#### Introduction to Computer Networks

# **Computer Networks: Performance Analysis**

#### https://pages.cs.wisc.edu/~mgliu/CS640/S25/index.html

Ming Liu mgliu@cs.wisc.edu

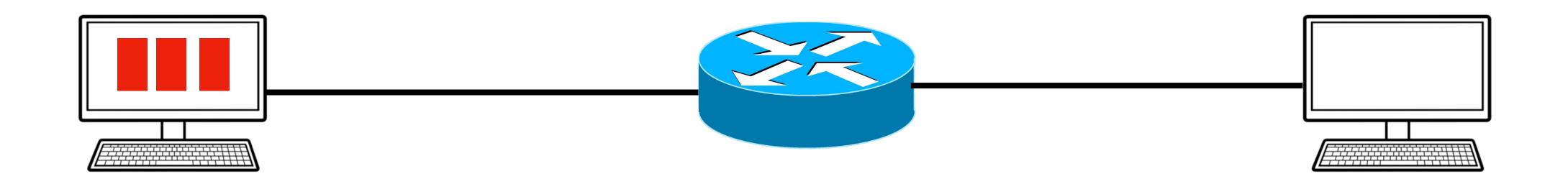
- Last
  - Computer networks: hardware infrastructure
  - Computer networks: software system
- Today
  - Delay
  - Throughput
  - RTT and BDP
- Announcements
  - Lab1 will be released this Thursday
  - No class on Thursday

#### Outline



#### **Packet-Switched Networks**

- Packet communication path
  - Host 1 —> Router A —> Router B —> ... —> Host 2

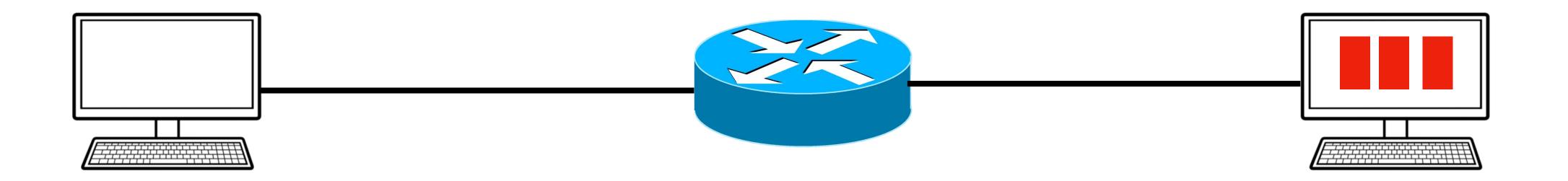




#### Packet-Switched Networks

- Packet communication path
  - Host 1 —> Router A —> Router B —> ... —> Host 2
- Delay

