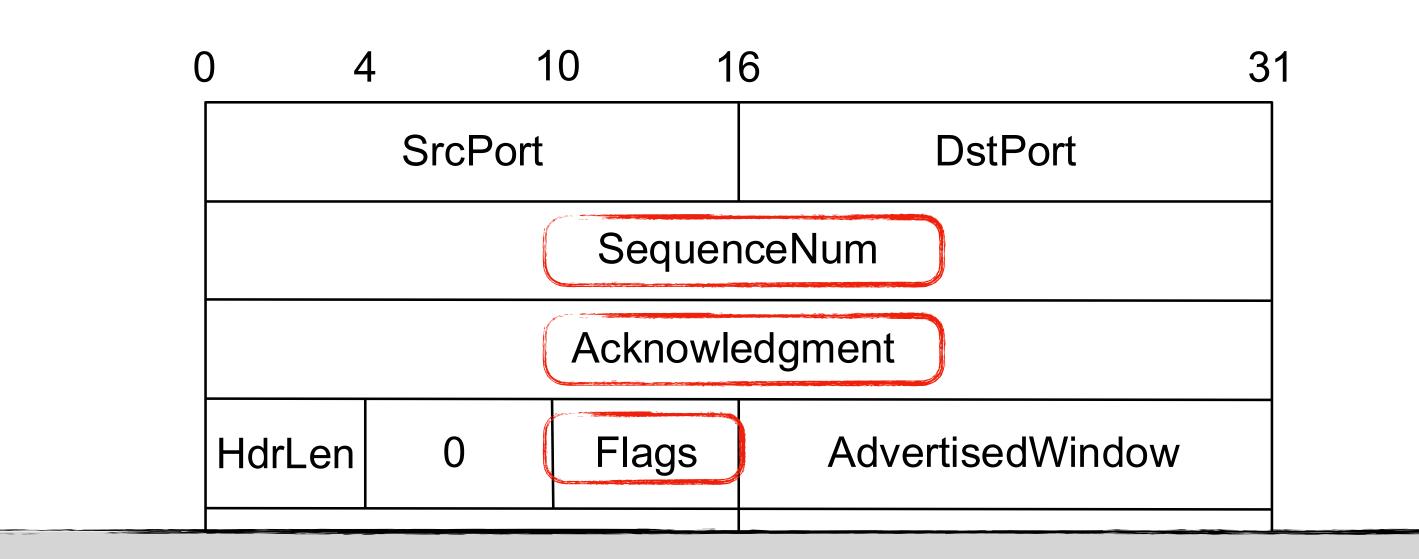
State Machine (Step 3)