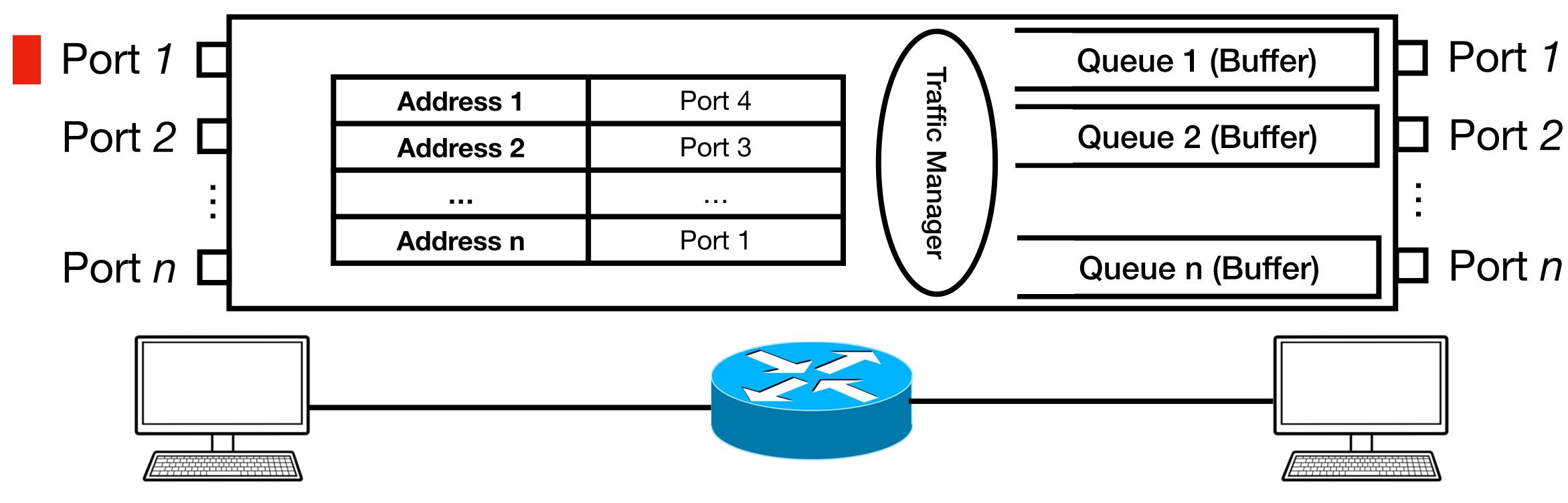
  - Four types of delay



#### The total amount of time transfers N bits across the packet path



### Processing Delay (Tproc)



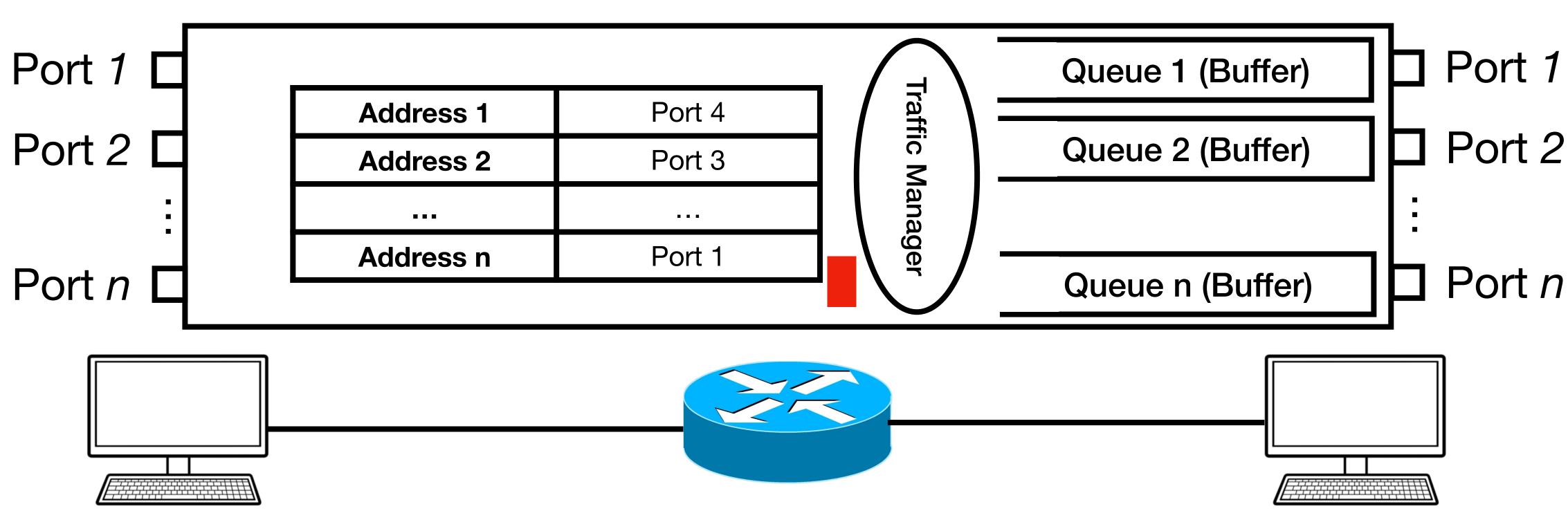




## Processing Delay (Tproc)

#### The time required to examine the packet header and determine where to forward

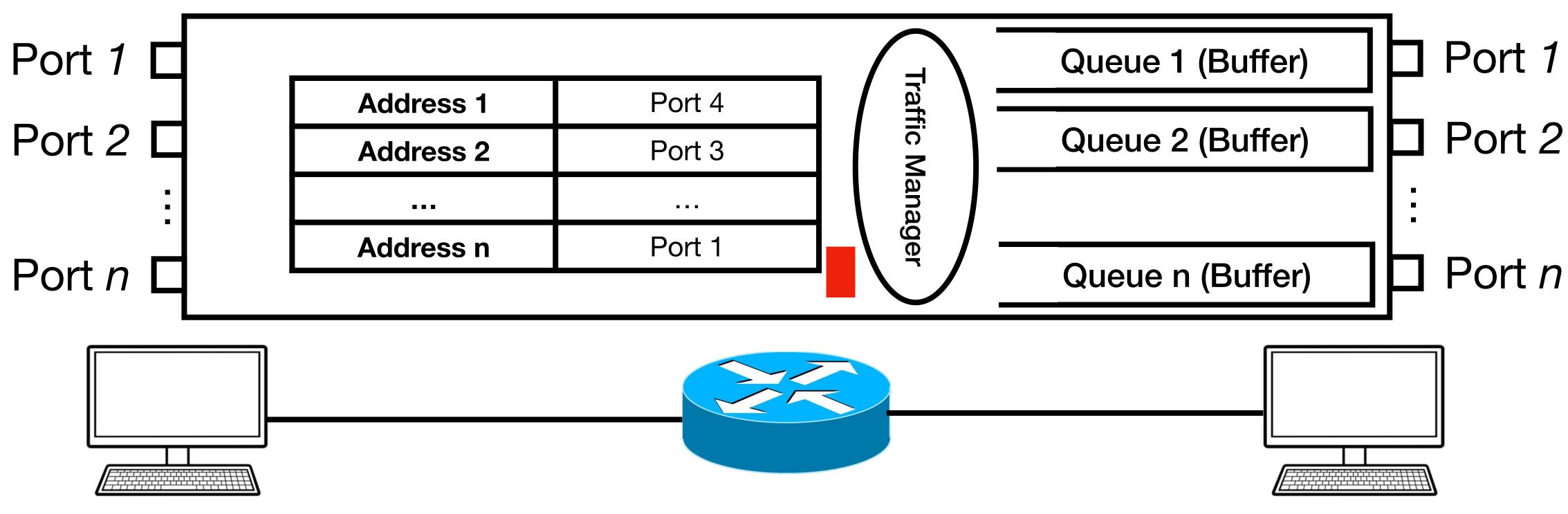
- Including checking bit-level errors
- Microsecond







### Queueing Delay (Tqueue)

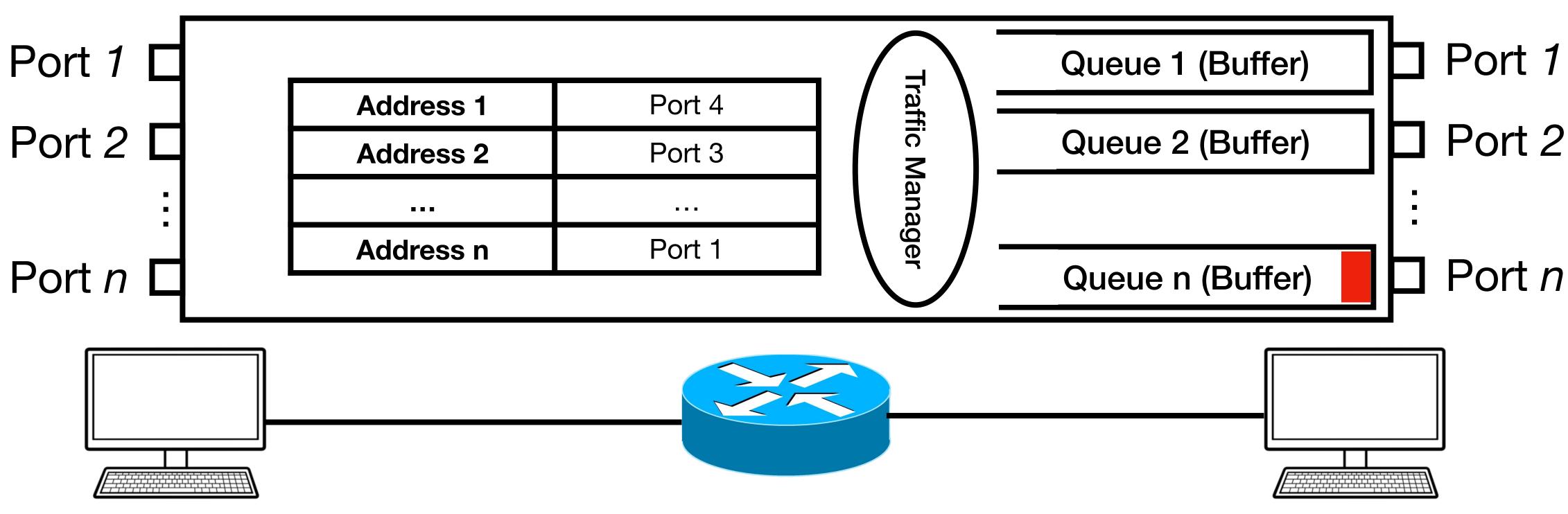






## Queueing Delay (Tqueue)

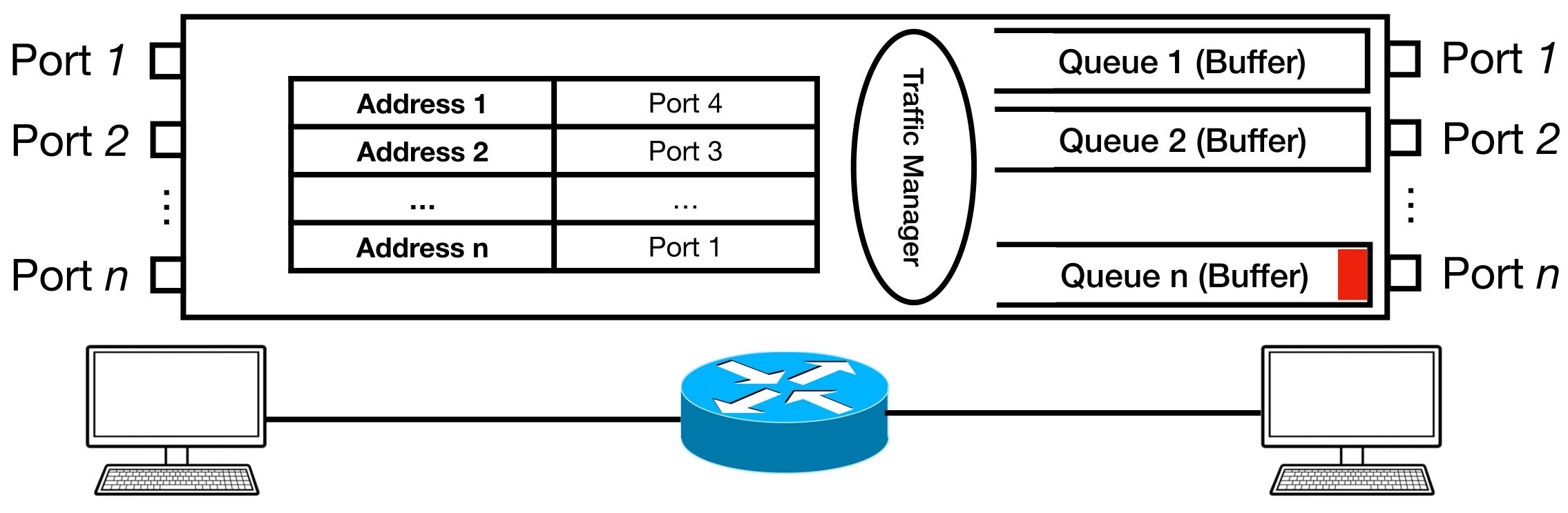
- The time it takes to wait to be transmitted
  - Depend on the number of earlier-arriving packets
  - Microsecond~Millisecond







### Transmission Delay (Ttrans)

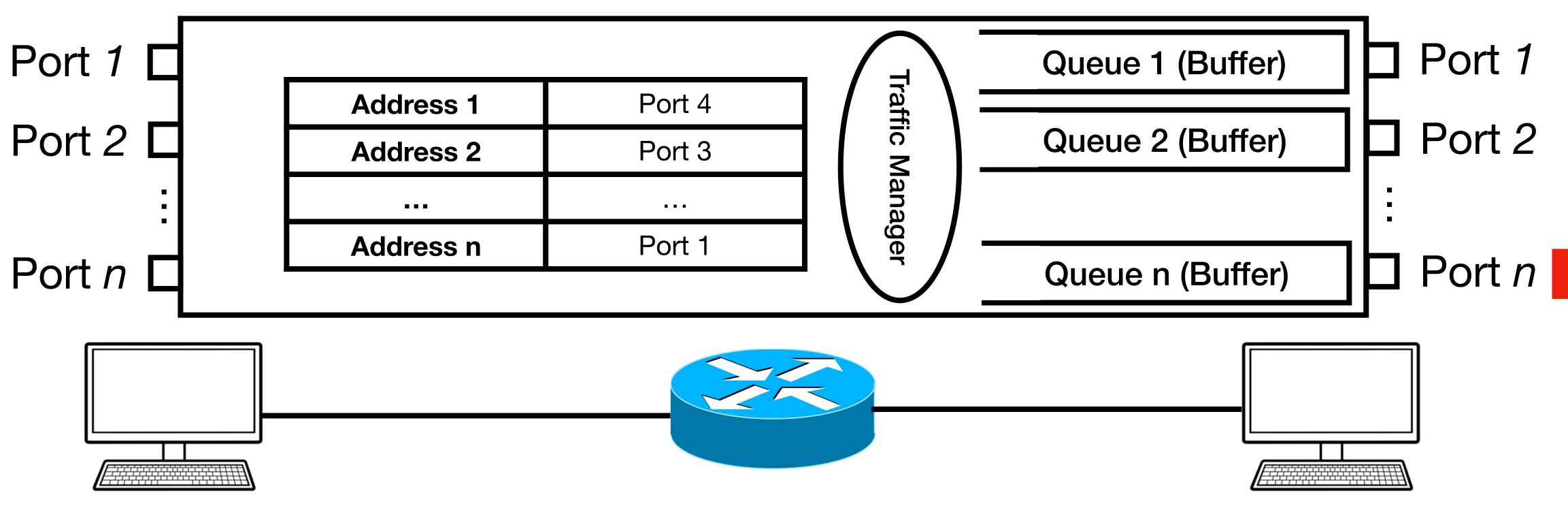






# Transmission Delay (Ttrans)

- The time it takes to transmit through the communication ports
  - First-come-first-served manner
  - Depend on the packet length and transmission rate



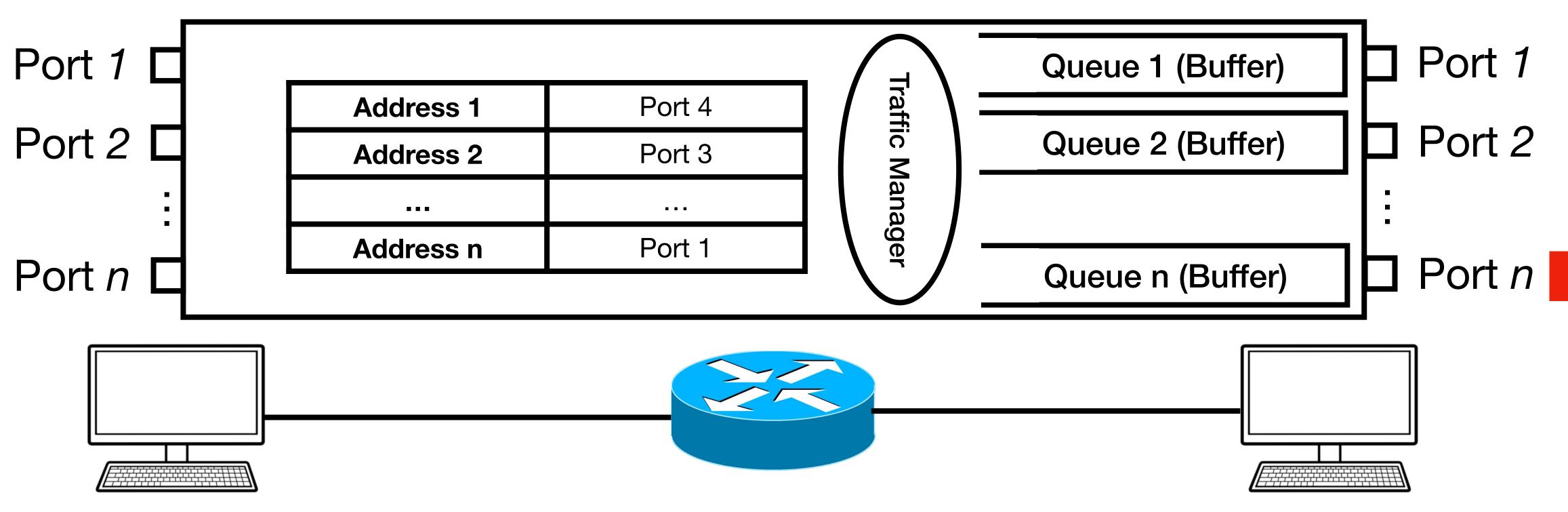






### **Transmission Delay Calculation**

- Packet length (L bits) and Transmission Rate (R bits/sec)
  - Transmission Delay = L/R
  - Suppose L=1KB, R = 10Mbps/100Mbps, what is the transmission delay?