TCP Connection Establishment Summary



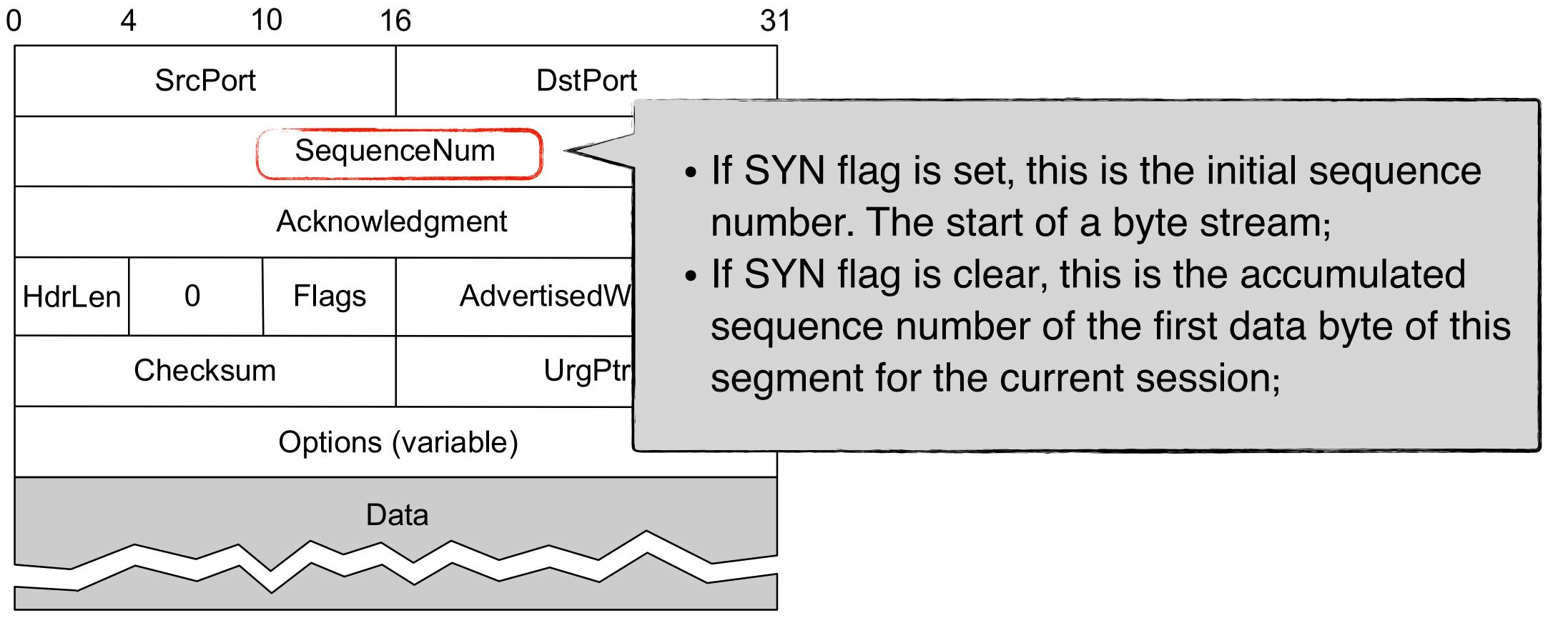
Revisit the TCP Header



- SYN/FIN -> TCP connection establishment and teardown
- ACK -> Acknowledgement is valid \bullet
- URG -> The segment contains urgent data. UrgPtr will be setup
- PUSH -> Notify the receiving process \bullet
- RESET -> The receiving side gets confused information