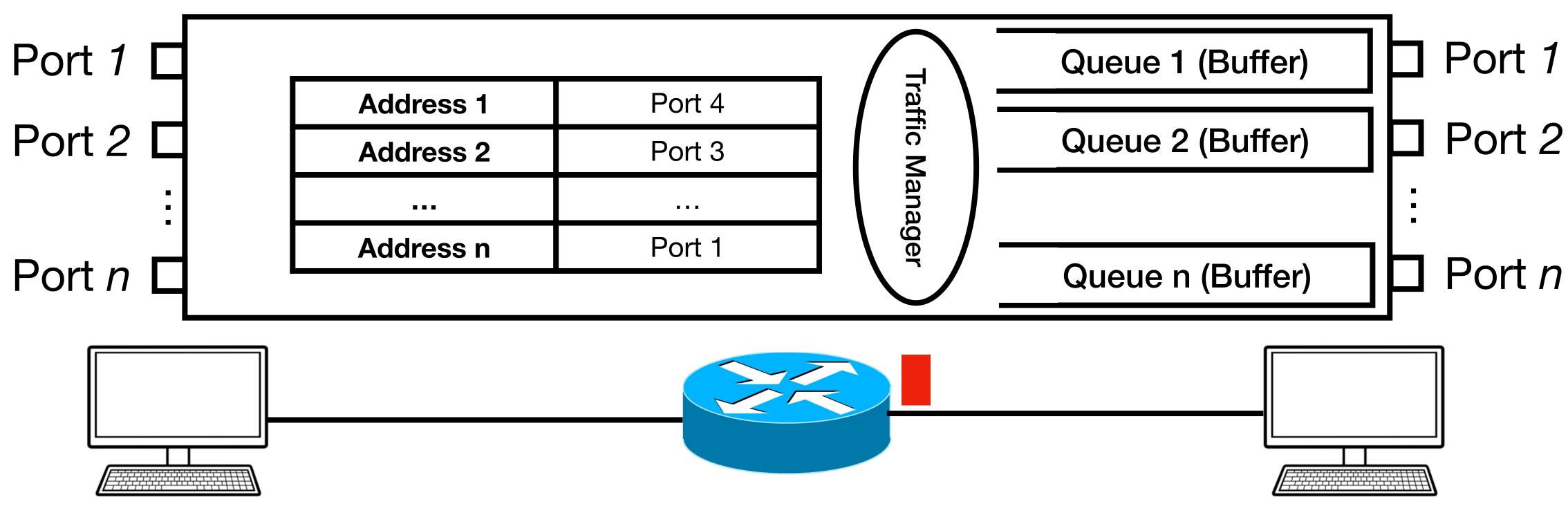








## **Propagation Delay (T**prop)

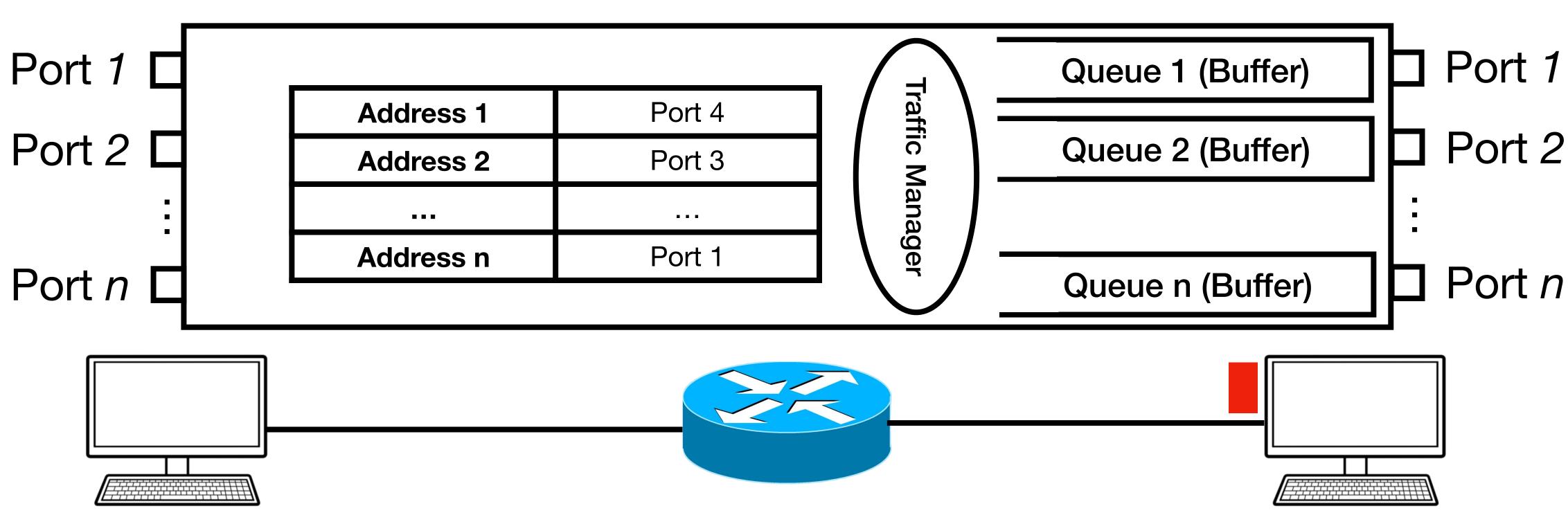






# **Propagation Delay (T**prop)

- - 2x10^8 ~ 3x10^8 meters/second

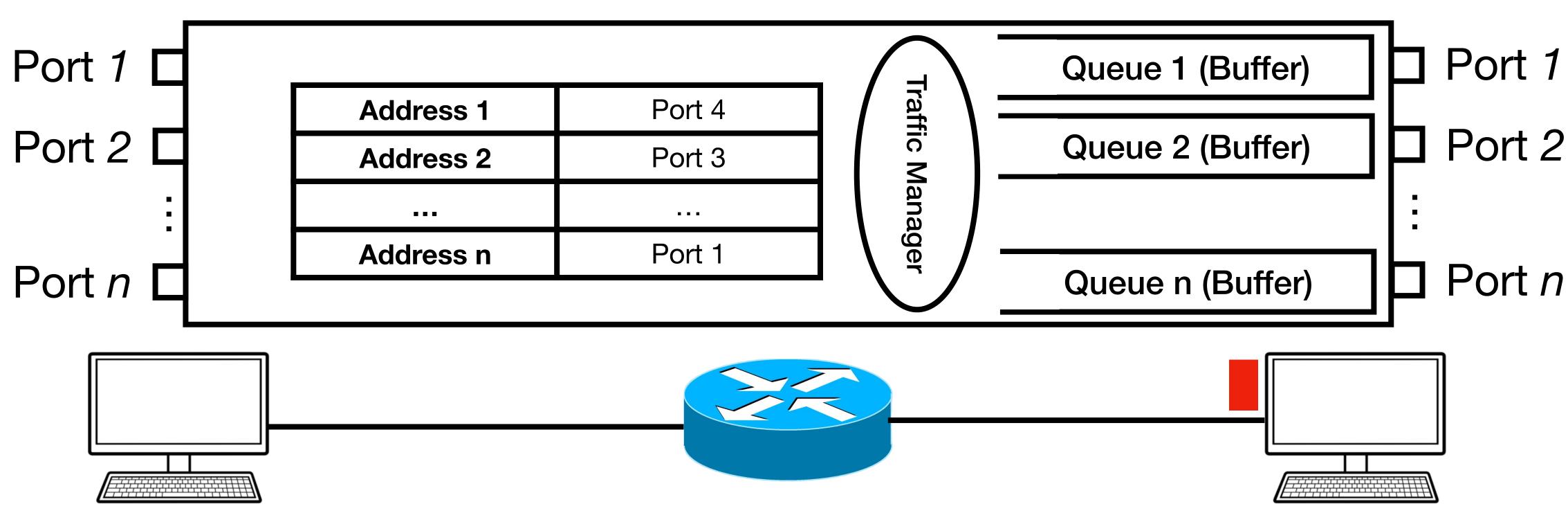


 The time required to propagate over the communication link • Depend on the physical media (e.g., fiber optics, copper wire, etc.)





- Distance (d) and propagation speed (s)
  - Propagation delay = d/s
  - Distance matters!



#### **Propagation Delay Calculation**





- Difference
  - Transmission delay: time taken to push out the packet
  - Propagation delay: time taken to traverse the link

to push out the packet to traverse the link



- Difference

  - Transmission delay: time taken to push out the packet Propagation delay: time taken to traverse the link
  - Suppose

    - Packet length (L) = 1.5KB and transmission rate (R) = 1Mbps • Distance (d) = 3m and propagation speed (s) =  $3^8 m/s$ What is the transmission delay and propagation delay?



- Difference

  - Transmission delay: time taken to push out the packet • Propagation delay: time taken to traverse the link
  - Suppose

    - Packet length (L) = 1.5 KB and transmission rate (R) = 1 Mbps• Distance (d) = 3m and propagation speed (s) =  $3^8 m/s$  What is the transmission delay and propagation delay? • If we increase the transmission rate to 1Gbps, what happens?