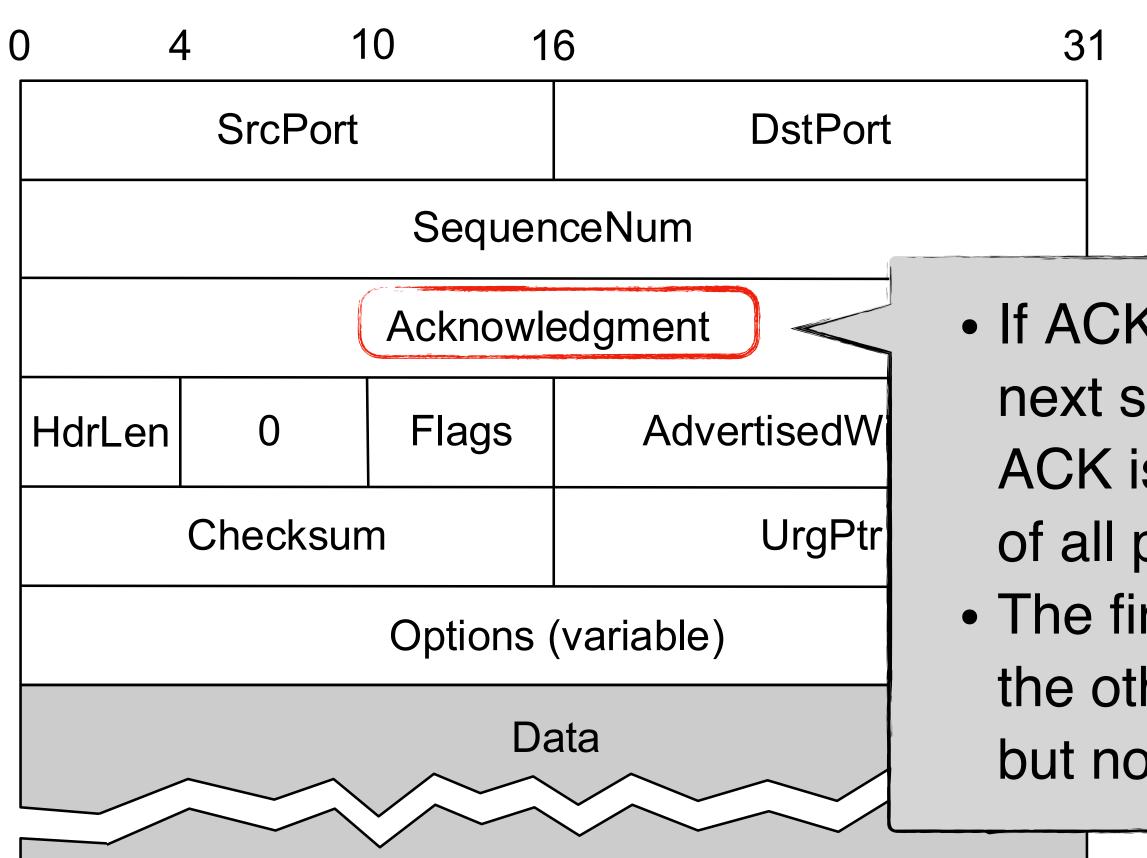


Revisit the TCP header





Revisit the TCP header



 If ACK flag is set, the value of this field is the next sequence number that the sender of the ACK is expecting. This acknowledges receipt of all prior bytes (if any)
The first ACK cent by each and acknowledges

 The first ACK sent by each end acknowledges the other ends's initial sequence number itself, but no data



Terminology

- 1. Host
- 2. NIC
- 3. Multi-port I/O bridge 19. Timeout
- 4. Protocol
- 5. RTT
- 6. Packet
- 7. Header
- 8. Payload
- 9. BDP
- 10. Baud rate
- 11. Frame/Framing
- 12. Parity bit
- 13. Checksum
- 14. Ethernet
- 15. MAC
- 16. (L2) Switch

- 17. Broadcast
- 18. Acknowledgement
- - 20. Datagram
 - 21. TTL
 - 22. MTU
 - 23. Best effort
 - 24. (L3) Router
 - 25. Subnet mask
 - 26. CIDR
 - 27. Converge
 - 28. Count-to-infinity
 - 29. Line card
 - 30. Network processor
 - 31. Gateway
 - 32. Private network

- 33. IPv6
- 34. Multicast
- 35. IGMP
- 36. SDN
- 37. (Transport) port
- 38. Pseudo header
- 39. SYN/ACK
- 40. Incarnation

Principle

- 1. Layering
- 2. Minimal States
- 3. Hierarchy

Technique

- 1. NRZ Encoding
- 2. NRZI Encoding
- 3. Manchester Encoding
- 4. 4B/5B Encoding
- 5. Byte Stuffing
- 6. Byte Counting
- 7. Bit Stuffing
- 8. 2-D Parity
- 9. CRC
- 10. MAC Learning
- 11. Store-and-Forward
- 12. Cut-through
- 13. Spanning Tree
- 14. CSMA/CD
- 15. Stop-and-Wait
- 16. Sliding Window

- 16. Fragmentation and Reassembly
- 17. Path MTU discovery
- 18. DHCP
- 19. Subnetting
- 20. Supernetting
- 21. Longest prefix match
- 22. Distance vector routing (RIP)
- 23. Link state routing (OSPF)
- 24. Boarder gateway protocol (BGP)
- 25. Network address translation (NAT)
- 26. User Datagram Protocol (UDP)
- 27. Transmission Control Protocol (TCP)
- 28. Three-way Handshake



Summary

Today's takeaways

#1: TCP connection setup uses a three-way handshake to build a duplex channel#2: The three-way handshake can be implemented via a state machine

Next lecture

TCP connection teardown