- Difference

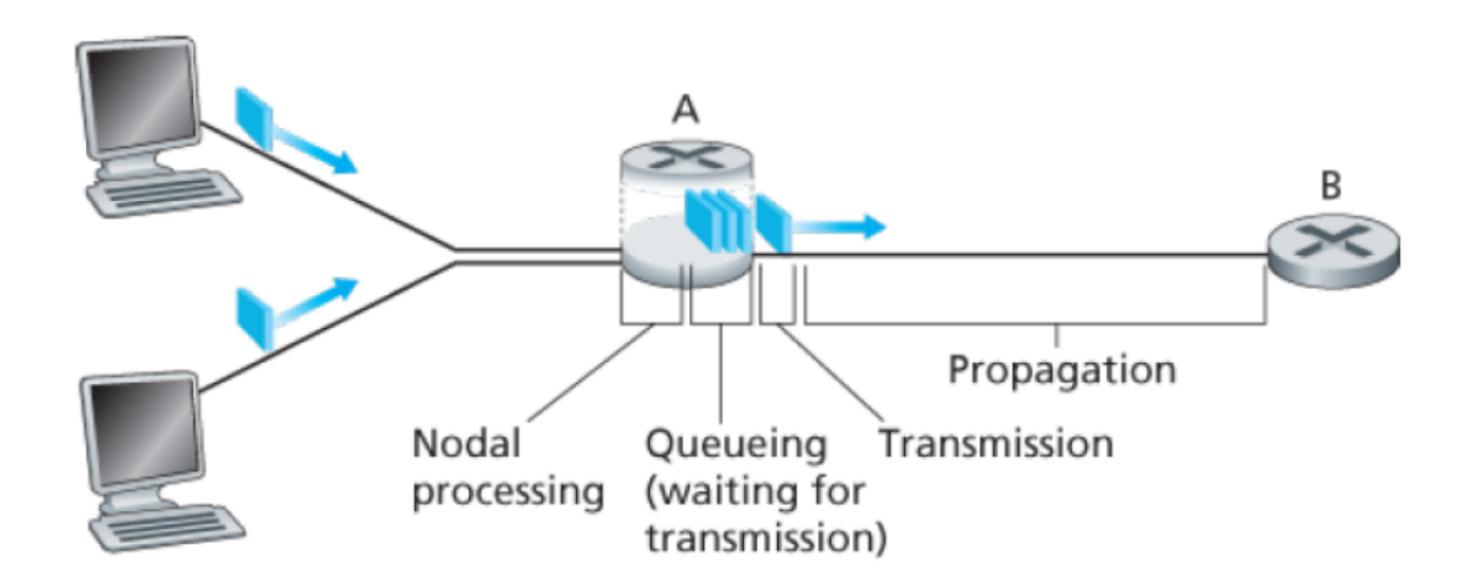
  - Transmission delay: time taken to push out the packet • Propagation delay: time taken to traverse the link
  - Suppose

    - Packet length (L) = 1.5 KB and transmission rate (R) = 1 Mbps• Distance (d) = 3m and propagation speed (s) =  $3^8 m/s$  What is the transmission delay and propagation delay? • If we increase the transmission rate to 1Gbps, what happens? • If we increase the distance to 3km, what happens?



# Total (Nodal) Delay

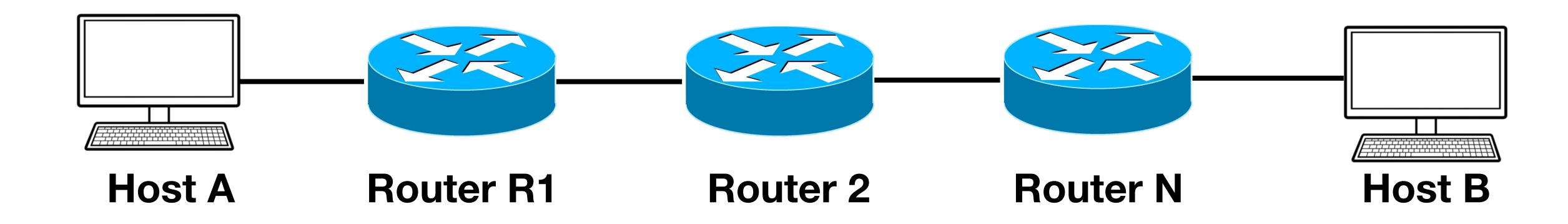
- Ttotal = Tproc + Tqueue + Ttrans + Tprop
  - Per-node, also called total nodal delay





#### End-to-End Delay

- Suppose
  - Host A can send a packet infinitely fast
  - Host B can receive a packet infinitely fast
- What is the end-to-end delay for a packet from host A to host B? • Tend-to-end = N (Tproc + Tqueue + Ttrans + Tprop) + Tprop



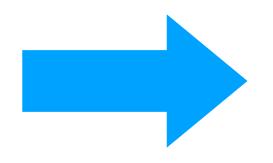


## Queueing Delay

- The most complicated and interesting one
  - Varying packet to packet and scenario to scenario
- Statistical metrics
  - Average queueing delay
  - Variance of queueing delay
  - Tail queueing delay
  - The probability that the queuing delay exceeds some value

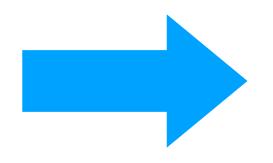
11



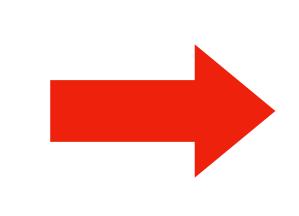


- Packet arrival rate a
- Unit: packets/second





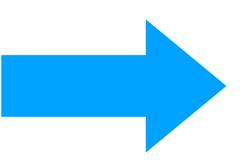
- Packet arrival rate a
- Unit: packets/second



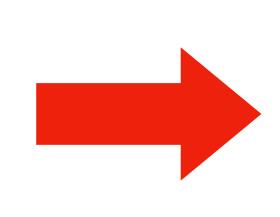
- Transmission rate R
- Unit: bits/second



- Given time T, incoming traffic load> outgoing traffic load
  - Incoming traffic load =  $\sum_{i=1}^{N} pkt\_size_i$  Out going traffic load = R \* T



- Packet arrival rate a
- Unit: packets/second



- Transmission rate R
- Unit: bits/second



- We define the ratio (La/R) as the traffic intensity Suppose all packets consist of L bits



- We define the ratio (La/R) as the traffic intensity Suppose all packets consist of L bits
- La/R > 1: queue built up!
  - The average rate at which bits arrive at the queue exceeds the rate at which the bits can be transmitted from the queue



- We define the ratio (La/R) as the traffic intensity Suppose all packets consist of L bits
- La/R > 1: queue built up!
  - The average rate at which bits arrive at the queue exceeds the rate at which the bits can be transmitted from the queue

## Design your system so that the traffic intensity is no greater than 1!





- We define the ratio (La/R) as the traffic intensity Suppose all packets consist of L bits
- La/R > 1: queue built up!
  - The average rate at which bits arrive at the queue exceeds the rate at which the bits can be transmitted from the queue

#### Can we see queueing when La/R <=1?





#### Traffic Arrival process

Suppose one packet arrives every L/R seconds



#### Traffic Arrival process

- Suppose one packet arrives every L/R seconds
  - No queueing
- Suppose N packets arrive every (L/R)N seconds



#### Traffic Arrival process

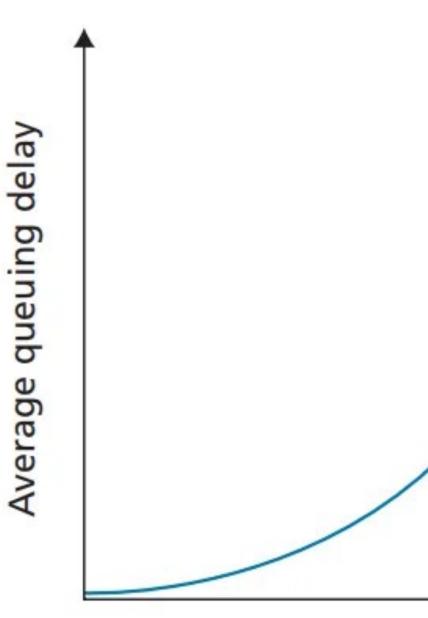
- Suppose one packet arrives every L/R seconds
  - No queueing
- Suppose N packets arrive every (L/R)N seconds • The 1st packet has no queueing delay, Tqueue = 0 • The 2nd packet has to wait for the 1st one,  $T_{queue} = L/R \times 1$ • The 3nd packet has to wait for the 1st and 2nd ones,  $T_{queue} = L/R \times 2$ • The nth packet has to wait for the (n-1) ones,  $T_{queue} = L/R \times (n-1)$



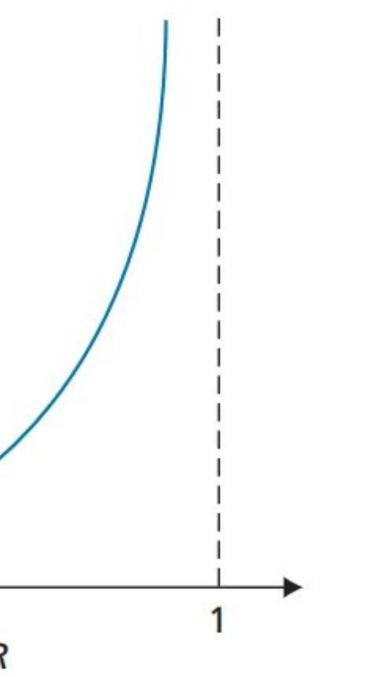
## Traffic Intensity Curve

- Close to 1, the average queueing delay increases

  - In reality, queue is fixed-sized ==> Packet Loss



# • A small percentage increase causes a significant delay increase





### Throughput

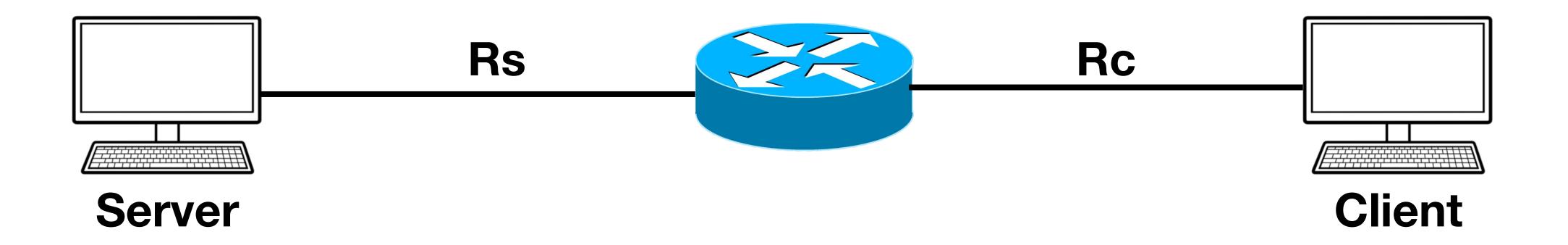
- Bandwidth: the number of bits transmitted per second at a communication port and link
  - bps, Kbps, Mbps, Gbps, Tbps, ...
- Throughput: the number of bits transmitted from A to B • A and B can be host, switch, etc.

  - bps, Kbps, Mbps, Gbps, Tbps, ...



#### A Simple Throughput Example

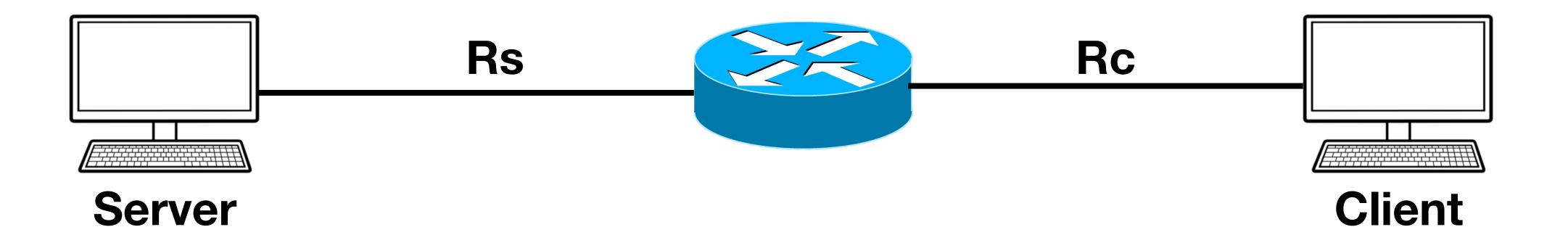
#### • A file transferring from a server to a client





### A Simple Throughput Example

- A file transferring from a server to a client
  - If Rs < Rc ==> The client receives the file at Rs

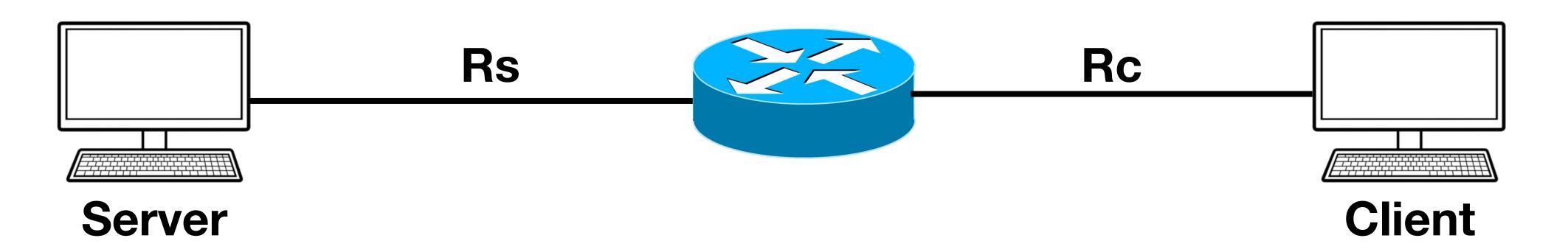


• If Rs > Rc ==> The client receives the file at Rc, but router is queued



## A Simple Throughput Example

- A file transferring from a server to a client
  - If Rs < Rc ==> The client receives the file at Rs
- Throughput = min {Rs, Rc}
  - Depend on the bottleneck link
- A log file = 32M bits, Rs = 2Mbps, Rc = 1Mbps
  - Transfer time is 32 seconds



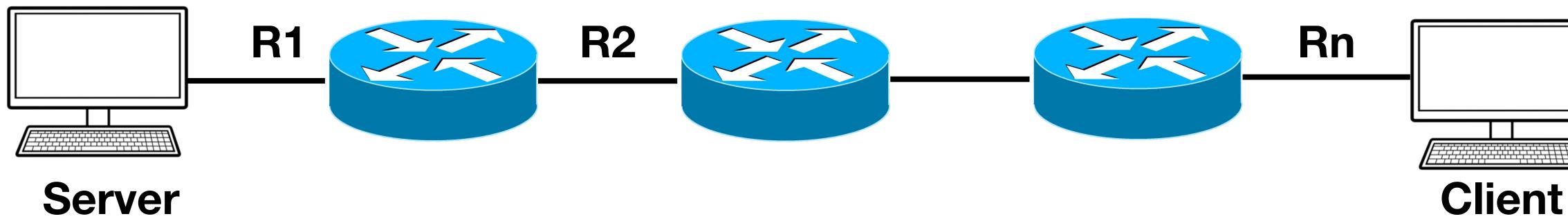
• If Rs > Rc ==> The client receives the file at Rc, but router is queued



## Another Throughput Example

Communication path: server -> a list of routers -> Client

- Throughput = min {R1, R2, R3, ..., Rn}
  - Depend on the bottleneck link



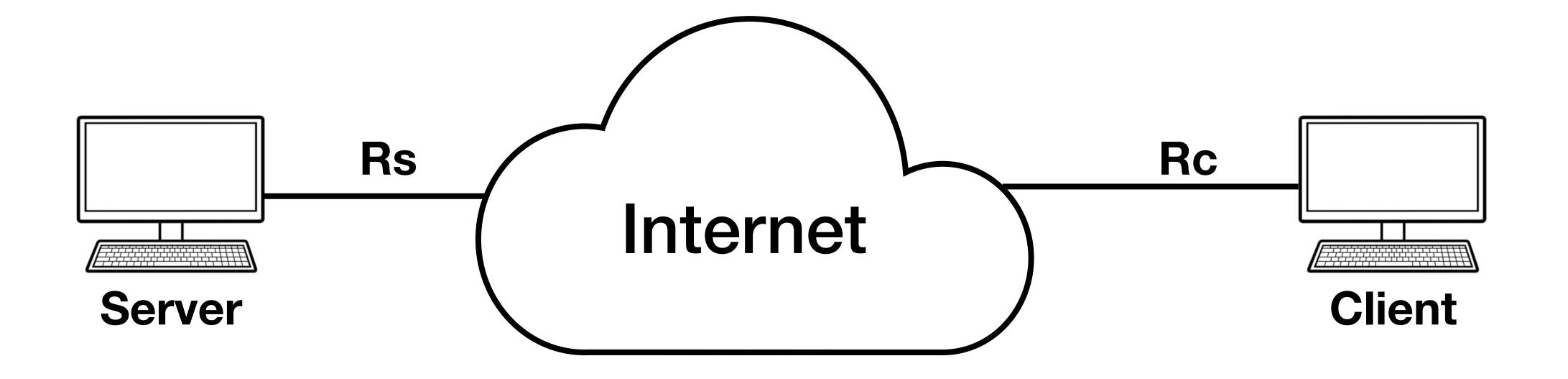






## Throughput in a Shared Network

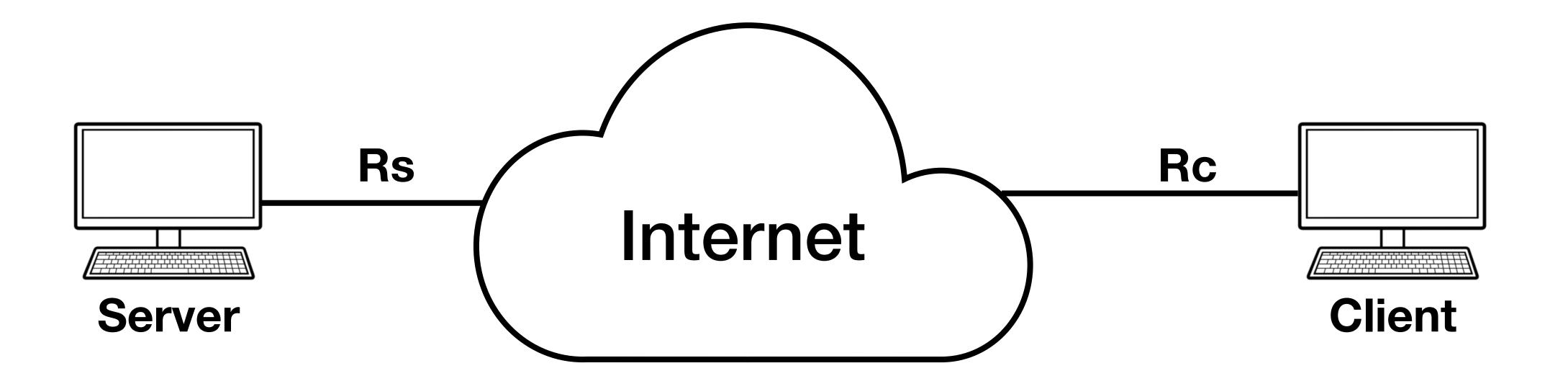
- Impossible to know the communication path details
- What is the throughput to transfer a file from a server to a client?





## Throughput in a Shared Network

- Impossible to know the communication path details
- - Actual Throughput = File Size / Total Transfer Time
  - Actual Throughput <= min {Rs, Rc}

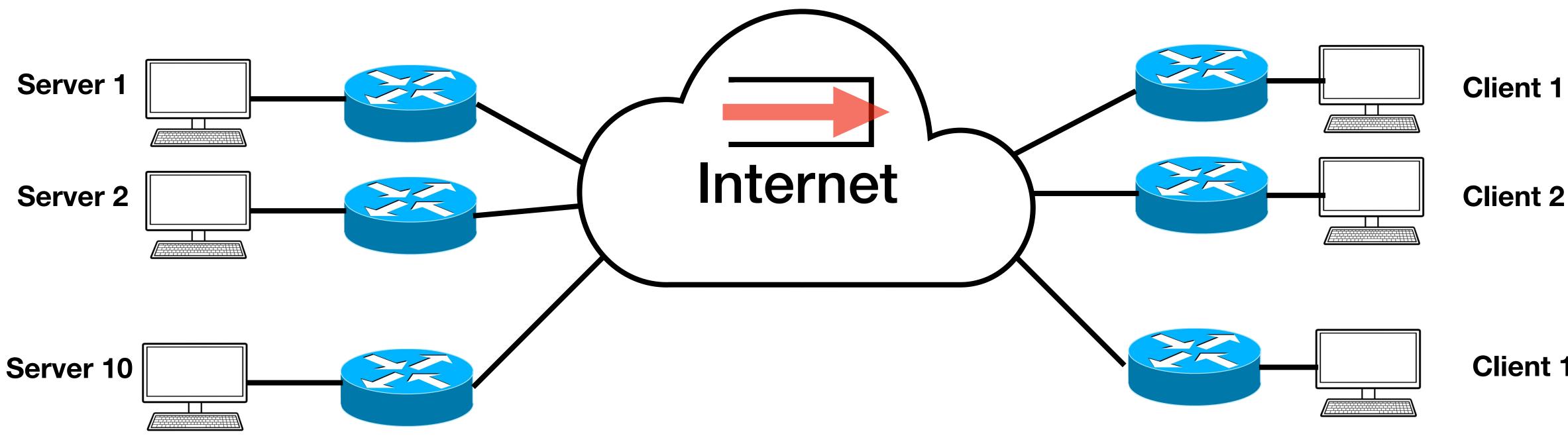


What is the throughput to transfer a file from a server to a client?



## Throughput under Concurrent Transmissions

Throughput also depends on intervening traffic

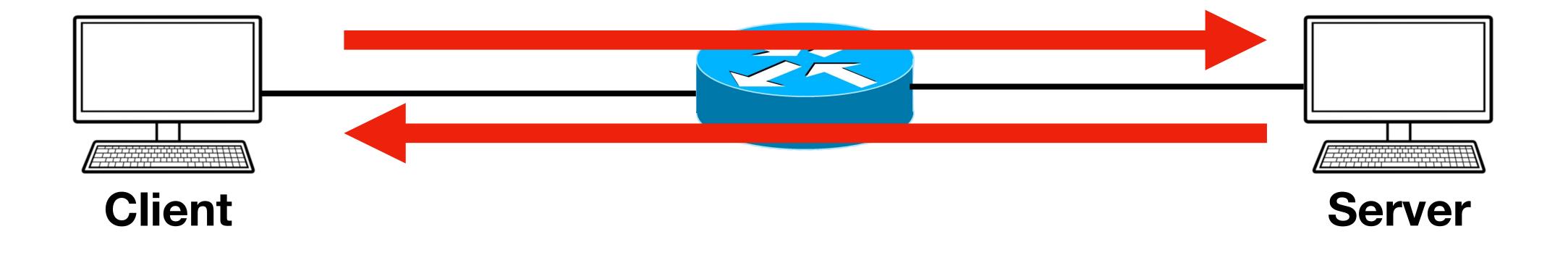






## RTT (Round-Trip Time)

# receive the response back



The time it takes to send a request from a client to a server and

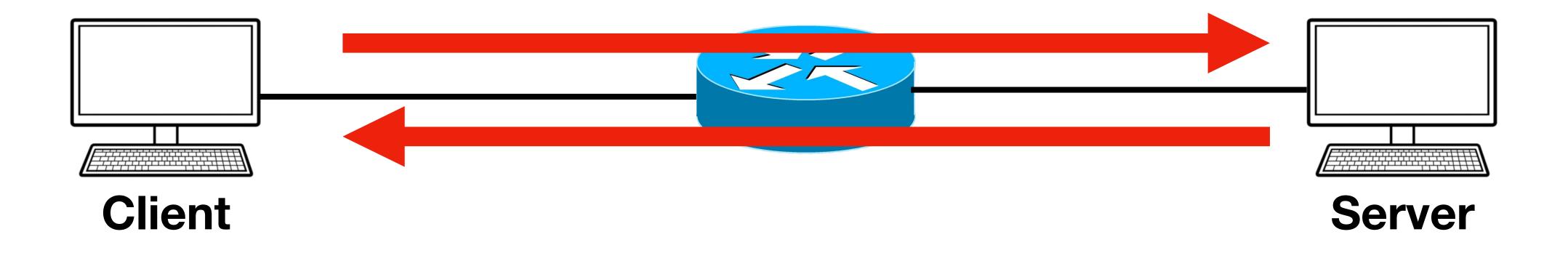
• RTT = Total delay (Client -> Server) + Total Delay (Server -> Client)



## RTT (Round-Trip Time)

- The time it takes to send a request from a client to a server and receive the response back
  - RTT = Total delay (Client -> Server) + Total delay (Server -> Client)
- RTT is application-dependent

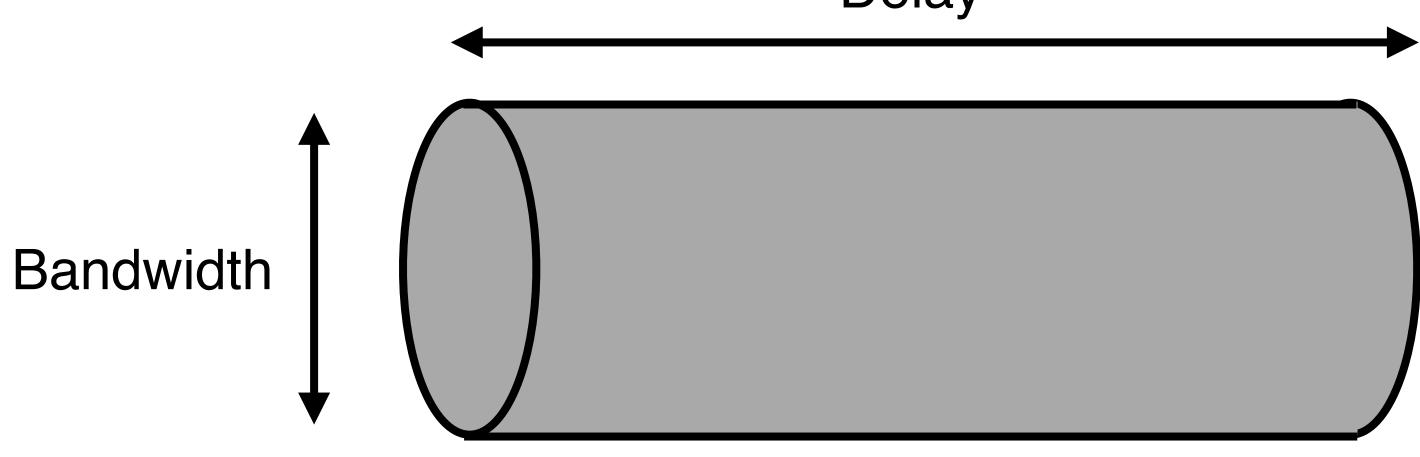
  - Web browsing: page download time (time to retrieve the first object) Cloud gaming: interactive latency
  - Video conferencing and streaming: Time to First Frame
  - LLM: Time to first token





## BDP (Bandwidth-Delay Product)

- The volume of a data pipe for one transmission
  - Bandwidth \* Delay
  - The number of bits have left the sender and are yet to reach the receiver

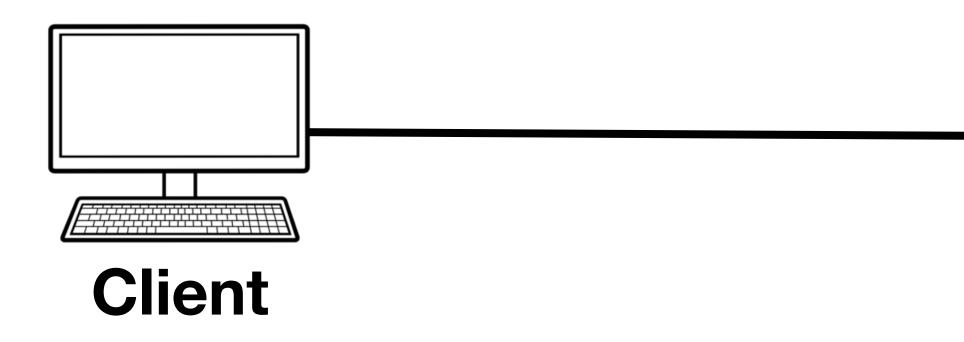


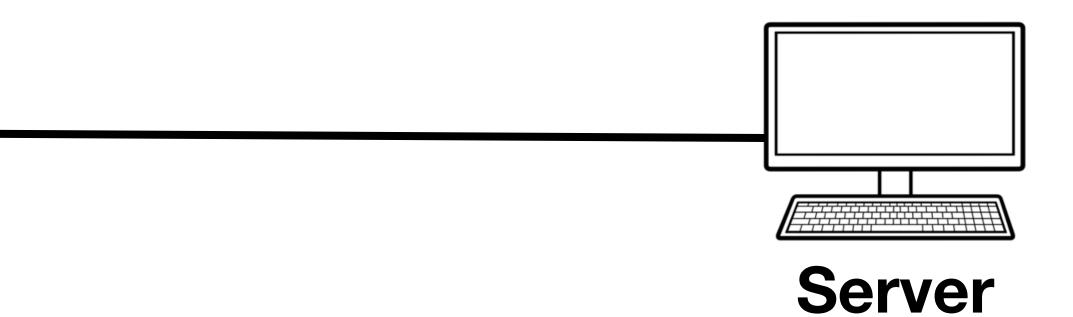
Delay



## **BDP of a Communication Link**

- Delay=Propagation delay ==> Link BDP
  - The number of bits over the wire

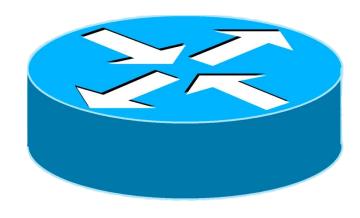






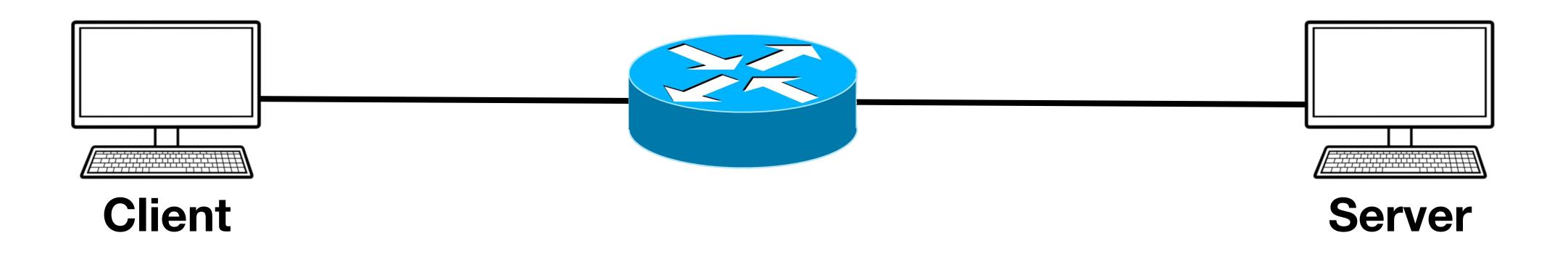
## BDP of a Router

- Delay=Propagation delay ==> Link BDP • The number of bits over the wire
- Delay=Processing + Queueing + Transmission ==> Router BDP The number of bits a router can hold





- Delay=Propagation delay ==> Link BDP
  - The number of bits over the wire
- Delay=Processing + Queueing + Transmission ==> Router BDP The number of bits a router can hold
- Delay=Total delay ==> End-to-End BDP The number of bits that stay in-flight between two hosts



- BDP of Client -> Server



- Today
  - Computer networks: performance analysis

- Next lecture
  - Physical Layer: Encoding

## Summary

